

Computer Networks

COL 334/672

TCP Reliability and Connection Establishment

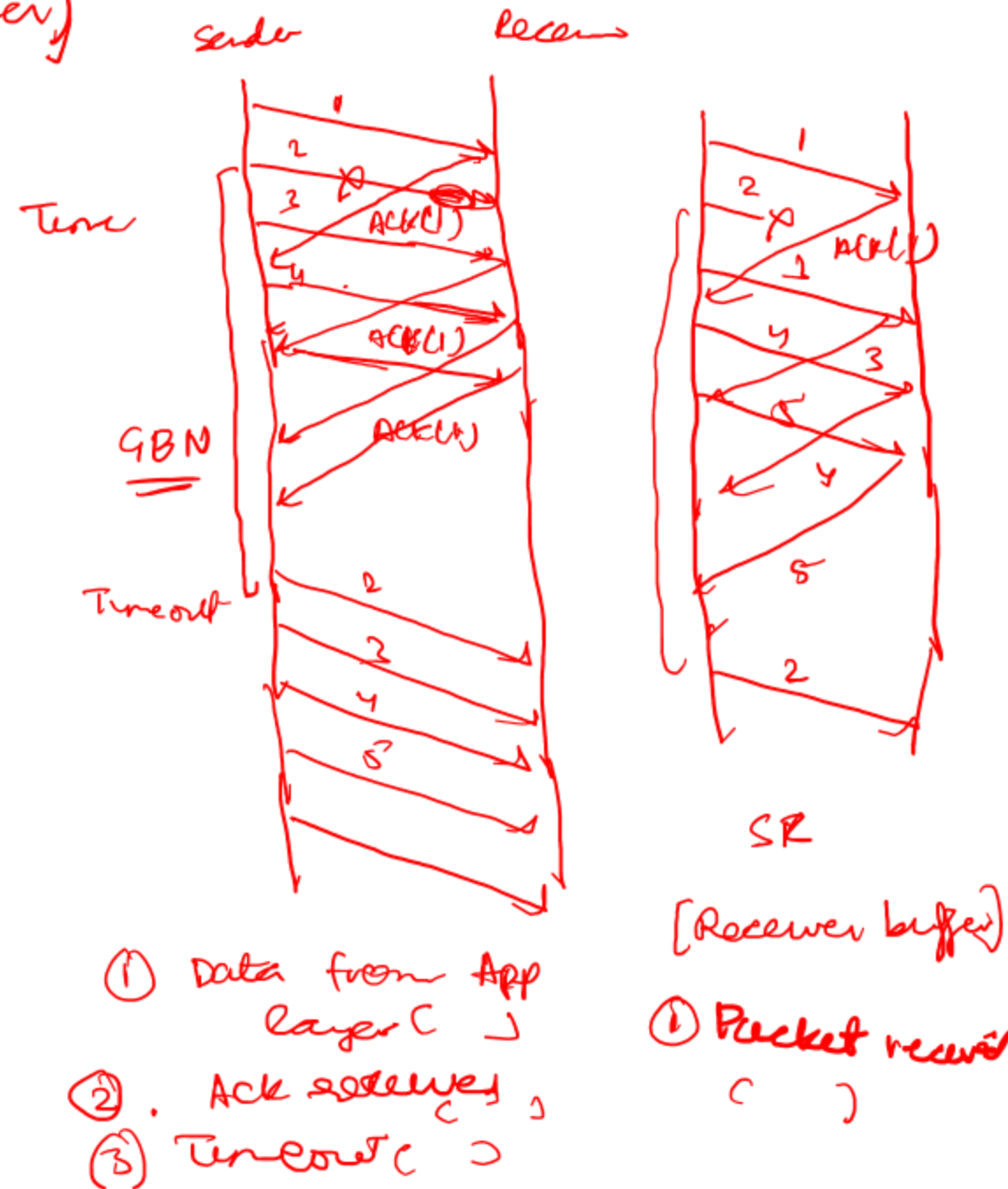
Slides adapted from KR

Sem 1, 2025-26

Recap [Reliability in Transport layer]

(Negative Acknowledgement)

