

COL215 Assignment 1 - 2x1 and 4x1 Multiplexers

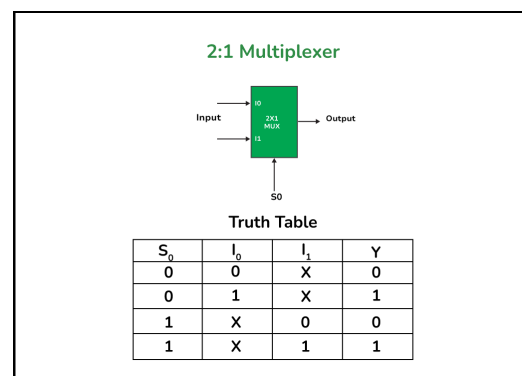
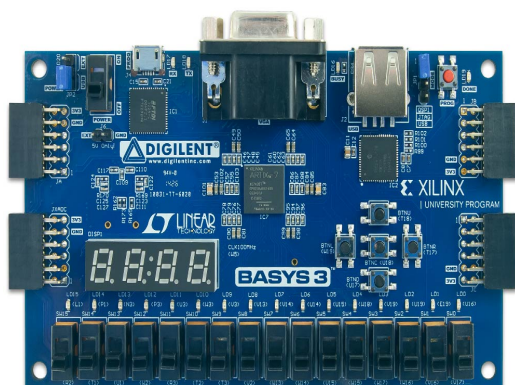
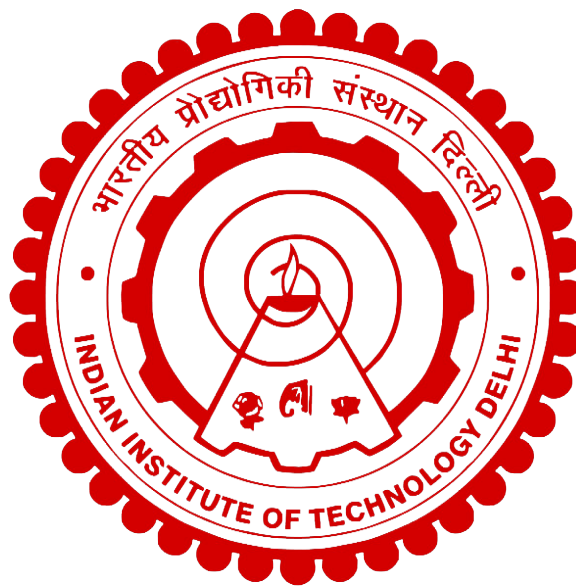
Submission By :

Yash Rawat
2023CS50334

Priyanshi Gupta
2023CS10106

Department of Computer Science and Engineering
Indian Institute of Technology, Delhi

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1 Working Principle of Multiplexers

1.1 What is a Multiplexer ?

A multiplexer is a device that selects one of several input signals and forwards the chosen input into a single output signal. An $N \times 1$ mux (abbreviation for multiplexer) has the ability to choose between N distinct inputs, for which it requires atleast $\lceil \log_2 N \rceil$ selection bits. Hence a 2×1 mux requires 1 input bit whereas 4×1 mux requires 2 input bits.

The truth table for a 2×1 Mux and a 4×1 Mux Look as following (Note that X/Y/Z means that Output is invariant of the respective bit, be it 0 or 1) :

D_0	D_1	S	Out
1	X	0	1
0	X	0	0
X	1	1	1
X	0	1	0

Table 1: Truth Table 2×1 Mux

D_0	D_1	D_2	D_3	S_0	S_1	Out
1	X	Y	Z	0	0	1
0	X	Y	Z	0	0	0
X	1	Y	Z	1	0	1
X	0	Y	Z	1	0	0
X	Y	1	Z	0	1	1
X	Y	0	Z	0	1	0
X	Y	Z	1	1	1	1
X	Y	Z	0	1	1	0

Table 2: Truth Table of a 4×1 Mux

1.2 Expression of 2×1 , 4×1 Mux and Recursive Definition of 4×1 Mux using 4×1 Mux

From the above truth table we get the formula for the 2×1 Mux as:

$$MUX_2(D_0, D_1, S) = S' \cdot D_0 + S \cdot D_1$$

whereas the formula for 4×1 Mux is :

$$MUX_4(D_0, D_1, D_2, D_3, S_0, S_1) = S'_0 \cdot (S'_1 \cdot D_0 + S_1 \cdot D_2) + S_0 \cdot (S'_1 \cdot D_1 + S_1 \cdot D_3)$$

