COL:750/7250

Foundations of Automatic Verification

Instructor: Priyanka Golia

Course Webpage

https://priyanka-golia.github.io/teaching/COL-750-COL7250/index.html

LTL Syntax

F = True

= p (atomic proposition)

$= F_1 \wedge F_2, F_1 \vee F_2, F_1 \rightarrow F_2, F_1 \leftrightarrow F_2$

$= \neg F_1$

= **N** F_1 **N** is "Next". F_1 is True at next step. Often represented as **O**, **X** .

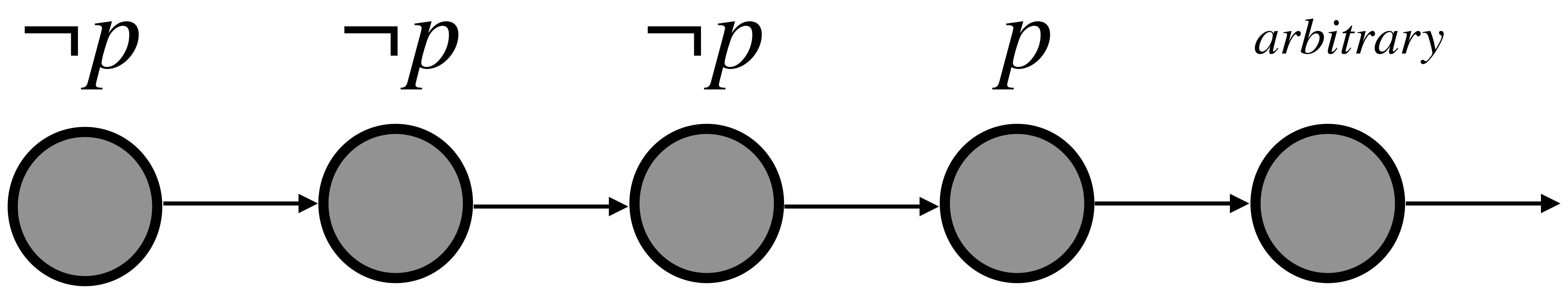
= F_1 **U** F_2 **U** is "Until". F_2 is True at "some point, say t", and until then F_1 is True.
At "t", F_1 doesn't have to hold any more!

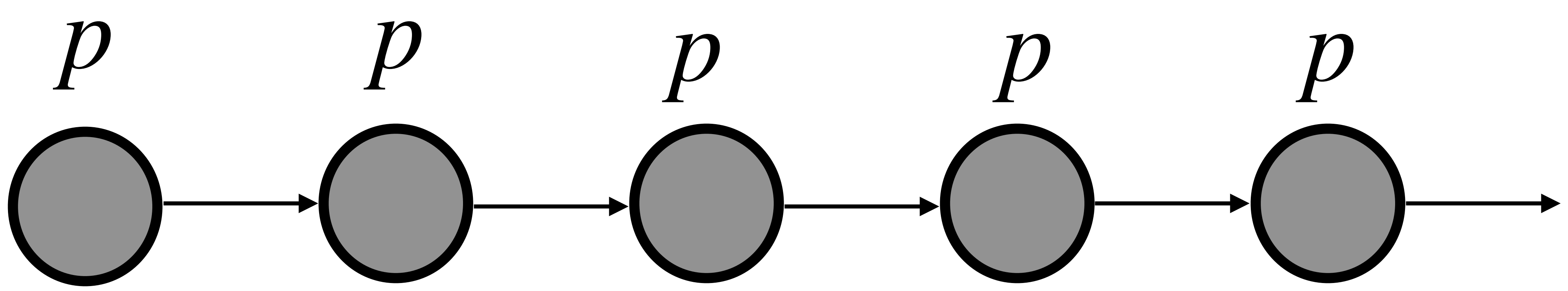
LTL Syntax

Primary temporal operators: **N U**

Eventually $\Diamond$ F F will become true at some point in the future.