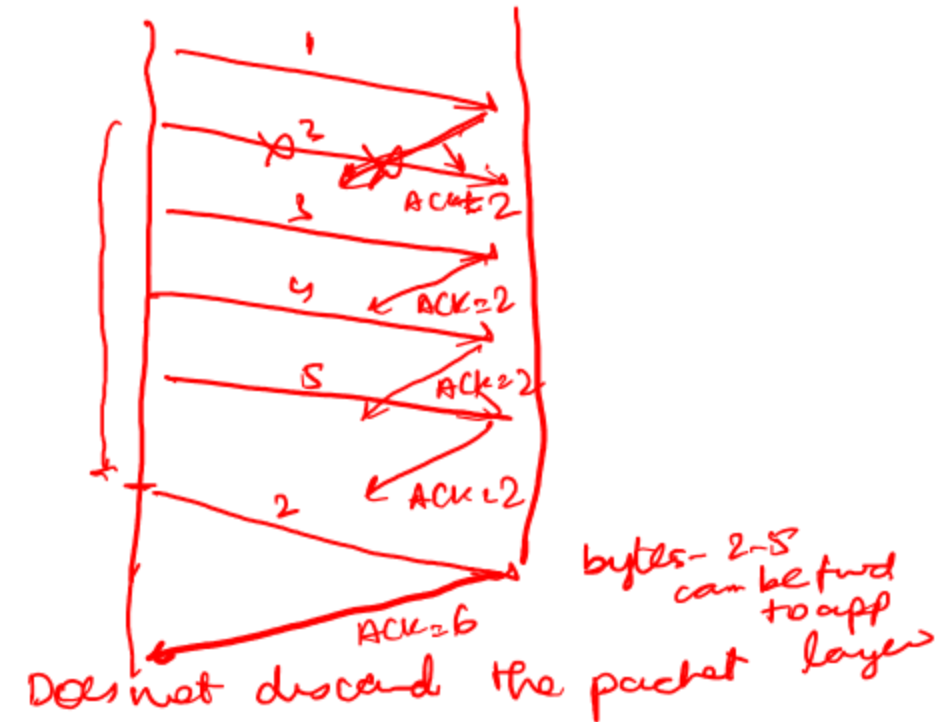
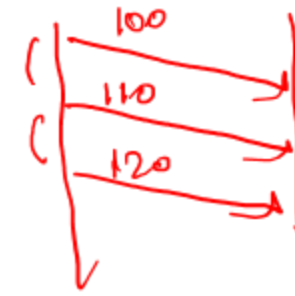
- Sliding window protocols are more efficient than simple stop-and-wait
 - Go-back-N -- cumulative ACKs
 - simpler mechanism at the receiver, but unnecessary retransmissions
 - Selective Repeat – each packet is acked individually
- Event-driven paradigm
- This class: How does TCP implement reliability



Reliability in TCP

- Uses a sliding window protocol
- ACKs are cumulative
- Specification vary for handling out-of-order delivery
 - Generally, out of order delivery packets within the window are not discarded

