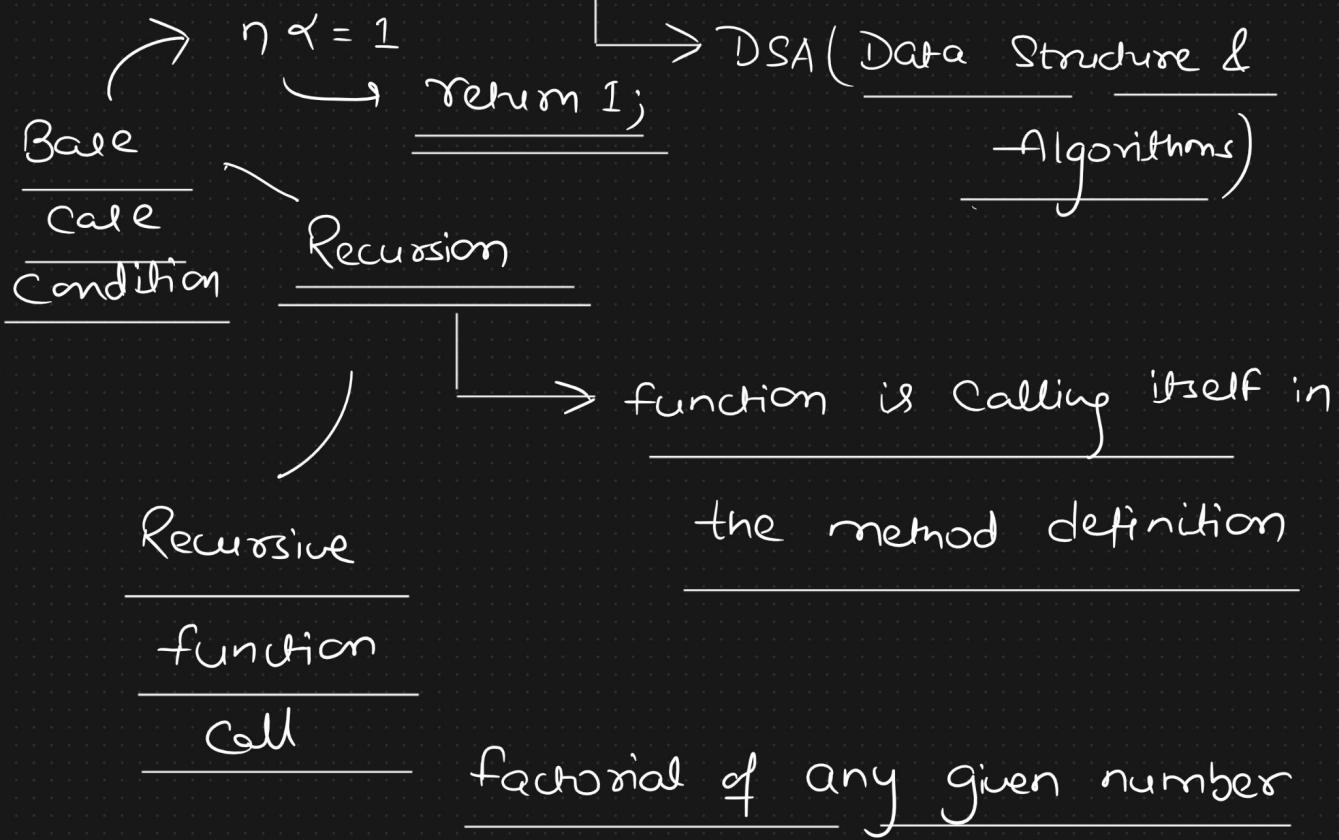


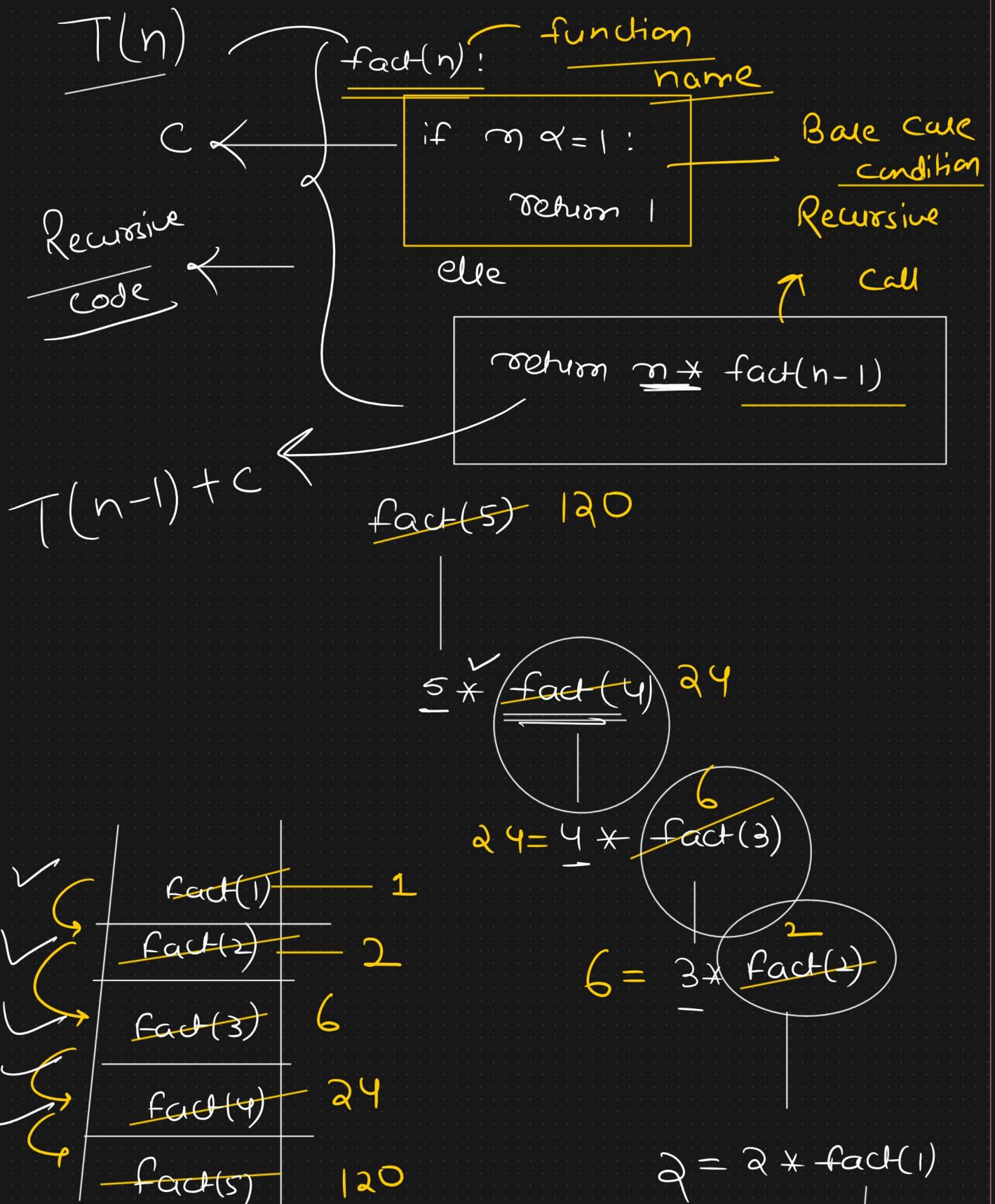
Recursion in Python



$$\begin{aligned}5! &= 5 \times 4 \times 3 \times 2 \times 1 \\&= 120\end{aligned}$$

$$\left\{ \begin{array}{l} 0! = 1 \\ 1! = 1 \end{array} \right.$$

$$\begin{aligned}5! &= 5 \times 4! \\&\quad \downarrow 4 \times 3! \\&\quad \downarrow 3 \times 2!\end{aligned}$$



Stack Data Structure

\hookrightarrow LIFO (Last In First Out)

Note:

Every Recursive code requires
extra Stack Space to
store the function
calls.

