- 1. Update Graphics Drivers
- 2. Install OpenCL SDK
- 3. Git clone/download the folder <a href="https://github.com/premsasidharan/GpgpuOpenCL">https://github.com/premsasidharan/GpgpuOpenCL</a>
- 4. Open GpGpuOpenCL/vs2015/GpGpuOpenCL.sln
- 5. Try to build VectorAdd

# HETEROGENEOUS / GPGPU COMPUTING WITH OPENCL

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#### From Wikipedia ---

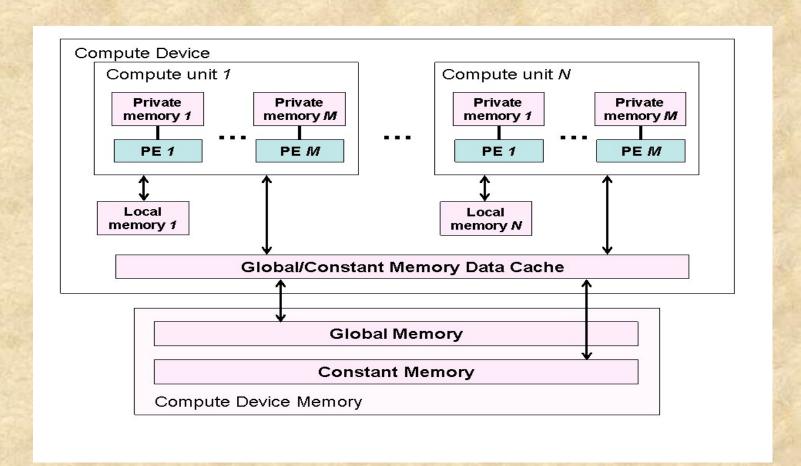
"Open Computing Language (OpenCL) is a framework for writing programs that execute across heterogeneous platforms consisting of CPUs, GPUs, DSPs, FPGAs and other processors or hardware accelerators"

	Platform	Silicon Vendor
12.00	x86/x64 SoC's (Laptop, Desktop)	Intel, AMD
	Mobile SoC's	Qualcomm, Mediatek, Intel Atom, Nvidia Tegra, Samsung Exynos, etc
	Discrete GPUs	Nvidia, AMD
	FPGA	Xilinix, Altera(Intel)
Carlotte Services	DSP	TI

### Agenda

- Abstract Hardware Architecture
- OpenCL Programming Model
- Coding Example 1: Vector Addition
- Coding Example 2: Color Inversion
- Coding Example 3: Image Gradient
- Coding Example 4: Canny Edge Detection
- Coding Example 5: Image Histogram
- Parallel Processing Patterns
- Coding Example 6: Reduce Sum
- Coding Example 7: Prefix Sum
- Coding Example 8: Harris Corner detection (Stream Compaction)
- OpenCL 2.0
- Coding Example 9: Prefix Sum using OpenCL 2.0
- Coding Example 10: Harris Corner detection using OpenCL 2.0
- Coding Example 11: Radix Sort
- Coding Example 12: Matrix Multiplication
- Reference
- Exercise: RGB Histogram

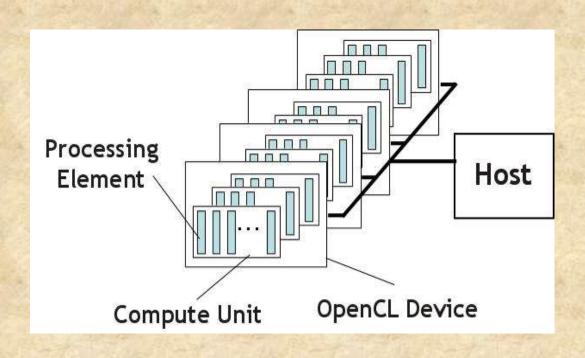
#### Abstract Hardware Architecture



#### **Processor Architecture**

SIMD (Single Instruction Multiple Data)	SIMT (Single Instruction Multiple Threads)	SMT (Simultaneous Multithreading)
<ul> <li>Vector processing instructions</li> <li>Performance issues with branch divergence</li> </ul>	<ul> <li>A hybrid between vector processing and hardware threading</li> <li>Not affected by branch divergence</li> </ul>	Instructions of several threads run in parallel
Example: Old GPU architectures	Example: New GPU architectures	Example: Modern multicore CPUs

## OpenCL Platform Model



#### Data Parallelism

is a form of parallelization across multiple processors in parallel computing environments. It focuses on distributing the data across different nodes, which operate on the data in parallel. It can be applied on regular data structures like arrays and matrices by working on each element in parallel.

#### Task Parallelism

is a form of parallelization of computer code across multiple processors in parallel computing environments. Task parallelism focuses on distributing tasks concurrently performed by processes or threads across different processors.

Courtesy: Wikipedia

## OpenCL Memory Model

- Private Memory (registers)
  - Per work-item
- Local Memory
  - Shared within a work-group
- Global Memory
  - OpenCL Buffers
  - OpenCL Images (1D, 2D and 3D)
- Host Memory

# OpenCL Application

- Host Code
  - Written in a High level language (C, C++).
  - Runs on the Host machine
  - Host code issues data transfer commands
  - Host code issues commands for OpenCL device execution
- OpenCL Device Code
  - Written in OpenCL C
  - Runs on OpenCL Device
  - Parallel code executes on many PEs

Work Item (thread): The basic unit of work on an OpenCL device

**Kernel**: The code for a work Item (like a C Function)

**Program**: Collection of kernels and other functions

**Context**: The environment within which work items execute

CommandQueue: Queue used by the host application to submit work to an OpenCL device. Work is queued in order. Work can be executed in-order or out-of-order.

# OpenCL Execution Model

#### Tradition C Code (vector addition)

```
void add(const float* a, const float* b, float* c, int N)
{
    for (int i = 0; i < N; i++)
    {
        c[i] = a[i]+b[i];
    }
}</pre>
```

#### OpenCL Kernel (vector addition)

```
kernel void add(global const float* a, global const float* b, global float* c)
{
   int i = get_global_id(0);
   c[i] = a[i]+b[i]
}
```

- Kernels are executed across a global domain of work items (threads)
- Work items (threads) can have 1D, 2D or 3D id
- Work item (thread) ID dimension is configurable from the host side depending on the data
- Work items (threads) are grouped into local work-groups (blocks)

- cl::Platform
- cl::Context
- cl::Device
- cl::CommandQueue
- Cl::Event
- cl::Program
- cl::Kernel
- cl::Buffer, cl::BufferGL
- cl::lmage1D, cl::lmage2D, cl::lmage3D, cl::lmageGL
- cl::Sampler

#### How to schedule an OpenCL kernel on a GPU (OpenCL Device)?

kernel void add(global const int\* pA, global const int\* pB,

global int\* pC)

const int id = get\_global\_id(0);

pC[id] = (pA[id] + pB[id]);"

Create a OpenCL context instance

```
cl::Context context(CL_DEVICE_TYPE_GPU);
std::vector<cl::Device> devices = context.getInfo<CL_CONTEXT_DEVICES>();
```

2. Create a CommandQueue instance using the OpenCL context and on the desired OpenCL device.

cl::CommandQueue queue(context, devices[0], CL\_QUEUE\_PROFILING\_ENABLE);

- Create the required OpenCL input, output buffers using OpenCL context cl::Buffer bufferA(context, CL\_MEM\_READ\_WRITE, count\*sizeof(cl\_int));
- 4. Build the OpenCL program

```
cl::Program program(context, sSource);
program.build();
```

Create the OpenCL kernel instance

cl::Kernel kernel(program, "add");