*TCP is bursty
w/ maximum # of
unacked bytes*



Optimization #1: Delayed ACKs

48 bytes

- Reduce the number of ACKs by sending 1 ACK every 2 packets

Maximum Segment Size

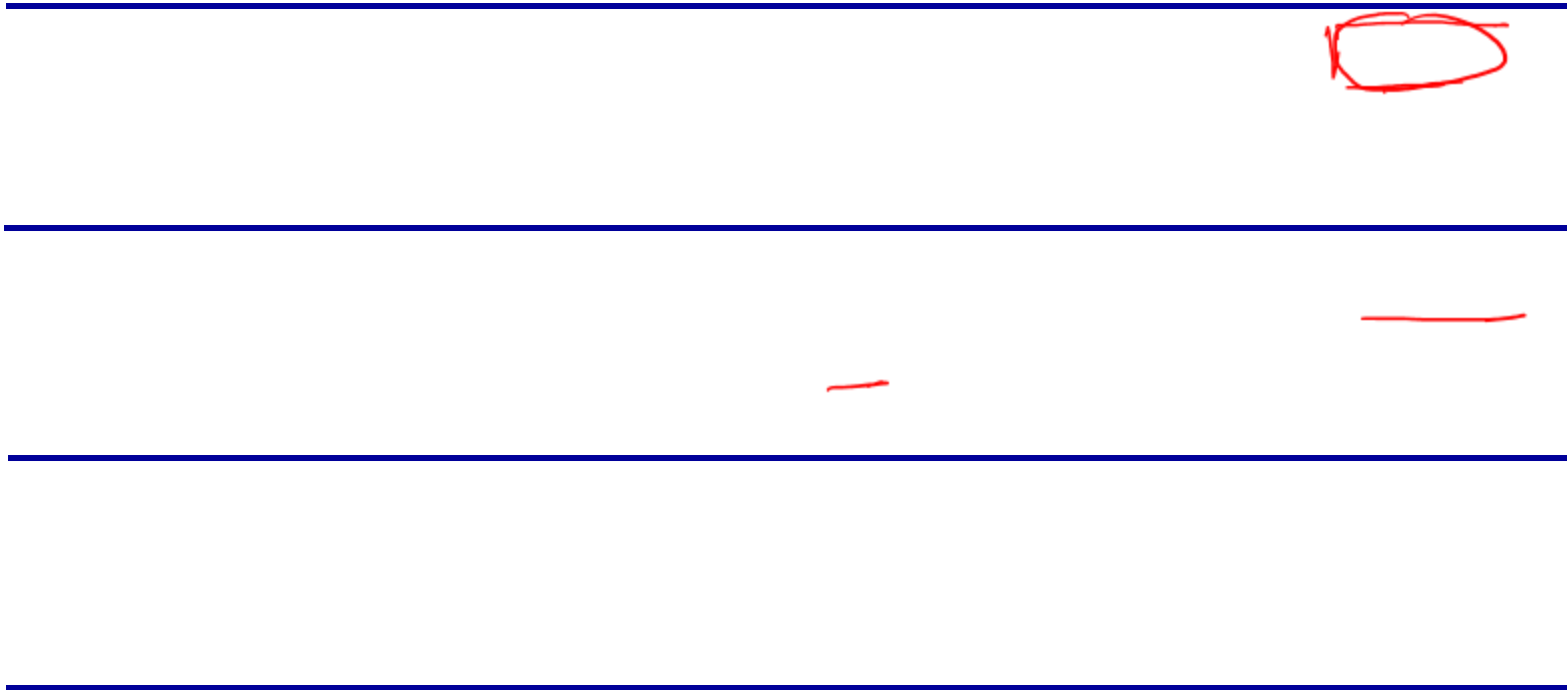
11 MSS - size

6 ACKs

Event at receiver

TCP receiver action

n packets



How to Set Timeout?