In the case of 4X1 Mux ,by collecting the sum of minterms in the manner mentioned above, it provides insight into the recursive nature of 4X1 Mux and how we can make it from 3 2X1 Mux as shown below :

$$MUX_4(D_0, D_1, D_2, D_3, S_0, S_1) = S'_0 \cdot (S'_1 \cdot D_0 + S_1 \cdot D_2) + S_0 \cdot (S'_1 \cdot D_1 + S_1 \cdot D_3)$$

$$MUX_4(D_0, D_1, D_2, D_3, S_0, S_1) = S'_0 \cdot MUX_2(D_0, D_2, S_1) + S_0 \cdot MUX_2(D_1, D_3, S_1)$$

$$MUX_4(D_0, D_1, D_2, D_3, S_0, S_1) = MUX_2(MUX_2(D_0, D_2, S_1), MUX_2(D_1, D_3, S_1), S_0)$$

We will be implementing the 4X1 Mux based on the above principle instead of implementing the actual formula. Below attached are the circuit diagrams of our implementation of 2X1 and 4X1 Multiplexers :

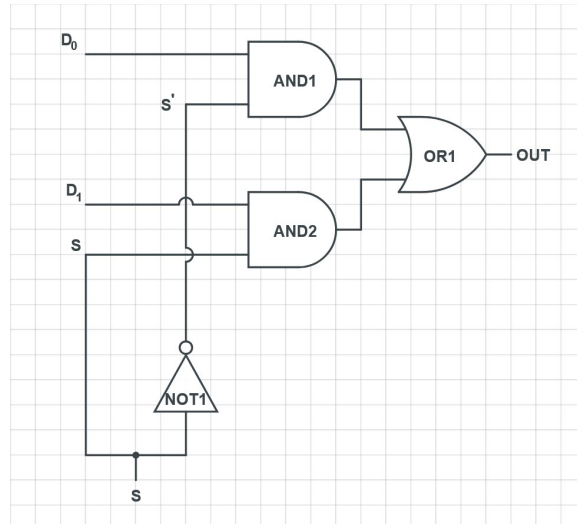


Figure 1: Implementation of 2X1 Mux using AND, NOT and OR Gates

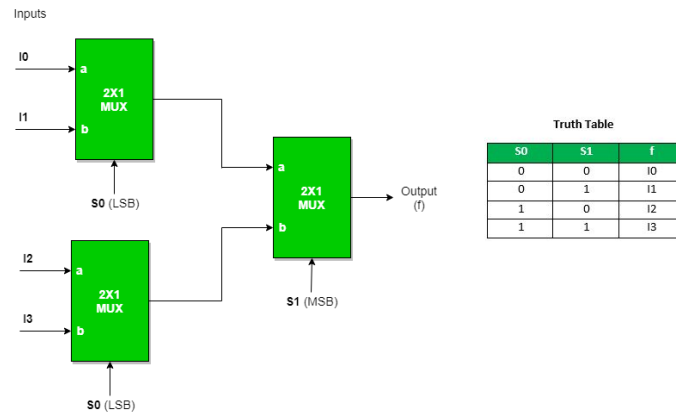
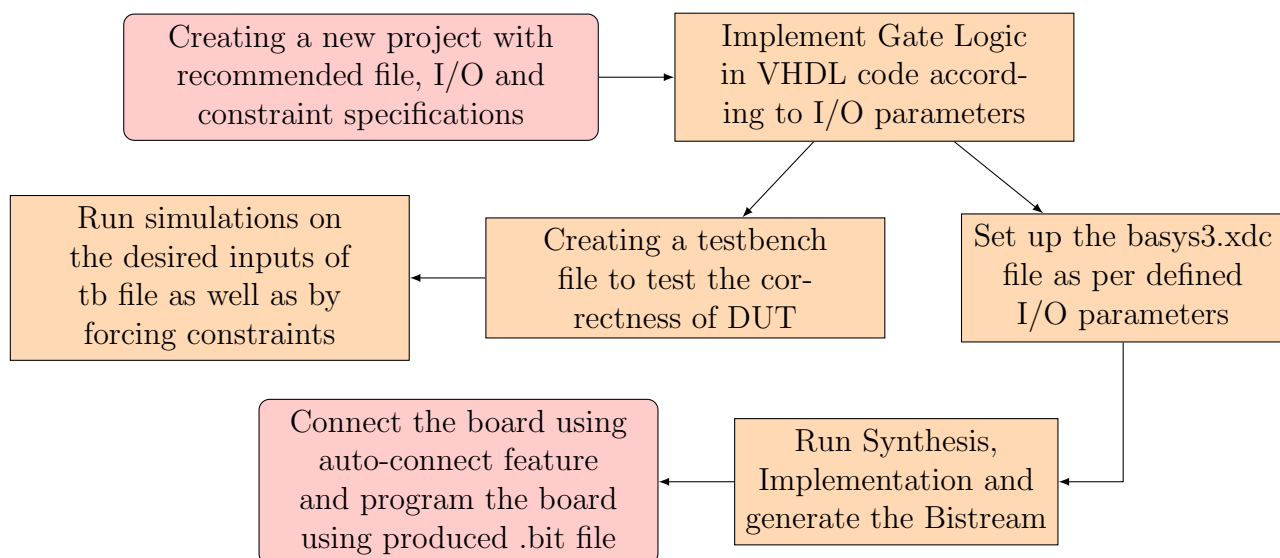


