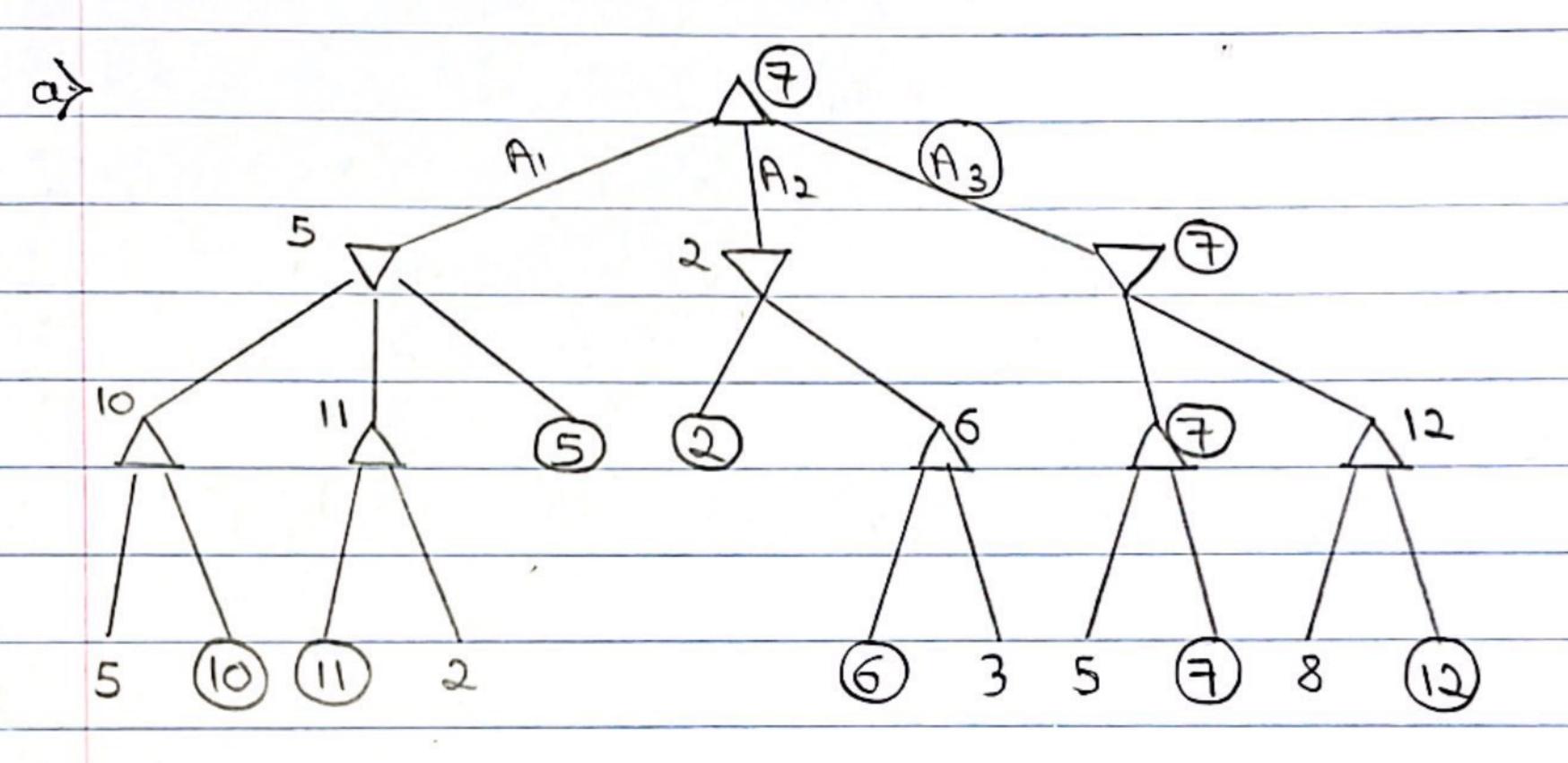
