

Data Structures Using C

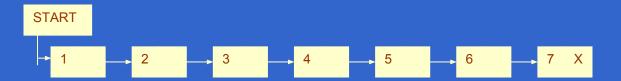
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CHAPTER 8

<u>INTRODUCTION</u>

- A linked list in simple terms is a linear collection of data elements. These data elements are called nodes.
- Linked list is a data structure which in turn can be used to implement other data structures. Thus, it acts as building block to implement data structures like stacks, queues and their variations.
- A linked list can be perceived as a train or a sequence of nodes in which each node contain one or more data fields and a pointer to the next node.

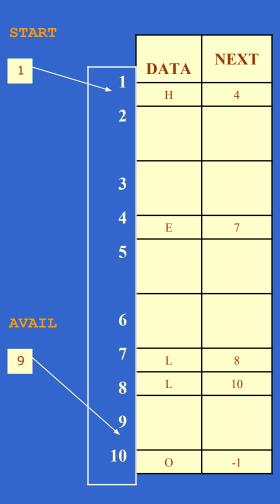


In the above linked list, every node contains two parts- one integer and the other a pointer to the next node. The left part of the node which contains data may include a simple data type, an array or a structure. The right part of the node contains a pointer to the next node (or address of the next node in sequence). The last node will have no next node connected to it, so it will store a special value called NULL.

INTRODUCTION contd.

- Linked list contains a pointer variable, START which stores the address of the first node in the list.
- We can traverse the entire list using a single pointer variable START.
 The START node will contain the address of the first node; the next part
 of the first node will in turn store the address of its succeeding node.
- Using this technique the individual nodes of the list will form a chain of nodes. If START = NULL, this means that the linked list is empty and contains no nodes.
- In C, we will implement a linked list using the following code:

```
struct node{int data;struct node *next;};
```



START pointing to the first element of the linked list in memory If we want to add a node to an already existing linked list in the memory, we will first find any free space in the memory and then use it to store the information.

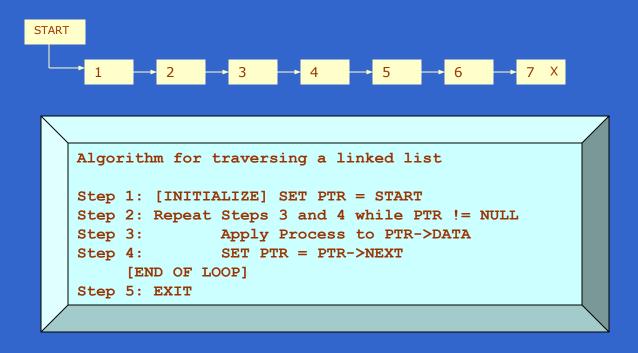
The operating system maintains a free pool which is a linked list of a of all free memory cells. It maintains a pointer variable AVAIL which stores the address of the first free space.

When you delete a node from the linked list, the operating system adds the freed memory to the free pool.

The operating system will perform this operation whenever it finds the CPU idle or whenever the programs are falling short of memory. The operating system scans through all the memory cells and mark the cells that are being used by some or the other program. Then, it collects all those cells which are not being used and add their address to the free pool so that it can be reused by the programs. This process is called garbage collection. The whole process of collecting unused memory cells (garbage collection) is transparent to the programmer.

Singly Linked List

 A singly linked list is the simplest type of linked list in which every node contains some data and a pointer to the next node of the same data type. By saying that the node contains a pointer to the next node we mean that the node stores the address of the next node in sequence.



```
Algorithm to print the information stored in each node of the linked list

Step 1: [INITIALIZE] SET PTR = START

Step 2: Repeat Steps 3 and 4 while PTR != NULL

Step 3: Write PTR->DATA

Step 4: SET PTR = PTR->NEXT

[END OF LOOP]

Step 5: EXIT
```

