

//KTR//

## # Infrastructure:

Chapter - 01

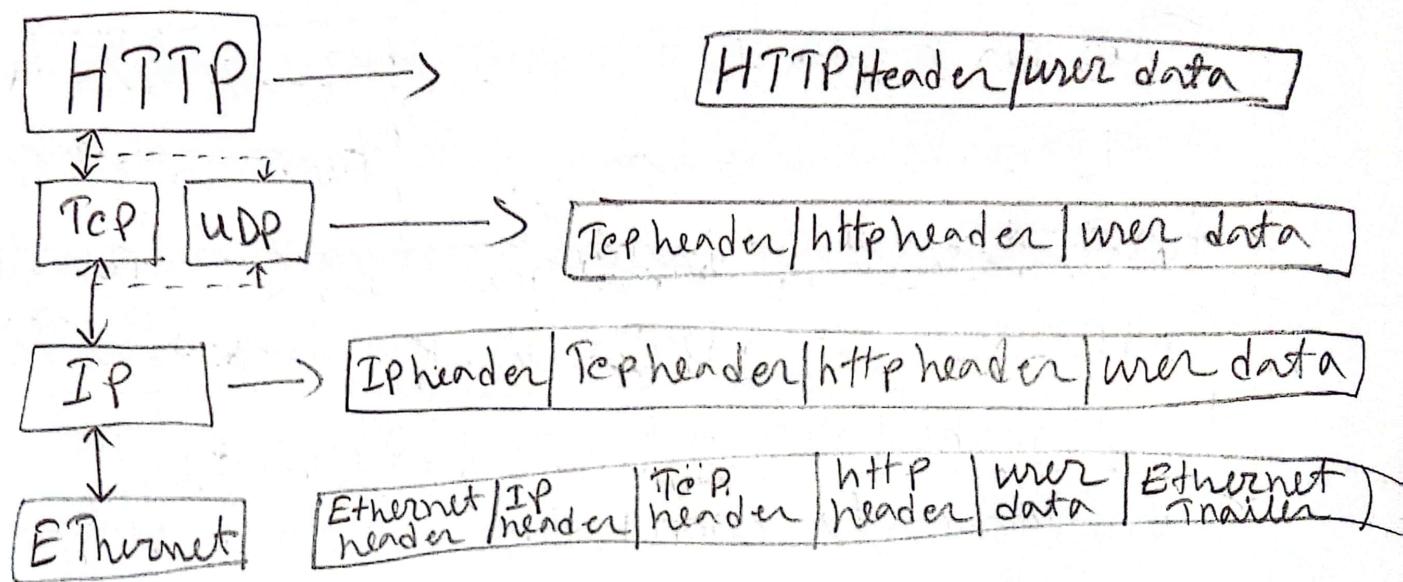
This refers to the physical and virtual components that enable data transmission over the internet. It includes network infrastructure, protocols, ISPs, DNS, CDN, peering, cloud computing, mobile networks and IoT infrastructure. These elements work together to facilitate global connectivity and data exchange.

