## Organización de Computadoras TP 3: Punto Flotante

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a) M\begin{cases} fraccionaria & E \begin{cases} BCS \\ 4 \ bits \end{cases}
\underbrace{010111}_{M} \underbrace{0110}_{E} \Rightarrow 0,010111 \ x \ 2^{0110} = (2^{-2} + 2^{-4} + 2^{-5} + 2^{-6}) \ x \ 2^{6} = 2^{4} + 2^{2} + 2^{1} + 2^{0} = 16 + 4 + 2 + 1 = 23 \quad \Leftarrow
\underbrace{000001}_{M} \underbrace{0000}_{E} \Rightarrow 0,000001 \ x \ 2^{0000} = (2^{-6}) \ x \ 2^{0} = 2^{-6} \quad \Leftarrow
\underbrace{000011}_{M} \underbrace{1001}_{E} \Rightarrow 0,0000011 \ x \ 2^{1001} = (2^{-5} + 2^{-6}) \ x \ 2^{-1} = 2^{-6} + 2^{-7} \quad \Leftarrow
\underbrace{111111}_{M} \underbrace{1111}_{E} \Rightarrow 0,1111111 \ x \ 2^{1111} = (2^{-1} + 2^{-2} + 2^{-3} + 2^{-4} + 2^{-5} + 2^{-6}) \ x \ 2^{-7} = 2^{-8} + 2^{-9} + 2^{-10} + 2^{-11} + 2^{-12} + 2^{-13} \quad \Leftarrow
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$$\underbrace{000000}_{M} \underbrace{0000}_{E} \Rightarrow 0,0000000 \ x \ 2^{0000} = 0 \ x \ 2^{0} = 0 \iff$$

$$\underbrace{0000000}_{M} \underbrace{1111}_{E} \Rightarrow 0,0000000 \ x \ 2^{1111} = 0 \ x \ 2^{-7} = 0 \iff$$

$$\underbrace{111111}_{M} \underbrace{0000}_{E} \Rightarrow 0,1111111 \ x \ 2^{0000} = (2^{-1} + 2^{-2} + 2^{-3} + 2^{-4} + 2^{-5} + 2^{-6}) \ x \ 2^{0} =$$

$$= 2^{-1} + 2^{-2} + 2^{-3} + 2^{-4} + 2^{-5} + 2^{-6} \iff$$

$$\underbrace{1000000}_{M} \underbrace{00000}_{E} \Rightarrow 0,1000000 \ x \ 2^{0000} = 2^{-1} \ x \ 2^{0} = 0,5 \iff$$

$$\underbrace{0000001}_{M} \underbrace{1111}_{E} \Rightarrow 0,0000001 \ x \ 2^{1111} = 2^{-6} \ x \ 2^{-7} = 2^{-13} \iff$$

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a) M\begin{cases} fraccionaria & E\begin{cases} BCS \\ 4 \ bits \end{cases}
0 \underbrace{10111}_{S} \underbrace{0110}_{M} \underbrace{0110}_{E} \Rightarrow 0 \ 0,10111 \ x \ 2^{0110} = +(2^{-1} + 2^{-3} + 2^{-4} + 2^{-5}) \ x \ 2^{6} = +(2^{5} + 2^{3} + 2^{2} + 2^{1}) = 
= 32 + 8 + 4 + 2 = 46 \quad \Leftarrow 
0 \underbrace{00001}_{S} \underbrace{00001}_{M} \underbrace{0000}_{E} \Rightarrow 0 \ 0,00001 \ x \ 2^{0000} = +(2^{-5}) \ x \ 2^{0} = +2^{-5} \quad \Leftarrow 
0 \underbrace{00011}_{S} \underbrace{1001}_{M} \Rightarrow 0 \ 0,000011 \ x \ 2^{1001} = +(2^{-4} + 2^{-5}) \ x \ 2^{-1} = +(2^{-5} + 2^{-6}) \quad \Leftarrow 
1 \underbrace{11111}_{S} \underbrace{11111}_{E} \underbrace{1111}_{E} \Rightarrow 1 \ 0,111111 \ x \ 2^{1111} = -(2^{-1} + 2^{-2} + 2^{-3} + 2^{-4} + 2^{-5}) \ x \ 2^{-7} = 
= -(2^{-8} + 2^{-9} + 2^{-10} + 2^{-11} + 2^{-12}) \quad \Leftarrow
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$$\underbrace{00000}_{S} \underbrace{00000}_{M} \underbrace{0000}_{E} \Rightarrow 0 \ 0,00000 \ x \ 2^{0000} = +0 \ x \ 2^{0} = 0 \quad \Leftarrow$$

$$\underbrace{000000}_{S} \underbrace{1111}_{M} \Rightarrow 0 \ 0,000000 \ x \ 2^{1111} = +0 \ x \ 2^{-7} = 0 \quad \Leftarrow$$

$$\underbrace{11111}_{M} \underbrace{0000}_{E} \Rightarrow 1 \ 0,111111 \ x \ 2^{0000} = -(2^{-1} + 2^{-2} + 2^{-3} + 2^{-4} + 2^{-5}) \ x \ 2^{0} =$$

