

0-1 Knapsack

$$W=8$$

$$n=4$$

$$P = \{1, 2, 5, 6\}$$

$$\omega = \{2, 3, 4, 5\}$$

$$\sum \omega_i x_i \leq W$$

$$x_i = \{1/0, 1/0, 1/0, 1/0\}$$

constraint

$$\downarrow 2^n$$

Maximize Profit

$$\max \sum P_i x_i$$

Memoization → Recursion

→ Storing the recursive values

$$\text{Net Pending Profit} = 8 - 6 = 2$$

Tabulation → No Recursion

$$W=8$$



$$2-2=0$$

$$W \longrightarrow$$



		0	1	2	3	4	5	6	7	8
P _i	0	0	0	0	0	0	0	0	0	0
ω _i	0	0	0	0	0	0	0	0	0	0
Obj1	1	2	1	0	0	1	1	1	1	1
Obj2	2	3	2	0	0	1	2	3	3	3
Obj3	5	4	3	0	0	1	2	5	6	7
Obj4	6	5	4	0	0	1	2	5	6	7
x ₁	2	6	1	0	1	1	1	1	1	1
x ₂	1	0	1	0	1	1	1	1	1	1
x ₃	0	1	0	1	1	1	1	1	1	1
x ₄	1	1	1	1	1	1	1	1	1	1

final Result ← Max Profit

Generalized formula

→ V[i, W] = max { V[i-1, k], V[i-1, W - ω_i] + P_i }

including current object

not including current object

$$V(4,1) = \max \{ \underline{V(3,1)}, V(3,1-s) + 6 \}$$

$$\max \{ 0, \underline{V(3,-4)} + 6 \}$$

$$5 \quad s \\ W - \omega t(i)$$

$$V(4,s) = \max \{ \underline{V(3,s)}, V(3,0) + 6 \}$$

$$= \underline{6}$$

$$\omega t(i-1) > w$$

\hookrightarrow skip that $i=0$

problem 2

$N=3, W=4$ object \hookrightarrow filled value

with 0

$$P = \{ 1, 2, 3 \}$$

$$\uparrow \{ 0, 1, 2 \}$$

$$\omega t = \{ 4, 5, 1 \}$$

$$\uparrow 6$$

$$i=1$$

\hookrightarrow picking values
of $P \& \omega t$.

$$x_1 \quad x_2 \quad x_3$$

$$\max P_{\text{Profit}} =$$

$$V(3,4) = \underline{\underline{\underline{}}$$

$$0 \quad 0 \quad 1$$

$$W \rightarrow$$

$$\downarrow$$

$$\text{Net Pending} \\ \text{Profit} =$$

$$3 - 3 = 0$$

P	ωt	0	0	1	2	3	4
0	0	0	0	0	0	0	0
1	4	1	0	0	0	0	1
2	5	2	0	0	0	0	1
3	1	3	0	3	3	3	3

$$V(1,3) = \max \{ \underline{V(0,3)}, V(0,-1) + 1 \}$$

$$V(1, 4) = \max\{ V(0, 4), V(0, 0) + 1 \}$$

$$= \max\{ 0, 1 \}$$

$$V(2, 4) = \max\left\{ \frac{V(1, 4)}{1}, V(1, -1) + 2 \right\}$$

$$= 1$$

$$V(3, 4) = \max\{ V(2, 4), V(2, 3) + 3 \}$$

$$= \max\{ 1, 0 + 3 \}$$

$$= 3$$

→ Display the object names which contribute to maximum profit.

$$\begin{cases} \text{Time complexity} \rightarrow \Theta(m * k) \\ \text{Space complexity} \rightarrow \Theta(m * k) \end{cases}$$