## # PDU (Protocol data unit):

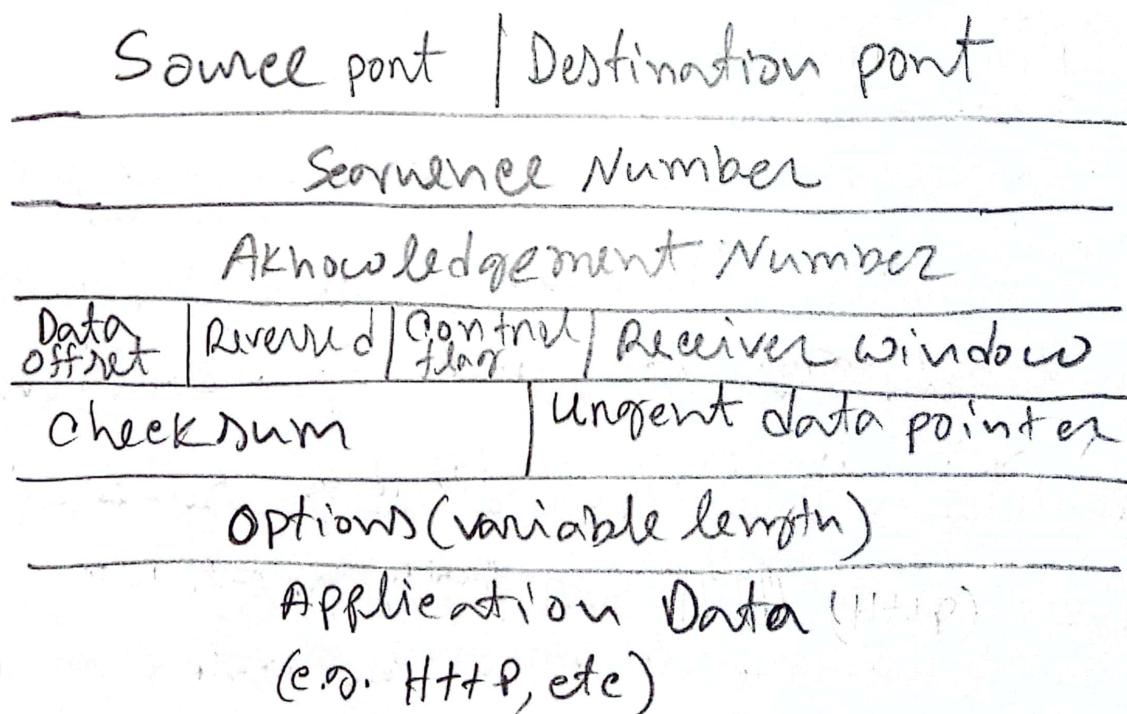
It is the basic unit of exchange between entities that communicate using a specified networking protocol. There are several protocols using PDU transfer in their protocol stack, such as TCP, UDP, IP, Ethernet, HTTP, SMTP and DNS.

To transport user data using layer specific Protocol data units (PDU) with header, trailer and

transported service data units (SDU),



\* Format of a TCP PDU / TCP Segment:



## \* IPv4 - packet format:

Version (4 bit)	Header length (4 bit)	Type of service (8 bit)	Total length
Identifier (16 bit)	Flags	Fragment offset	
Time to live (8 bit)	Protocol (8 bit) TCP, UDP, ICMP	Header checksum (16 bit)	
Source IP Address (32 bit)			
Destination IP address (32 bit)			
options			
Application data			

## → Purpose of header in IPv4 packet:

Header in an IPv4 packet, provides essential information for routing and delivering the packet. It helps to identify the packet, determine its priority, specify its size, set a maximum hop limit, encapsulated protocol and ensure header integrity.

## # Components of web architecture in the internet:

- ① Client (web browser, mobile app)
- ② Web servers    ③ Protocols (HTTP)
- ④ DNS    ⑤ ISP    ⑥ CDN    ⑦ Database
- ⑧ Web application frameworks and programming languages.
- ⑨ Security mechanism (SSL/TLS)
- ⑩ Standards and specifications defined by organizations like W3C (www consortium)

## # Basic components of a web architecture:

- ① Clients: user device, like computer.
- ② Servers: stores and manages web resources and responds to requests.
- ③ Web browser: displays web content to user.
- ④ HTTP Protocol: defines how data is transmitted over the web.

- ⑤ Web servers: Handles incoming requests and serves files.
- ⑥ Application / Database server
- ⑦ Front-end      ⑧ Back-end
- ⑨ Caching: Stores frequently used data for quick retrieval
- ⑩ load balancer: Distributes traffic among multiple servers
- ⑪ CDN (Content Delivery Network): Optimizes content delivery for speed.

## # HTTP protocol and its evolution:

### \* Basic notions:

→ WWW - world wide web

→ HTML - Hypertext markup language

→ HTTP -          u      transfer protocol

→ URL - uniform resource locator

#HTTP: Standard application protocol between <sup>web</sup> server and web client to data access and transport.

→ It applies TCP as a reliable transport protocol with port 80 at a server.

→ It is a readable ASCII protocol.

#URL: It is a unique <sup>web</sup> address of each web page or object in the web. URL denotes the logical location of an instance in a network context.

Eg: `http://example.com/books/book-039`

#URI:

uniform resource identifier, is a 7-bit encoded ASCII character string to identify objects in a unique manner.