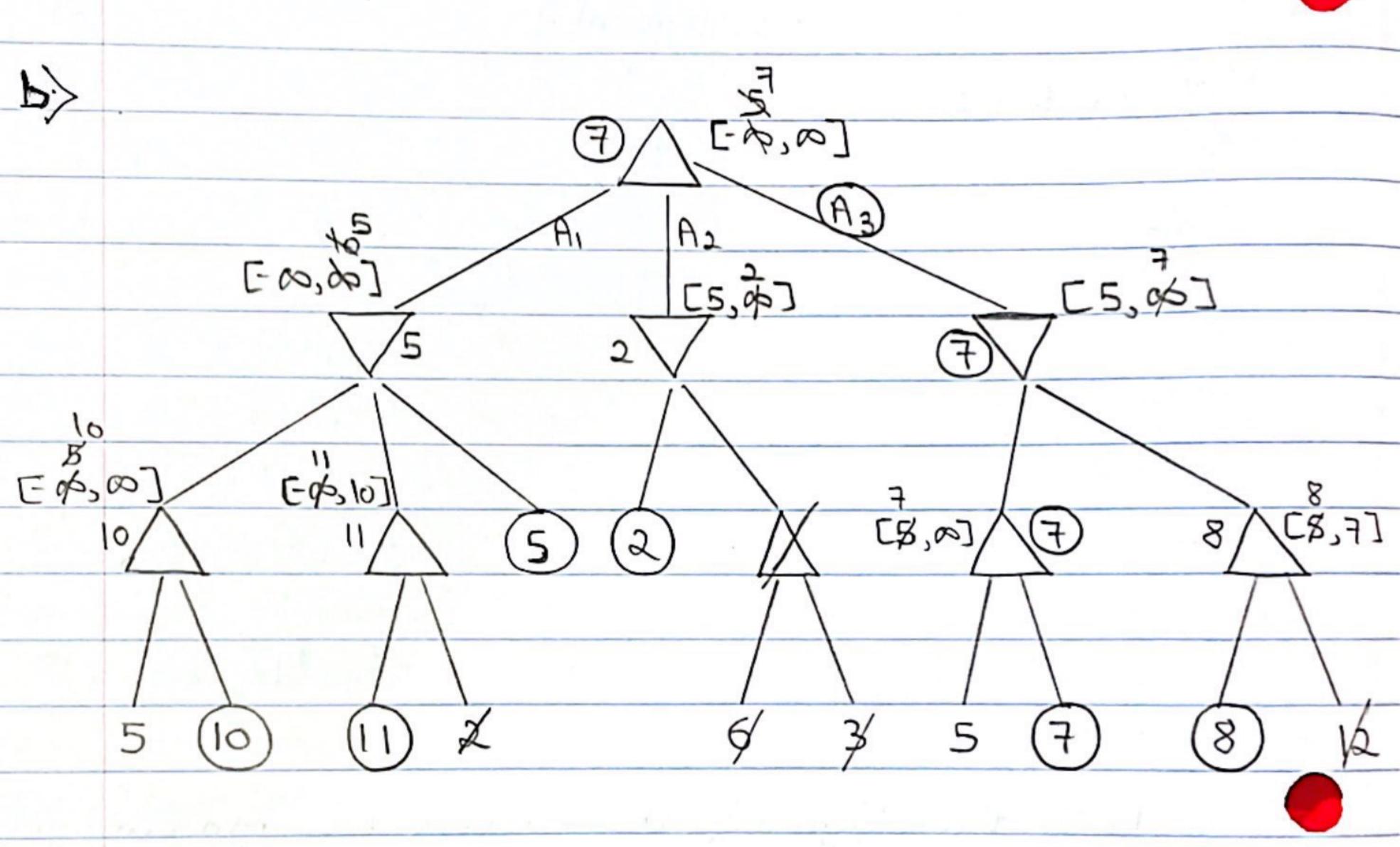
-Assignment 2-

