Brit Doseii 1002170533. Correctness of Selection Sort: Algorithm: Find the smallest element is the unsorted part of the array and Swap it with the first element in that part of the array. For (=0 to n-18 minimum = ( for y= i+1 Jon & If (orr [i] < orr [min]) { minimum 2 j 3 (i!= mirymum) {. Swap (ith element, minimum (tramele

Z

Corrections: The correctness of the selection sort Algorithm can be proven by wing a loop inveriant and demonstrating that it holds true during the execution of the algorithm. Initialization:-Invariant At the start of each iteration of the outer doop, the suburray before the arrent index "i' is sorted, and all elements in this subarray are Smaller than or squal to any selement in other unsorted nortion years ant fo for i=0 ton-1 &.

minimum : i

Mointenance:

(0

Before the execution of the inner look, all the elements before i are sorted. The inner look identifies the sorted and he minimum element in the array and super the values are

for j=1+1 to n g. Af cour [j] < our [minimum ] { minimum = } Evap (arr[i], arr[minimum]) The algorithm forts the element at rebrue ett to sprobed deitu, 9 rebru hur your reture ent. mumining Agriand Dato n - 1, iterating through the array. Termination: Once the outer fool rature ant sono algorithm terminales . The entre array is sorted, and there are no tent enaitibres to expeal stingfue would prevent the completion.