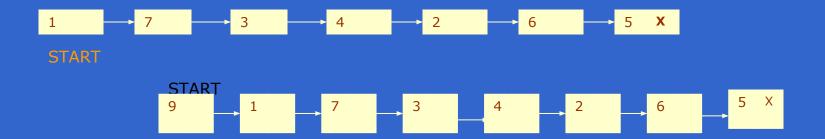
```
Step 1: [INITIALIZE] SET PTR = START
Step 2: Repeat Steps 3 while PTR != NULL
Step 3:
              IF VAL = PTR->DATA
               SET POS = PTR
              Go To Step 5
          ELSE
               SET PTR = PTR->NEXT
          [END OF IF]
     [END OF LOOP]
Step 4: SET POS = NULL
Step 5: EXIT
                                 2 -
   PTR
                                 2
           PTR
                                2
                                       6 -
```

 $7 \longrightarrow 3 \longrightarrow 4 \longrightarrow 2 \longrightarrow 6 \longrightarrow 5 X$

PTR

PTR

Algorithm to search an unsorted linked list







PTR

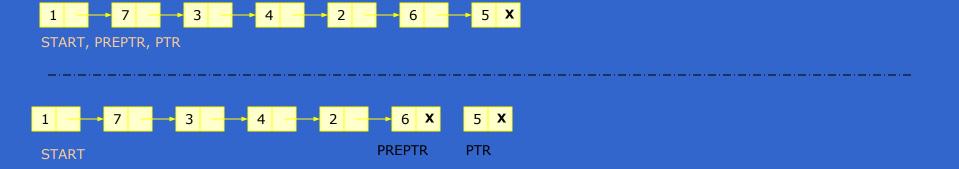
```
Algorithm to insert a new node at the end of the linked list
Step 1: IF AVAIL = NULL, then
          Write OVERFLOW
          Go to Step 10
     [END OF IF]
Step 2: SET New Node = AVAIL
Step 3: SET AVAIL = AVAIL->NEXT
Step 4: SET New Node->DATA = VAL
Step 5: SET New Node->Next = NULL
Step 6: SET PTR = START
Step 7: Repeat Step 8 while PTR->NEXT != NULL
Step 8: SET PTR = PTR ->NEXT
     [END OF LOOP]
Step 9: SET PTR->NEXT = New Node
Step 10: EXIT
```

```
Algorithm to insert a new node after a node that has value NUM
          Step 1: IF AVAIL = NULL, then
                    Write OVERFLOW
                    Go to Step 12
               [END OF IF]
          Step 2: SET New Node = AVAIL
          Step 3: SET AVAIL = AVAIL->NEXT
          Step 4: SET New Node->DATA = VAL
          Step 5: SET PTR = START
          Step 6: SET PREPTR = PTR
          Step 7: Repeat Step 8 and 9 while PREPTR->DATA != NUM
                         SET PREPTR = PTR
          Step 8:
                  SET PTR = PTR->NEXT
          Step 9:
               [END OF LOOP]
          Step 10: PREPTR->NEXT = New Node
          Step 11: SET New Node->NEXT = PTR
          Step 12: EXIT
             3 4 2 5 X
START, PTR, PREPTR
                                                         PREPTR
                                                                   PTR
                                             START
                                4 - 2 -
                       9 –
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```

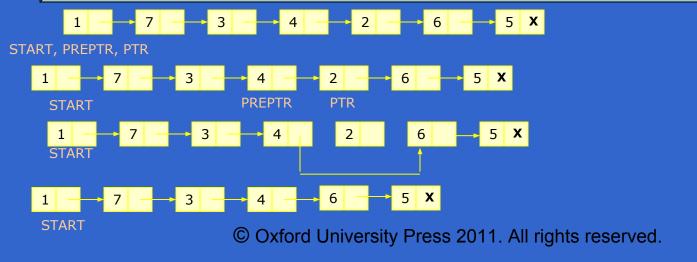
START





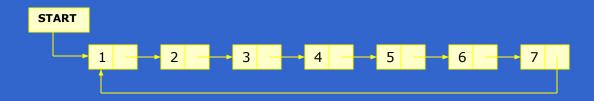