* Jask 1



the min values to 7. 30, we can say that Minimax algorithm will first action A3 to execute.



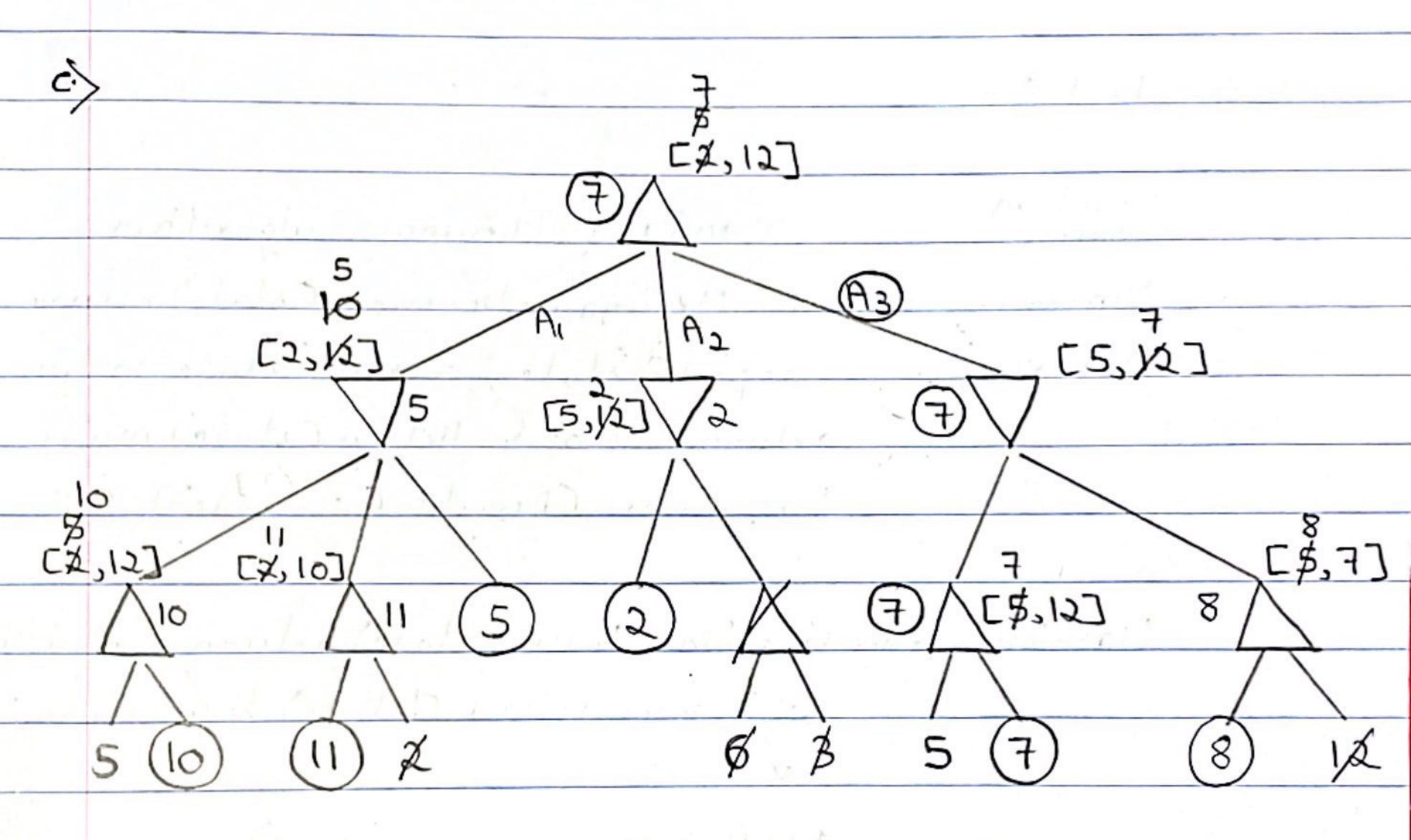
Fran the figure, Mirimax algorithm will pick action As to execute.
Tes, the answer is some as part a.

c> For MIN node:

if
$$v <= 2$$
 or $v == 2$ then prune

if $v < B$ then $B = V$

For MAX node: if V > = B or V = = 12 then pune if V > Q then Q = V



If we know that maximum possible value can 12 and minimum possible value is 2 then we can assign the \approx 3 B values accordingly. \approx 2 = 2 \Rightarrow Highest possible score till now. \approx B = 12 \Rightarrow Least possible score till now.