Figure 2: Recursive implementation of 4X1 Mux using 2X1 Mux

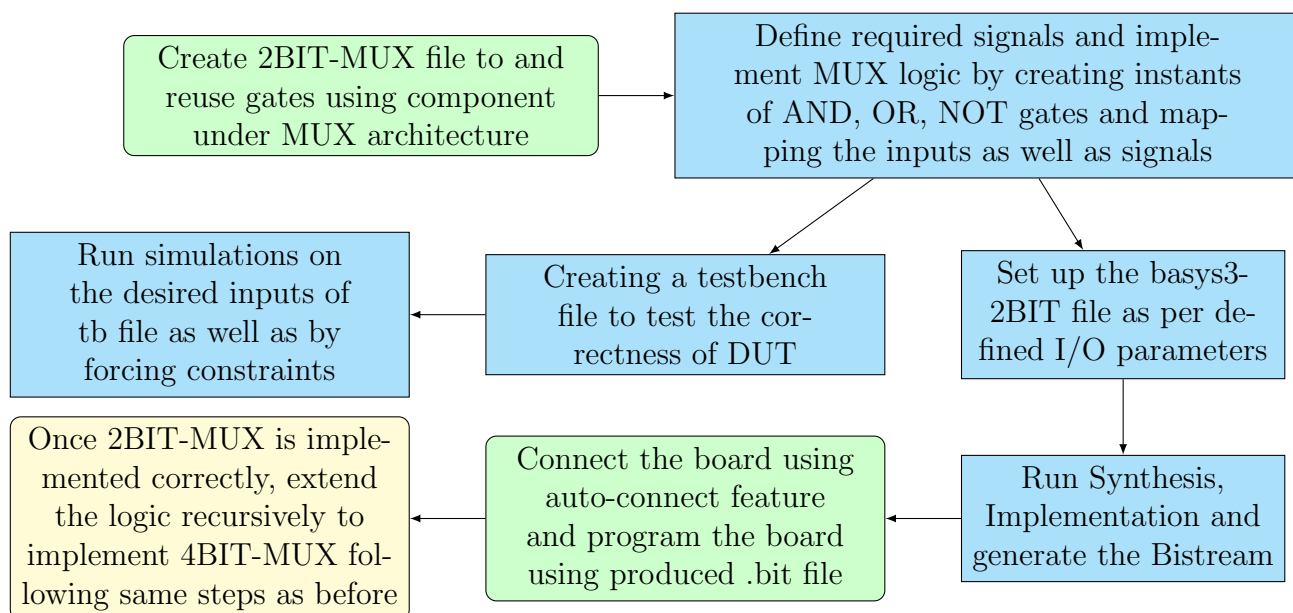
2 Design Philosophy

2.1 Design Process Flow Chart

2.1.1 Week 1 - Implementing AND, OR, NOT Gate Logic



2.1.2 Week 2 - Implementing 2x1 MUX and by extension 4x1 MUX



2.2 I/O Terminals

2.2.1 Week 1 - Basic Setup

Implementing a 2 input AND gate, 2 input OR gate and 1 input NOT gate required us to define 5 inputs and three outputs. The mapping of ports to the constraint file are as give below :