```
Algorithm to delete the node after a given node from the linked list
Step 1: IF START = NULL, then
         Write UNDERFLOW
         Go to Step 10
     [END OF IF]
Step 2: SET PTR = START
Step 3: SET PREPTR = PTR
Step 4: Repeat Step 5 and 6 while PRETR->DATA != NUM
Step 5:
              SET PREPTR = PTR
Step 6:
              SET PTR = PTR->NEXT
     [END OF LOOP]
Step7: SET TEMP = PTR->NEXT
Step 8: SET PREPTR->NEXT = TEMP->NEXT
Step 9: FREE TEMP
Step 10: EXIT
```

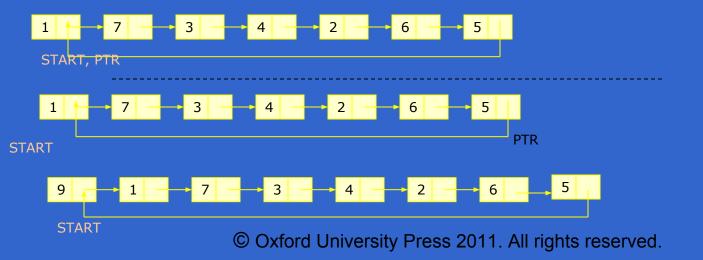


Circular Linked List

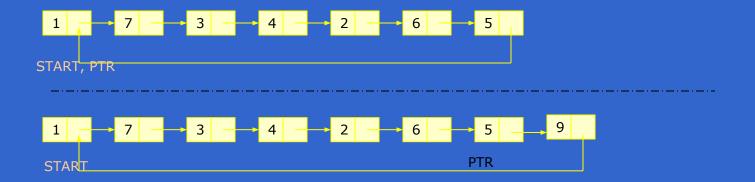
In a circular linked list, the last node contains a pointer to the
first node of the list. We can have a circular singly listed list as
well as circular doubly linked list. While traversing a circular
linked list, we can begin at any node and traverse the list in
any direction forward or backward until we reach the same
node where we had started. Thus, a circular linked list has no
beginning and no ending.



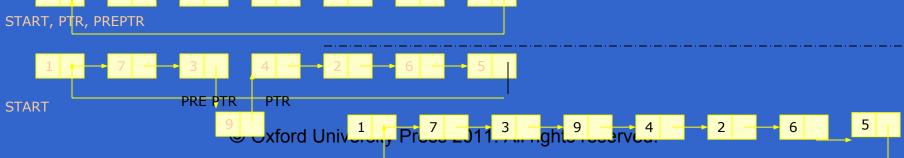
```
Algorithm to insert a new node in the beginning of circular the linked list
Step 1: IF AVAIL = NULL, then
         Write OVERFLOW
         Go to Step 7
     [END OF IF]
Step 2: SET New Node = AVAIL
Step 3: SET AVAIL = AVAIL->NEXT
Step 4: SET New Node->DATA = VAL
Step 5: SET PTR = START
Step 6: Repeat Step 7 while PTR->NEXT != START
Step 7:
              PTR = PTR->NEXT
Step 8: SET New Node->Next = START
Step 8: SET PTR->NEXT = New Node
Step 6: SET START = New Node
Step 7: EXIT
```