Also, the α - β priving minimax algorithm can be improved as follows to prime the node if any successor of max node has value = 12 or any successor of min node has value = 2.

Task 2-	
Function	Standard Minimax algorithm
Minimax-	
Decision	enpert : tente : tente
	return the a in Action Costate) maximizing
	Min-value (Result (a. state)
Max-Value	Sunction Max-Value (state) returns a utility verb
	function Max-Value (state) returns a utility veil if Terminal-Test (state) then return Utility CState
	1 C 5 to 5
	for as in ourcerson (state)
	do V < Max (V, Min-value (S))
	return V.
Min-Value	function Min-Value Cotate) returns a utility value
15 1 7 1	tilitel venter redt (state) test-haring je
	$v \leftarrow \infty$ (State
	for as 5 in Successor (state)
,	do V & Min (V, Max-value (5))
	return.

	Function	Modified Minimax Algorithm
	Minimax-	fronction Minimax-Decision (State) returns an action
	Decision	enog ni etata trerus, etata: alugrei
	1. 64	prigraixam (elata) eraited rie a ente convertier
1	1 1 1 2 2 3 2	Min-Value CResult Castates
1	Max-Value	Junetian Max-Value Catata acturrs a utility value
1	2	estisted results root (state) test-locaines fi Cestistes
,	A sond	v ~ - ∞.
		for as 5 in Euccesson (state)
	2.6 11.1.	do V = Max (V, Min-Value (S))
	211 21	return V.
	Min-Value	function Min-Value (state) returns a utility value
	, ,	if Terminal - Test Catale) Then return Utility
		V < Max-value (Deepareen More (State) Catales
		return v.
	DeepGreen	function Deep Green Move (state)
	Move	returns state from min players more
	y ,	retpo etate reuter at printenas ab #
	*	# min players mare without using
		# minimax algarithm.
		return state.

Here, the only required change is in the Min-Value function.