SNo.	Port Name	Board Contstraint	I/O
1	a_and_in	V17	IN
2	b_and_in	V16	IN
3	c_and_out	U16	OUT
4	d_or_in	W16	IN
5	e_or_in	W17	IN
6	f_or_out	E19	OUT
7	g_not_in	W15	IN
8	h_not_out	U19	OUT

Table 3: Mapping of Ports and Constraints in Basys3.xdc for BASIC-gates

2.2.2 Week 2 - 2x1-MUX and 4x1-MUX

Implementing a 2x1 MUX using 2 AND gates, 1 OR gate and 1 NOT gate required us to define three inputs and one output. The mapping of ports to the constraint file are as give below :

SNo.	Port Name	Board Contstraint	I/O
1	D1	V17	IN
2	D2	V16	IN
3	S	W16	IN
4	O	U16	OUT

Table 4: Mapping of Ports and Constraints in Basys3.2BIT.xdc for 2x1 MUX

Implementing a 4x1 MUX using 4 AND gates, 3 OR gates and 2 NOT gates required us to define six inputs and one output. The mapping of ports to the constraint file are as give below :

SNo.	Port Name	Board Contstraint	I/O
1	d1	V17	IN
2	d2	V16	IN
3	d3	W16	IN
4	d4	W17	IN
5	s1	W15	IN
6	s2	V16	IN
7	o	U16	OUT