$$\underbrace{-(2^{-1} + 2^{-2} + 2^{-3} + 2^{-4} + 2^{-5})}_{S} \iff \underbrace{-(2^{-1} + 2^{-2} + 2^{-3} + 2^{-4} + 2^{-5}) \ x \ 2^{0} =$$

$$\underbrace{100000}_{S} \underbrace{00000}_{M} \underbrace{0000}_{E} \Rightarrow 1 \ 0,000000 \ x \ 2^{0000} = -(0 \ x \ 2^{0}) = -0 \ x \ 1 = -0 \iff$$

$$\underbrace{000001}_{S} \underbrace{1111}_{E} \Rightarrow 0 \ 0,000001 \ x \ 2^{1111} = +(2^{-5} \ x \ 2^{-7}) = +2^{-12} \iff$$

2) 
$$M \begin{cases} fraccionaria \\ BCS 5 bits \end{cases}$$
  $E \begin{cases} BSS \\ 3 bits \end{cases}$   
a)  $Sin \ normalizar$   

$$0 \underbrace{1000}_{S} \underbrace{111}_{E} \Rightarrow 0 \ 0,1000 \ x \ 2^{111} = + (2^{-1}) \ x \ 2^{7} = + (2^{6}) = 64 \iff$$

$$Normalizada$$

$$0 \underbrace{1000}_{S} \underbrace{111}_{E} \Rightarrow 0 \ 0,1000 \ x \ 2^{111} = + (2^{-1}) \ x \ 2^{7} = + (2^{6}) = 64 \iff$$

$$Con \ bit \ implícito$$

$$0 \underbrace{1000}_{S} \underbrace{111}_{E} \Rightarrow 0 \ 0,11000 \ x \ 2^{111} = + (2^{-1} + 2^{-2}) \ x \ 2^{7} =$$

$$= + (2^{6} + 2^{5}) = 64 + 32 = 96 \iff$$

b) Sin normalizar  

$$\frac{1}{S} \underbrace{1000}_{M} \underbrace{011}_{E} \Rightarrow 10,1000 \ x \ 2^{011} = -(2^{-1}) \ x \ 2^{3} = -(2^{2}) = -4 \quad \Leftarrow$$

Normalizada

$$\underbrace{1}_{S} \underbrace{1000}_{M} \underbrace{011}_{E} \Rightarrow 10,1000 \ x \ 2^{011} = -(2^{-1}) \ x \ 2^{3} = -(2^{2}) = -4 \quad \Leftarrow$$

Con bit implícito

$$\frac{1}{\tilde{s}} \underbrace{1000}_{M} \underbrace{011}_{E} \Rightarrow 10,11000 \ x \ 2^{011} = -(2^{-1} + 2^{-2}) \ x \ 2^{3} = \\
= -(2^{2} + 2^{1}) = -(4 + 2) - 6 \quad \Leftarrow$$

c) Sin normalizar

$$\underbrace{0}_{S} \underbrace{0000}_{M} \underbrace{000}_{E} \Rightarrow 0 \ 0,0000 \ x \ 2^{000} = +0 \quad \Leftarrow$$

Normalizada

No se puede. La mantisa no empieza con 0,1...

Con bit implícito

$$\underbrace{0000}_{S} \underbrace{0000}_{M} \underbrace{000}_{E} \Rightarrow 00,10000 \ x \ 2^{000} = +(2^{-1}) \ x \ 2^{0} =$$

$$=+(2^{-1}) x 1 = +0,5 \Leftarrow$$

d) Sin normalizar

$$\frac{1}{\tilde{s}} \underbrace{1111}_{\tilde{M}} \underbrace{111}_{\tilde{E}} \Rightarrow 10,1111 \times 2^{111} = -(2^{-1} + 2^{-2} + 2^{-3} + 2^{-4}) \times 2^{7} = \\
= -(2^{6} + 2^{5} + 2^{4} + 2^{3}) = -(64 + 32 + 16 + 8) = -120 \quad \Leftarrow$$

Normalizada

$$\frac{1}{S} \underbrace{1111}_{M} \underbrace{111}_{E} \Rightarrow 10,1111 \times 2^{111} = -(2^{-1} + 2^{-2} + 2^{-3} + 2^{-4}) \times 2^{7} = \\
= -(2^{6} + 2^{5} + 2^{4} + 2^{3}) = -(64 + 32 + 16 + 8) = -120 \quad \Leftarrow$$

Con bit implícito

$$\frac{1}{s} \underbrace{1111}_{M} \underbrace{111}_{E} \Rightarrow 10,11111 x 2^{111} = -(2^{-1} + 2^{-2} + 2^{-3} + 2^{-4} + 2^{-5}) x 2^{7} = \\
= -(2^{6} + 2^{5} + 2^{4} + 2^{3} + 2^{2}) = -(64 + 32 + 16 + 8 + 4) = -124 \quad \Leftarrow$$

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