6. Transfer data from host to OpenCL input buffers(map, or copy) using CommandQueue

```
cl_int* pDataA = (cl_int *)queue.enqueueMapBuffer(bufferC, CL_TRUE, CL_MAP_WRITE, 0, count*sizeof(cl_int)); for (size_t i = 0; i < count; i++) pDataA[i] = rand()/100; queue.enqueueUnmapMemObject(bufferA, pDataA);
```

7. Set Input, Output arguments to kernel instance

kernel.setArg(0, bufferA);

8. Enqueue the OpenCL kernel on the Command Queue

```
cl::Event event;
queue.enqueueNDRangeKernel(kernel, cl::NullRange, cl::NDRange(count), cl::NullRange, NULL, &event);
```

9. Wait for the kernel to finish execution

event.wait();

## Kernel work item ID/Thread ID

- Global ID (Global ID for the work Item)
  - size\_t get\_global\_id(size\_t dim)
- Local ID (Work Item ID within the Block)
  - size\_t get\_local\_id(size\_t dim)
- Work group ID (Work Group ID for the work Item)
  - size\_t get\_group\_id(size\_t dim)
- size\_t get\_global\_size(size\_t dim)
- size\_t get\_local\_size(size\_t dim)
- size\_t get\_num\_groups(size\_t dim)

# Kernel Work Item ID (1D data)

```
        Data (0-255)
        Data (256-511)
        Data (512 – 767)
        Data (N-255, N)

        Block 0
        Block 1
        Block 2
        ...
        Block M
```

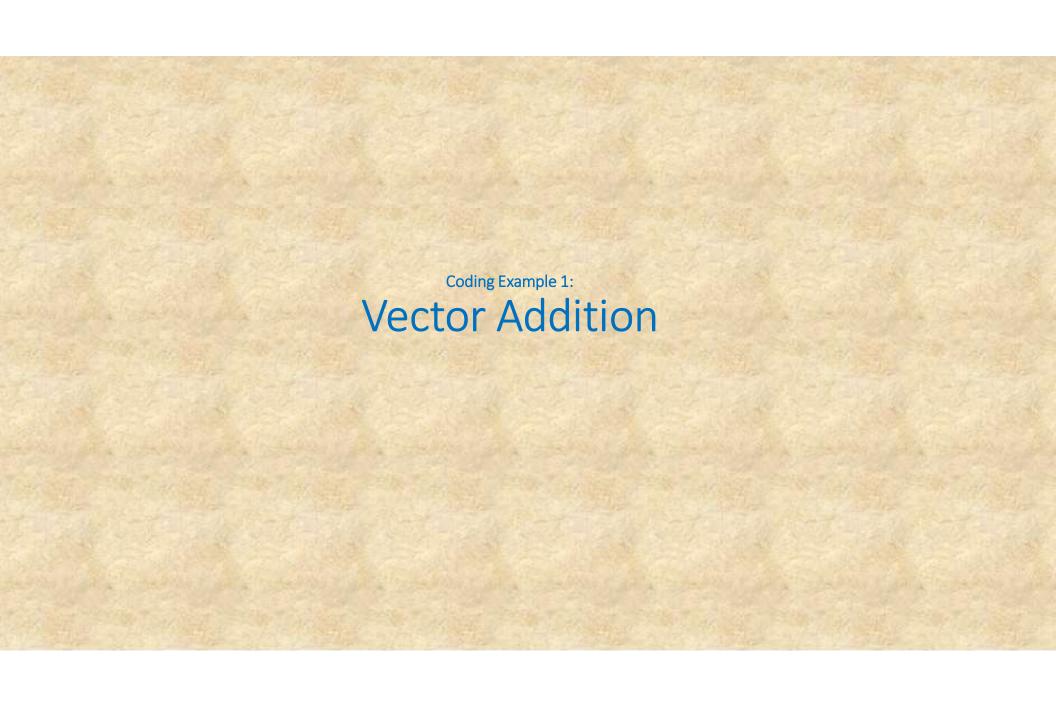
```
size_t get_global_id(0)
size_t get_local_id(0)
size_t get_group_id(0)
size_t get_num_groups(0)
```

## Kernel Work Item ID (2D Data)

Ex:- Image with 32x24 pixels organized into blocks of size 8x8

Block (0, 0)	Block (0, 0) Block			0)			Blo	ock	(2, 0	Block(3, 0)
Block(0, 1) Block(1, 1)  Block(0, 2) Block(1, 2)			k(1,	1, 1)			BI	ock(	) Block(3, 1)	
			Block(2, 2)				) Block(3, 2)			
								题		
		8, 8	9, 8	10, 8	11, 8	12, 8	13, 8	14, 8	15, 8	
		8, 9	9, 9	10, 9	11, 9	12, 9	13, 9	14, 9	15, 9	
		8, 10	9, 10	10, 10	11, 10	12, 10	13, 10	14, 10	15, 10	
		8, 11	9, 11	10, 11	11, 11	12, 11	13, 11	14, 11	15, 11	
		8, 12	9, 12	10, 12	11, 12	12, 12	13, 12	14, 12	15, 12	
		8, 13	9, 13	10, 13	11, 13	12, 13	13, 13	14, 13	15, 13	
E BURNEY		8,	9, 14	10, 14	11, 14	12, 14	13, 14	14, 14	15, 14	
		8, 15	9, 15	10, 15	11, 15	12, 15	13, 15	14, 15	15, 15	

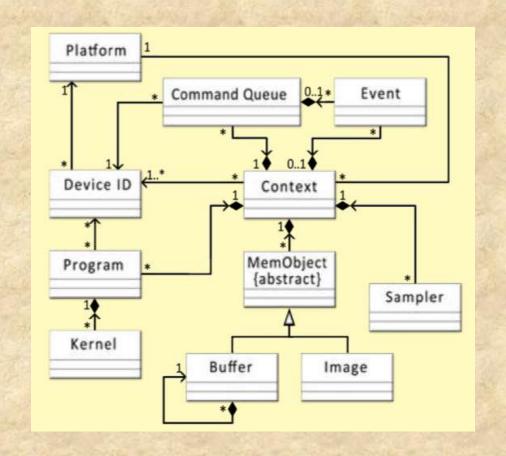
```
size_t get_global_id(0)
size t get global id(1)
size t get local id(0)
size_t get_local_id(1)
size_t get_group_id(0)
size t get group id(1)
size t get global size(0)
size t get global size (1)
size t get local size(0)
size t get local size (1)
size t get_num_groups(0)
size_t get_num_groups(1)
```



# Summary (Vector Addition)

- Learned to launch a basic kernel on the GPU
- Optimization tips
  - Utilize the available memory band width per processing element
  - Use int4 instead of int and process the 4 elements within the thread