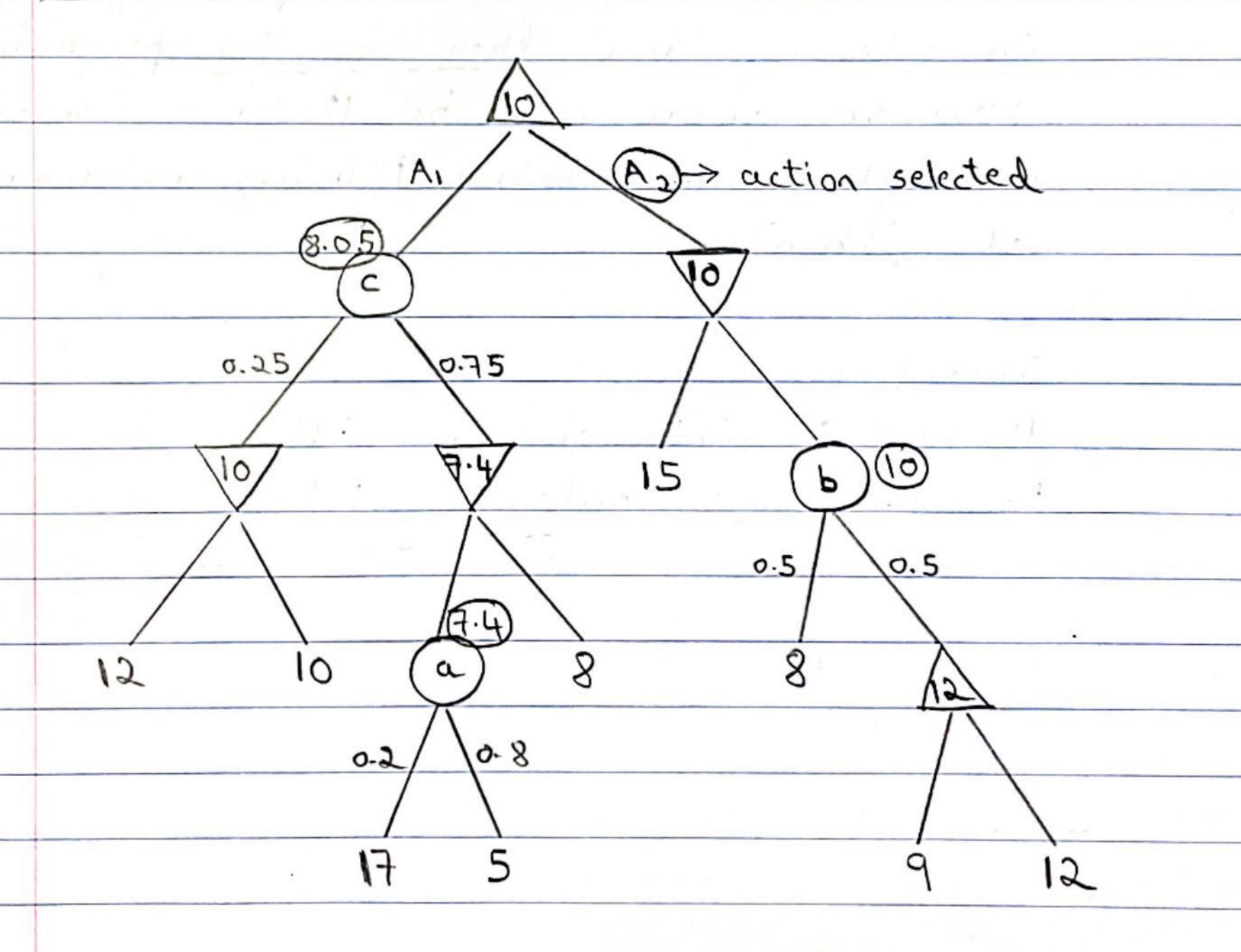
As given in the question, the function
Deeplereen Move returns the state after the
Min player plays. So now we can use this
function in the Min-Value function to
predict the Min player's moves and prune
the rest of the unreceasing branches in
the tree instead of iterating through every
possible move that the Min player can
play. This can reduce the number of
nades in the tree to almost half.

* Advantage of modified Minimax Algorithm using DeephreenMove function.

As mentioned above this will result in fewer branches ie pruning the branches corresponding to the moves that the Min player will never play and reducing the tree size to approx. half . For an optimal player this algorithm will return the same action as the standard Minimax algorithm using less memory. And for non-optimal player

This algorithm will gives the best action resulting in a state with maximum utility value.

* Task 3:



