

Computer Networks

COL 334/672

Multi-access rules

Slides adapted from KR

Sem 1, 2025-26

Quiz on Moodlenew

Password: manchester



Link Layer: Services

- Encoding
- Framing
- Error detection
- Addressing
- Link access

Cyclic Redundancy Check (CRC)

Algorithm

1. Multiply $M(x)$ by x^k ; that is, add k zeros at the end of the message. Call this zero-extended message $T(x)$.
2. Divide $T(x)$ by $G(x)$ and find the remainder.
3. Subtract the remainder from $T(x)$ → get $P(x)$

■ **Claim:** If $G(x)$ has $x+1$ as one of its factors, all single-bit errors can be detected

■ **Ethernet** protocol uses a 32-bit error check

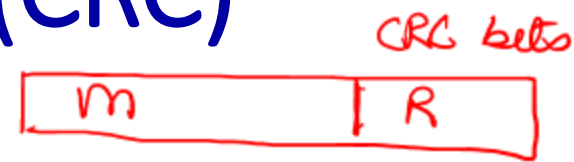
→ **CRC-32** = $x^{32} + x^{26} + x^{23} + x^{22} + x^{16} + x^{12} + x^{11} + x^{10} + x^8 + x^7 + x^5 + x^4 + x^2 + x + 1$

■ Where is CRC implemented?

Hardware (NIC)

End-to-end principle & Error detection

↳ Link layer shouldn't implement ED but done for performance



$$G(x) \mid P'(x)$$

$$(x+1) \times G'(x) \mid P'(x)$$

$$(x+1) \mid P(x) + E(x)$$

$$x+1 \mid E(x)$$

$$x+1 \mid x^k$$

Not possible

Handwritten long division of 1101 by 11. The quotient is 101 and the remainder is 0.

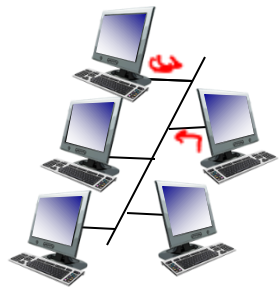
Link Layer: Services

- Encoding
- Framing
- Error detection
- Addressing and link access
multiple

Multiple access links, protocols

two types of “links”:

- point-to-point
 - point-to-point link between Ethernet switch, host
 - PPP for dial-up access
- broadcast (shared wire or medium)
 - old-school Ethernet
 - upstream HFC in cable-based access network
 - 802.11 wireless LAN, 4G/4G. satellite



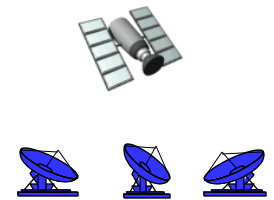
shared wire (e.g.,
cabled Ethernet)