# OpenCL Programming Model



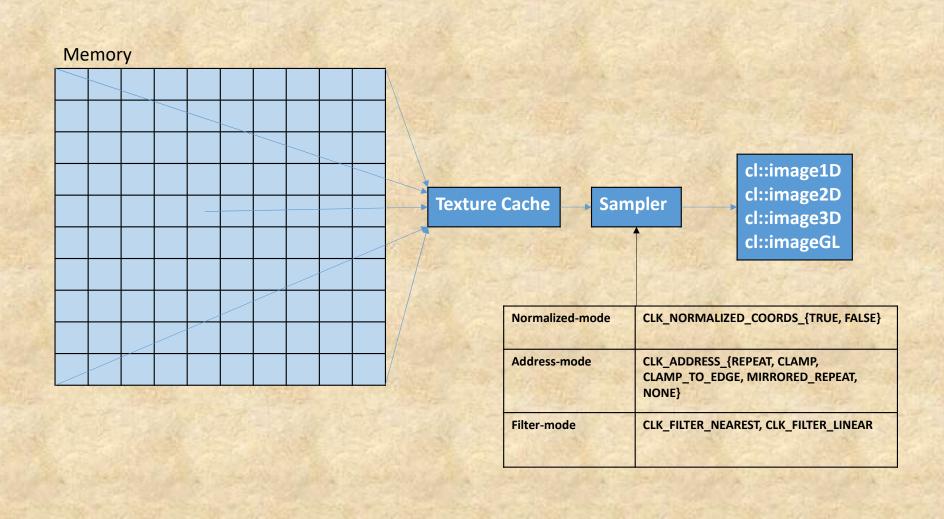
```
Host
                                                      Kernel
                                                      kernel void test(write_only global int* pData, int count)
cl::Buffer, cl::BufferGL
size t count = 1024;
                                                          int id = get global id(0);
cl::Buffer buffA(context, CL_MEM_READ_WRITE,
count*sizeof(cl int));
                                                          pData[id] = value;
kernel.setArg(0, buffA);
kernel.setArg(1, (cl_int)count);
cl::Image2D, cl::ImageGL
                                                      kernel void avg(read only image2d t inp, write only image2d t out)
cl::Image2D inplmg(context, CL MEM READ ONLY,
                                                        const int x = get global id(0);
cl::ImageFormat(CL_RGBA, CL_FLOAT), w, h);
                                                        const int y = get_global_id(1);
cl::Image2D outImg(context, CL MEM READ WRITE,
cl::ImageFormat(CL_R, CL_FLOAT), w, h);
                                                        float4 color = read imagef(inp, (int2)(x, y));
                                                        float idata = (color.x+color.y+color.z)/3.0;
kernel.setArg(0, inplmg);
kernel.setArg(1, outImg);
                                                        write imagef(out, (int2)(x, y), idata)
```

### Inter-op with OpenGL

OpenCL Context Creation with OpenGL on Windows

- OpenGL Buffers, Textures can be converted to OpenCL buffers and Images.
  - · cl::BufferGL, cl::ImageGL
  - cl::BufferGL buff gl(ContextCL, CL MEM READ WRITE, glBufferID)
  - cl::ImageGL img\_gl(ContextCL, CL\_MEM\_READ\_ONLY, GL\_TEXTURE\_2D, 0, glTextureID)
  - std::vector<cl::Memory> gl objs = { buff gl, img gl };
  - queueCL.enqueueAcquireGLObjects(&gl\_objs)
    - ....
  - queueCL.enqueueReleaseGLObjects(&gl objs);
- For some GPUs (Intel HD graphics) interop works only if OpenGL texture internal format is float32 or float16 (GL\_R32F, GL\_R16F, GL\_RG32F, GL\_RGBA32F etc)

void glTexImage2D(GLenum target, GLint level, GLint internalformat, GLsizei width, GLsizei height, GLint border, GLenum format, GLenum type, const GLvoid \* data)



#### Coding Example 2:

#### Color Inversion

```
kernel void invert color(read only image2d t inpImg, write only image2d t outImg)
        int2 coord = (int2)(get global id(0), get global id(1));
        float3 idata = read imagef(inpImg, coord).xyz;
        float4 odata = (float4)(1.0f-idata.x, 1.0f-idata.y, 1.0f-idata.z, 1.0f);
        write imagef(outImg, coord, odata);
  cl::ImageGL inplmgGL(mCtxtCL, CL MEM READ ONLY, GL TEXTURE 2D, 0, mBgrImg());
  cl::ImageGL outImgGL(mCtxtCL, CL_MEM_WRITE_ONLY, GL_TEXTURE_2D, 0, mInvImg());
Н
  std::vector<cl::Memory> gl objs = { inplmgGL, outlmgGL };
  cl::Event event;
  mKernel.setArg(0, inplmgGL);
C | mKernel.setArg(1, outImgGL);
O | mQueueCL.enqueueAcquireGLObjects(&gl_objs);
  mQueueCL.enqueueNDRangeKernel(mKernel, cl::NullRange, cl::NDRange(mBgrImg.width(), mBgrImg.height()), cl::NDRange(8, 8), NULL, &event);
  event.wait();
  mQueueCL.enqueueReleaseGLObjects(&gl objs);
```

#### Coding Example 3:

# Image Gradient

	MaskX	
-1.0f	+0.0f	+1.0f
-2.0f	+0.0f	+2.0f
-1.0f	+0.0f	+1.0f

	MaskY	
+1.0f	+2.0f	+1.0f
+0.0f	+0.0f	+0.0f
-1.0f	-2.0f	-1.0f

- 1. Apply MaskX on Every Pixel and Compute Ix
- 2. Apply MaskY on Every Pixel and Compute ly
- 3. I=sqrt((Ix\*Ix)+(Iy\*Iy))

## Summary: Image Gradient

```
kernel void gradient(read_only image2d_t inplmg, write_only image2d_t outlmg, global const float* plx, global const float* ply)
{
    const int x = get_global_id(0);
    const int y = get_global_id(1);
    const sampler_t sampler = CLK_NORMALIZED_COORDS_FALSE|CLK_ADDRESS_CLAMP|CLK_FILTER_LINEAR;

float ix = 0.0f;
    for (int i = -1; i <= +1; i++)
    {
        for (int i = -1; i <= +1; j++)
        {
            int index = (3*(i+1))+(j+1);
            ix += (plx[index]*read_imagef(inplmg, sampler, (int2)(x+j, y+i)).x);
            iy += (ply[index]*read_imagef(inplmg, sampler, (int2)(x+j, y+i)).x);
        }
        float odata = sqrt((ix*ix)+(iy*iy));
        write_imagef(outlmg, (int2)(x, y), odata);
    }
}</pre>
```

- OpenGL OpenCL interop
- Texture cache (cl::Image)
- Samplers
  - Normalized-mode: CLK\_NORMALIZED\_COORDS\_{TRUE, FALSE}
  - Address-mode: CLK\_ADDRESS\_{REPEAT, CLAMP, CLAMP\_TO\_EDGE, MIRRORED\_REPEAT, NONE}
  - Filter-mode: CLK\_FILTER\_NEAREST, CLK\_FILTER\_LINEAR
- Constant Memory (Accessed through cache)

# Coding Example 4: Canny Edge Detection

- 1. Apply Gaussian filter to smooth the image in order to remove the noise
- 2. Find the intensity gradients of the image
- 3. Apply non-maximum suppression to get rid of spurious response to edge detection
- 4. Apply double threshold to determine potential edges
- 5. Finalize the detection of edges by suppressing all the other edges that are weak and not connected to strong edges.

## Summary: Canny Edge Detection

- Event wait lists
- Chaining kernel's using Event wait list
- Measuring Kernel Execution time

```
cl::CommandQueue queue(context, devices[0], FL_GUEUE_NECEPTERG_ENABLE);
...

bool kernelExecTime(const cl::CommandQueue& queue, const cl::Event& event, int& time)
{
    cl_command_queue_properties qProp;
    queue.getInfo<cl_command_queue_properties>(CL_QUEUE_PROPERTIES, &qProp);

if (qProp & CL_QUEUE_PROFILING_ENABLE)
{
    int start_time = event.getProfilingInfo<CL_PROFILING_COMMAND_START>();
    int end_time = event.getProfilingInfo<CL_PROFILING_COMMAND_END>();
    time = (end_time-start_time);
    return true
}
time = 0;
return false;
}
```

#### Barrier

- All work-items in a work-group executing the kernel on a processor must execute this
  function before any are allowed to continue execution beyond the barrier. This
  function must be encountered by all work-items in a work-group executing the
  kernel.
- If barrier is inside a conditional statement, then all work-items must enter the conditional if any work-item enters the conditional statement and executes the barrier.
- If barrier is inside a loop, all work-items must execute the barrier for each iteration of the loop before any are allowed to continue execution beyond the barrier
- CLK\_LOCAL\_MEM\_FENCE, CLK\_GLOBAL\_MEM\_FENCE.