Recursion code



Recurrence Relation

Substitution

Master's
Theorem

Recursive
Tree

$$T(n) = T(n-1) + c$$

$$T(n) = \mathcal{O}(n)$$

Linear
Time
Complexity

Stack Space

Space Complexity = $O(n)$

Fibonacci Series

0	1	2	3	4	5	6	7
0	1	1	2	3	5	8	13



$T(n)$

$\text{fib}(n)$:

if $n \neq 1$



Bare case

condition

c \leftarrow

return n

else



Recursive

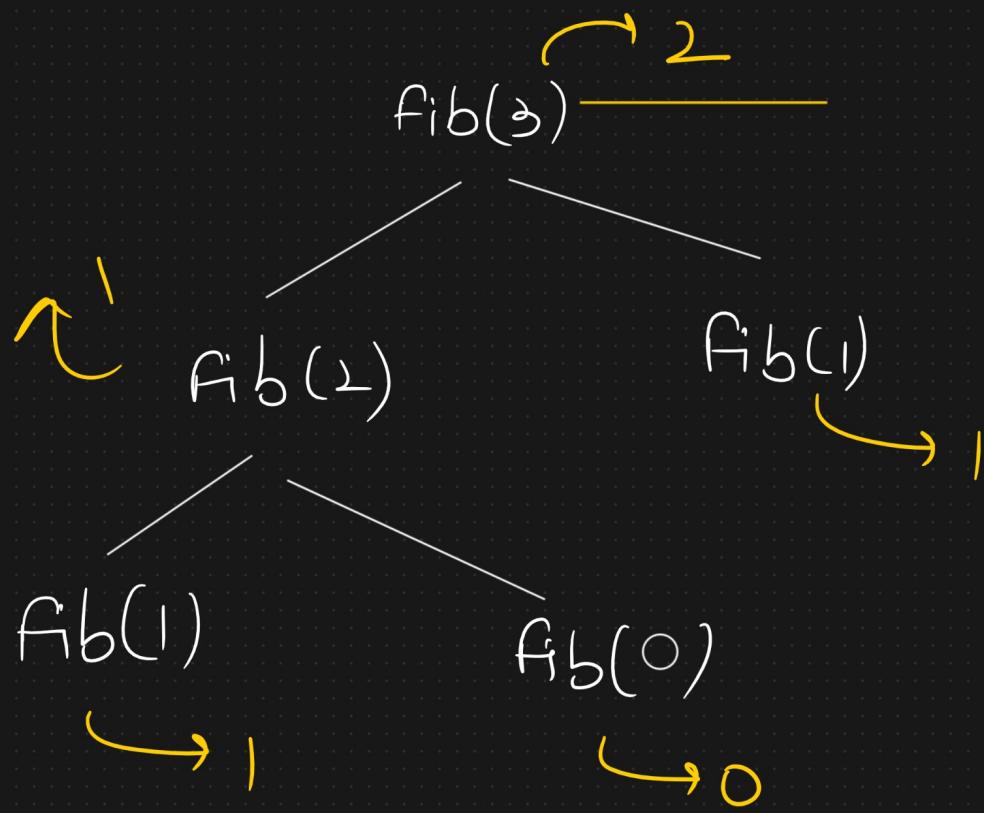
