

Tumhara Approach: Priyansh ka Logic

1. Array `arr` me elements `0` se `n` ke beech hain, `length = n`.
2. Ideal position har element ka `arr[i]` ka index = `arr[i]` hona chahiye (agar value < n).
3. Tum kya kar rahe ho:
 - Jab tak element apni jagah par nahi hai, swap karte raho.
 - Jab sahi position aa jaye, next index par move.
4. End me, jo index aur value mismatch kare, wahi missing number hoga.
5. Agar koi mismatch nahi mila, missing number = `n`.

Tumhara Code ka Explanation Line by Line:

```
int missingNumber(vector<int>& arr) {
    int i = 0;
    int n = arr.size();

    // Step 1: Place every element at its correct index
    while(i < n) {
        if(arr[i] < n && arr[i] != i) {
            swap(arr[i], arr[arr[i]]);
        } else {
            i++;
        }
    }

    // Step 2: Check which index is incorrect
    for(int i = 0; i < n; i++) {
        if(arr[i] != i) return i; // Found the missing number
    }
}
```

```

    }

    // Step 3: If all are correct, missing number is n
    return n;
}

```

✅ Dry Run Priyansh ke Perspective se:

Example: `nums = [3, 0, 1]`

`n = 3`

- `i = 0`: `arr[0] = 3` → condition `(arr[i] < n)` fail → `i++`
- `i = 1`: `arr[1] = 0` → condition true → `swap(arr[1], arr[arr[1]])`
Swap(0th and 1st) → `[0, 3, 1]`
- `i` stays 1: `arr[1] = 3` → fail → `i++`
- `i = 2`: `arr[2] = 1` → true → `swap(arr[2], arr[arr[2]])`
Swap(2nd and 1st) → `[0, 1, 3]`
- `i` stays 2: `arr[2] = 3` → fail → `i++`
- Done.

Now second loop:

- `i=0` → `arr[0] = 0` ✅
- `i=1` → `arr[1] = 1` ✅
- `i=2` → `arr[2] = 3` ❌ → return 2 ✅

Answer = 2.

✅ Complexity:

- **Time:** $O(n)$ (kyunki har element max 1-2 swaps lega).
- **Space:** $O(1)$.

🔥 Priyansh ka tareeqa kyun OP hai?

- Tumne **extra space use nahi kiya**, unlike boolean array method.
 - Tumne **sum method** ya **XOR trick** avoid kiya, but ye approach **more generalized** hai (ye logic aur problems me bhi kaam aata hai like "First Missing Positive").
 - Tumhara code neat hai aur interview friendly hai.
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