

Computer Networks

COL 334/672

Transport Layer

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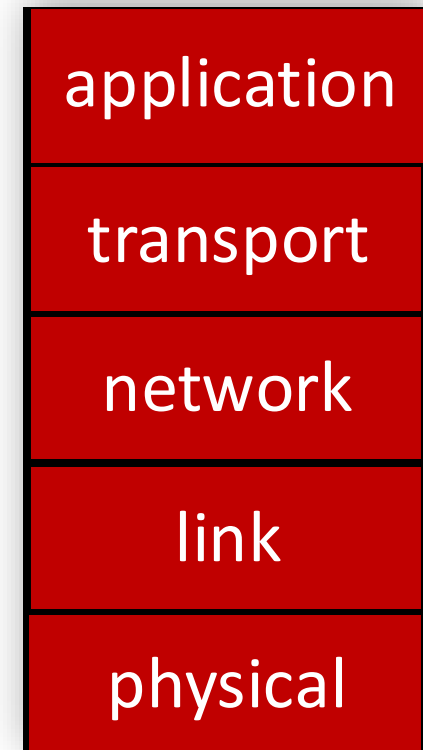
Sem 1, 2024-25

Quiz on Moodle

- Password: sdn

Course Recap

- *application*: supporting network applications
 - HTTP, IMAP, SMTP, DNS
- *transport*: process-process data transfer
 - TCP, UDP
- *network*: routing of datagrams from source to destination
 - IP, routing protocols
- *link*: data transfer between neighboring network elements
 - Ethernet, 802.11 (WiFi), PPP
- *physical*: bits “on the wire”