#URN: uniform resource name, an info

unit in the namespace having a designated identity.

Example: urn:isbn: 087222156002 > URI  
                          URN

# URC: uniform resource characteristics, it is a meta data of a resource.

# Task and feature of Http:

① Http uses tcp transport service to receive the data transportation. A client never initiates a Tcp connection and that will be closed after the data transfer in http 1.0 and kept open for further transfers in http 1.1.

② HTTP is stateless. Because, state management is too expensive for a protocol. It will make the operation slow and complex.

## # Persistent and non-persistent TCP connection:

### ① Non-persistent:

- Each request opens its own TCP connection for the object data transport, of the replies.
- The server closes the connections after a transfer.
- Used in HTTP 1.0

### ② Persistent:

- An opened connection remains open.
- All further object requests use that open TCP connection.
- Used in HTTP 1.1 (in sequential order)
- New request issued without waiting for a reply of previous one, in the same TCP packet.
- Better performance with small objects.

## # Limitations of HTTP 1.1, TLS and TCP

- TCP connection in HTTP uses 3-way handshake and several flow starts for upto 4-8 parallel TCP connections in the browser.
- Head of line blocking due to a slow process of the first request.
- High TLS complexity

## # HTTP 1.0 vs 2.1 vs 2.0 :

Every http protocol version is stated as,

### \* HTTP 1.0 :

- Each http request establishes a new connection to the server and is closed after receiving the response.
- Connections are non-persistent.
- Headers are often lengthy and inefficient.
- Resource loading is slow because of multi-

-ple connection required.

→ Caching: Weak caching management.

## ⑧ HTTP(1.1):

→ Persistent connection; which allows multi-  
-ple HTTP request and response to be sent over  
a single TCP connection.

→ Host header; It allows multiple site to  
be hosted in a single ip address, which  
helps to adopt the virtual hosting feature.

→ Chunked Transfer-Encoding; It helps  
to transfer large amount of data efficiently.

→ Caching mechanisms were improved,

→ compression; Added content compression  
to reduce data transfer size.

## ④ HTTP 2.0:

- Multiplexing requests: HTTP 2.0 can send multiple requests for data in parallel over a single TCP connection. Helps to faster loading times.
- Header compression: Header fields are compressed and reduced the size to save bandwidth.
- Server push: Here server can proactively push resources to client, like a server can push CSS or JavaScript files after the initial HTML request.
- Stream prioritization: Requests can be prioritize by assigning priority level.
- Binary protocol: Uses binary protocol instead of text, which is efficient to data parsing.
- Backward compatible: It designed to be compatible with HTTP 1.1, allowing older clients or servers to work with it.

## #QUIC protocol:

QUIC (Quick UDP Internet connection) is a <sup>transport</sup> protocol developed by Google.

- It is a generally applicable, secure, connection oriented version layer protocol with multiplexing functionality for data streams.
- It uses UDP connection instead of TCP in HTTP2.0. And combined with TLS 1.3.
- Multiplexing of parallel uni- or bi-directional data streams with or without reliability and security requirements.
- Uses encrypted, authenticated packet header and payload elements for data streams.
- It prevents the HOL (head of line) blocking of data stream during the multiplexing along TCP transport channel.

→ This protocol can reconstruct lost data using a redundant encoding of messages, e.g. by fountain codes.