```
Algorithm to insert a new node at the end of the circular linked list
Step 1: IF AVAIL = NULL, then
          Write OVERFLOW
          Go to Step 7
     [END OF IF]
Step 2: SET New Node = AVAIL
Step 3: SET AVAIL = AVAIL->NEXT
Step 4: SET New Node->DATA = VAL
Step 5: SET New Node->Next = START
Step 6: SET PTR = START
Step 7: Repeat Step 8 while PTR->NEXT != START
Step 8:
               SET PTR = PTR ->NEXT
     [END OF LOOP]
Step 9: SET PTR ->NEXT = New Node
Step 10: EXIT
```

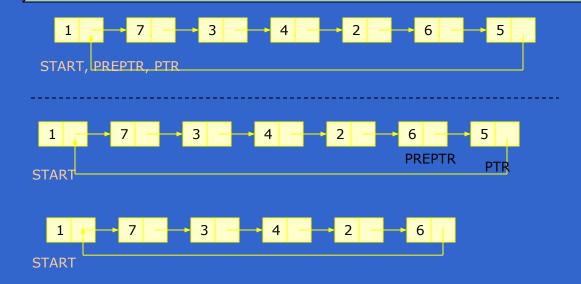


```
Algorithm to insert a new node after a node that has value NUM
Step 1: IF AVAIL = NULL, then
         Write OVERFLOW
         Go to Step 12
     [END OF IF]
Step 2: SET New Node = AVAIL
Step 3: SET AVAIL = AVAIL->NEXT
Step 4: SET New Node->DATA = VAL
Step 5: SET PTR = START
Step 6: SET PREPTR = PTR
Step 7: Repeat Step 8 and 9 while PTR->DATA != NUM
Step 8:
              SET PREPTR = PTR
Step 9: SET PTR = PTR->NEXT
     [END OF LOOP]
Step 10: PREPTR->NEXT = New Node
Step 11: SET New Node->NEXT = PTR
Step 12: EXIT
```



Algorithm to delete the first node from the circular linked list **Step 1: IF START = NULL, then** Write UNDERFLOW Go to Step 8 [END OF IF] **Step 2: SET PTR = START** Step 3: Repeat Step 4 while PTR->NEXT != START Step 4: **SET PTR = PTR->NEXT** [END OF IF] **Step 5: SET PTR->NEXT = START->NEXT Step 6: FREE START Step 7: SET START = PTR->NEXT** Step 8: EXIT START, PTR START START

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Algorithm to delete the node after a given node from the circular linked list

Step 1: IF START = NULL, then
Write UNDERFLOW
Go to Step 9

[END OF IF]

Step 2: SET PTR = START Step 3: SET PREPTR = PTR

Step 4: Repeat Step 5 and 6 while PREPTR->DATA != NUM

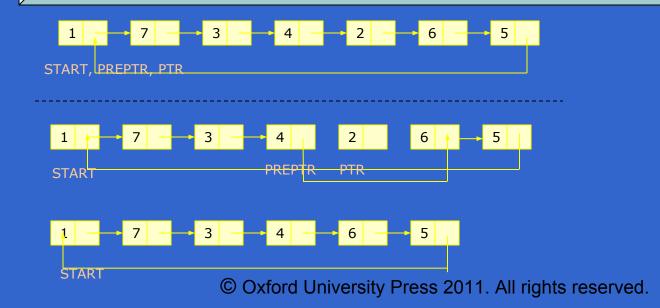
Step 5: SET PREPTR = PTR
Step 6: SET PTR = PTR->NEXT

[END OF LOOP]

Step 7: SET PREPTR->NEXT = PTR->NEXT

Step 8: FREE PTR

Step 9: EXIT



Doubly Linked List

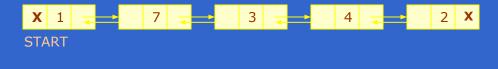
A doubly linked list or a two way linked list is a more complex type of linked list which contains a pointer to the next as well as previous node in the sequence. Therefore, it consists of three parts and not just two. The three parts are data, a pointer to the next

node and a pointer to the previous node

```
In C language, the structure of a doubly linked list is given as, struct node
{ struct node *prev; int data; struct node *next; };

