Computer Networks
COL 334/672

Application Layer: DNS

Sem 1, 2025-26

Recap: Application Layer

- HTTP
- Email
- **DNS** or Domain Name System
- P2P
- Video streaming

DNS: Domain Name System

DHCP : Application-layer protocol

- Humans understand names (google.com), Naming
- Internet hosts, routers understand IP address (12.123.12.12) Addressing
- *Q:* how to map between IP address and name, and vice versa ?

Two questions:

- How to design the database?
- How to retrieve the IP for a given domain name?

Domain Name System (DNS):

- *phone book* of the Internet
- *application-layer protocol:* hosts, DNS servers communicate to *resolve* names (address/name translation)
  - *note:* core Internet function, implemented as application-layer protocol
  - complexity at network's "edge"

End-to-end design principle

How would you design the DNS database?

- **Primary goal**: a database of domain to IP mappings
- What are other design constraints/challenges?

a single machine storing all mappings?

↳ Scalable : # of queries to the table & size of the table

↳ Fast (Performance)

↳ Secure

↳ Resilience

AWS outage

↳ Dynamic

Approach 1: Centralized DNS

- single point of failure
- traffic volume
- maintenance
- ..

Replicate the database across N servers

↳ (1) Scalability & Performance issues

(2) Expensive to maintain & update

Alternate

Split the database and store on multiple servers

Distributed and Hierarchical System

split the database?

- Q: On what basis to decentralize?

Hierarchical domain name space

E.g. cse.iitd.ac.in - IP mapping

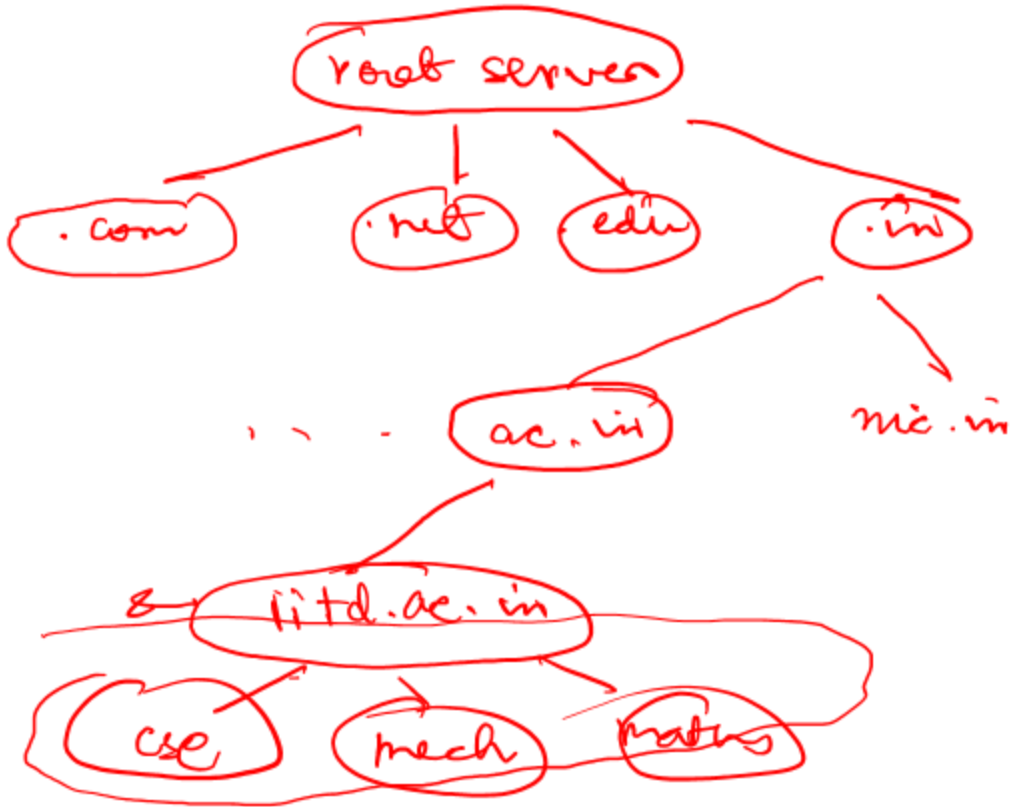
① Split it alphabetically

iitd.ac.in

cse.iitd.ac.in → should be on the same server

(Autonomy engraved into this)