$\Diamond F \equiv True\ \mathbf{U}\ F$





Always (valid) $\Box F$ F is always True.

$\Box F \equiv \neg \Diamond \neg F$ (Never (Eventually ($\neg F$))).

LTL Syntax Sequence of states (paths).

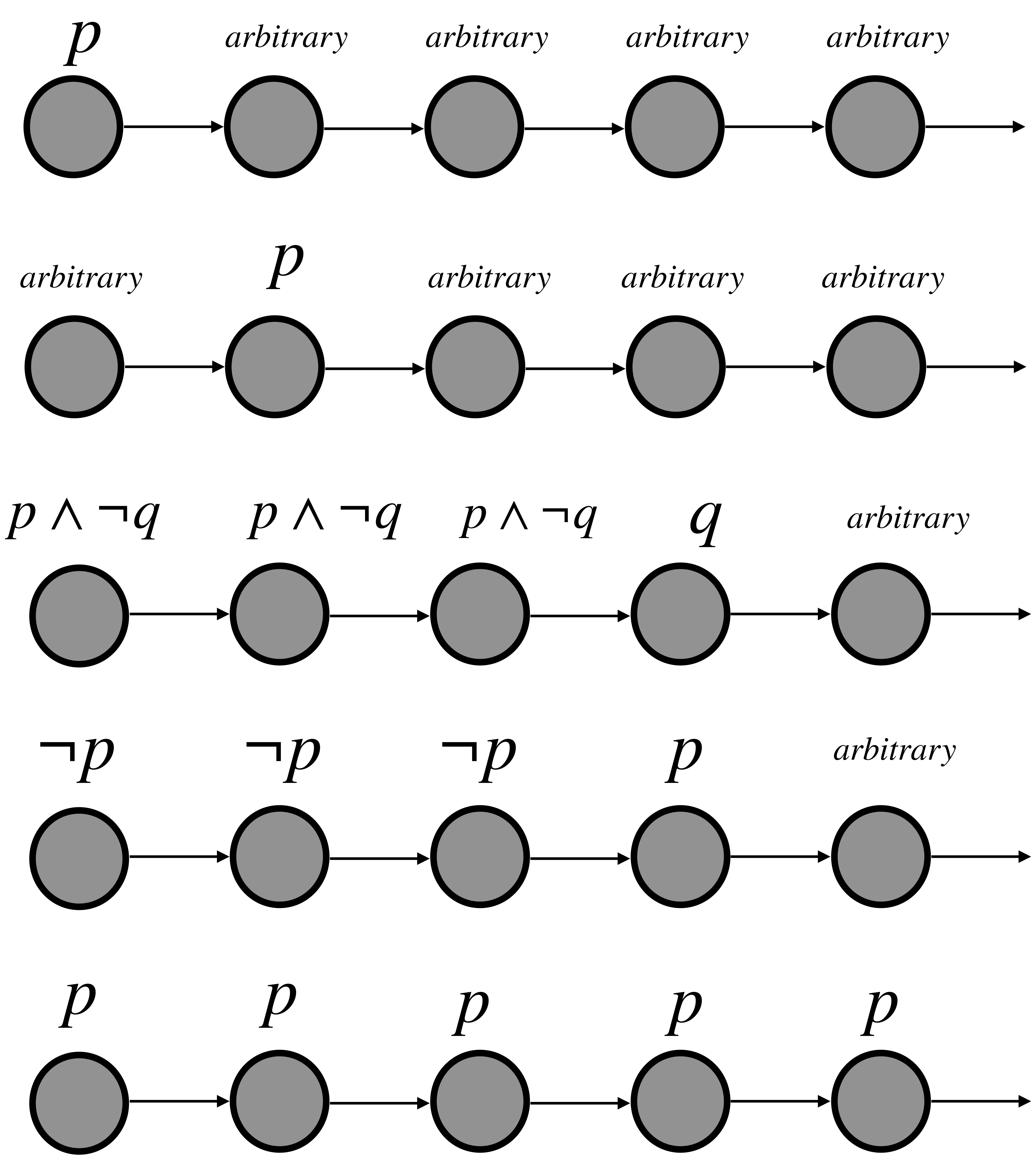
Atomic prop. P

$\mathbf{N}\ p$

$p\ \mathbf{U}\ q$

$\Diamond\ p$

$\Box\ p$



Once red, the light always becomes green eventually after being yellow for some time.