The prov field of the first node and the post field of the last node.
```

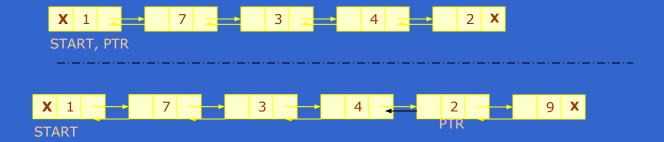
The prev field of the first node and the next field of the last node will contain NULL. The prev field is used to store the address of the preceding node. This would enable to traverse the list in the backward direction as well. Oxford University Press 2011. All rights reserved.



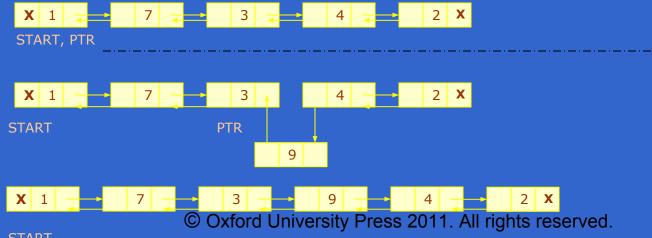


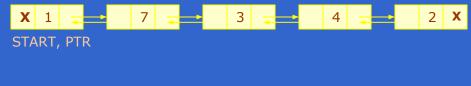
START

```
Algorithm to insert a new node at the end of the doubly linked list
Step 1: IF AVAIL = NULL, then
          Write OVERFLOW
          Go to Step 11
     [END OF IF]
Step 2: SET New Node = AVAIL
Step 3: SET AVAIL = AVAIL->NEXT
Step 4: SET New Node->DATA = VAL
Step 5: SET New Node->Next = NULL
Step 6: SET PTR = START
Step 7: Repeat Step 8 while PTR->NEXT != NULL
Step 8:
               SET PTR = PTR->NEXT
     [END OF LOOP]
Step 9: SET PTR->NEXT = New Node
Step 10: New Node->PREV = PTR
Step 11: EXIT
```

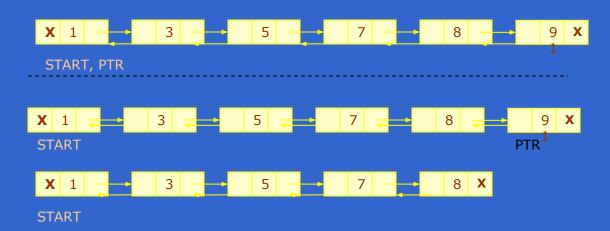