premature timeout : Unnecessary Retx

long timeout : Higher latency

$$\text{Timeout} \geq \underline{\underline{\text{RTT}}}$$

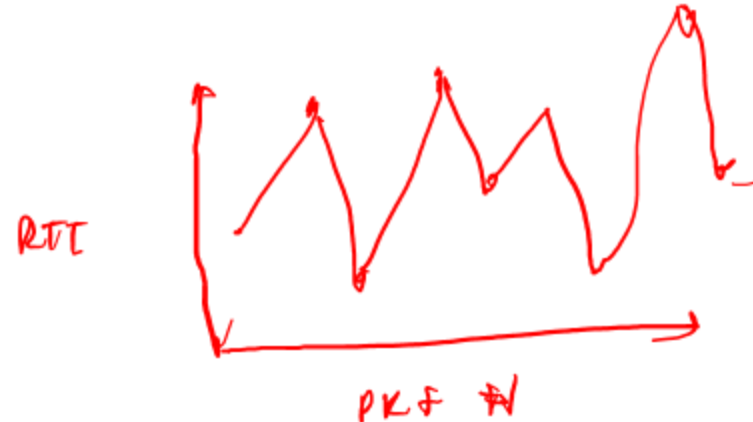
last 10 packet

last packet RTT

$$\boxed{T_o = 1 \text{ Second}}$$



$$T_o = \underline{\underline{\text{RTT}_c}} + \text{safety margin}$$

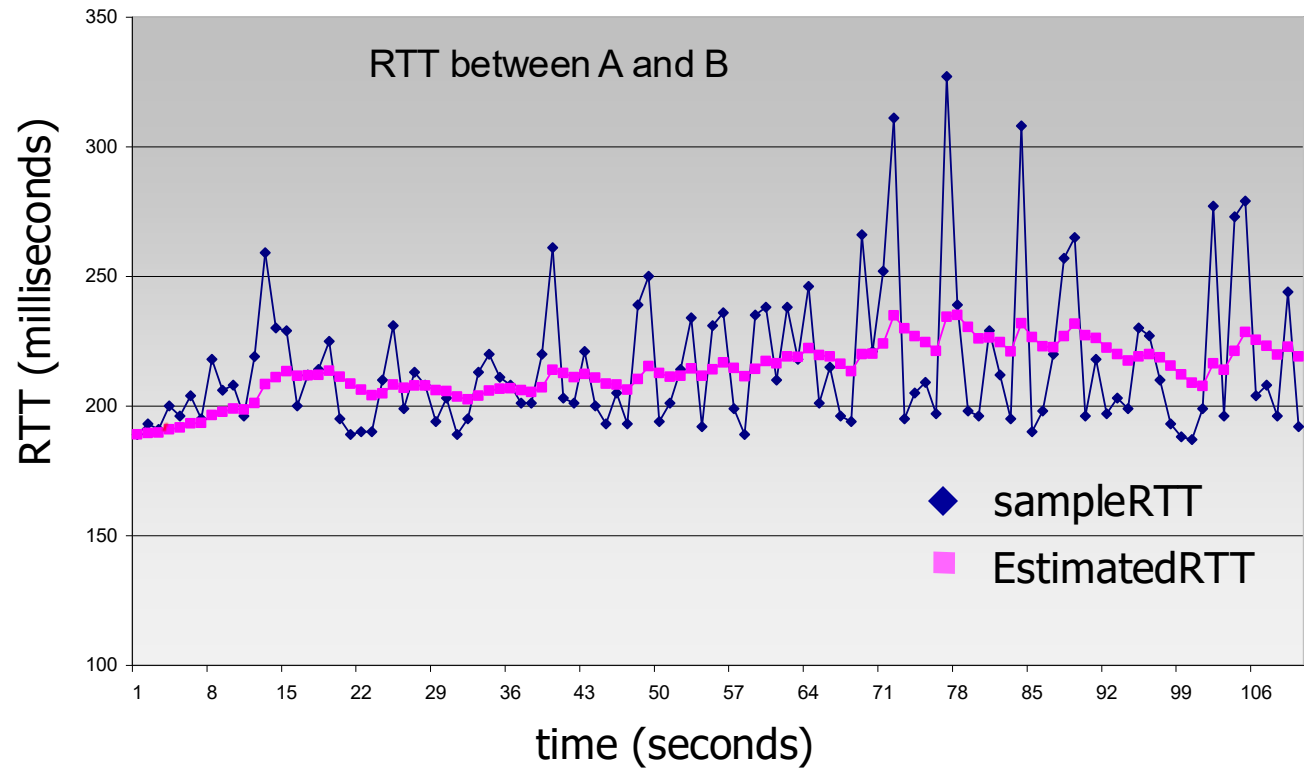


TCP round trip time, timeout

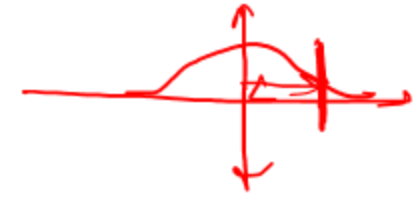
$$\text{EstimatedRTT} = (1 - \alpha) * \text{EstimatedRTT} + \alpha * \text{SampleRTT}$$

- exponential weighted moving average (EWMA)
- influence of past sample decreases exponentially fast
- typical value: $\alpha = 0.125$

$$\text{NewRTT}_E = \text{oldRTT}_E (1 - \alpha) + \alpha \text{ SampleRTT}$$



TCP round trip time, timeout



- timeout interval: **EstimatedRTT** plus “safety margin”
 - large variation in **EstimatedRTT**: want a larger safety margin

$$\text{TimeoutInterval} = \text{EstimatedRTT} + 4 * \text{DevRTT}$$



↑
estimated RTT

↑
“safety margin”

- **DevRTT**: EWMA of (**SampleRTT** deviation from **EstimatedRTT**) :

$$\text{DevRTT} = (1-\beta) * \text{DevRTT} + \beta * |\text{SampleRTT} - \text{EstimatedRTT}|$$

↑
EWMA

(typically, $\beta = 0.25$)

Timeouts take too long!

- Optimization #2:

3 Duplicate all
some data has readed

3 duplicate ACKs

↳ Refractive

Fast Retransmit

TCP fast retransmit

if sender receives 3 additional ACKs for same data ("triple duplicate ACKs"), resend unACKed segment with smallest seq #

- likely that unACKed segment lost,

TCP options: allow selective ACKs



Once there is a timeout
; Double Timeout value (Binary Exp Backoff)

Timeouts take too long!