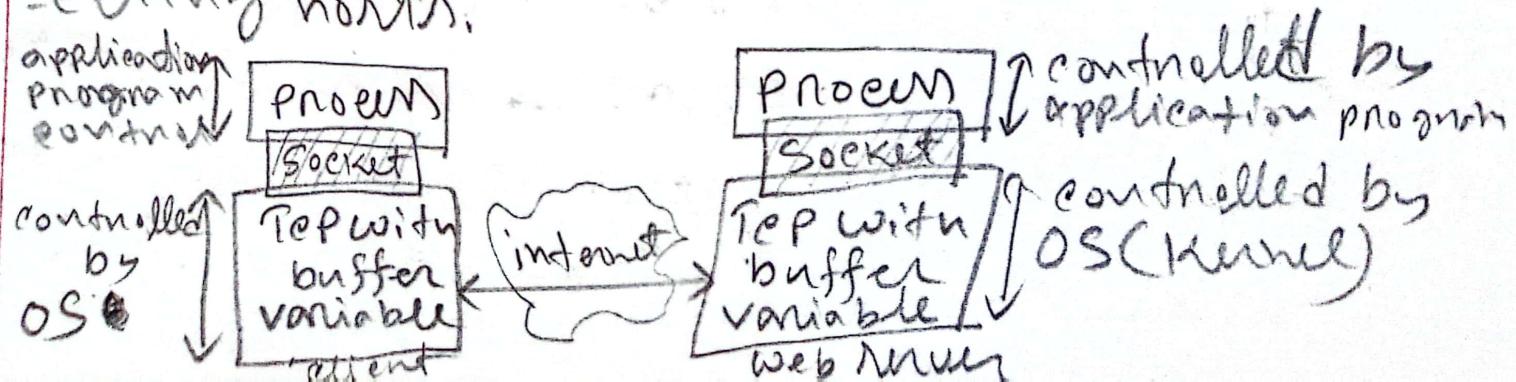
## Chapter 1.2.2 [TCP/IP protocol stack]

### # TCP/IP protocol suite:

- DNS: Domain name system
- HTTP: Hypertext transfer protocol
- FTP: File Transfer protocol
- SMTP: Simple mail transfer protocol
- IMAP: Internet message access protocol
- Telent: Remote call of programs

### # Transport connection setup by application program

- Socket interface: It is a buffer and API for the data exchange between two application process in the user space and an end-to-end transport protocol in the kernel space.
- Tcp transport service: This ensures a reliable transport of a byte stream from one process to another peerings process on two interconnecting hosts.



## # Transport protocol concept:

- ① Link by link, perspective of the data link layer on a physical transmission segment.
- ② End to end, perspective of the transport layer on interconnected transmission segments called transmission path.

## # TCP features:

### ✳ 3-way handshake:

This is a fundamental process of TCP that establishes a reliable connection between a client and a server. It ensures both parties permission before data exchange. The 3 way handshake described below:

#### ① Synchronize (SYN):

Here client initiates the connection and sends a TCP segment with the SYN flag set to 1. Also a sequence number includes, which is an initial sequence number chosen

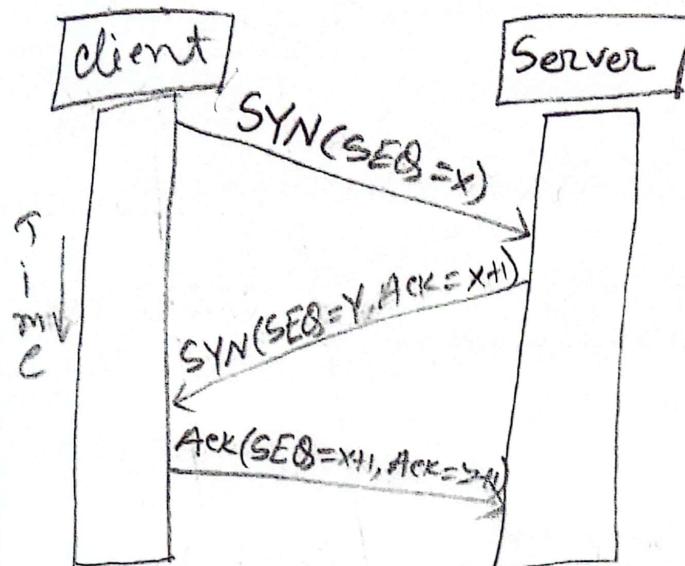
by the client. This SEQ Number represents the first byte of data in the client message.

② Synchronize-Acknowledgement (SYN ACK):

Upon receiving the client's SYN segment, the server responds by sending a TCP segment with both SYN and ACK flags set to 1. This segment includes its own ISN (initial sequence number) and an acknowledgement number indicating the next expecting sequence number from the client.