② Split based on the domain name structure



Reverse DNS lookup

IP → Hostname

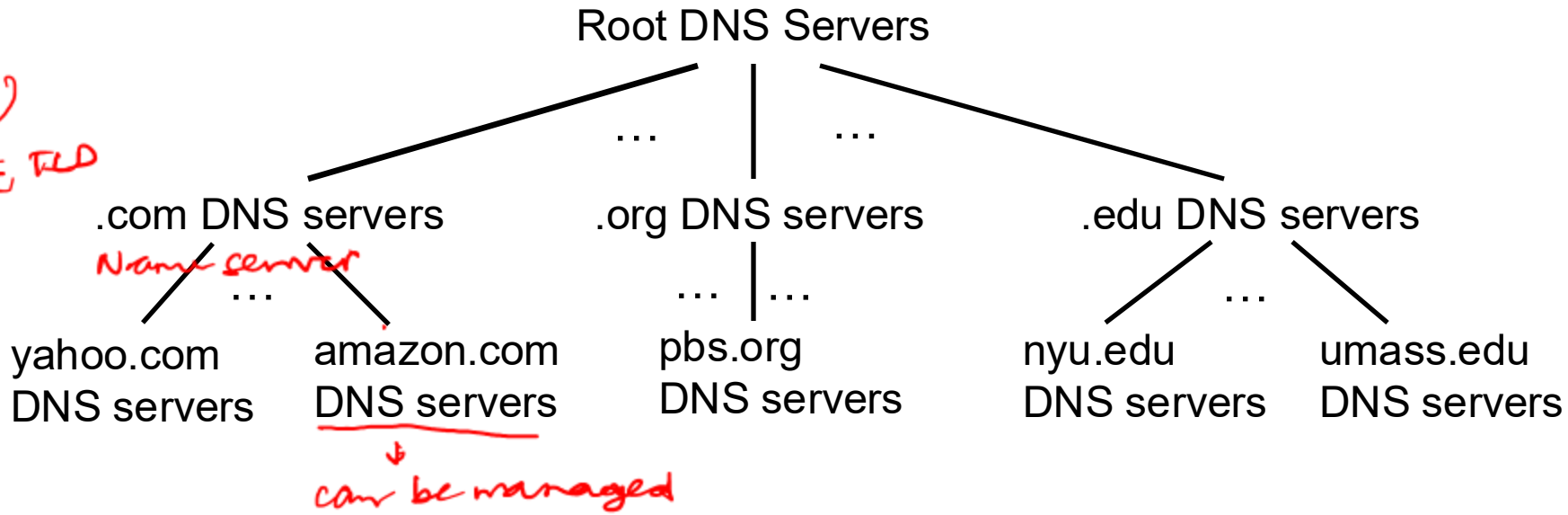
manage

Decentralized and distributed system

(cool334.com , IPaddress)

- Q: On what basis to decentralize? Hierarchical domain name space
- Partition domain name hierarchy into zones managed by some authority → DNS registrars
  - E.g., ICANN is responsible for storing information about top-level domains
- Each zone corresponds to a name server