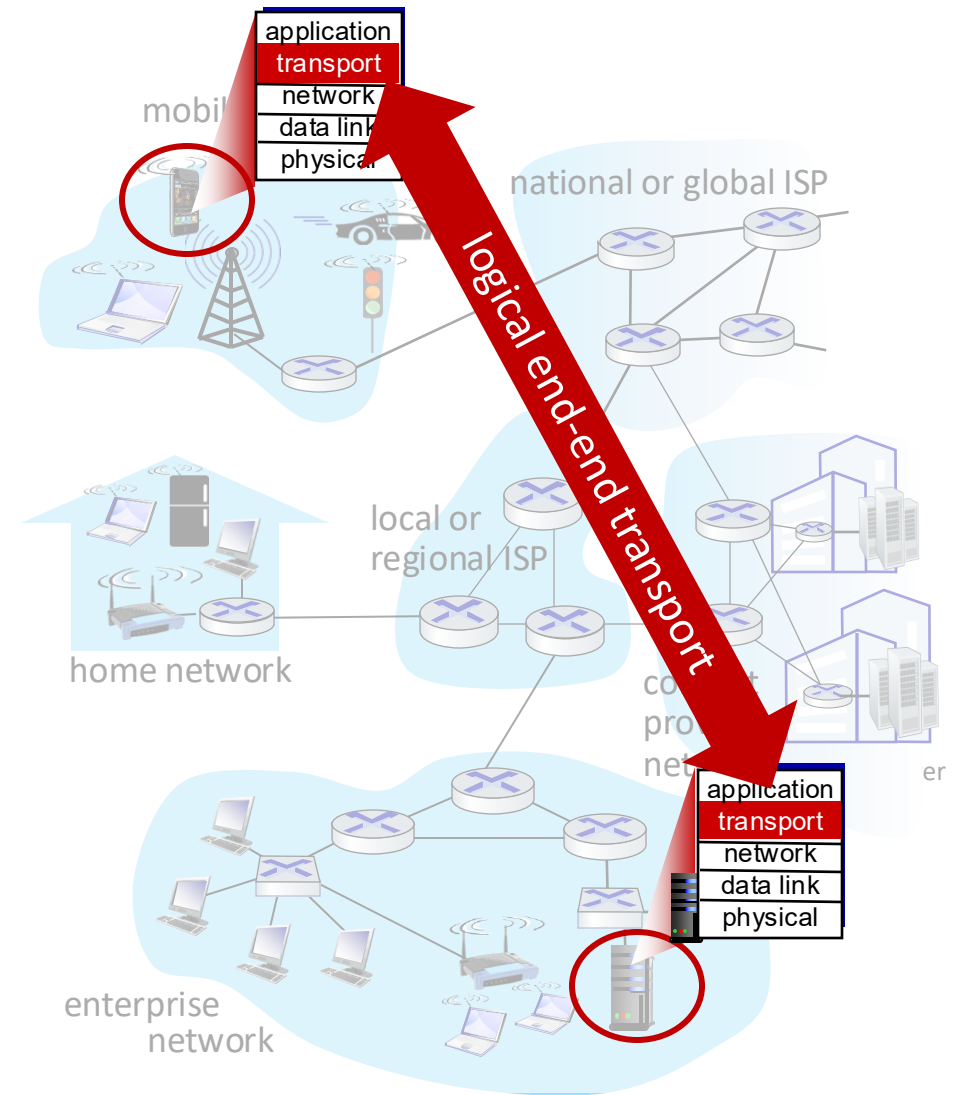
Transport layer: overview

Our goal:

- Understand principles behind transport layer services:
 - multiplexing, demultiplexing
 - reliable data transfer
 - flow control
 - congestion control
- Learn about Internet transport layer protocols:
 - UDP:
 - connectionless transport
 - TCP:
 - connection-oriented reliable transport
 - TCP congestion control

Transport services and protocols

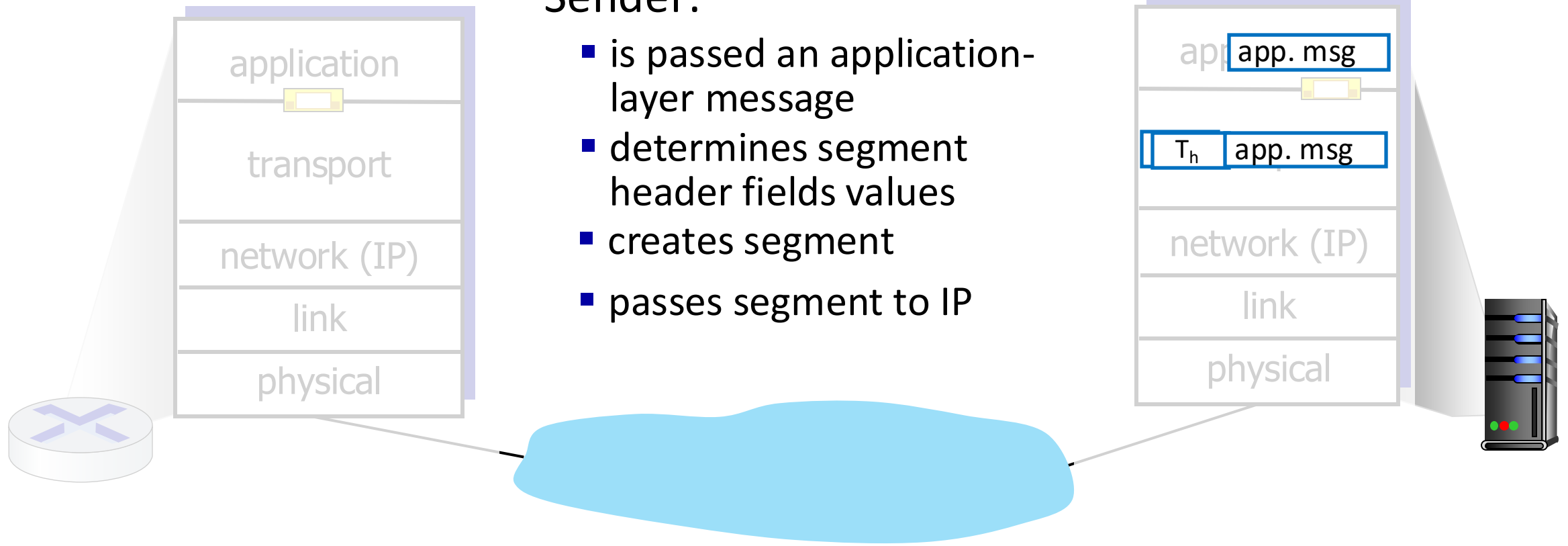
- provide *logical communication* between application processes running on different hosts
- transport protocols actions in end systems:
 - sender: breaks application messages into *segments*, passes to network layer
 - receiver: reassembles segments into messages, passes to application layer
- two transport protocols available to Internet applications
 - TCP, UDP