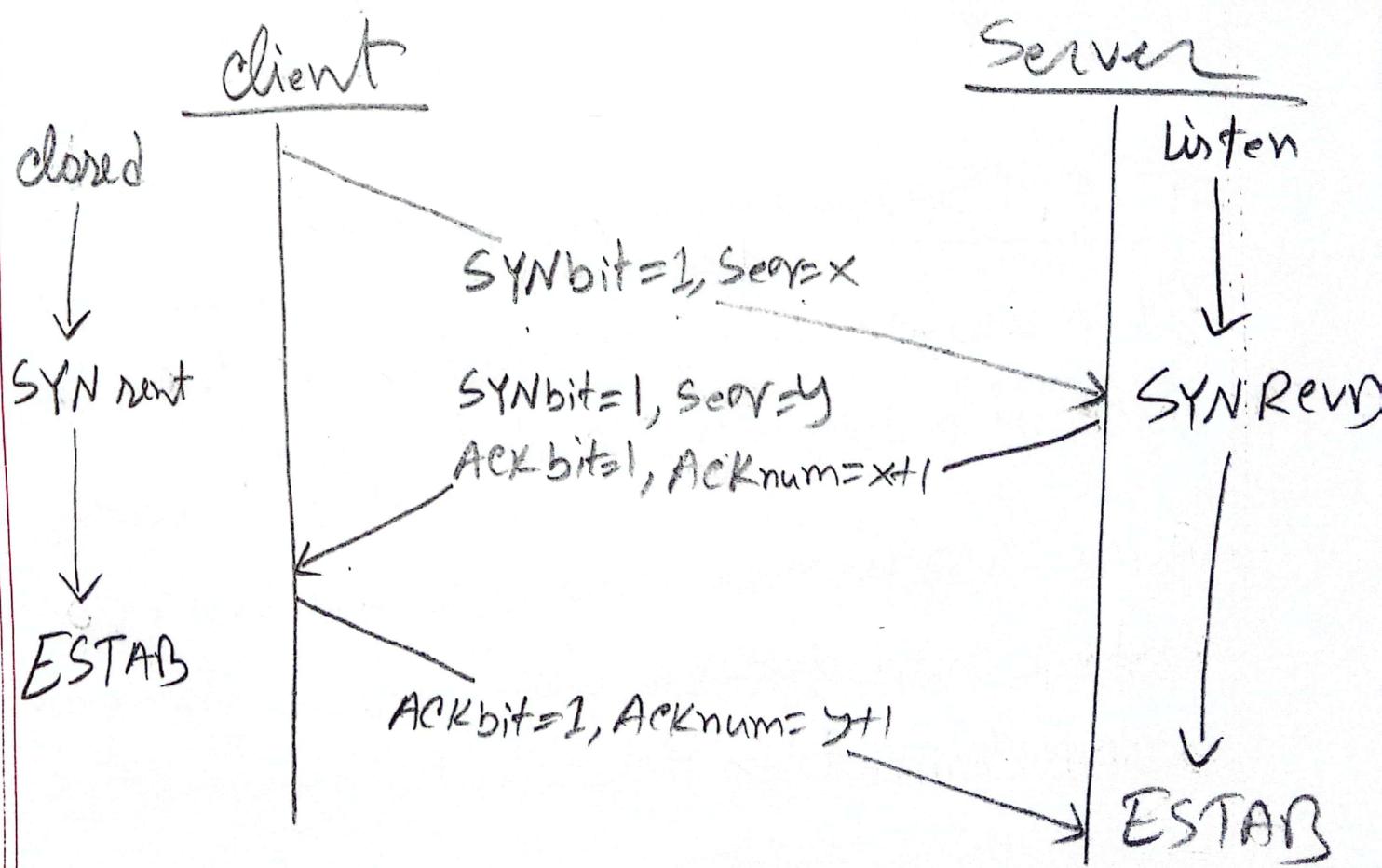
③ Acknowledgement (ACK): Here client acknowledges the server's SYN-ACK segment by sending a TCP segment with the ACK flag set to 1. The ACK number is set to the server's ISN + 1, indicating that the client is ready to start the data transmission.