Name servers



Root

Top Level Domain

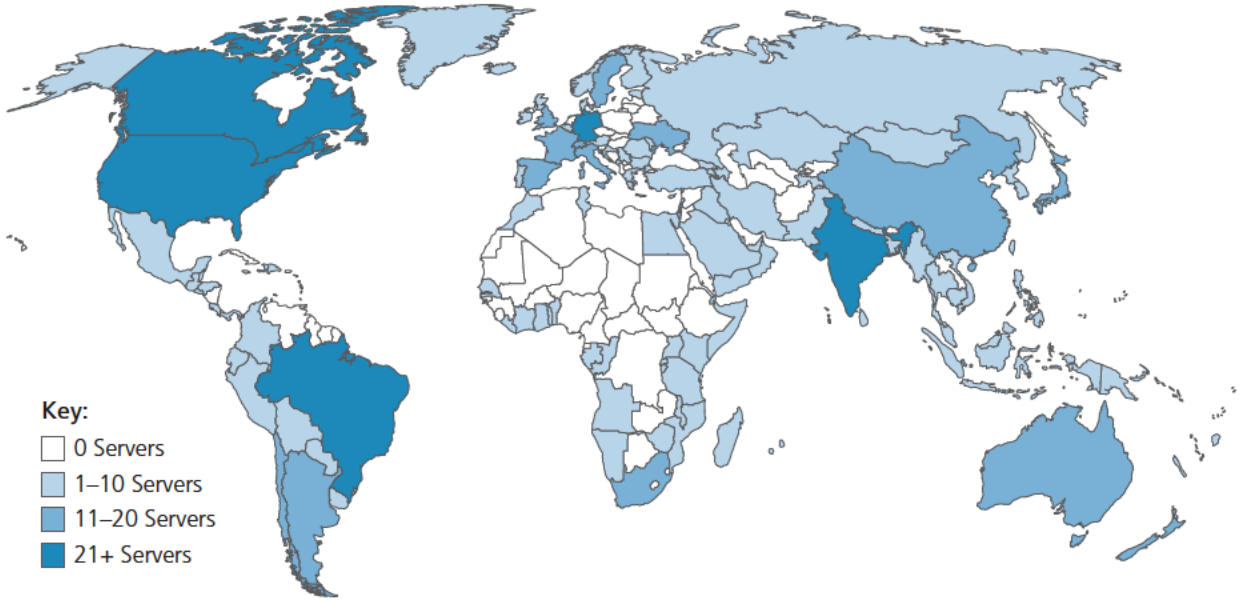
Authoritative

- Name servers are replicated and may be geographically distributed for reliability

DNS: root name servers

- official, contact-of-last-resort by name servers that can not resolve name
- *incredibly important* Internet function
  - Internet couldn't function without it!
  - DNSSEC – provides security (authentication, message integrity)
- ICANN (Internet Corporation for Assigned Names and Numbers) manages root DNS domain

13 logical root name "servers" worldwide each "server" replicated many times (~200 servers in US)



DNS records

DNS: distributed database storing resource records (RR)

RR format: (`name, value, type, ttl`) → used for caching

type=A

- `name` is hostname
- `value` is IP address

type=NS

- `name` is domain (e.g., foo.com)
- `value` is hostname of authoritative name server for this domain ns.foo.com

type=CNAME

- `name` is alias name for some "canonical" (the real) name
- www.ibm.com is really servereast.backup2.ibm.com
- `value` is canonical name

type=MX

- `value` is name of SMTP mail server associated with `name`

Example

AAAA type record - IPv6

Glue record

Root Name server (Top-level domain)

① . in , dns.xyz.in , NS, TTL

② dns.xyz.in , 112.112.112.112, A,

③ .com , dns.xyz.com , NS

④ dns.xyz.com , 113.113.113.113, A

- edu

. net .

(who has authority where to ut?)

IN name Server

① .ernet.in , dns.ernet.in , NS, TTL

② dns.ernet.in , x.y.z.w, A

example.com → ns.cloudfare.com → $IP_x \rightarrow IP_y$

ns.example.com

cd334.com

COM Name server

① example.com , ns.example.com , NS, TTL

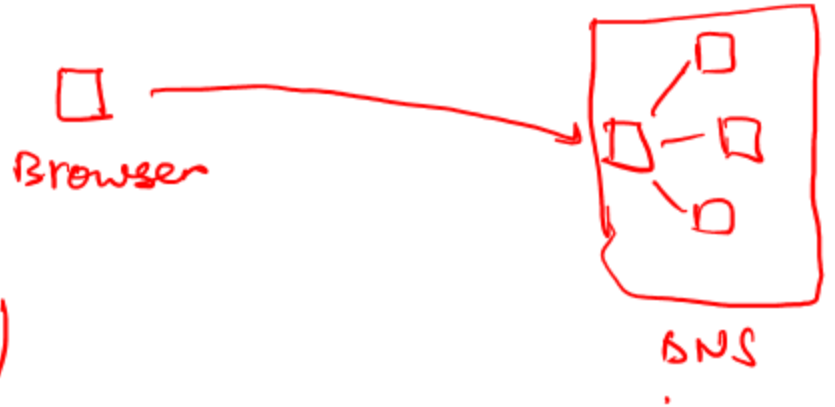
② ns.example.com , IP_x , A, TTL

① example.com , ns.cloudfare.com , NS, TTL

② cd334.com , ns.cloudflare.com , NS, TTL

How to do name resolution?