$$\Box(red \rightarrow (\Diamond\, green \;\wedge (\neg green\ \mathbf{U}\ yellow)))$$ What about — <{red}, {yellow and green}>

$$\Box(red \rightarrow \mathbf{N}\,(red\ \mathbf{U}\ (yellow \;\wedge\; \mathbf{N}\,(yellow\ \mathbf{U}\ green))))$$

$$\Box(red \rightarrow (\mathbf{N}\,(\neg green \wedge yellow) \wedge ((\neg green \wedge yellow)\mathbf{U}\ green))))$$

Suggestions from today's class:

$$\Box(red \rightarrow (\neg\mathbf{N}\ green \wedge (yellow\mathbf{U}\ green)))$$

$$\Box(red \rightarrow (\neg\mathbf{N}\ green \wedge ((\mathbf{N}\ yellow)\mathbf{U}\ green)))$$

$$\Box(red \rightarrow \mathbf{N}\ red \vee (\ \mathbf{N}\ yellow \wedge (yellow\ \mathbf{U}\ green)))$$

$$\Box(red \rightarrow ((\neg\mathbf{N}\ green) \wedge \Diamond yellow \wedge (yellow\ \mathbf{U}\ green)))$$

LTL Syntax

Primary temporal operators: **N U**

Weak Until $- F_1$ **W** F_2, F_1 must remain true until F_2 becomes true, but F_2 doesn't necessarily need to become true at any point.

$F_1 \mathbf{W} F_2 \equiv (F_1 \ \mathbf{U} \ F_2) \vee (\Box F_1)$

It is considered weaker version of **U**, which requires F_2 to eventuallyTrue.

System is in safe mode **W** system is ready

LTL Syntax

Primary temporal operators: **N U**

Release $- F_1 \mathbf{R} F_2$, F_2 must remain true until and including the point
where F_1 first becomes true, but F_1 doesn't necessarily need
to become true at any point.

$$F_1 \mathbf{R}\, F_2 \equiv ((F_2 \wedge \neg F_1)\ \mathbf{W}\ (F_2 \wedge F_1))$$

LTL: Formulas

Duality Law $\neg \mathbf{N}\, p \equiv \mathbf{N} \neg p$ $\neg \Diamond p \equiv \Box \neg p$ $\neg \Box p \equiv \Diamond \neg p$

Absorption Law $\Diamond \Box \Diamond P \equiv \Box \Diamond p$ $\Box \Diamond \Box P \equiv \Diamond \Box p$

Distributive Law $\mathbf{N}(p\ \mathbf{U}\ q) \equiv ((\mathbf{N}\ p)\ \mathbf{U}\ (\mathbf{N}\ q))$ $\Diamond(p \vee q) \equiv \Diamond p \vee \Diamond q$

$\Diamond(p \wedge q) \not\equiv \Diamond p \wedge \Diamond q$ $\Box(p \wedge q) \equiv \Box p \wedge \Box q$

Expansion Law $p\ \mathbf{U}\ q \equiv q \vee (p \wedge (\mathbf{N}\ (p\ \mathbf{U}\ q))$ $\Box p \equiv p \wedge (\mathbf{N}\ (\Box\ p))$

$\Diamond p \equiv p \vee (\mathbf{N}\ (\Diamond\ p))$

LTL: Examples

If an intruder is detected, then an alert must be raised at the 3 step.