function

call

return $\text{fib}(n-1) +$

$\text{fib}(n-2)$

$T(n-1) + T(n-2)$



$$T(n) = T(n-1) + T(n-2) + c$$

Recursive tree

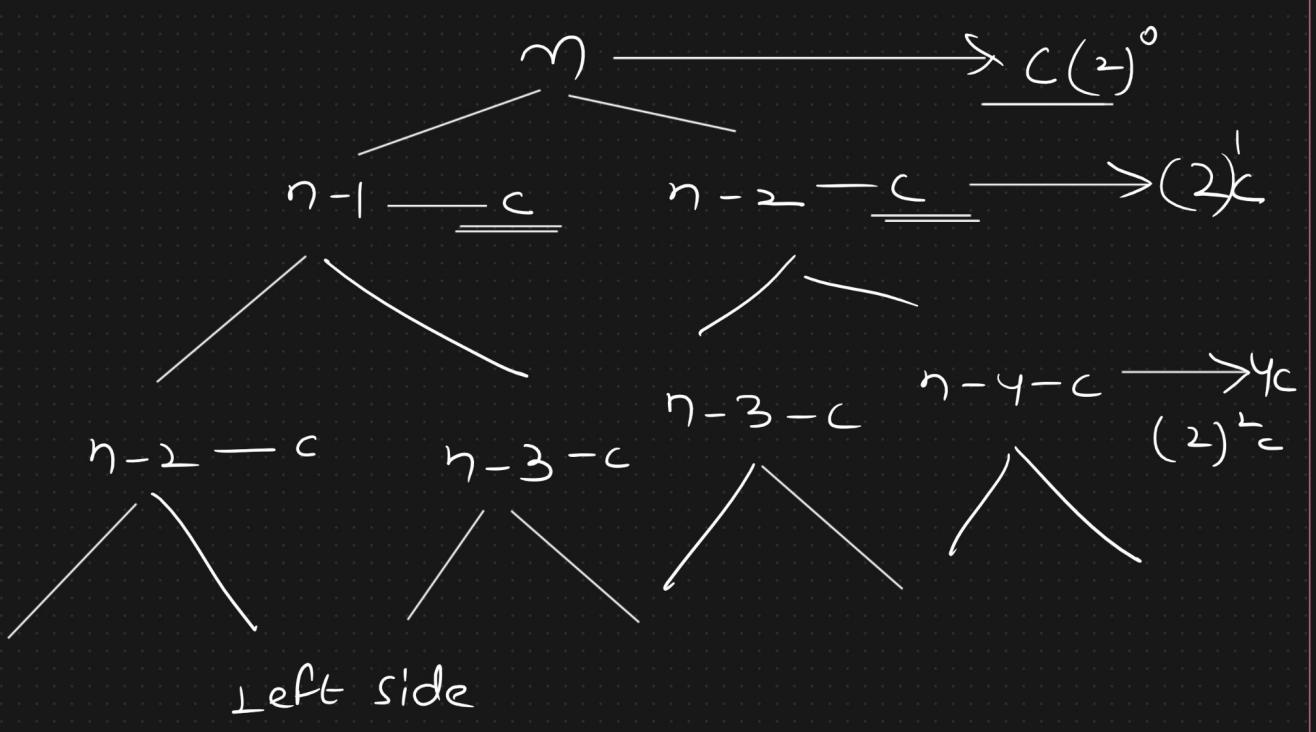
approach

Dynamic
Programming

$O(2^n)$

Drawback of
using
recursion

Exponential
time
complexity



$$n - k = 1$$

$$\frac{k = (n+1)}{}$$

Level

$$c \left[2^0 + 2^1 + 2^2 + 2^3 + \dots + 2^k \right]$$