For node a value = (17) (0.2) + (5) (0.8) = 7.4

For node b, value = (8) (0.5) + (12) (0.5) = 10

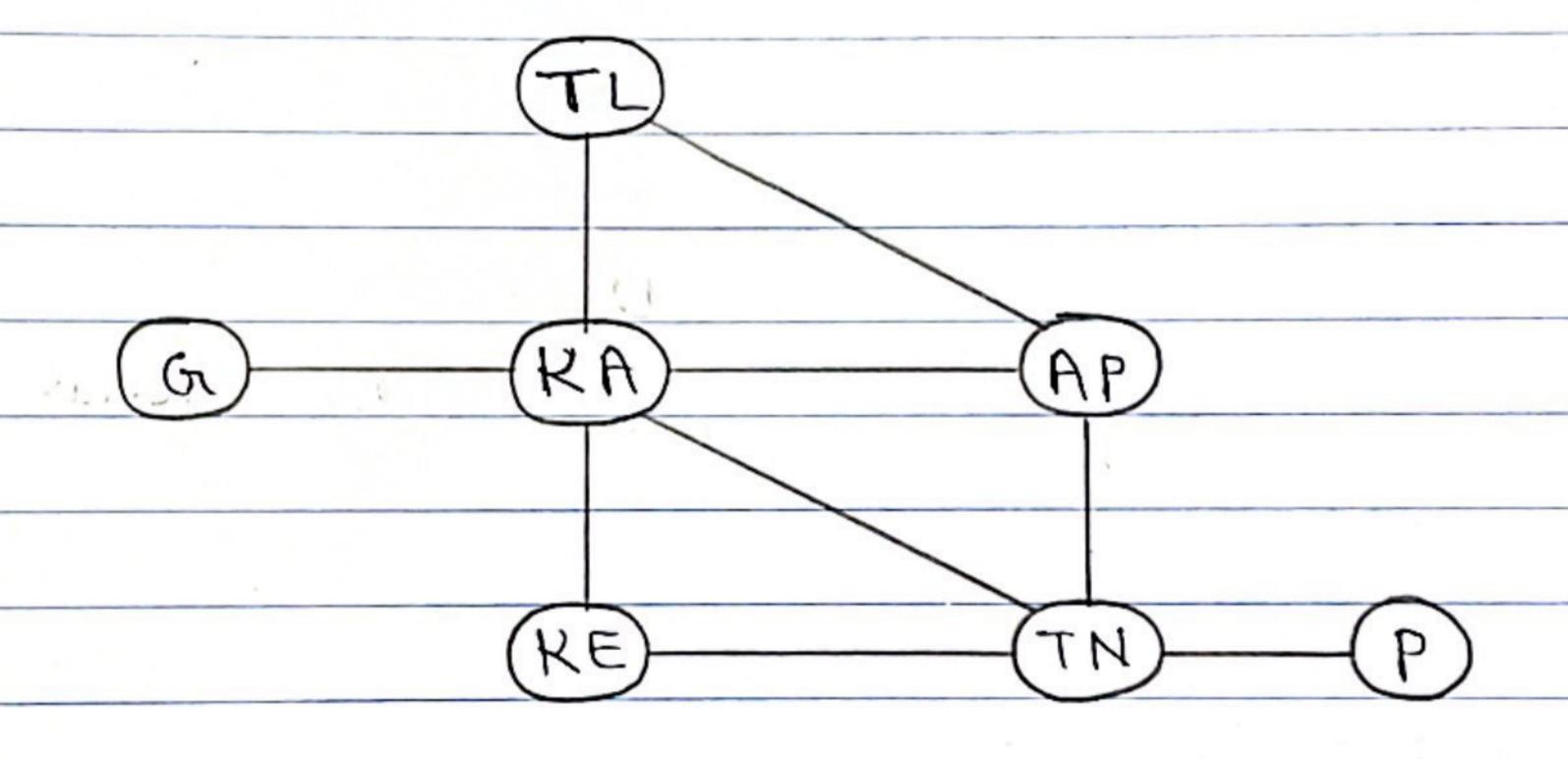
Ear nade C, value = (10) (0.25) + (7.4) (0.75) = 8.05

The action selected by the minimax algorithm is Az as it maximizes the minimum values. After selecting action Az, an optimal opponent will select the more corresponding to node b. On selecting that there is 0.5 probability that the score will be 12 or 8 (since the nax player will select 8 using minimax · Condinable Therefore, Highest passible outcome: 12 Courest passible outcome: 8