$$\Box(IntruderDettected \rightarrow (\mathbf{N} \neg alert \wedge \ \mathbf{N} \ \mathbf{N} \neg alert \wedge \ \mathbf{N} \ \mathbf{N} \ \mathbf{N}\, alert))$$

A robot must keep moving until it reaches the charging station, and once charged, it must always eventually move again.

$$\Box(Move \ \mathbf{U} \ AtChargeStation) \wedge \Box(Charged \rightarrow \Diamond Move)$$

LTL: Semantics

We interpret our temporal formulae in a discrete, linear model of time.

$< N, I >$, where N is a set of Natural number and $I : N \mapsto 2^{\Sigma}$

I maps each Natural number (representing a moment in time) to a set of propositions

Let $\pi = a_0, a_1, a_2, \ldots$ $\pi(i) = a_i$ AP at i^{th} level.

$\pi^i = a_i, a_{i+1}, a_{i+2}, \ldots$ Suffix of π

LTL: Semantics

Semantics with respect to a given Trace (or Path) π

Let $\pi = a_0, a_1, a_2, \ldots$ $\pi(i) = a_i$ AP at i^{th} level. $\pi^i = a_i, a_{i+1}, a_{i+2}, \ldots$ Suffix of π

$\pi \models p$ Iff $p \in \pi(0)$ $\pi^i \models p$ Iff $p \in \pi(i)$

$\pi \models \mathbf{N}\ F_1$ Iff $\pi^1 \models F_1$ $\pi^i \models \mathbf{N}\ F$ Iff $\pi^{i+1} \models F_1$

$\pi \models F_1\ \mathbf{U}\ F_2$ Iff $\exists j \geq 0,\ \pi^j \models F_2$, and $\pi^i \models F_1$ for all $0 \leq i < j$

$\pi \models \Diamond\ F_1$ Iff $\exists j \geq 0,\ \pi^j \models F_1$

$\pi \models \Box\ F_1$ Iff $\forall j \geq 0,\ \pi^j \models F_1$

$\pi \models \Box \Diamond\ F_1$ Iff $\exists^\infty j \geq 0,\ \pi^j \models F_1$ $\exists^\infty = \forall i \geq 0, \exists j \geq i$

$\pi \models \Diamond \Box\ F_1$ Iff $\forall^\infty j \geq 0,\ \pi^j \models F_1$ $\exists^\infty = \exists i \geq 0, \forall j \geq i$

LTL: Semantics Kripke Structure

AP — is a set of atomic propositions (Boolean valued variables, predicates)

Kripke structure over AP as a 4-tuple $M = (S, I, R, L)$

S = a finite set of states.