- Optimization #2:

— *TCP fast retransmit* —

if sender receives 3 additional ACKs for same data (“triple duplicate ACKs”), resend unACKed segment with smallest seq #

- likely that unACKed segment lost, so don't wait for timeout

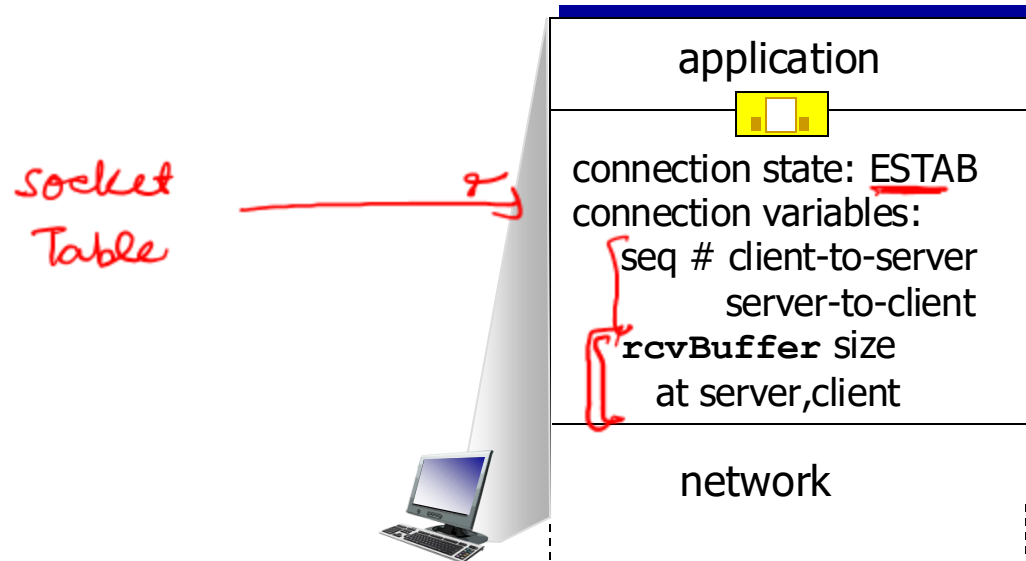
Does TCP implement GBN or SR?

Hybrid

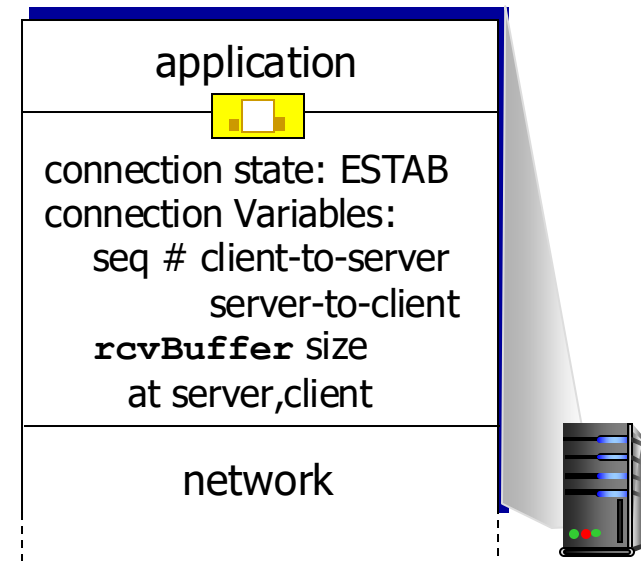
TCP connection management

before exchanging data, sender/receiver “handshake”:

- agree to establish connection (each knowing the other willing to establish connection)
- agree on connection parameters (e.g., starting seq #s)



```
Socket clientSocket =  
    newSocket("hostname", "port number");
```



```
Socket connectionSocket =  
    welcomeSocket.accept();
```


TCP 3-way handshake

Client state

```
clientSocket = socket(AF_INET, SOCK_STREAM)
```

LISTEN

```
clientSocket.connect((serverName, serverPort))
```

SYNSENT

ESTAB

choose init seq num, x
send TCP SYN msg

received SYNACK(x)
indicates server is live;
send ACK for SYNACK;
this segment may contain
client-to-server data



SYNbit=1, Seq=x

SYNbit=1, Seq=y
ACKbit=1; ACKnum=x+1

ACKbit=1, ACKnum=y+1



choose init seq num, y
send TCP SYNACK
msg, acking SYN

received ACK(y)
indicates client is live

Server state

```
serverSocket = socket(AF_INET, SOCK_STREAM)  
serverSocket.bind(('', serverPort))  
serverSocket.listen(1)  
connectionSocket, addr = serverSocket.accept()
```

LISTEN

SYN RCVD

ESTAB