How To Do DNS Resolution?



- What is the transport protocol?

- Does the browser directly query the root server?

Aggregation → DNS resolver

Local DNS name servers

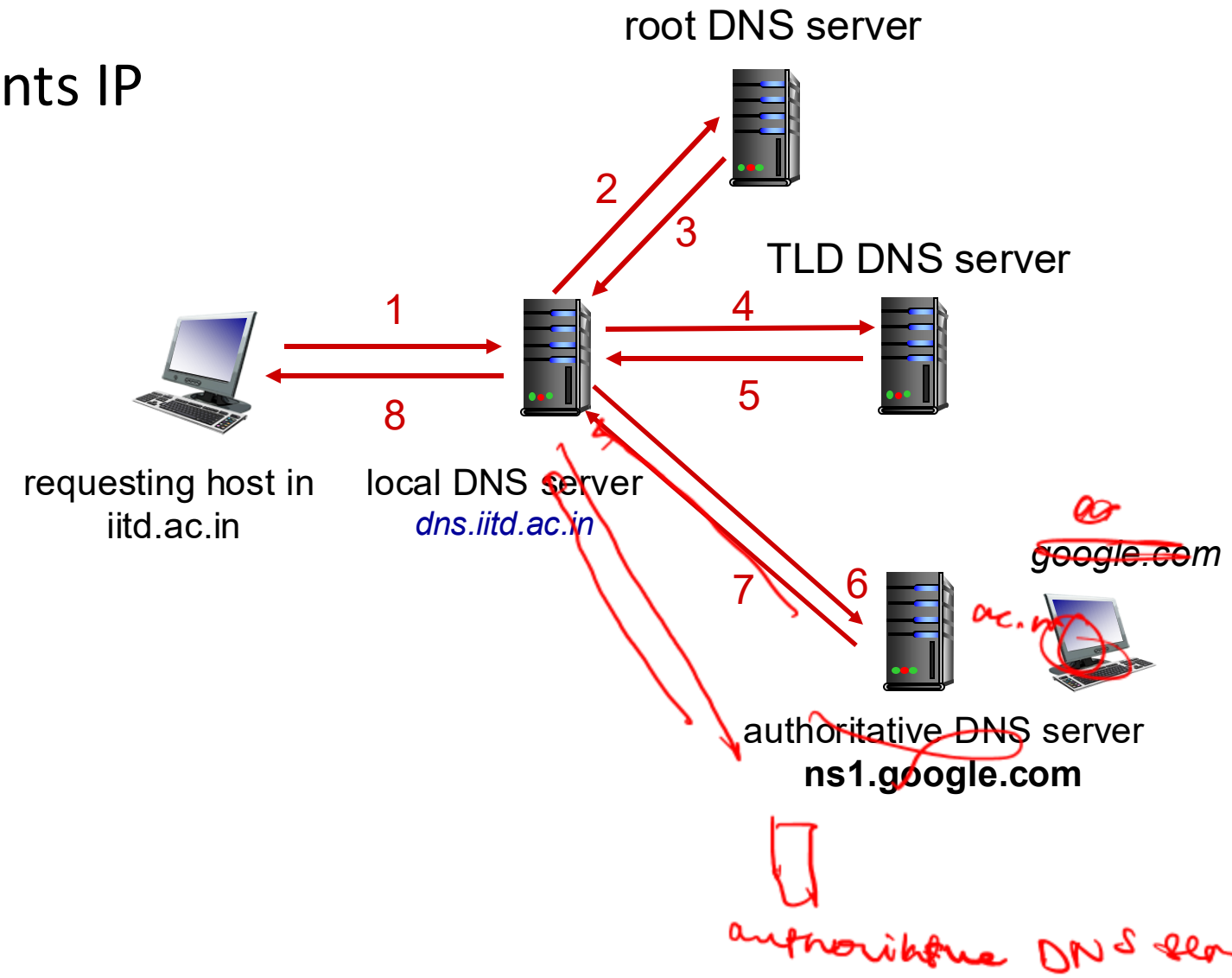
- when host makes DNS query, it is sent to its *local* DNS server
  - Local DNS server returns reply, answering:
    - from its local cache of recent name-to-address translation pairs (possibly out of date!)
    - forwarding request into DNS hierarchy for resolution
  - each ISP has local DNS name server; to find yours:
    - MacOS: `% scutil --dns`
    - Windows: `>ipconfig /all`
- local DNS server doesn't strictly belong to hierarchy