#### Atomic Operations

 atomic\_add, atomic\_sub, atomic\_xchg, atomic\_inc, atomic\_dec, atomic\_cmpxchg, atomic\_min, atomic\_max, atomic\_and, atomic\_or, atomic\_xor

#### Coding Example 5:

## Image Histogram computation

```
kernel void init(global int* pHistData)
{
    pHistData[get_global_id(0)] = 0;
}
kernel void histogram(read_only image2d_t inplmg, global int* pHistData)
{
    const int x = get_global_id(0);
    const int y = get_global_id(1);
    int idata = (int)ceil(255.0f*read_imagef(inplmg, (int2)(x, y)).x);
    atomic_inc(&pHistData[idata]);
}
```

## How to improve performance?

- Atomic operation on global memory is very slow
- Intermediate Image Histogram using local (shared) memory and atomic ops
- Accumulate the intermediate histograms using atomic addition to global memory
- Histogram without atomics ??? [Exercise]

```
kernel void init(global int* pHistData)
 pHistData[get_global_id(0)] = 0;
kernel void histogram(read only image2d t inplmg, global int* pHistData)
 local int shHistData[256];
 const int x = get_global_id(0);
  const int y = get_global_id(1);
 int index = (get_local_size(0)*get_local_id(1)) + get_local_id(0);
 if (index < 256) shHistData[index] = 0;</pre>
 barrier(CLK LOCAL MEM FENCE);
 int idata = (int)ceil(255.0f*read_imagef(inplmg, (int2)(x, y)).x);
  atomic_inc(&shHistData[idata]);
 barrier(CLK_LOCAL_MEM_FENCE);
 if (index < 256) atomic_add(&pHistData[index], shHistData[index]);</pre>
```

## More performance improvement

- Write the intermediate histogram's into global memory
- Accumulate the intermediate histogram data using a second kernel
  - Accumulation using Simple addition
  - Accumulation using reduce sum

							100			The second second
100		Bin: 0	Bin: 1	Bin: 2					Bin:254	Bin:255
-	Block 0:									
		+	+	+	+	+	+	+	+	+
-	Block 1:									
		+	+	+	+	+	+	+	+	+
-										
		+	+	+	+	+	+	+	+	+
_	Block N:									
100	Final Histogram									

# Parallel Processing Patterns

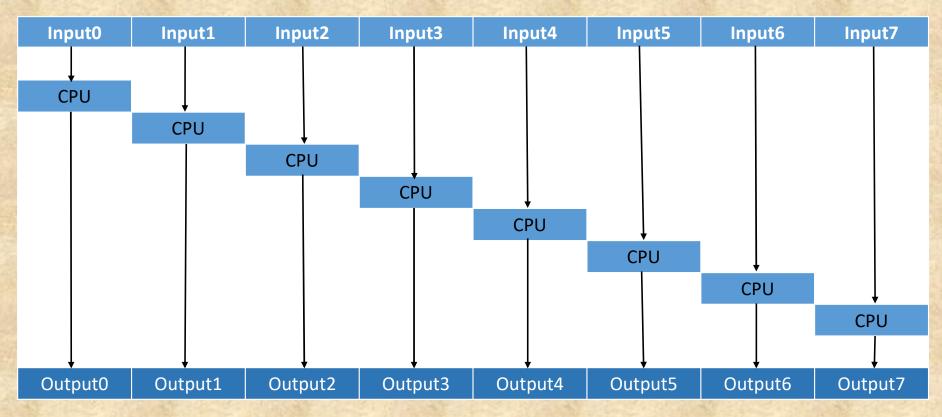
- Map
- Stencil
- Gather
- Scatter
- Reduction (Eg:- sum, min, max etc)
- Scan
  - Inclusive Sum (Prefix Sum)
  - Exclusive Sum
- Stream Compaction
- Histogram

## Map

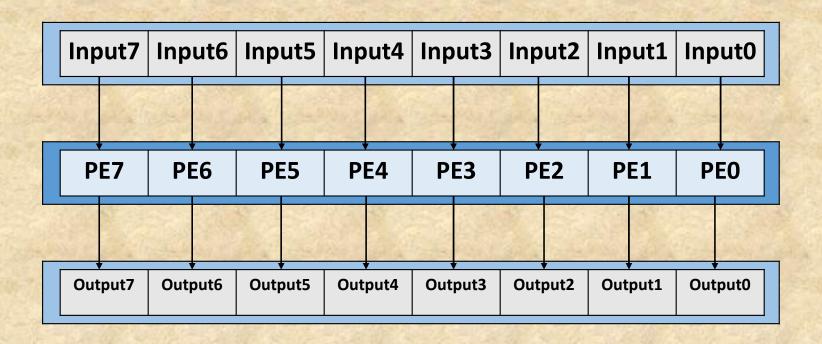
• In the map pattern a function, which we call an elemental function, is applied to different data.

Ex:- vector addition, color inversion, image gradient

# Map (Serial)



# Map (Parallel)



### Gather

 Given a collection of locations (addresses or array indices) and a source array, gather collects all the data from the source array at the given locations and places them into an output collection.

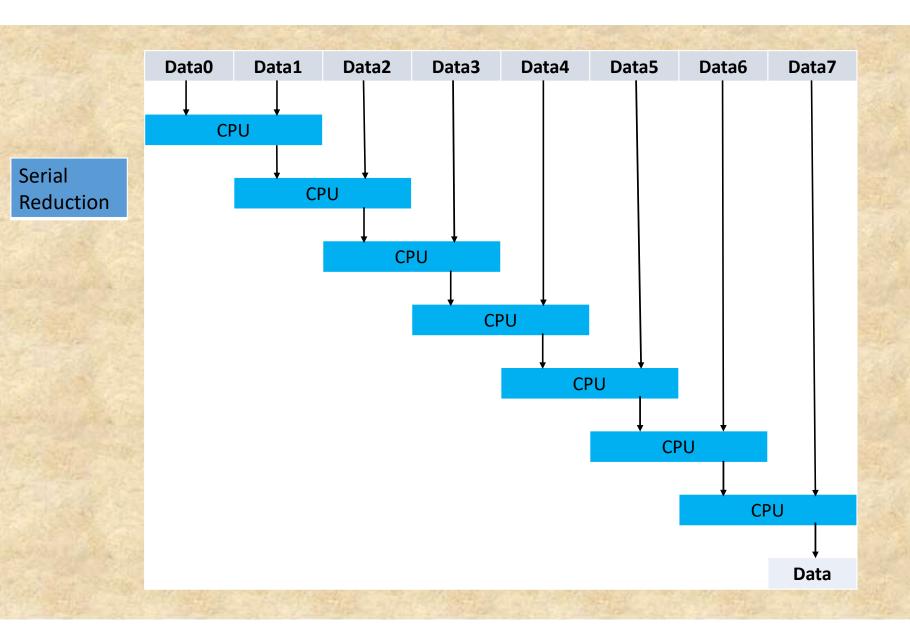
### Scatter

 Scatter is similar to gather, but write locations rather than read locations are provided as input.