G.P Series

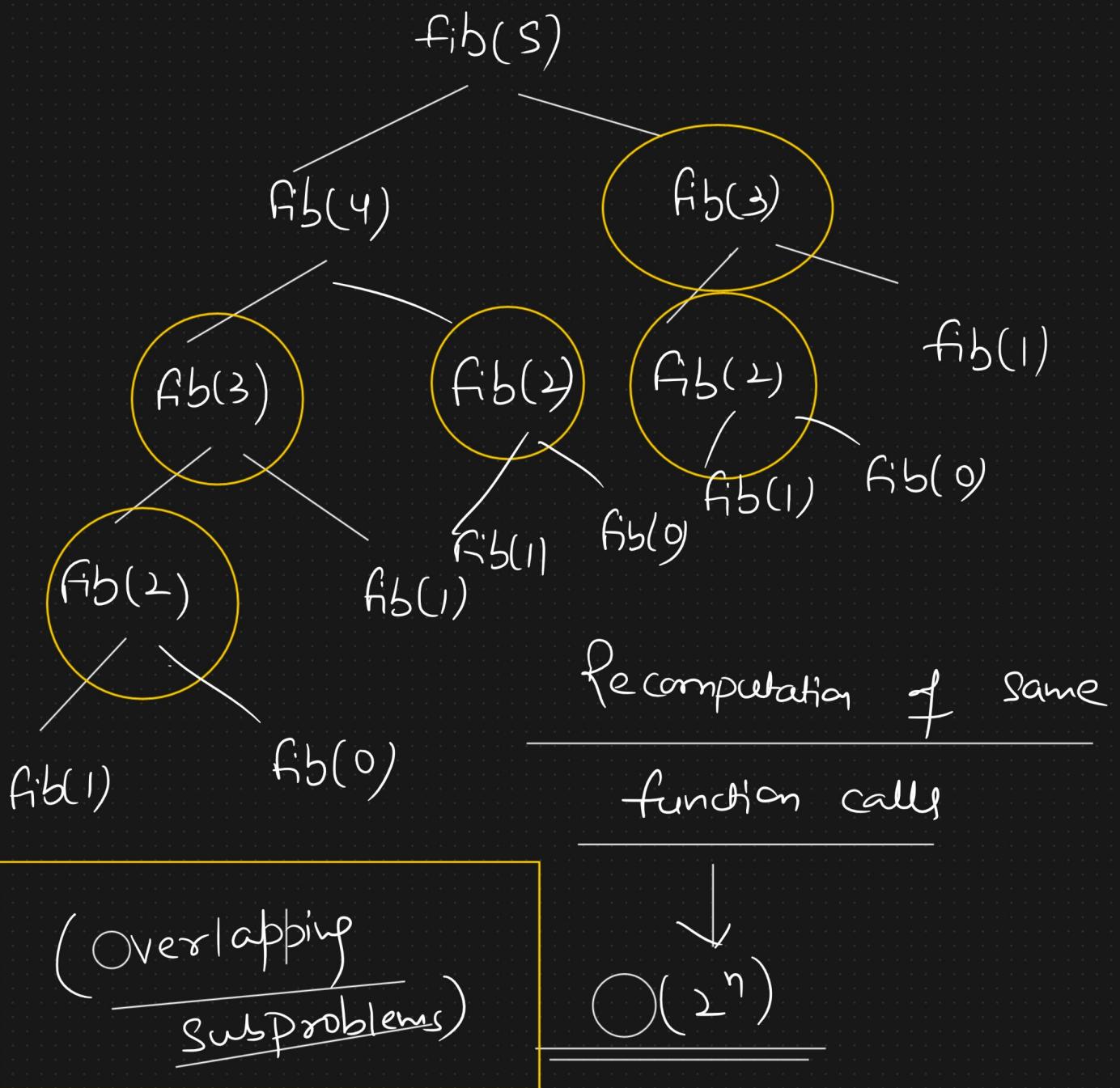
$$a = 1$$

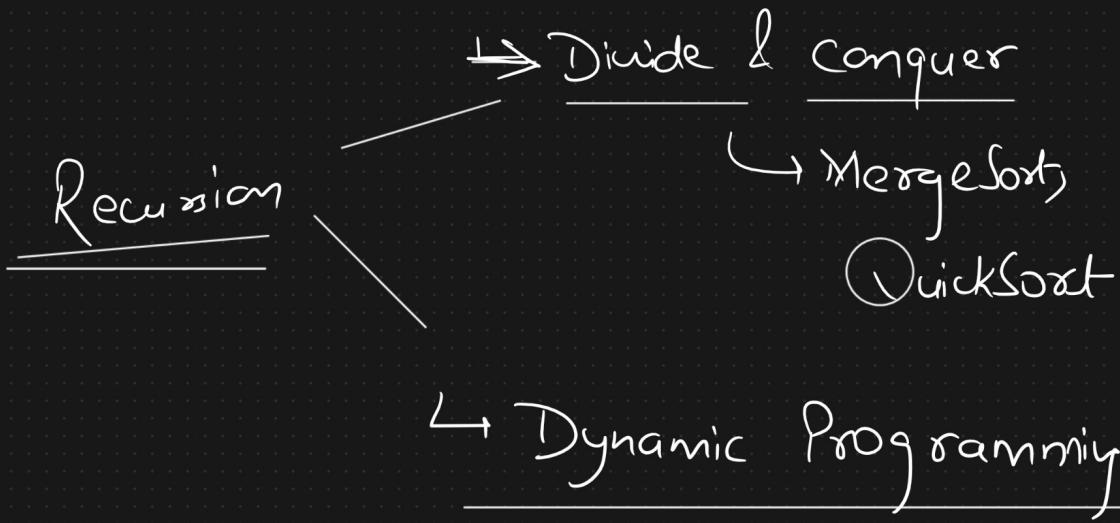
$$\frac{\tau = 2}{}$$

$$\tau = 2$$

$$\frac{\tau > 1 \rightarrow a (\tau^n - 1)}{\tau - 1}$$

$$\frac{1(2^n - 1)}{1} = O(2^n)$$





Sum of Given number

$\hookrightarrow \text{num} = 1234$

$$\begin{array}{r} \\ \hookrightarrow \\ 1+2+3+4=10 \end{array}$$