```
Algorithm to insert a new node after a node that has value NUM
Step 1: IF AVAIL = NULL, then
          Write OVERFLOW
          Go to Step 11
     [END OF IF]
Step 2: SET New Node = AVAIL
Step 3: SET AVAIL = AVAIL->NEXT
Step 4: SET New Node->DATA = VAL
Step 5: SET PTR = START
Step 6: Repeat Step 8 while PTR->DATA != NUM
Step 7:
              SET PTR = PTR->NEXT
     [END OF LOOP]
Step 8: New Node->NEXT = PTR->NEXT
Step 9: SET New Node->PREV = PTR
Step 10: SET PTR->NEXT = New Node
Step 11: EXIT
```

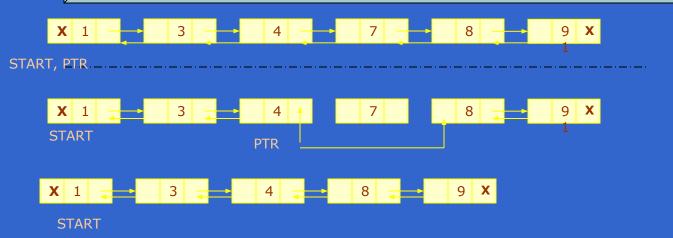








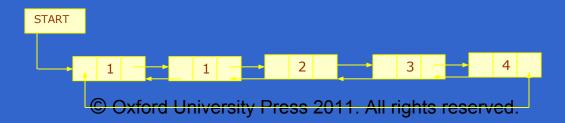
```
Algorithm to delete the node after a given node from the doubly linked list
Step 1: IF START = NULL, then
          Write UNDERFLOW
          Go to Step 9
     [END OF IF]
Step 2: SET PTR = START
Step 3: Repeat Step 4 while PTR->DATA != NUM
Step 4: SET PTR = PTR->NEXT
     [END OF LOOP]
Step 5: SET TEMP = PTR->NEXT
Step 6: SET PTR->NEXT = TEMP->NEXT
Step 7: SET TEMP->NEXT->PREV = PTR
Step 8: FREE TEMP
Step 9: EXIT
```



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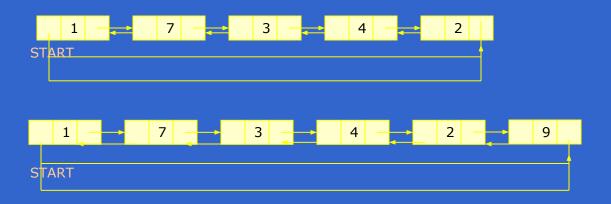
Circular Doubly Linked List

- A circular doubly linked list or a circular two way linked list is a more complex type of linked list which contains a pointer to the next as well as previous node in the sequence.
- The difference between a doubly linked and a circular doubly linked list is same
 as that exists between a singly linked list and a circular linked list. The circular
 doubly linked list does not contain NULL in the previous field of the first node
 and the next field of the last node. Rather, the next field of the last node stores
 the address of the first node of the list, i.e; START. Similarly, the previous field of
 the first field stores the address of the last node.
- Since a circular doubly linked list contains three parts in its structure, it calls for more space per node and for more expensive basic operations. However, it provides the ease to manipulate the elements of the list as it maintains pointers to nodes in both the directions. The main advantage of using a circular doubly linked list is that it makes searches twice as efficient.



```
Algorithm to insert a new node in the beginning of the circular doubly
      linked list
      Step 1: IF AVAIL = NULL, then
               Write OVERFLOW
               Go to Step 12
           [END OF IF]
      Step 2: SET New Node = AVAIL
      Step 3: SET AVAIL = AVAIL->NEXT
      Step 4: SET New_Node->DATA = VAL
      Step 6: SET START->PREV->NEXT = new node;
      Step 7: SET New_Node->PREV = START->PREV;
      Step 8: SET START->PREV= new_Node;
      Step 9: SET new_node->next = START;
      Step 10: SET START = New_Node
      Step 11: EXIT
START
START
```

```
Algorithm to insert a new node at the end of the circular
doubly linked list
Step 1: IF AVAIL = NULL, then
          Write OVERFLOW
          Go to Step 11
     [END OF IF]
Step 2: SET New Node = AVAIL
Step 3: SET AVAIL = AVAIL->NEXT
Step 4: SET New Node->DATA = VAL
Step 5: SET New Node->Next = START
Step 6: SET New Node->PREV = START->PREV
Step 7: EXIT
```



```
Step 1: IF AVAIL = NULL, then
                   Write OVERFLOW
                   Go to Step 11
              [END OF IF]
         Step 2: SET New Node = AVAIL
         Step 3: SET AVAIL = AVAIL->NEXT
         Step 4: SET New Node->DATA = VAL
         Step 5: SET PTR = START
         Step 6: Repeat Step 8 while PTR->DATA != NUM
         Step 7:
                        SET PTR = PTR->NEXT
              [END OF LOOP]
         Step 8: New Node->NEXT = PTR->NEXT
         Step 9: SET PTR->NEXT->PREV = New Node
         Step 9: SET New Node->PREV = PTR
         Step 10: SET PTR->NEXT = New Node
         Step 11: EXIT
    1
START, PTR
                     PTR
                              9
  START
                    © Oxford University Fiess 2011. All rights reserved.
                                     START
```

Algorithm to insert a new node after a node that has value NUM