DNS name resolution: iterated query

Example: host at iitd.ac.in wants IP address for google.com

Iterated query:

- contacted server replies with name of server to contact
- "I don't know this name, but ask this server"

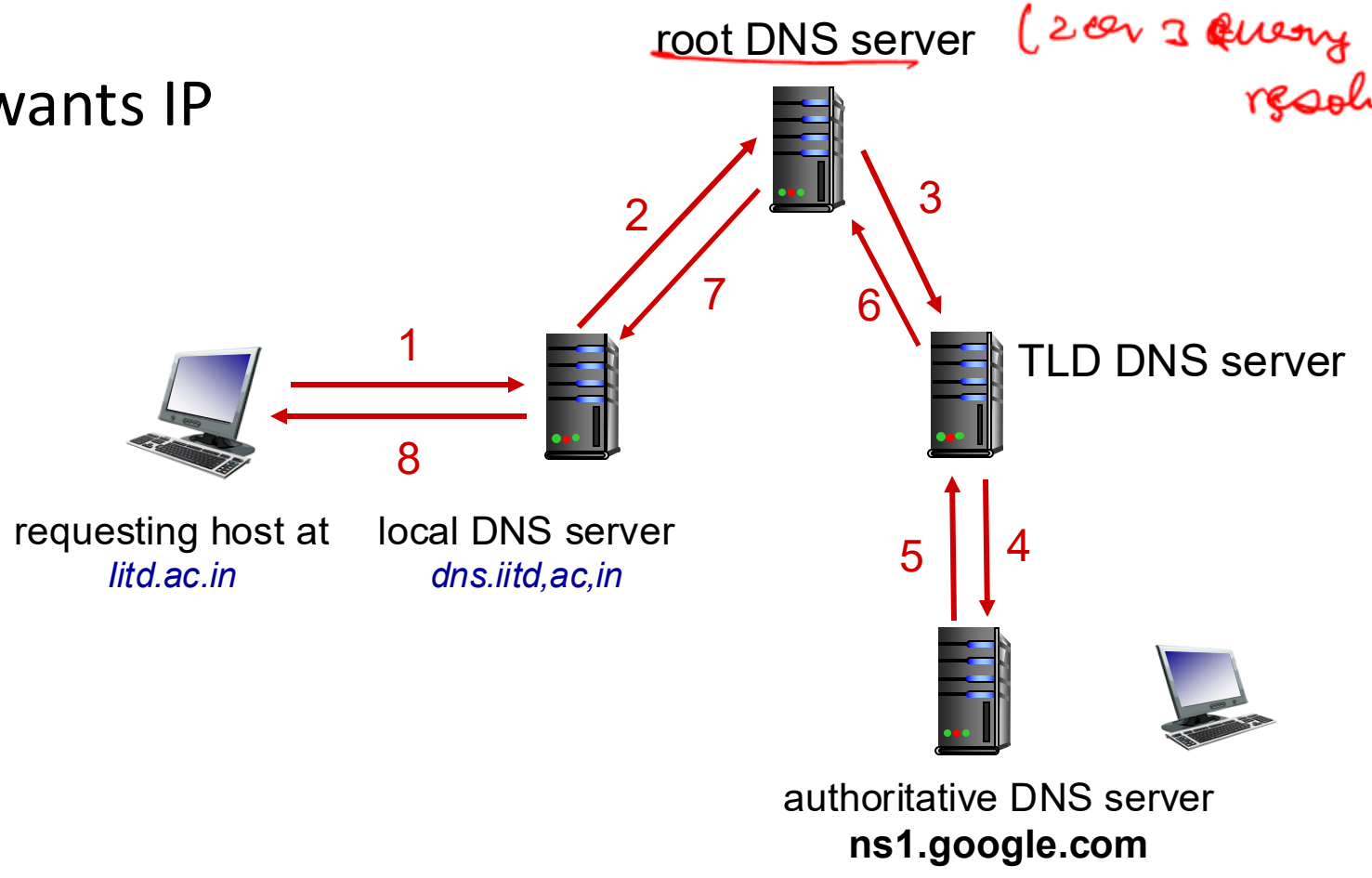


DNS name resolution: recursive query

Example: host at iitd.ac.in wants IP address for google.com

Recursive query:

- puts burden of name resolution on contacted name server
- heavy load at upper levels of hierarchy?

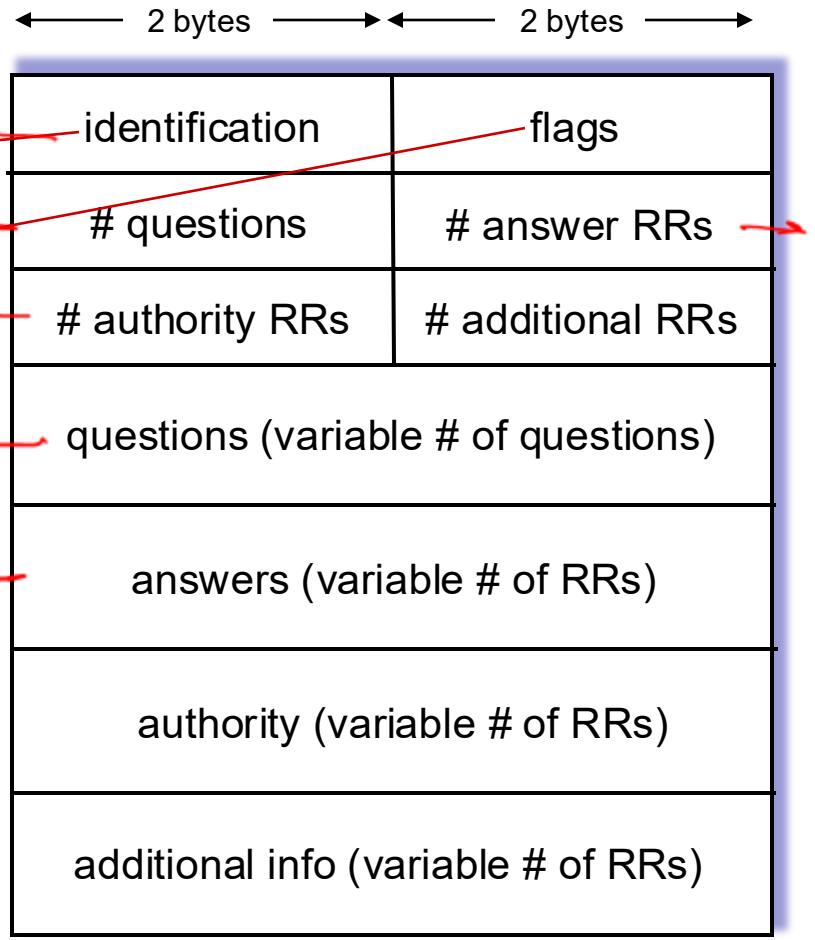


DNS protocol messages

DNS *query* and *reply* messages, both have same *format:*

message header:

- identification: 16 bit # for query, reply to query uses same #
- flags:
  - query or reply
  - recursion desired
  - recursion available
  - reply is authoritative