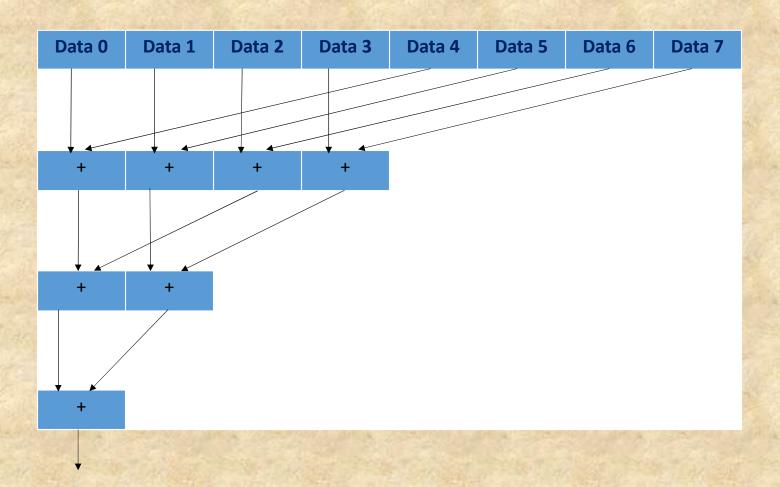
## Reduce

The reduce pattern allows data to be summarized; it combines all the elements in a collection into a single element using some associative combiner operator.

Ex:- min, max, sum, etc



#### **Parallel Reduction**



### Parallel Reduction (Sum)

	Index	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
	Data	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	Initial Data in shared Memory
Ų	Step 1									10	12	14	16	18	20	22	24	If (id < 8) data[id] += data[id+8]
	Step 2													28	32	36	40	If (id < 4) data[id] += data[id+4]
g	Step 3															64	72	If (id < 2) data[id] += data[id+2]
	Step 4																136	If (id < 1) data[id] += data[id+1]

# Block(workgroup) Reduction

Load the Elements into shared memory. For example if the Block contains 256 work items, Bring 512 data items into shared memory and do the reduction as described in the next block.

```
int bSize = get_local_size(0);
int i = get_local_id(0);
While (bSize > 0)
{
    if (i < bSize) data[i] += data[bSize+i];
    barrier(CLK_LOCAL_MEM_FENCE);
    bSize /= 2;
}</pre>
```

#### Coding Example 6:

### Reduce Sum

```
kernel void reduceSum(global const int* pInput, global int* pOutput, const int count)
                                                                                                                inline void breduce sum(local volatile int* shData)
 local int shData[SH_MEM_SIZE];
                                                                                                                  int bSize = get_local_size(0);
                                                                                                                  while (bSize > 0)
 int i = get_local_id(0);
 int id = (SH MEM SIZE*get group id(0))+i;
 shData[i] = (id < count)?pInput[id]:0;</pre>
                                                                                                                    int i = get_local_id(0);
                                                                                                                    if (i < bSize) shData[i] += shData[bSize+i];</pre>
 i += get_local_size(0);
 id += get_local_size(0);
                                                                                                                    barrier(CLK_LOCAL_MEM_FENCE);
                                                                                                                    bSize /= 2;
 shData[i] = (id < count)?pInput[id]:0;</pre>
 barrier(CLK LOCAL MEM FENCE)
 breduce sum(shData);
 if (get_local_id(0) == 0) pOutput[get_group_id(0)] = shData[0];
```

```
size t count = 1000000;
size_t gSize = count/256;
if ((count % 256)) ++gSize;
cl::Buffer buffInp(context, CL MEM READ ONLY, count*sizeof(cl int));
cl::Buffer buffOut(context, CL_MEM_READ_WRITE, gSize*sizeof(cl_int));
reduce.setArg(0, buffInp);
reduce.setArg(1, buffOut);
reduce.setArg(2, (int)count);
queue.enqueueNDRangeKernel(reduce, cl::NullRange, cl::NDRange(256*gSize), cl::NDRange(128), NULL, &events[evtCount]);
while (gSize > 1)
    reduce.setArg(0, buffOut);
    reduce.setArg(1, buffOut);
    reduce.setArg(2, (int)gSize);
    wEvents[wEvtCount].push_back(events[evtCount++]);
    size_t ngSize = (gSize/256)+(((gSize%256)!=0)?1:0);
    queue.enqueueNDRangeKernel(reduce, cl::NDRange, cl::NDRange(ngSize*256), cl::NDRange(128), &wEvents[wEvtCount++], &events[evtCount]);
    gSize = ngSize;
events[evtCount].wait();
```

## Optimized Reduce Sum

```
inline void breduce sum(local volatile int* sh data)
  if (get local id(0) < 256) sh data[get local id(0)] += sh data[256+get local id(0)];
  barrier(CLK LOCAL MEM FENCE);
  if (get local id(0) < 128) sh data[get local id(0)] += sh data[128+get local id(0)];
  barrier(CLK LOCAL MEM FENCE);
  if (get local id(0) < 64) sh data[get local id(0)] += sh data[64+get local id(0)];
  barrier(CLK LOCAL MEM FENCE);
  if (get local id(0) < 32) sh data[get local id(0)] += sh data[32 + get local id(0)];
                                                                                                 No need to place
  if (get_local_id(0) < 16) sh_data[get_local_id(0)] += sh_data[16 + get_local_id(0)];
                                                                                                 barrier instruction
  if (get_local_id(0) < 8) sh_data[get_local_id(0)] += sh_data[8 + get_local_id(0)];
                                                                                                 within WARP
  if (get_local_id(0) < 4) sh_data[get_local_id(0)] += sh_data[4 + get_local_id(0)];
  if (get local id(0) < 2) sh data[get local id(0)] += sh data[2 + get local id(0)];
  if (get local id(0) < 1) sh data[get local id(0)] += sh data[1 + get local id(0)];
  barrier(CLK LOCAL MEM FENCE);
```

#### Scan

• The scan collective operation produces all partial reductions of an input sequence, resulting in a new output sequence. There are two variants: inclusive scan and exclusive scan. For inclusive scan, the nth output value is a reduction over the first n input values. For exclusive scan, the nth output value is a reduction over the first n-1 input values.

#### **Prefix Sum**

• The prefix sum, scan, or cumulative sum of a sequence of numbers  $x_0$ ,  $x_1$ ,  $x_2$ , ... is a second sequence of numbers  $y_0$ ,  $y_1$ ,  $y_2$ , ..., the sums of prefixes of the input sequence:  $y_0 = x_0 y_1 = x_0 + x_1 y_2 = x_0 + x_1 + x_2$  ...

	Input	1	2	3	4	5	6	7	8
1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1	Prefix Sum(Inclusive Sum)	1	3	6	10	15	21	28	36

Input	1	2	3	4	5	6	7	8
Exclusive Sum	0	1	3	6	10	15	21	28

#### Prefix Sum

	ID	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	PERSONAL PROPERTY.
8	Data	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	
	Step 1	2	2	2	2	2	2	2	2	2	2	2	2	2	2	2		if (id >= 1) data[id] += data[id-1]
S	Step 2	4	4	4	4	4	4	4	4	4	4	4	4	4	3	2		if (id >= 2) data[id] += data[id-2]
	Step 3	8	8	8	8	8	8	8	8	8	7	6	5					if (id >= 4) data[id] += data[id-4]
	Step 4	16	15	14	13	12	11	10	9									if (id >= 8) data[id] += data[id-8]

```
inline void bscan(local volatile int* shData)
{
   int level = 1;
   const int id = get_local_id(0);
   while (level < SH_MEM_SIZE)
   {
      if (id >= level) shData[id] += shData[id-level];
      barrier(CLK_LOCAL_MEM_FENCE);
      level *= 2;
   }
}
```

## Efficient Prefix Sum on a GPU workgroup

SI	IMD7	SIMD6	SIMD5	SIMD4	SIMD3	SIMD2	SIMD1	SIMD0
PE (2	255-224)	PE (223-192)	PE (191-160)	PE (159-128)	PE (127-96)	PE (95-64)	PE (63-32)	PE (31-0)

For most of the GPUs the workgroup contains eight 32-wide SIMD processors. (Nvidia uses the term WARP to denote these SIMD processors, AMD calls them WAVEFRONT).