Brute force approach

→ $\Theta(2^m)$ → Exponential Time Complexity

→ $\Theta(m * k)$

→ Polynomial Time Complexity

Matrix Chain Multiplication

$$M_1 = \begin{bmatrix} \overrightarrow{a_{11} \ a_{12} \ a_{13}} \\ a_{21} \ a_{22} \ a_{23} \end{bmatrix}_{2 \times 3} \quad M_2 = \begin{bmatrix} b_{11} & b_{12} \\ b_{21} & b_{22} \\ b_{31} & b_{32} \end{bmatrix}_{3 \times 2}$$

\downarrow

$$\begin{array}{c}
 \begin{array}{cc}
 \underline{2 \times 3} & \underline{3 \times 2}
 \end{array} \\
 M_1 * M_2 = \frac{\text{Output Matrix}}{2 \times 2} \quad \text{Total } \frac{a * b * c}{2 * 3 * 2} = 12 \\
 \text{Num of multiplication} = \frac{12}{(3 * 4)}
 \end{array}$$

$$\begin{array}{c}
 0 \left[\begin{array}{c}
 \underline{a_{11} * b_{11} + a_{12} * b_{21} + a_{13} * b_{31}} \\
 \underline{a_{11} * b_{21} +} \\
 \underline{a_{12} * b_{22} + a_{13} * b_{32}}
 \end{array} \right] \\
 1 \left[\begin{array}{c}
 \underline{a_{21} * b_{11} + a_{22} * b_{21} + a_{23} * b_{31}} \\
 \underline{a_{21} * b_{21} +} \\
 \underline{a_{22} * b_{22} +} \\
 \underline{a_{23} * b_{32}}
 \end{array} \right]
 \end{array}$$

$$M_1 = 2 \times 3$$

$$M_2 = 3 \times 4$$

$$M_3 = 4 \times 2$$

$$P_2 \quad P_3$$

↳ Matrix multiplication

(Associative)

Property

$$M_1 * M_2 * M_3$$

I

$$(M_1 * M_2) * M_3$$

II

$$M_1 * (M_2 * M_3)$$

Min number

of multip.

$$\left. \begin{array}{l} M_1 = 2 \times 3 \\ M_2 = 3 \times 4 \end{array} \right\} = 2 \times 3 \times 4 = 24 \quad (\xrightarrow{2 \times 4})$$

$$\left. \begin{array}{l} O_1 = 2 \times 4 \\ M_3 = 4 \times 2 \end{array} \right\} = \frac{2 \times 4 \times 2}{P_0 \quad P_2 \quad P_3} = 16$$

$$\begin{array}{rcl} \text{Total num of} & & \\ \hline \text{multiplications} & = 24 + 16 \\ & \hline & = 40 \end{array}$$

$$\left. \begin{array}{l} M_2 = 3 \times 4 \\ M_3 = 4 \times 2 \end{array} \right\} = 3 \times 4 \times 2 = 24$$

$$O_1 = M_2 * M_3 \quad (3 \times 2)$$

$$M_1 = 2 \times 3$$

$$\begin{array}{rcl} O_1 = 3 \times 2 \\ P_0 \quad P_1 \quad P_3 \\ 2 \times 3 \times 2 \end{array}$$

$$= 12$$

$$24 + 12 = 36$$

☞ commutative

$$M_1 * O_1 \neq O_1 * M_1$$

Property

$M_1 \ M_2 \ M_3 \ - \ - \ - \ M_{10}$

↳ Generalized formula

$$C(1,3) = \min_{i \leq k \leq j} \begin{cases} C(1,2) + C(3,3) + p_{i-1} p_k p_j \\ C(1,1) + C(2,3) + p_0 p_1 p_3 \end{cases}$$

$$k=1, \ k=2, \ k=3$$

$$C(1,4) = \min \left\{ \begin{array}{l} \dots \\ \dots \end{array} \right\}$$

Task

$$A = 1 \times 3, B = 3 \times 1, C = 1 \times 2, D = 2 \times 3$$

$$E = ABCD$$

1) How many min number of multiplications required to get an output matrix $E ??$

2) In what sequence ??

3) Try to do same thing using above generalization & verify that in both ways answers are same or not??