# CS118 Homework 5

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#### Problem 1

Since A is sending data faster than B can remove it, the reciever buffer will fill up. When the buffer is full, B will tell A to stop sending data until B can clear space in the buffer. Once B is ready to recieve more data, it sends a segment to A. A will then send data (likely filling the buffer again) until all data has been sent from A to B.

s port: 5670

d port: 2008

seq\_no: 3120

ack\_no: 981

control bits (all 6 bits): 010001

а

$$\begin{split} LossRate &= W/2 + (W/2+1) + L + W \\ &= \sum_{n=0}^{W/2} (W/2+n) \\ &= (W/2+1)W/2 + \sum_{n=0}^{W/2} n \\ &= (W/2+1)W/2 + \frac{W/2(W/2+1)}{2} \\ &= W^2/4 + W/2 + W^2/8 + W/4 \end{split}$$

$$LossRate = 3W^2/8 + 3W/4$$

b. For large W, 
$$3W^2/8 >> 3W/4$$
 so  $L pprox \frac{8}{3W^2}$ 

$$W=\sqrt{rac{8}{3L}}$$

$$avgRate = 3/4*W/RTT = 3/4*\sqrt{\frac{8}{3L}} \frac{MSS}{RTT} = 1.22 \frac{MSS}{RTT*\sqrt{L}}$$

a. To fully use the network, the rwnd must be larger than delay\*bandwidth.

$$delay = \frac{800Mb}{s} * \frac{10^{6}b}{1Mb} * \frac{1s}{10^{3}ms} = 8 * 10^{5}b/ms$$

$$RecieveWindow \geq 8*10^5 b/ms*400 ms = 3.2*10^8 b = 40 MB$$

$$bits = log_2(40MB) = 26$$
 for RecieveWindow

$$Sequence Num \geq maxLifetime*bandwidth = 25s*800Mbps = 20Gbit = 2.5GB$$

$$bits = log_2(2.5GB) = 32 \ {
m for \ Sequence Num}$$

b. If RecieveWindow of 16 bits then it would be smaller than delay\*bandwidth

$$effective = RecieveWindow/RTT = 2^{16}B/400ms = 163.84KB/s$$