## # TCP connection setup:

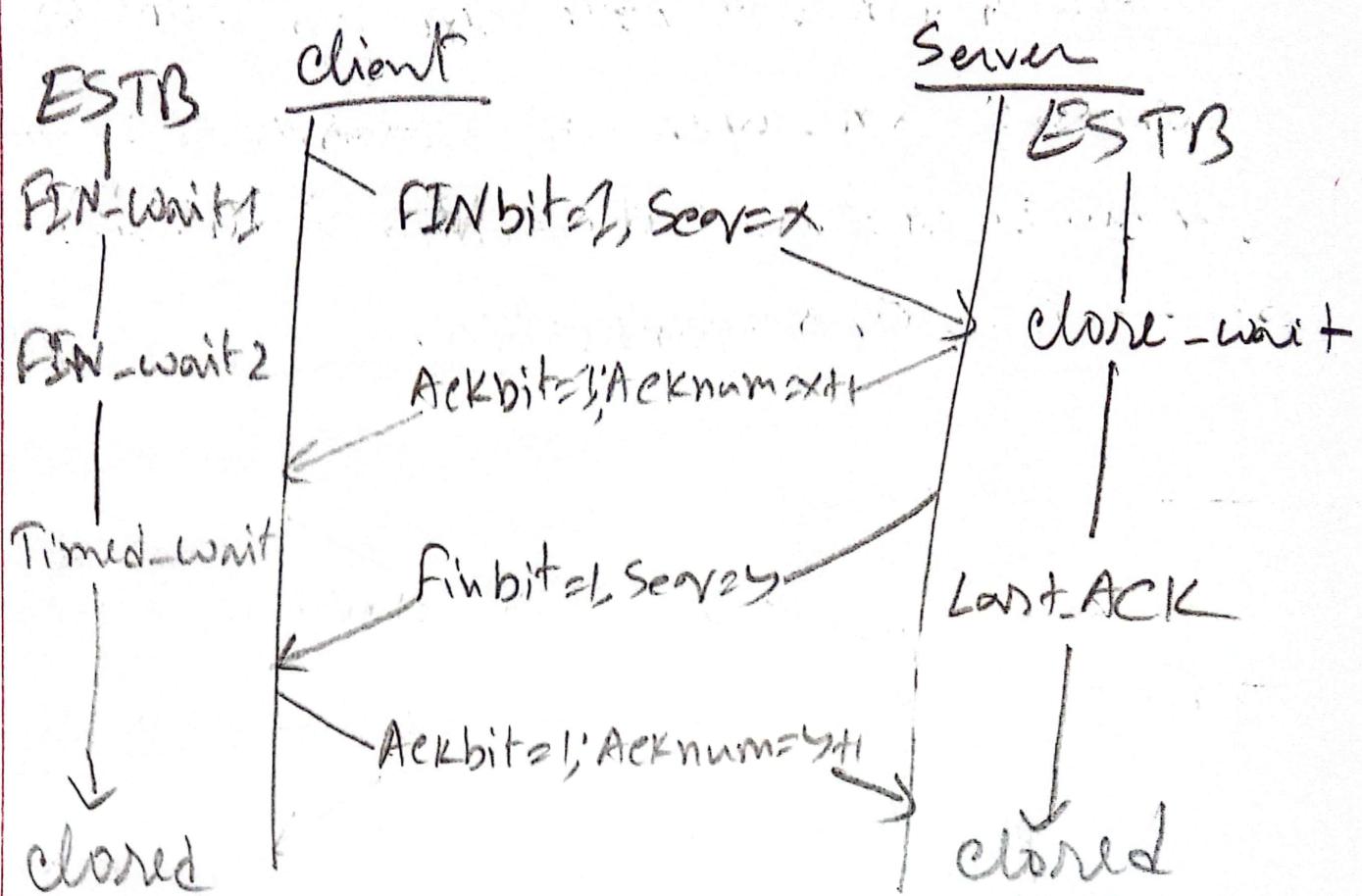


- ① Setting connection by a SYN message.
- ② Assigning numbering segments by a random sequence number.
- ③ The request for expected new segment by positive acknowledgement (cumulative ACK).

→ Tcp 3-way Handshake (Session establishment).



## #TCP Session closing:-



## #TCP Sequence numbering:-

If it is a byte stream number of the first byte of a segment's data part.

These sequence numbers are used to track and manage data segments within a TCP connection to maintain data integrity and correctness.