Transport Layer Actions

Sender:

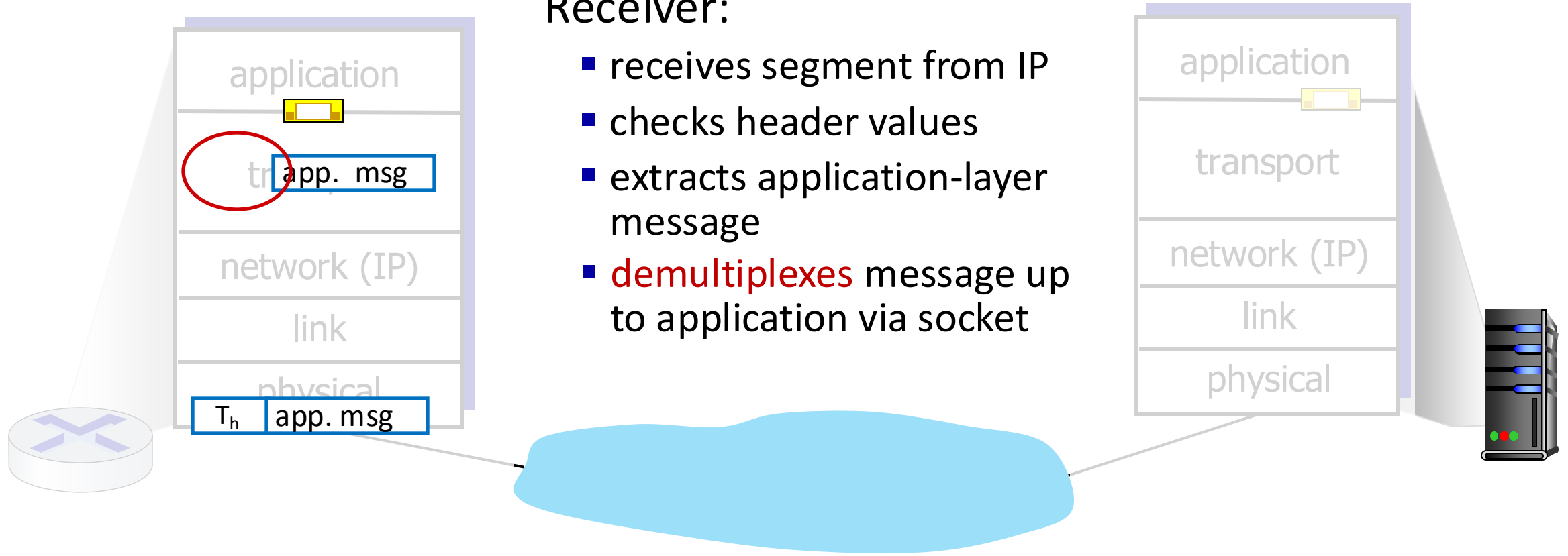
- is passed an application-layer message
- determines segment header fields values
- creates segment
- passes segment to IP