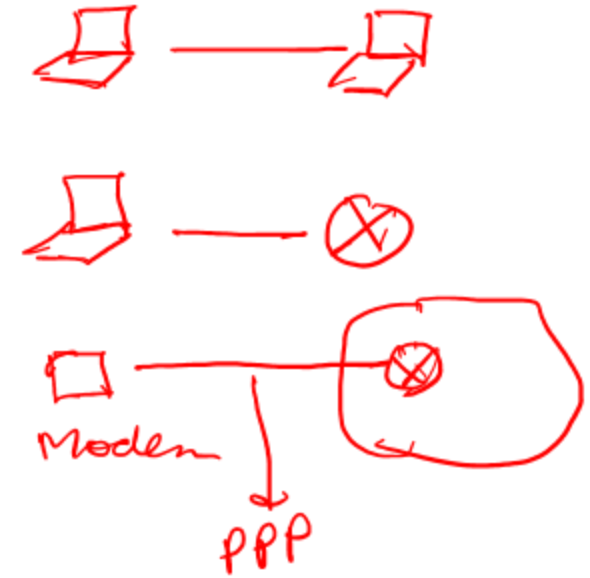
shared radio: 4G/5G



shared radio: WiFi

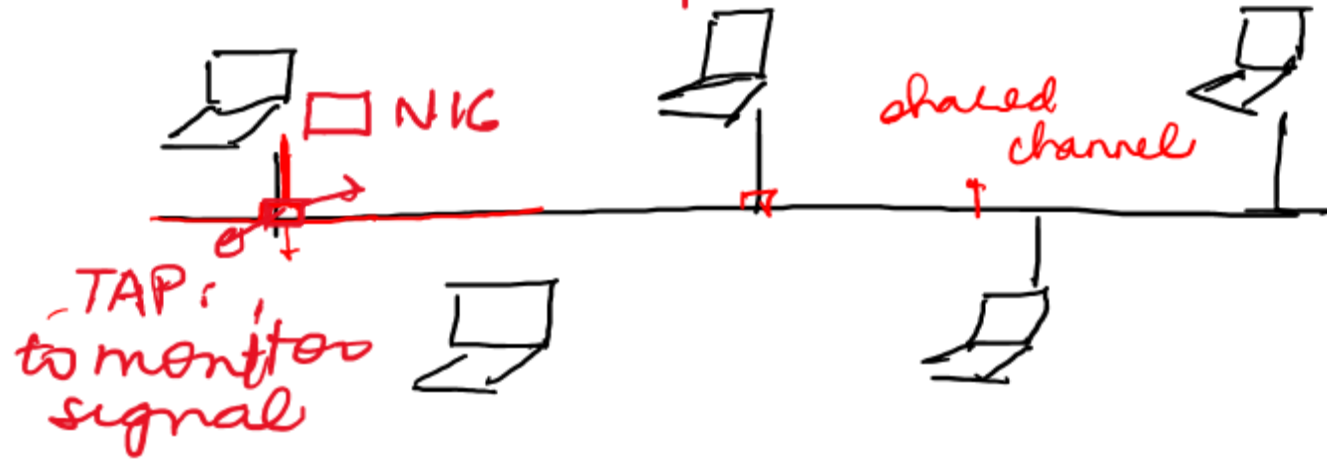


shared radio: satellite



Medium Access Control (MAC) protocol

Transceiver to interpret



Requirement 1: Addressing

- Need a way to identify end nodes
- Ethernet uses a 48-bit MAC address burned in NIC ROM, also sometimes software settable
- E.g., 1A-2F-BB-76-09-AD

Find out your host's MAC using
ipconfig / ifconfig.

MAC protocol

- **algorithm** that determines how nodes share channel, i.e., determine when node can transmit

- communication about channel sharing must use channel itself
- no out-of-band channel for coordination

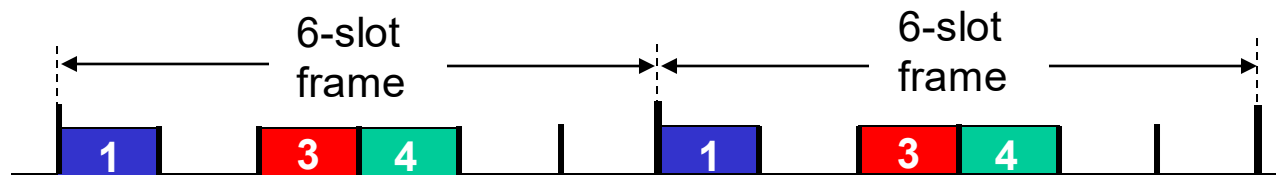
DESIGN GOALS

- ① Efficiency
- ② Fairness
- ③ Resilient
- ④ Simple

Channel partitioning MAC protocols: TDMA

TDMA: time division multiple access

- access to channel in “rounds”
- each station gets fixed length slot (length = packet transmission time) in each round
- unused slots go idle
- example: 6-station LAN, 1,3,4 have packets to send, slots 2,5,6 idle