#Acknowledgement Number: This is a sequence number of the next byte expected from other side. It ensures a reliable data transmission, maintains sequence order and ~~control~~ supports flow control in TCP connection.

### #Flow control:

This is a mechanism that ensures the sender does not overwhelm the receiver by adapting the rate of data transmission based on receiver capacity.

- Receiver will "advertise" free buffer space by including a value of free buffer variable "window".
- If the sender's data transmission rate is too high or receiver's buffer is close to full, the receiver will reduce the advertised window size in TCP acknowledgement. And will signal to slow down.
- If the receiver has more space in buffer and can handle more data, it will increase

the advertised window size to allow sender for transmitting more data.

→ This dynamic window size adjustment ensures comfortable data transmission rate by a receiver, avoiding congestion and potential data loss.

### # IP address class:

<u>Class</u>	<u>Range</u>
A	1.0.0.0 to 126.0.0.0
B	128.0.0.0 to 191.255.0.0
C	192.0.0.0 to 223.255.255.0
(multicast) D	224.0.0.0 to 239.255.255.255
(experimental) E	240.0.0.0 to 255.255.255.255

### # IPv4 address:

Logical name "wellington.cs.virginia.edu"

the IPv4 address is 128.143.132.144

IPv4 class B address: 128.143.0.0

host identification: 138.144

network mask : 255.255.0.0

→ prefix notation : 128.143.138.144/16

→ network prefix of a class B network is 16 bits long.

→ Residual 16 bits defines the address part used for host identification.

## # Multiplexing/Demultiplexing:

→ Multiplexing (at sender):

It is a process of handling data from multiple sockets, add transport header of segment, which, later used for demultiplexing on receiver side.

→ De-multiplexing (at receiver):

It is process to deliver received segments to the correct socket in an application program, using header information.

## # Connectionless Demultiplexing:

In this process, network protocols like UDP, direct incoming data packets to the correct application on the receiver side based on the destination port in the packet header. This allows multiple application to share network connection simultaneously without pre-established connection.

## # Port Numbers in TCP:

A port number identifies the endpoint of a TCP transport connection at the sending or receiving host of a traffic relationship. Some well known ports:

FTP - 21	Finger - 79	NNTP - 119
Telnet - 23	HTTP - 80	
SMTP - 25	POP - 110	
TFTP - 69		

## # Name Resolution:

### → HTTP Name resolution:

TCP does not operate with hostnames, that's why name must be translated into a IPv4 / IPv6 address.

### → DNS name resolution:

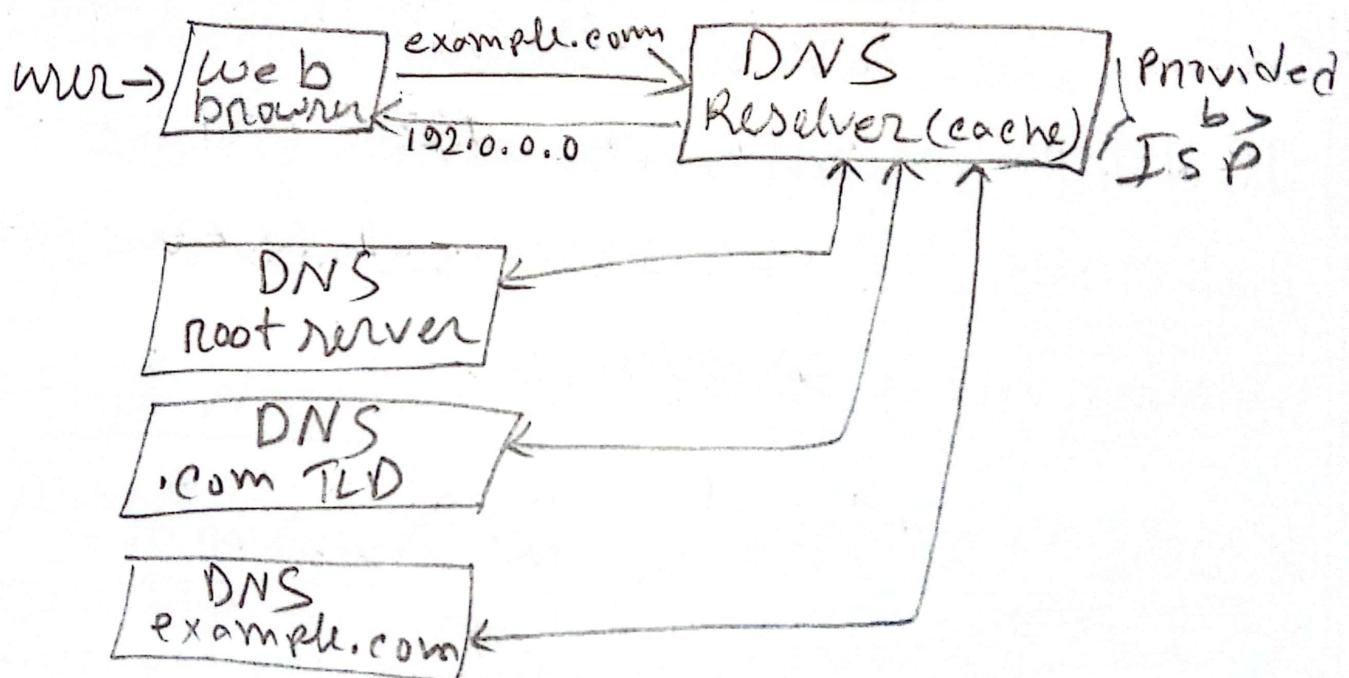
Name resolution of a logical host name to an IP address performed by sending a request to a distributed database system, which called Domain Name System.