Transport Layer Actions

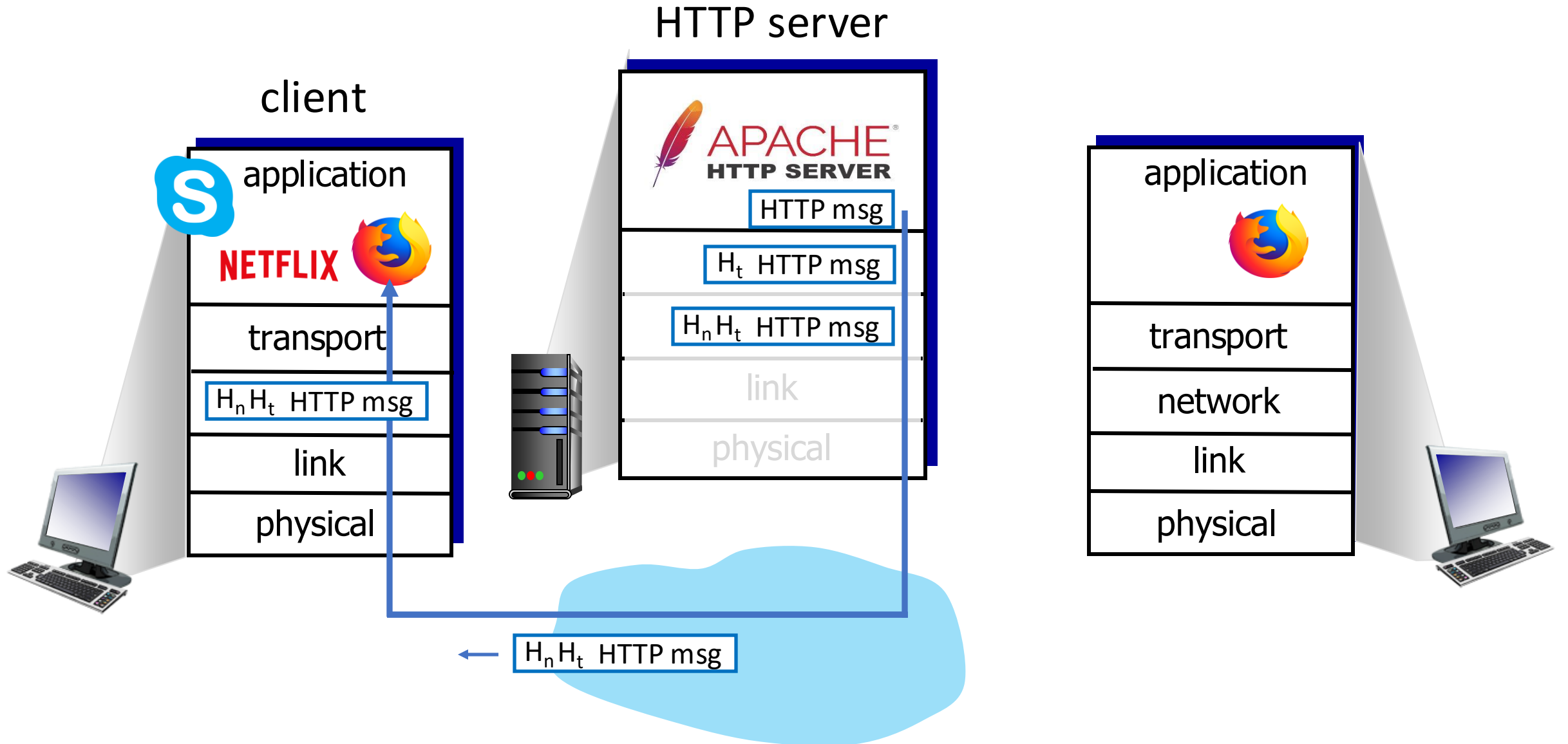
Receiver:

- receives segment from IP
- checks header values
- extracts application-layer message
- **demultiplexes** message up to application via socket