Table 5: Mapping of Ports and Constraints in Basys3_4BIT.xdc for 4x1 MUX

3 Simulation and Schematic Snapshots

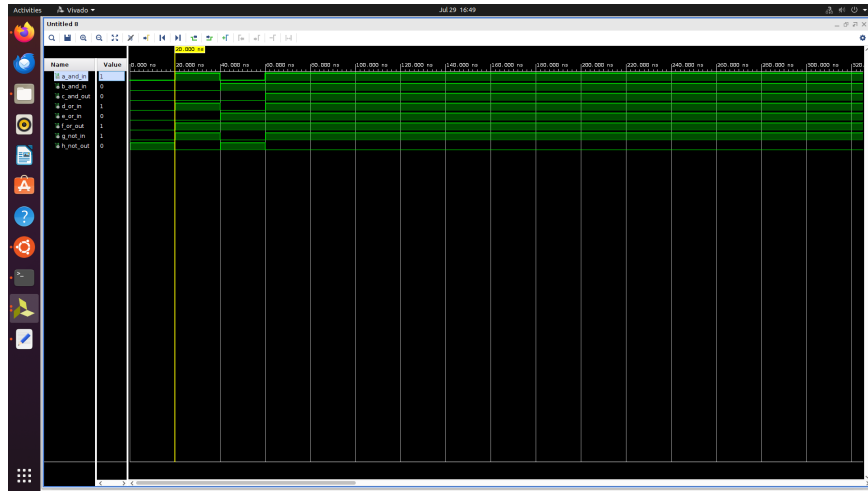


Figure 3: Simulation for AND,OR and NOT gates

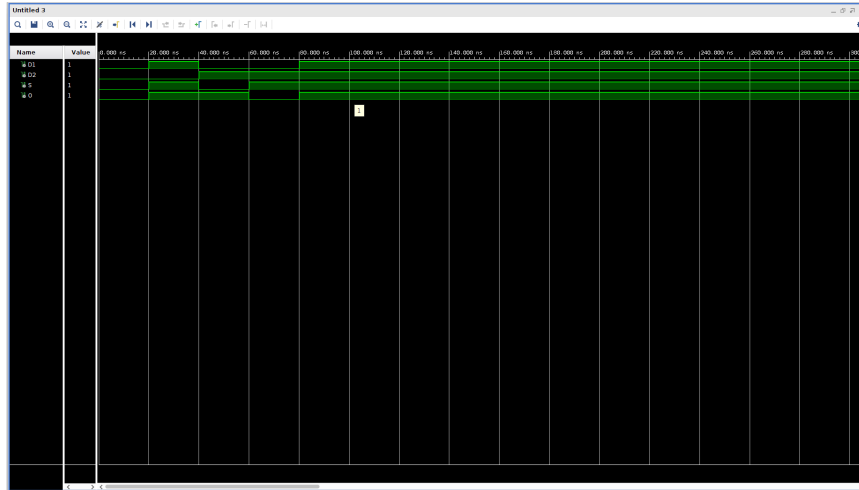


Figure 4: Simulation for 2x1 MUX

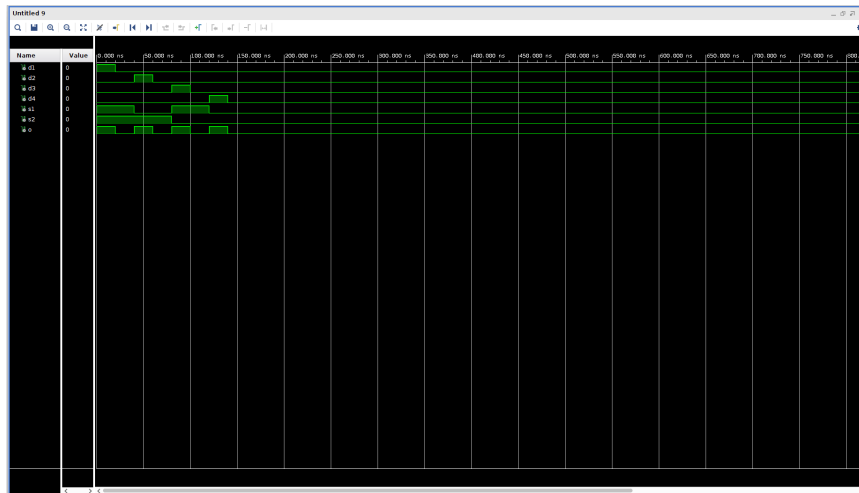


Figure 5: Simulation for 4x1 MUX

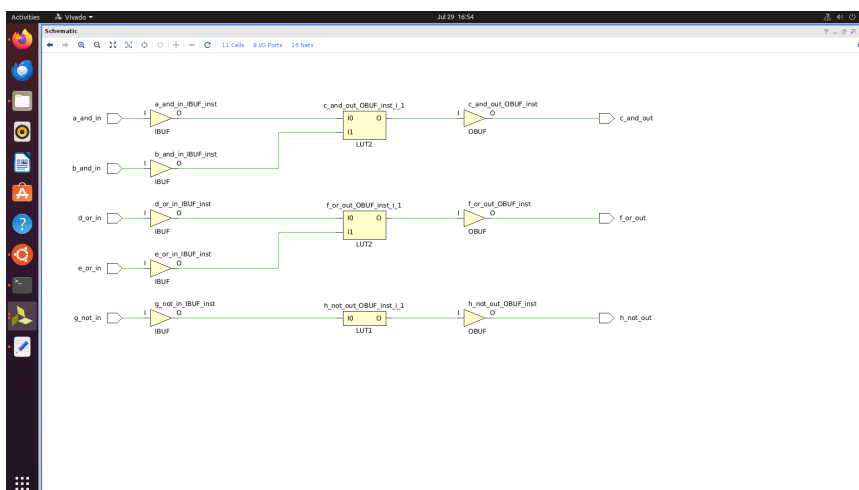


Figure 6: Schematic of AND,OR and NOT gates

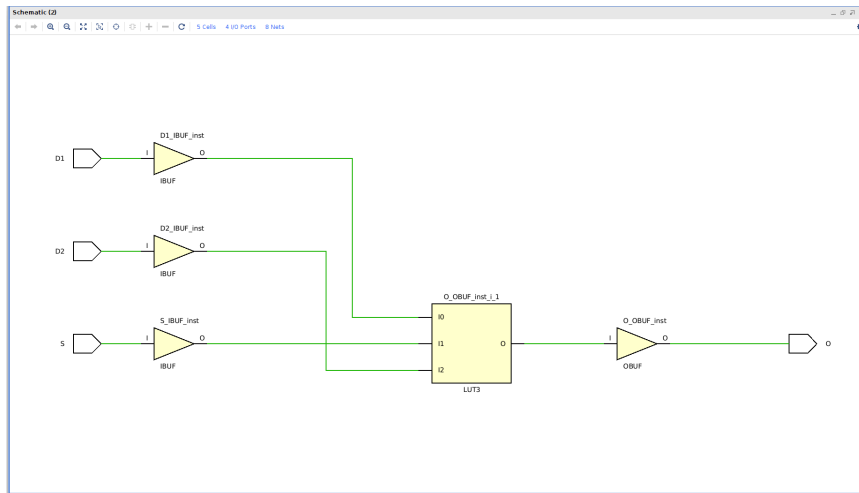


Figure 7: Schematic for 2x1 MUX

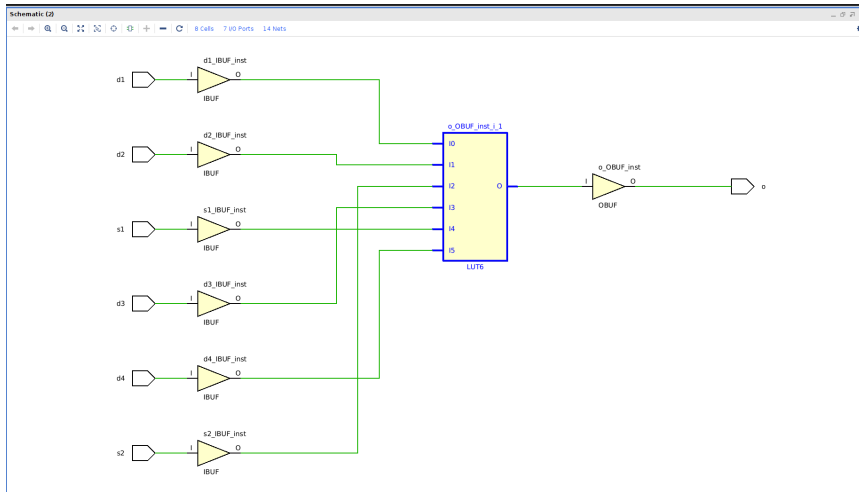


Figure 8: Schematic for 4x1 MUX

1. Slice Logic

Site Type	Used	Fixed	Prohibited	Available	Util%
Slice LUTs*	1	0	0	20800	<0.01
LUT as Logic	1	0	0	20800	<0.01
LUT as Memory	0	0	0	9600	0.00
Slice Registers	0	0	0	41600	0.00
Register as Flip Flop	0	0	0	41600	0.00
Register as Latch	0	0	0	41600	0.00
F7 Muxes	0	0	0	16300	0.00
F8 Muxes	0	0	0	8150	0.00

* Warning! The Final LUT count, after physical optimizations and full implementation, is typically lower. Run opt_design after synthesis, if not already completed, for a more realistic count.

2. Memory

Site Type	Used	Fixed	Prohibited	Available	Util%
Block RAM Tile	0	0	0	50	0.00
RAMB36/FIFO*	0	0	0	50	0.00
RAMB18	0	0	0	100	0.00

* Note: Each Block RAM Tile only has one FIFO logic available and therefore can accommodate only one FIFO36E1 or one FIFO18E1. However, if a FIFO18E1 occupies a Block RAM Tile, that tile can still accommodate a RAMB18E1

3. DSP

Site Type	Used	Fixed	Prohibited	Available	Util%
DSPs	0	0	0	90	0.00

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Slice Registers	0	0	0	41600	0.00
Register as Flip Flop	0	0	0	41600	0.00
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F7 Muxes	0	0	0	16300	0.00
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* Warning! The Final LUT count, after physical optimizations and full implementation, is typically lower. Run opt_design after synthesis, if not already completed, for a more realistic count.

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RAMB36/FIFO*	0	0	0	50	0.00
RAMB18	0	0	0	100	0.00

* Note: Each Block RAM Tile only has one FIFO logic available and therefore can accommodate only one FIFO36E1 or one FIFO18E1. However, if a FIFO18E1 occupies a Block RAM Tile, that tile can still accommodate a RAMB18E1

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