* Jask 4-

7

Carstraint Yraph -

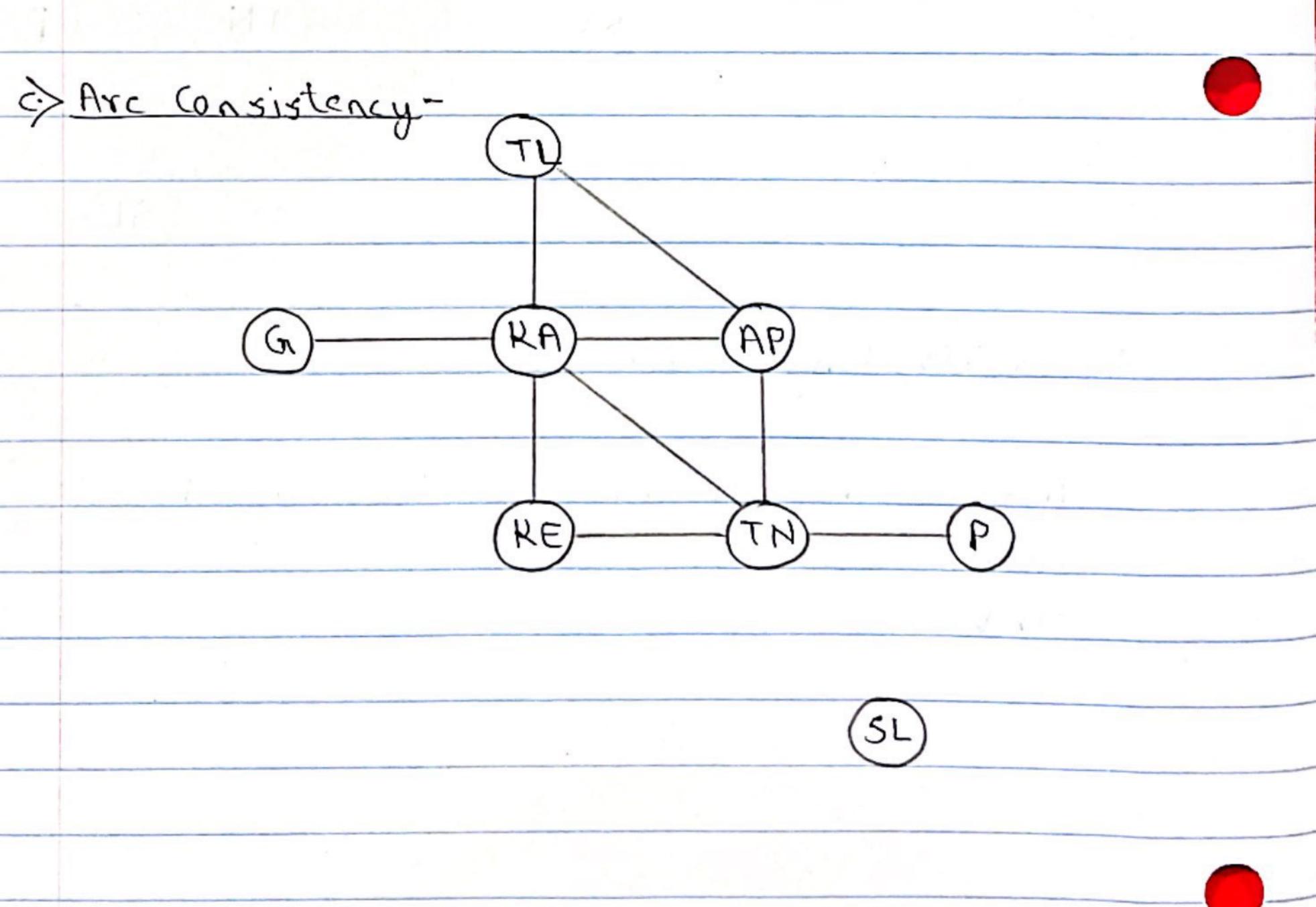


(5L)

B) Backtracking Search -

Node	RA	TN	AP	KE	TL	G	P	5L	_
Degrée	5	₩3	3×1	ZX0	210	XO	YO	0	
MRV	3	22	221	321	371	35	32	3	

Level	MRV	Degree Values	Node Variable Selecte
1	3	5	'KA
2	2	3	TN
3	1		AP
4	1	0	KE
5	1	0	TL
6	2	0	G
F	2	A o	P
8	3	0	5 L