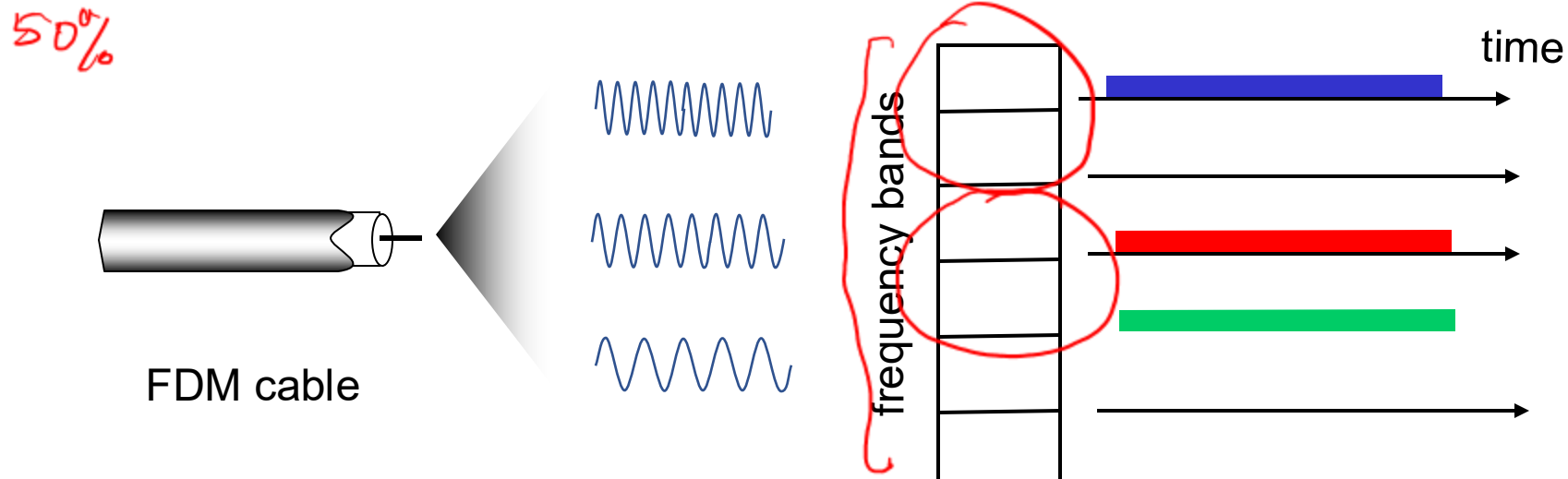


Channel partitioning MAC protocols: FDMA

FDMA: frequency division multiple access

- channel spectrum divided into frequency bands
- each station assigned fixed frequency band
- unused transmission time in frequency bands go idle
- example: 6-station LAN, 1,3,4 have packet to send, frequency bands 2,5,6 idle

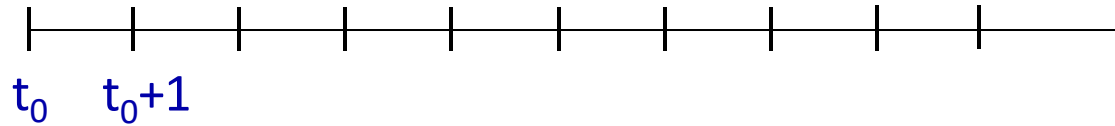
- ① Fair ✓
- ② Simple ✓
- ③ Resilient ✓
- ④ Efficiency X



Random access protocols

- when node has packet to send
 - transmit at full channel data rate R
 - no *a priori* coordination among nodes
- two or more transmitting nodes:
“collision”
- random access protocol specifies: *& how to (try to) avoid collisions*
 - ↳ • how to detect collisions
 - ↳ • how to recover from collisions (e.g., via delayed retransmissions) ¹
- examples of random access MAC protocols:
 - ALOHA, slotted ALOHA
 - CSMA, CSMA/CD, CSMA/CA

Slotted ALOHA



assumptions:

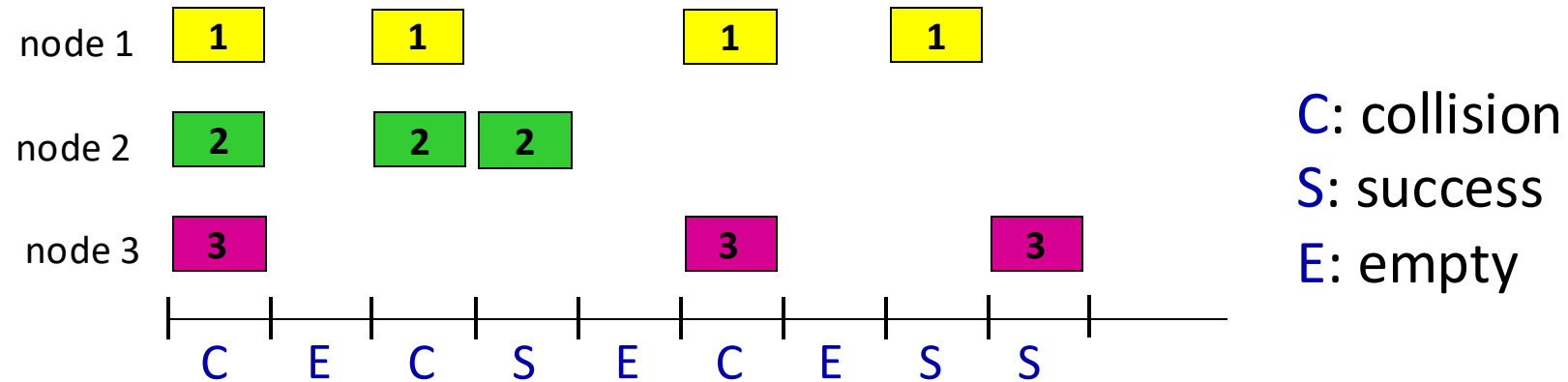
- all frames same size
- time divided into equal size slots (time to transmit 1 frame)
- nodes start to transmit only slot beginning
- nodes are synchronized
- if 2 or more nodes transmit in slot, all nodes detect collision

operation:

- when node obtains fresh frame, transmits in next slot
 - *if no collision*: node can send new frame in next slot
 - *if collision*: node retransmits frame in each subsequent slot with probability p until success

randomization – why?

Slotted ALOHA



Pros:

- single active node can continuously transmit at full rate of channel
- highly decentralized: only slots in nodes need to be in sync
- simple

Cons:

- collisions, wasting slots
- idle slots
- clock synchronization