I = a set of initial states $I \subseteq S$

R = a transition relation $R \subseteq S \times S$

L = a labelling function $L : S \rightarrow 2^{AP}$

LTL: Semantics Kripke Structure

Kripke structure over AP as a 4-tuple $M = (S, I, R, L)$

S = a finite set of states. $S = \{s_1, s_2, s_3\}$

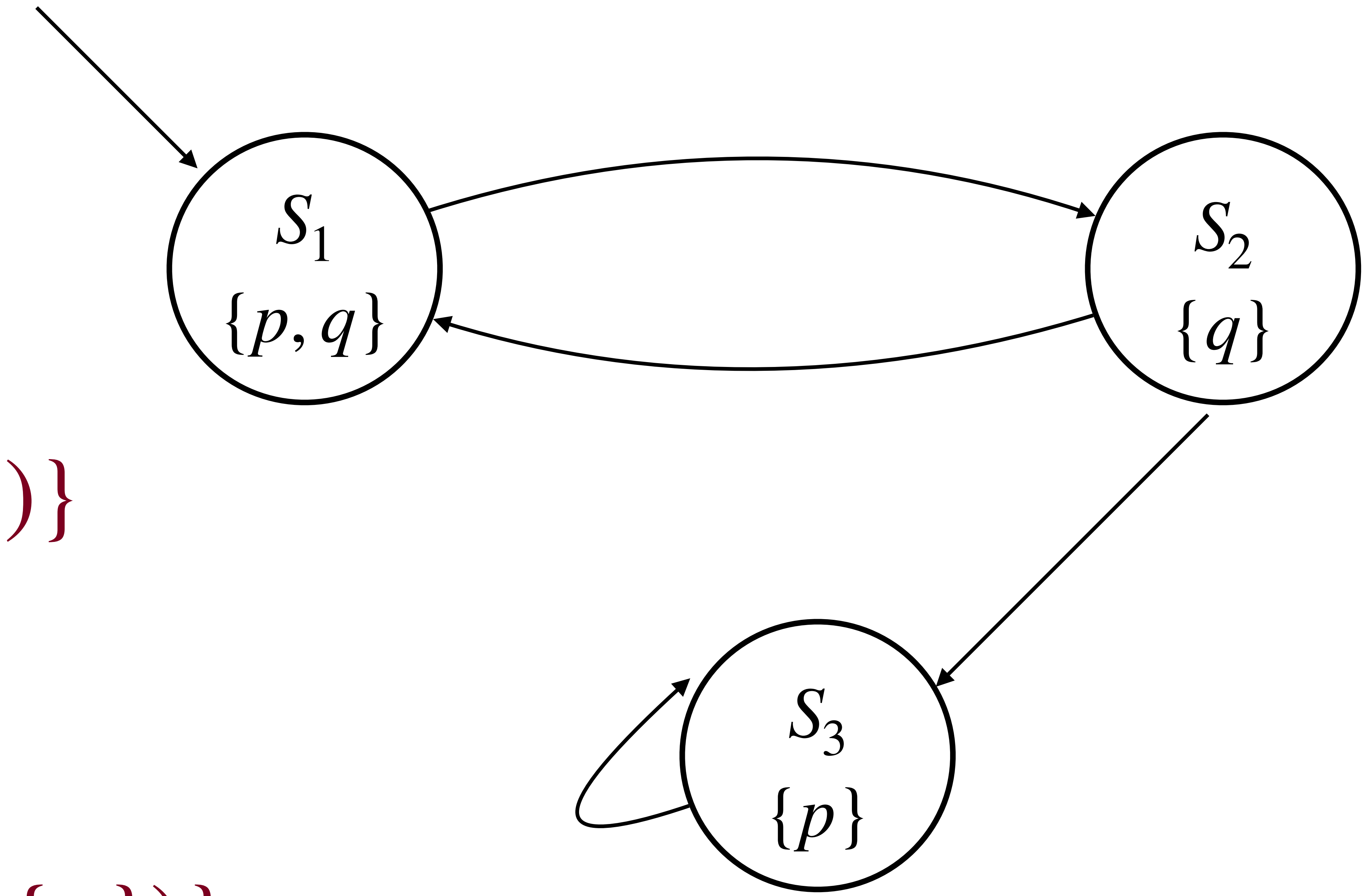
I = a set of initial states $I \subseteq S$ $I = \{s_1\}$