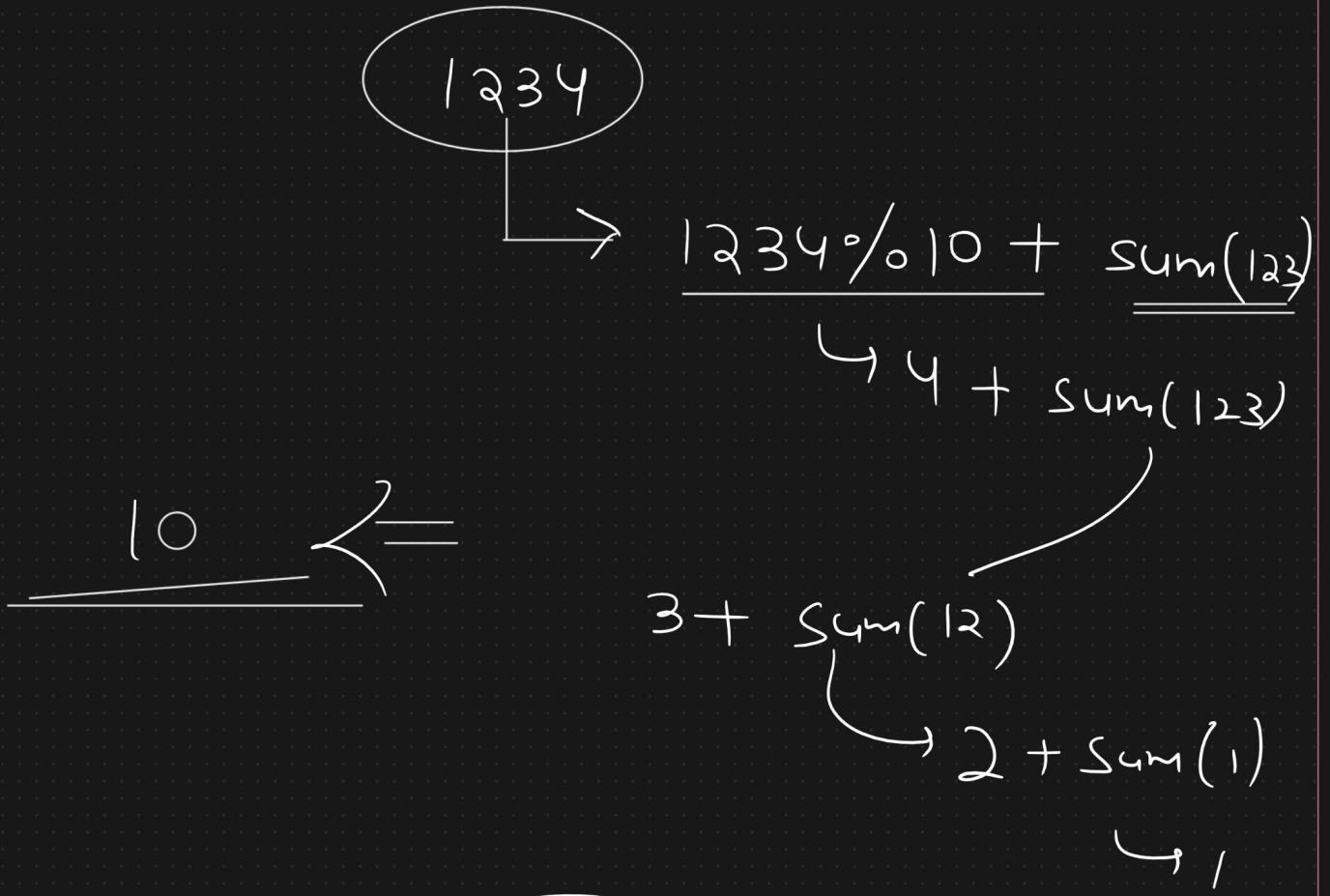
① Base case condition

$\text{sum}(n)$: if $n \leq 10$:
 $\hookrightarrow \text{return } n$

else:

$$n \% 10 + \text{sum}(n / 10)$$

4 remainder



	<u>Recursive</u>	<u>Iterative</u>
①	Bare case condition	for, while
②	Recursive function call	
③	Lesser lines of Code → Clarity	Higher num of lines of Code

③ Stack space
is required

No need of extra
stack space