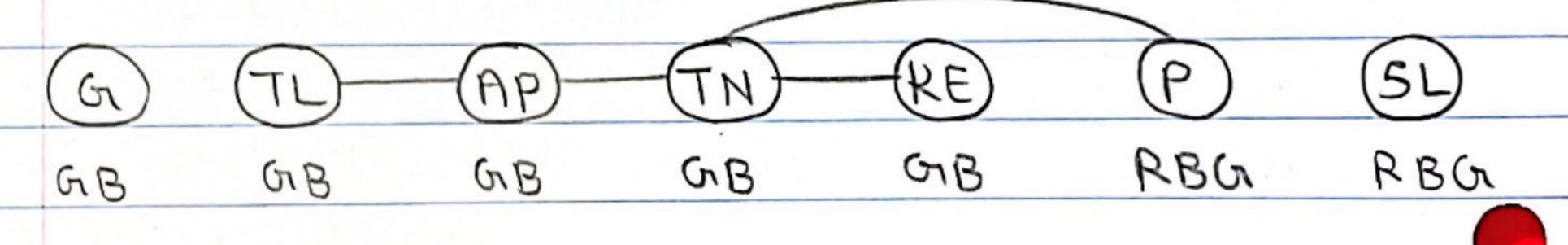
7,		, Red'de first var		Α'.
		AP RE TL'		<u>5</u> L
1	a deri			
	B BarB	RUB RUB RUB	RGB RGB	RGB
	1			
1	50->KA	2-KA>G	5 KA > KE	_
	TL->KA	3 -KA->-TL	THYE	•
	AP>KA	AP->TL	6 AP->TN-	
	RE>RA	4 TL >AP	-RA->TH-	
	TH->KA-	KA->AP-	HE-TH-	
		TN-AP-	-P→TH-	
1				
1	\			

de Yes, we can use the structure of the problem to make solving it more efficient.

There are 2 possible ways to do 50:

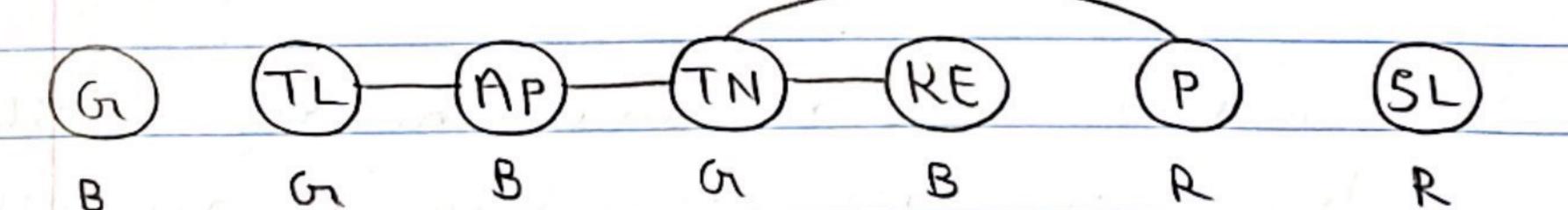
I First, we can see that the region '3L' is having no ares connecting to any other regions. So, we can separate it from the main region and consider to any other subfroblem and solve both the problems separately and concatenate the final result.

If we can see that the region graph is not cyclic and cannot be solved considering it as a tree (which would be efficient). But we can see that it is nearly tree and we can determine the cut-set here consist of node 'KA' and assigning it to 'Red'. And then solve these subproblem using trees as shown below.



The above sub-problem can be easily solved in $O(nd^2)$ time.

es alid solution -



Assign Blue to Gr
Assign Red to SL

Assign Green to TL

Assign Blue to AP

Possible Solution:

Assign Green to TN

Red > (RA, P, SL)

Assign Blue to RE

Green > (TL, TN)

Assign Red to P

Blue > (G, AP, RE)