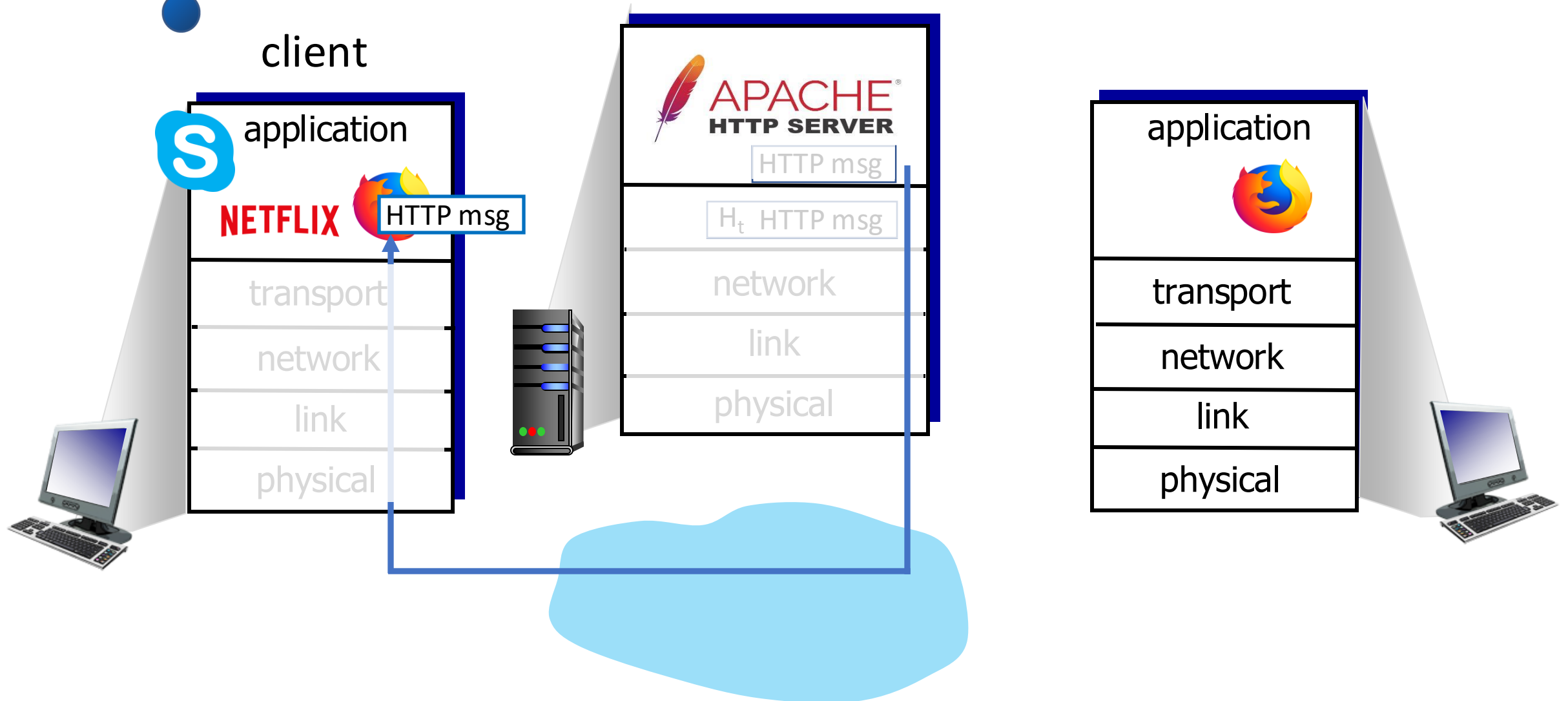
Transport-layer services

- Multiplexing/demultiplexing
- Reliable delivery
- Flow control
- Congestion control