R = a transition relation $R \subseteq S \times S$

$$R = \{(s_1, s_2), (s_2, s_1), (s_2, s_3), (s_3, s_3)\}$$

L = a labelling function $L : S \to 2^{AP}$

$$L = \{(s_1, \{p, q\}), (s_2, \{q\}), (s_3, \{p\})\}$$



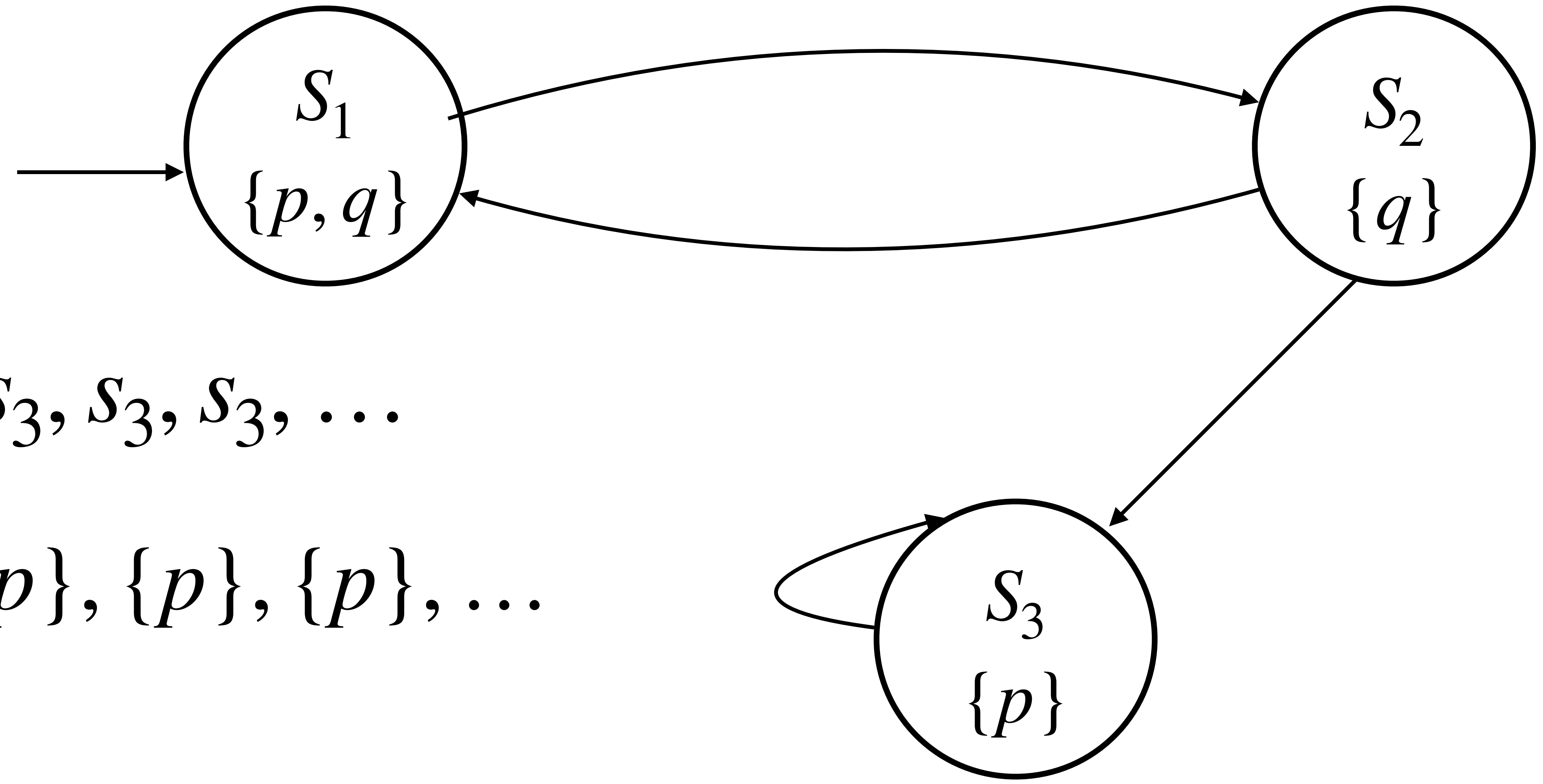
AP = $\{p, q\}$

LTL: Semantics Kripke Structure

Kripke structure over AP as a 4-tuple $M = (S, I, R, L)$ AP = $\{p, q\}$

$S = \{s_1, s_2, s_3\}$ $I = \{s_1\}$ $R = \{(s_1, s_2), (s_2, s_1), (s_2, s_3), (s_3, s_3)\}$