Follow the below steps for doing a block scan (Data is available in shared memory)

- 1. Perform Eight WARP/WAVEFRONT/32 wide prefix sum on the shared Memory (WARP wide sum). Eight 32-wide prefix sums are available in memory.
- 2. Grab the 31st, 63rd, 95th, 127th, 159th, 191th, 223rd and 255th elements and place them in shared memory variable

255th(7th) 2	223rd (6th)	191th(5th)	159th(4th)	127th(3rd)	95th(2nd)	63rd(1st)	31 <sup>st</sup> (0)
--------------	-------------	------------	------------	------------	-----------	-----------	----------------------

- 3. Do a warp prefix sum on the above shared memory variable
- 4. Add the above elements WARP wide on the WARP wide sum (add 0<sup>th</sup> element to elements 63-32, add 1<sup>st</sup> elements 95-64, add 2<sup>nd</sup> element to elements 127-96, add 3<sup>rd</sup> element to elements 159-128, add 4<sup>th</sup> element to elements 191-160, add 5<sup>th</sup> element to elements 223-192 and 6<sup>th</sup> element to elements 255-224)

D0	D1					D30	D31		SIMD/T 0		R0_0	R0_1				R0_30	R0_31	
D32	D33	••			••	D62	D63	M	SIMD/T1		R1_0	R1_1				R1_30	R1_31	
D64	D65					D94	D95		SIMD/T 2		R2_0	R2_1				R2_30	R2_31	
D96	D97					D126	D127	1	SIMD/T3	_	R3_0	R2_1				R3_30	R3_31	
D128	D129					D158	D159	7	SIMD/T 4	7	R4_0	R2_1				R4_30	R4_31	100
D160	D161					D190	D191		SIMD/T 5		R5_0	R2_1				R5_30	R5_31	
D192	D193					D232	D233		SIMD/T 6		R6_0	R2_1				R6_30	R6_31	
D234	D235	••				D254	D255		SIMD/T7		R7_0	R2_1				R7_30	R7_31	

Warp Level Prefix Sum

P0	P1			••	P30	P31	
P32	P33				P62	P63	
P64	P65				P94	P95	
P96	P97				P126	P127	1
P128	P129				P158	P159	1
P160	P161				P190	P191	
P192	P193				P232	P233	
P234	P235				P254	P255	

R0_0	R0_1					R0_30	R0_31			
R1_0	R1_1				••	R1_30	R1_31	SIMD/T 0	S0	Û
R2_0	R2_1					R2_30	R2_31	SIMD/T 1	<b>S1</b>	
R3_0	R2_1	-JL				R3_20	R3_31	SIMD/T 2	S2	
R4_0	R2_1					R4_30	R4_31	SIMD/T 3	S3	1
R5_0	R2_1	7				R5_30	R5_31	SIMD/T 4	<b>S4</b>	
R6_0	T	+				-		SIMD/T 5	S5	ij
*	R2_1	-				R6_30	R6_31			è
R7_0	R2_1	-				R7_30	R7_31	SIMD/T 6	S6	
								SIMD/T 7	<b>S7</b>	

SIMD/

(WARP LEVEL PREFIX SUM)

# Coding Example 7: Prefix Sum

```
inline void wscan(int i, local volatile int* sh data)
 const int wid = i&(WARP_SIZE-1);
 if (wid >= 1) sh_data[i] += sh_data[i-1];
 if (wid >= 2) sh_data[i] += sh_data[i-2];
 if (wid >= 4) sh data[i] += sh data[i-4];
 if (wid >= 8) sh_data[i] += sh_data[i-8];
 if (WARP_SIZE == 32)
   if (wid >= 16) sh data[i] += sh data[i-16];
inline void bscan(int i, local volatile int* shData, local volatile int* shDataT)
 wscan(i, shData);
 barrier(CLK_LOCAL_MEM_FENCE);
 if (i < (SH_MEM_SIZE/WARP_SIZE)) shDataT[i] = shData[(WARP_SIZE*i)+(WARP_SIZE-1)];</pre>
 barrier(CLK LOCAL MEM FENCE);
 if (i < WARP SIZE) wscan(i, shDataT);</pre>
 barrier(CLK_LOCAL_MEM_FENCE);
 if (i >= WARP_SIZE) shData[i] += shDataT[(i/WARP_SIZE)-1];
 barrier(CLK LOCAL MEM FENCE);
kernel void scan(global const int* plnput, global int* pOutput, const int count)
 local int shData[SH_MEM_SIZE+WARP_SIZE];
 int i = get_local_id(0);
 int id = get_global_id(0);
 shData[i] = (id < count)?pInput[id]:0;
 bscan(i, shData, &shData[SH_MEM_SIZE]);
 if (id < count) pOutput[id] = shData[i];</pre>
```

```
kernel void gather(global const int* plnput, global int* pOutput, const int start, const int count)
{
    local int shData[SH_MEM_SIZE+WARP_SIZE];
    int i = get_local_id(0);
    int id = start+(get_local_size(0)*get_local_id(0));
    shData[i] = (id < count)?plnput[id]:0;
    bscan(i, shData, &shData[SH_MEM_SIZE]);
    pOutput[i] = shData[i];
}

kernel void add(global const int* plnput, global int* pOutput, const int start, const int count)
{
    local int shData;
    if (get_local_id(0) == 0) shData = plnput[get_group_id(0)];
    barrier(CLK_LOCAL_MEM_FENCE);
    const int id = start+get_global_id(0);
    if (id < count) pOutput[id] += shData;
}</pre>
```

## Prefix Sum on a binary (0 or 1) sequence

ID	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19
Input	0	0	0	0	1	0	1	0	0	0	0	1	0	0	0	1	0	0	1	0
Inc. Scan	0	0	0	0	1	1	2	2	2	2	2	3	3	3	3	4	4	4	5	5
Ex. Scan	0	0	0	0	0	0	1	1	1	1	1	2	2	2	2	3	3	3	4	4

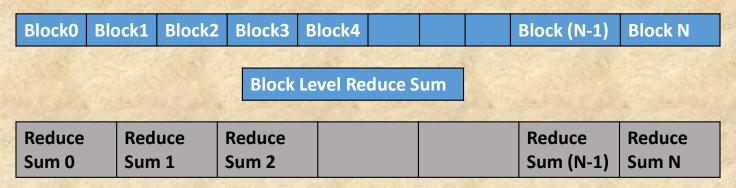
Indexes with 1's	Inc. Scan Value	Ex. Scan Value
4	1	0
6	2	1
11	3	2
15	4	3
18	5	4

# Prefix Sum Applications

- Stream Compaction
- Radix-sort
- Parallel Quick-sort
- Integral Image
- Polynomial Evaluation
- Histogram
- Lexical Analysis
- Solving recurrences

## **Faster Stream Compaction**

1. Do block level reduce sum's on the input and store the outputs on global memory.



- 1. Do an Exclusive Prefix Sum on the block level reduce sum's
- 2. Add the results in Step2 block wise to all elements in the block. (add the value at index 0 to all elements in block 0, add the value at index 1 to all elements in block1, add the value at index N to all elements in blockN)

#### Coding Example 8:

## Harris Corner Detection

- 1. Compute Ix (gradient x) and Iy (gradient y) at every pixel
- 2. Compute the product of derivative at every pixel
  - 1.  $|x^2| = |x^*|x$
  - 2.  $|y|^2 = |y^*|y|$
  - 3.  $|xy = |x^*|y$
- 3. Apply Gaussian smoothing on Ix2, Iy2, Ixy and compute G\_Ix2, G\_Iy2, and G\_Ixy
- 4. Define at each pixel Matrix H