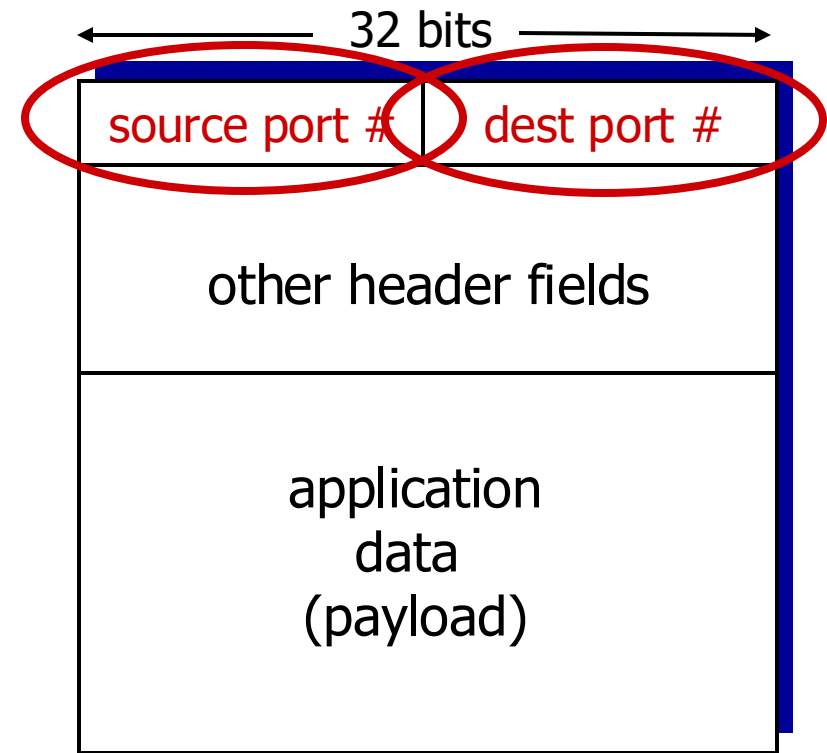


Q: how did transport layer know to deliver message to Firefox browser process rather than Netflix process or Skype process?



How demultiplexing works

- host receives IP datagrams
 - each datagram has source IP address, destination IP address
 - each datagram carries one transport-layer segment
 - each segment has source, destination port number
- host uses *IP addresses & port numbers* to direct segment to appropriate socket



TCP/UDP segment format

Multiplexing/Demultiplexing

- Implemented using **logical** ports
- Difference in UDP/TCP multiplexing
- UDP or connection-less: Uses 2-tuple
 - (Dst IP, Dst Port) of the incoming packet aka bind (IP, port)
- TCP or connection-oriented: Uses 4-tuple
 - (Src IP, Src Port, Dst IP, Dst Port)

UDP: User Datagram Protocol

- “no frills,” “bare bones”
Internet transport protocol
- “best effort” service, UDP
segments may be:
 - lost
 - delivered out-of-order to app
- *connectionless*:
 - no handshaking between UDP
sender, receiver
 - each UDP segment handled
independently of others

Why is there a UDP?

- No connection establishment delay. Makes it particularly suitable for applications involving short transactions:
 - DNS
 - SNMP
 - DHCP
- Provides applications with the flexibility to implement only a subset of features
 - E.g., reliability is not needed for video conferencing applications

UDP: User Datagram Protocol [RFC 768]

INTERNET STANDARD

RFC 768

J. Postel

ISI

28 August 1980

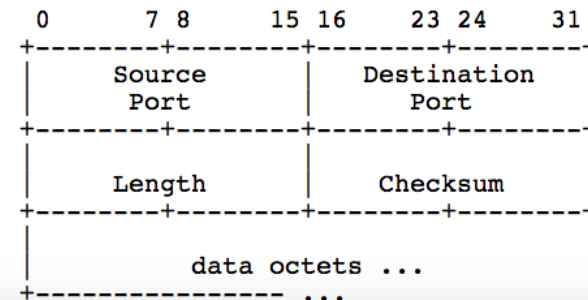
User Datagram Protocol

Introduction

This User Datagram Protocol (UDP) is defined to make available a datagram mode of packet-switched computer communication in the environment of an interconnected set of computer networks. This protocol assumes that the Internet Protocol (IP) [1] is used as the underlying protocol.

This protocol provides a procedure for application programs to send messages to other programs with a minimum of protocol mechanism. The protocol is transaction oriented, and delivery and duplicate protection are not guaranteed. Applications requiring ordered reliable delivery of streams of data should use the Transmission Control Protocol (TCP) [2].

Format



UDP segment header