```
Algorithm to delete the first node from the Circular doubly linked list

Step 1: IF START = NULL, then

Write UNDERFLOW

Go to Step 8

[END OF IF]

Step 2: SET PTR = START

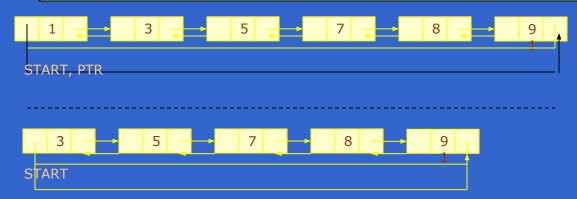
Step 3: SET PTR->PREV=>NEXT= PTR->NEXT

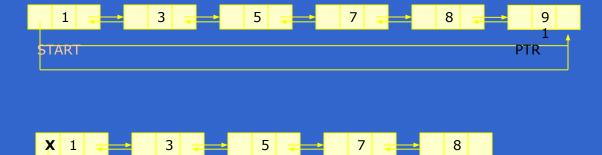
Step 4: SET PTR->NEXT->PREV = PTR->PREV

Step 5: SET START = START->NEXT

Step 6: FREE PTR

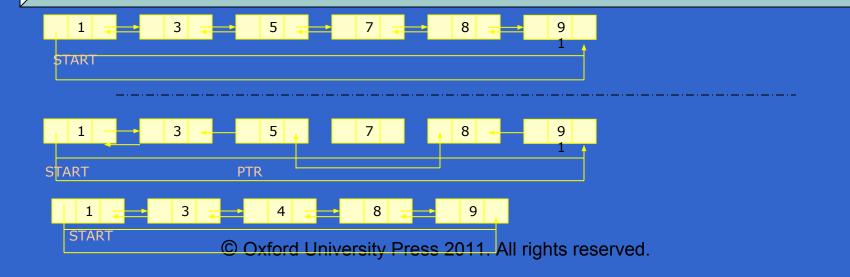
Step 7: EXIT
```





START

```
Algorithm to delete the node after a given node from the circular doubly
linked list
Step 1: IF START = NULL, then
         Write UNDERFLOW
         Go to Step 9
     [END OF IF]
Step 2: SET PTR = START
Step 3: Repeat Step 4 while PTR->DATA != NUM
Step 4:
              SET PTR = PTR->NEXT
     [END OF LOOP]
Step 5: SET TEMP = PTR->NEXT
Step 6: SET PTR->NEXT = TEMP->NEXT
Step 7: SET TEMP->NEXT->PREV = PTR
Step 8: FREE TEMP
Step 9: EXIT
```



Header Linked List

- A header linked list is a special type of linked list which contains a header node at the beginning of the list. So, in a header linked list START will not point to the first node of the list but START will contain the address of the header node. There are basically two variants of a header linked list-
- Grounded header linked list which stores NULL in the next field of the last node
- Circular header linked list which stores the address of the header node in the next field of the last node. Here, the header node will denote the end of the list.


```
Step 5: EXIT
  Algorithm to insert a new node after a given node
  Step 1: IF AVAIL = NULL, then
            Write OVERFLOW
            Go to Step 10
       [END OF IF]
  Step 2: SET New Node = AVAIL
  Step 3: SET AVAIL = AVAIL->NEXT
  Step 4: SET PTR = START->NEXT
  Step 5: SET New Node->DATA = VAL
  Step 6: Repeat step 4 while PTR->DATA != NUM
  Step 7: SET PTR = PTR->NEXT
       [END OF LOOP]
  Step 8: New Node->NEXT = PTR->NEXT
  Step 9: SET PTR->NEXT = New Node
  Step 10: EXIT
```

Algorithm to traverse a Circular Header Linked List

Step 2: Repeat Steps 3 and 4 while PTR != START

Step 3: Apply PROCESS to PTR->DATA

Step 1: SET PTR = START->NEXT

[END OF LOOP]

Step 4: SET PTR = PTR->NEXT