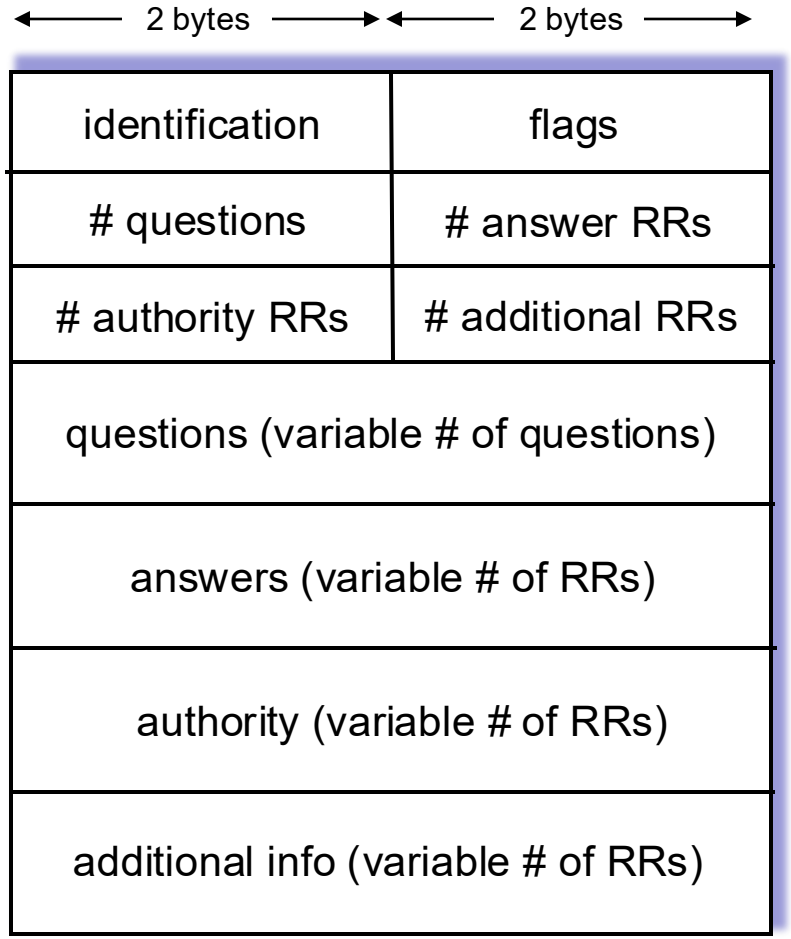


DNS protocol messages

DNS *query* and *reply* messages, both have same *format:*

dig google.com

```
▾ Queries
  ▾ google.com: type A, class IN
      Name: google.com
      [Name Length: 10]
      [Label Count: 2]
      Type: A (Host Address) (1)
      Class: IN (0x0001)
▾ Answers
  ▾ google.com: type A, class IN, addr 142.250.194.142
      Name: google.com
      Type: A (Host Address) (1)
      Class: IN (0x0001)
      Time to live: 227 (3 minutes, 47 seconds)
      Data length: 4
      Address: 142.250.194.142
```



Caching DNS Information

- once (any) name server learns mapping, it *caches* mapping, and i*mmediately* returns a cached mapping in response to a query
  - caching improves response time
  - cache entries timeout (disappear) after some time (TTL)
  - TLD servers typically cached in local name servers
- cached entries may be *out-of-date*
  - if named host changes IP address, may not be known Internet-wide until all TTLs expire!
  - *best-effort name-to-address translation!*

Getting your info into the DNS

example: new startup “Network Utopia”

- register name networkuptopia.com at *DNS registrar* (e.g., Network Solutions)
  - provide names, IP addresses of authoritative name server (primary and secondary)
  - registrar inserts NS, A RRs into .com TLD server:

```
(networkutopia.com, dns1.networkutopia.com, NS)
(dns1.networkutopia.com, 212.212.212.1, A)
```

- create authoritative server locally with IP address `212.212.212.1`
  - type A record for www.networkuptopia.com
  - type MX record for networkutopia.com

xyz.github.io Nameserver lukup
↳ Managed by Github

DNS observations

DDoS attacks

- bombard root servers with traffic
  - not successful to date
  - traffic filtering
  - local DNS servers cache IPs of TLD servers, allowing root server bypass

Spoofing attacks

- intercept DNS queries, returning bogus replies
  - DNS cache poisoning
  - RFC 4033: DNSSEC authentication services

Centralization of DNS

- Name servers hosted by third-party (e.g., cloudflare, amazon)
  - Why?
  - Single point of failure?