G_lx2	G_lxy
G_lxy	G_ly2

- 5. Calculate R = det(H)-k\*trace(H)\*trace(H)
- 6. Threshold on value of R and due non max suppression
- 7. Do stream compaction

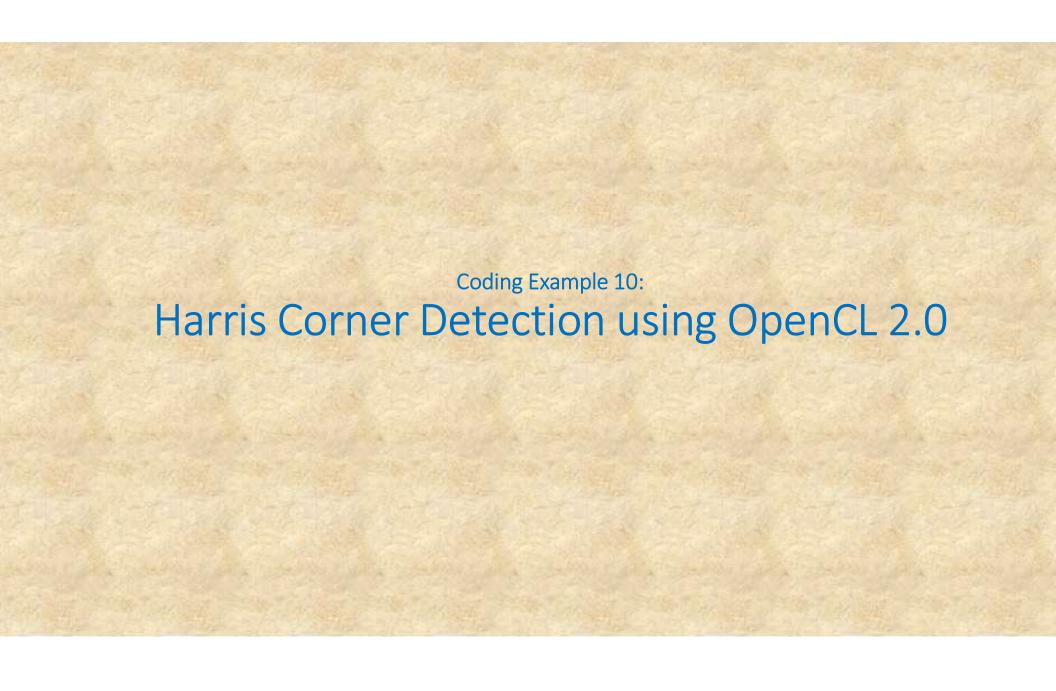
## OpenCL 2.0

- Dynamic parallelism (Nested parallelism)
  - Device side kernels capable of launching other kernels
- Built-in work group functions
  - Scan (Inclusive, Exclusive)
  - Reduce (Min, Max, Sum etc)
- Shared virtual memory
  - Map Free access
  - Fine Grained Coherent access
  - Fine Grained Synchronization
  - Complex Data structures with Pointers (eg:- BTree, Graphs etc)

```
cl::Context context(CL_DEVICE_TYPE_GPU);
std::vector<cl::Device> devices = context.getInfo<CL CONTEXT DEVICES>();
cl::CommandQueue queue(context, devices[0], CL QUEUE PROFILING ENABLE);
cl_queue_properties qprop[] = { CL_QUEUE_PROPERTIES,
(cl command queue properties)(CL QUEUE OUT OF ORDER EXEC MODE ENABLE|CL QUEUE ON DEVICE|
CL QUEUE ON DEVICE DEFAULT | CL QUEUE PROFILING ENABLE), 0 };
int err = 0;
cl_command_queue dev_q = clCreateCommandQueueWithProperties(context(), devices[0](), qprop, &err);
std::ostringstream options;
options << "-cl-std=CL2.0";
cl::Program program(context, sSource);
program.build(options.str().c_str());
```

# Coding Example 9: Prefix Sum using OpenCL 2.0

```
kernel void scan(global int* pdata, global int* tempData, int count)
  clk event tevent1;
  clk event tevent2;
  int ncount = count/4;
  int ncount4 = ncount*4;
  int gcount = ((ncount/BLK_SIZE)+(((ncount%)
  const int BLK SIZE = 256;
  const int BLK_SIZE2 = (BLK_SIZE*BLK_SIZE); BLK_SIZE) == 0)?0:1))*BLK_SIZE;
  enqueue_kernel(get_default_queue(), CLK_ENQUEUE_FLAGS_NO_WAIT,
       ndrange 1D(gcount, BLK SIZE), 0, 0, & event1, \( \) child scan(pdata, ncount); \( \);
  for (int i = (4*BLK_SIZE); i < ncount4; i += (16*BLK_SIZE2))
    enqueue kernel (get default queue(), CLK ENQUEUE FLAGS NO WAIT, ndrange 1D(BLK SIZE, BLK SIZE), 1,
           &event1, &event2, ^{ gather_scan(pdata, tempData, i, ncount4); });
    enqueue kernel(get default queue(), CLK ENQUEUE FLAGS NO WAIT, ndrange 1D((4*BLK SIZE2), BLK SIZE), 1,
          &event2, &event1, ^{ add_data(pdata, tempData, i, ncount); });
  enqueue kernel(get_default_queue(), CLK_ENQUEUE_FLAGS_NO_WAIT, ndrange_1D(1, 1), 1,
          &event1, &event2, ^{ add remainders(&pdata[ncount4], count%4); });
  release event(event2);
```



## Radix Sort

#### **Input Data**

Data	Binary
6	0110
13	1101
2	0010
3	0011
4	0100
5	0101
6	0110
7	0111
8	1000
9	1001

Re-ordering based on bit position 0 (Stream compaction with the mask b'0001)

Data	Binary
6	0110
2	0010
4	0100
6	0110
8	1000
13	1101
3	0011
5	0101
7	0111
9	1001

Re-ordering based on bit position 1 (Stream compaction with the mask b'0010)

Data	Binary
4	0100
8	1000
13	1101
5	0101
9	1001
6	0110
2	0010
6	0110
3	0011
7	0111

Re-ordering based on bit position 2 (Stream compaction with the mask b'0100)

Data	Binary
8	1000
9	1001
2	0010
3	0011
4	0100
13	1101
5	0101
6	0110
6	0110
7	0111

Re-ordering based on bit position 3 (Stream compaction with the mask b'1000)

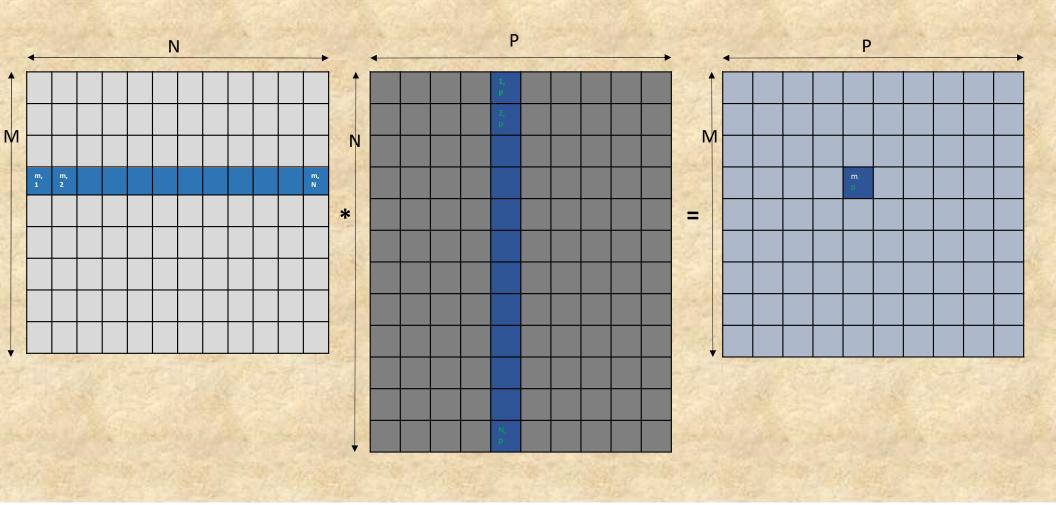
Data	Binary
2	0010
3	0011
4	0100
5	0101
6	0110
6	0110
7	0111
8	1000
9	1001
13	1101