$L = \{(s_1, \{p, q\}), (s_2, \{q\}), (s_3, \{p\})\}$



M may produce a path $w = s_1, s_2, s_1, s_2, s_3, s_3, s_3, s_3, \ldots$

π^{s_1} $\pi = \{p, q\}, \{q\}, \{p, q\}, \{q\}, \{p\}, \{p\}, \{p\}, \ldots$

M can produce words belonging to the language —

$(\{p, q\}\{q\})^*(\{p\})^\omega \cup (\{p, q\}\{q\})^\omega$

LTL: Semantics Kripke Structure

Given a kripke structure M and a path π in M, a state $s \in S$, and an LTL formula F:

1. $< M, \pi > \vDash F$ iff $\pi^{s_o} \vDash F$, where s_o is initial state of π
2. $< M, s > \vDash F$ iff $< M, \pi > \vDash F$ for all paths starting at s.
3. $< M > \vDash F$. iff $< M, s_o > \vDash F$ for every $s_o \in I$, where I initial states of M.

LTL: Semantics

A formula F is satisfiable if there exists at least one Kripke Structure M, and at least one initial state s_o such that:

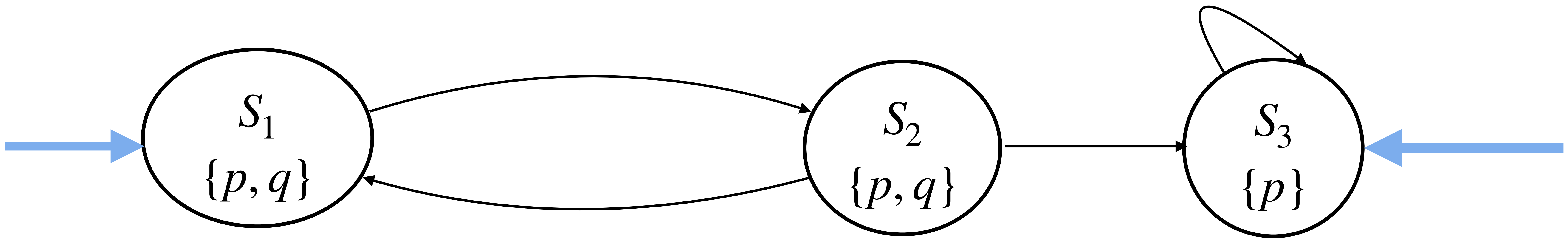
$$< M, s_o > \models F$$

A formula F is valid if for all Kripke Structures M, and for all initial states s_o:

$$< M, s_o > \models F$$

LTL model checking — Given formula F, and Kripke Structure M checks if $< M, s_o > \models F$ holds for every initial state $s_o \in I$

LTL: Semantics



Does $M \models \Box p$? Yes, $< M, s_1 > \models \Box p$ and $< M, s_3 > \models \Box p$

$\pi_1^{s_1} = < \{p, q\}\{p, q\}, \{p, q\}, \{p, q\}\ldots >$ $\pi_2^{S_1} = < \{p, q\}\{p, q\}, \{p, q\}, \{p, q\}, \{p\}, \{p\}\ldots >$ $\pi_3^{S_3} = < \{p\}, \{p\} \ldots. >$

Does $M \models \mathbf{N}(p \land q)$? No, $< M, s_1 > \models \mathbf{N}(p \land q)$, but $< M, s_3 > \not\models \mathbf{N}(p \land q)$

Does $M \models \Box(\neg q \to \Box(p \land \neg q))$? Yes

Does $M \models q\ \mathbf{U}(p \land \neg q)$? No, $< M, \pi_1 > \not\models q\ \mathbf{U}\ (p \land \neg q)$

Course Webpage

Thanks!