# DNS: It is a critical system that helps to translate human readable domain name to IP address to locate resources on the internet.

→ Domain: A human readable name that represents a group of devices/resources on the internet. It is organized hierarchically with sub-domains branching off the main domain.

→ DNS Zone: A collection of DNS records for a specific domain or subdomain in DNS namespace. Zones are used for administrative control and management of DNS settings.

## # DNS working mechanism:



\* It is also the procedure of dns lookup.

## # DNS Types:-

A → IPv4 address record, (ID = 1)

AAAA → IPv6 address record, (ID = 28)

NS = Name Server (id=2)

CNAME = Canonical name (id=5)

DNSKEY = DNS key resource record (id=48)

MX = Mail exchange record (id=15)

URI = uniform resource Identifier (id=255)

## # ARP (Address resolution protocol) :

It is a network protocol that helps devices to find the physical (MAC) address associated to the known IP address within the same network (local). It also allows devices to communicate on the same network segment by mapping ip address to MAC address through ARP request and responses.

## //Chapter: 2 //

#Cross layer communication of a client Server:

/Noted in Pg/ 2.1 slide-23

# IPv4 packet header :

Version (4 bit)	IHL (4 bit)	Type of Service (8 bit)	Total length (16 bit)
Identification (16 bit)		Flags (3 bit)	Fragment offset (13 bit)
Time to live (8 bit)	Protocol (8 bit)	Header checksum (16 bit)	
Source Address (32 bit)			
Destination Address (32 bit)			
Options (variable)		padding (variable)	

## # Private ip addresses;

Private ip address are reserved for use within private network and not routable on the public internet;

## IPv4 private ip addresses:-

→ class A: 10.0.0.0 to 10.255.255.255

→ class B: 172.16.0.0 to 172.31.255.255

→ class C: 192.168.0.0 to 192.168.255.255

## IPv6 private ip address prefix:

→ fc00::/8

→ fe80::/10

## #IPv6 Networking:

It is a 128 bit address system for large scale use. This address consists of 8 segments, every segment is 16 bit.

Ex: 9BF5:0000:0000:0000:BA5F:039A:000A:2126

Type prefix

Rest of address

→ IPv4 Embedding address:

8 bits	88 bits	32 bits
00000000	All 0s	IPv4 address

→ Header format:

4 bit	8 bit	20 bit
version	Priority	Flow label
Payload length	next Header	Hop limit

Source address

Destination address

## //KTR - Important Links for Lab //

# VyOS Docker hub → unibaktr  
for init.vbnsh bonie config and lab  
config instruction.

# VyOS interface command → for init files  
# VyOS next-hop ~~command~~<sup>interface</sup> → for static route  
in init.vbnsh

# VyOS dhcp command → for dhcp commands.