#### Coding Example 11:

## Radix Sort

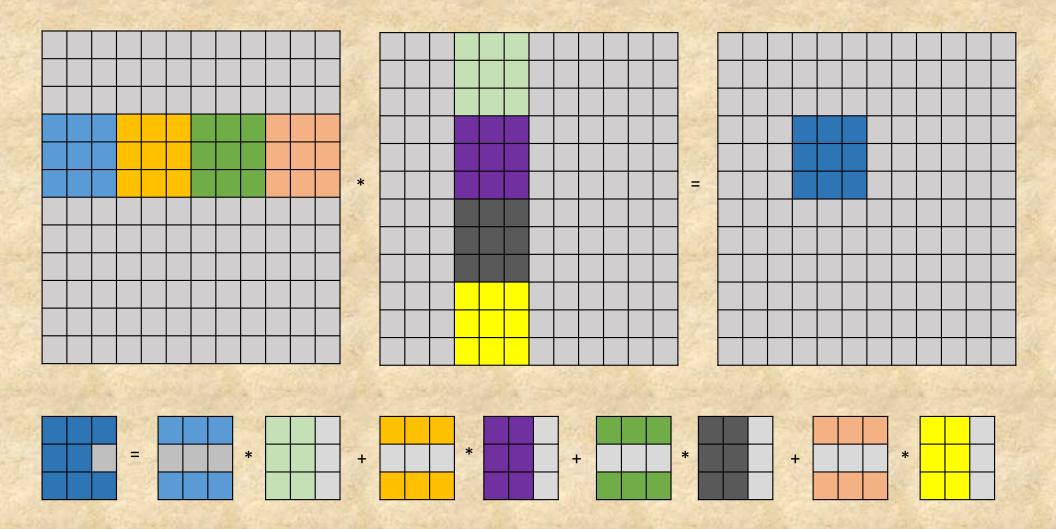
```
kernel void radix sort(global int* p in data, global int* p temp data, int count)
  clk_event_t event[4];
  const int BLK_SIZE = 256;
  int size = 4+(count/256);
  int temp buff size = 1024*((size/1024)+((size%1024)==0?0:1));
  global int* p_odd_blk_count = p_temp_data + temp_buff_size;
  global int* p_even_blk_count = p_odd_blk_count + size;
  global int* p_out_data = p_even_blk_count + size;
  unsigned int mask = 0x0001;
  for (int i = 0; mask != 0; i++)
    global int* p_input_data = ((i%2) == 0)?p_in_data:p_out_data;
    global int* p_output_data = ((i%2) == 0)?p_out_data:p_in_data;
    enqueue_kernel(get_default_queue(), CLK_ENQUEUE_FLAGS_NO_WAIT, ndrange_1D(count, BLK_SIZE), (i==0)?0:1,
                      (i==0)?0:&event[3], &event[0], ^{ block even odd count(p input data, mask, p even blk count, p odd blk count); });
    enqueue kernel(get_default_queue(), CLK_ENQUEUE_FLAGS_NO_WAIT, ndrange_1D(1, 1), 1,
                     &event[0], &event[1], ^{ scan(p_even_blk_count, p_temp_data, size); });
    enqueue kernel(get_default_queue(), CLK_ENQUEUE_FLAGS_NO_WAIT, ndrange_1D(1, 1), 1,
                     &event[0], &event[2], ^{ scan(p odd blk count, p temp data, size); });
    enqueue kernel(get default queue(), CLK ENQUEUE FLAGS NO WAIT, ndrange 1D(count, BLK SIZE), 2,
                     &event[1], &event[3], ^{ compact(p input data, mask, p even blk count, p odd blk count, p output data); });
    mask = mask << 1;
  release_event(event[3]);
```

# Matrix Multiplication



```
kernel void multiply(global float* pMatC, global const float* pMatA, global const float* pMatB, int M, int N, int P)
{
  int m = get_global_id(0);
  int p = get_global_id(1);
  float result = 0.0;
  for (int n = 0; n < N; n++)
  {
     result += (pMatA[(m*N)+n]*pMatB[(n*P)+p]);
  }
  pMatC[(P*m)+p] = result;
}</pre>
```

- Too much global memory access
- Is it possible to reduce the global memory acces?



#### Coding Example 12:

## Matrix Multiplication

```
kernel void multiply(global float* pMatC, global const float* pMatA, global const float* pMatB, int M, int N, int P)
 local float shA[16][16];
 local float shB[16][16];
 int m = get_global_id(0);
 int p = get_global_id(1);
 int pc = (get_group_id(1)<<4)+get_local_id(0);</pre>
 float result = 0.0;
 for (int n = get local id(1); n < N; n += 16)
    shA[get_local_id(0)][get_local_id(1)] = pMatA[(N*m)+n];
    shB[get_local_id(0)][get_local_id(1)] = pMatB[(P*n)+pc];
    barrier(CLK LOCAL MEM FENCE)
    for (int i = 0; i < 16; i++)
      result += (shA[get_local_id(0)][i]*shB[get_local_id(1)][i]);
    barrier(CLK LOCAL MEM FENCE);
  pMatC[(P*m)+p] = result;
```

# Optimization tips

- Algorithm (Parallelization ?)
- Memory
  - Cached access
  - Use of local memory
  - Bandwidth
- Parallel patterns
- Inline assembly ???
  - Many different GPU architectures (Nvidia, AMD, Intel, ARM, Qualcomm, Imagination)
  - Instruction set open for Nvidia and AMD

## Debugging

- Intel OpenCL SDK
  - Breakpoints (Remote)
  - Single stepping (Remote)
  - Profiling
- Nvidia Nsight
  - Profiling support for OpenCL
  - Breakpoints (Only for CUDA apps)
  - Single stepping (Only for CUDA apps)
- Visual Studio supports OpenCL CPU device debugging

# CUDA vs OPENCL

CUDA	OpenCL
- Proprietary Framework (Only Nvidia GPUs)	+ Open (Available on most of the Platforms) OpenCL 1.2 is supported by Intel, AMD, Nvidia GPUs
+ Supports C++ (Templates are very useful)	- Supports only C (C++ support available in OpenCL 2.1)
+ CUDA preprocessor (Allows mixing of device and host code)	- No preprocessor
+ CUDA compiler optimizations much better	- Not as good as CUDA
+ Lots of optimized 3 <sup>rd</sup> party libraries available	- Not as good as CUDA
+ Excellent tools for debugging, profiling etc.	- Not good

## Reference

- Books
  - OpenCL Programming Guide [Aftab Munshi, Benedict R. Gaster etc]
  - Structured Parallel Programming Patterns for Efficient Computation [Michael McCool, Arch D.Robinson, James Reinders]
  - The CUDA Handbook [Nicholas Wilt]
  - GPGPU Gems [Nvidia]
- Useful Libraries
  - Boost compute
  - clBLAS
  - Intel clDNN (Neural Network inference)

## Exercise

- 1. Compute and display RGB histogram for a real time camera feed?
  - a) with atomics
  - b) without atomics

Compare the performance of both options.

