



# 746. Min Cost Climbing Stairs (Iterative DP)



```
class Solution {  
public:  
    int minCostClimbingStairs(vector<int>& arr) {  
        for (int i = 2; i < arr.size(); i++) {  
            arr[i] += min(arr[i - 1], arr[i - 2]);  
        }  
        return min(arr[arr.size() - 1], arr[arr.size() - 2]);  
    }  
};
```



## Step-by-Step Explanation

Step	Description
1	Start from the <b>2nd index (<math>i = 2</math>)</b> because the first two steps (0 and 1) are your starting options.
2	For each step $i$ , compute the <b>minimum cost</b> to reach it by taking the cheaper of the two previous steps: $\text{arr}[i] += \min(\text{arr}[i-1], \text{arr}[i-2])$ .
3	This means at every index, you're storing the <b>minimum total cost to reach that step</b> .
4	Finally, you can reach the top either from the <b>last step (<math>n-1</math>)</b> or the <b>second last step (<math>n-2</math>)</b> , so return $\min(\text{arr}[n-1], \text{arr}[n-2])$ .



At each step, you decide whether it's cheaper to come from one step below or two steps below.

By the end of the loop, each element of `arr` represents the **minimum cumulative cost** to reach that step.

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## Dry Run Example

**Input:** [1,100,1,1,1,100,1,1,100,1]

i	arr before update	Operation	arr after update
2	[1,100,1,...]	<code>arr[2] += min(100,1)</code> → <code>arr[2]=1+1</code>	[1,100,2,...]
3	[1,100,2,...]	<code>arr[3] += min(100,2)</code> → <code>arr[3]=1+2</code>	[1,100,2,3,...]
4	[1,100,2,3,...]	<code>arr[4] += min(2,3)</code> → <code>arr[4]=1+2</code>	[1,100,2,3,3,...]
...	continue	...	...
Final	[1,100,2,3,3,103,4,4,104,5]	<code>min(5,104)</code> →  6	

## Time Complexity

- **O(n)** → Single pass through the array.

## Space Complexity

- **O(1)** → In-place update of the input array.

## Key Idea

- Modify the input array to act as a **DP table**.
- Each cell represents the **minimum cost** to reach that step.
- Answer = `min(last, second_last)` .

## Quick Summary

Concept	Description
<b>Approach</b>	Bottom-Up Dynamic Programming
<b>Transition</b>	<code>arr[i] += min(arr[i-1], arr[i-2])</code>
<b>Base Case</b>	Start from step 0 or 1

Concept	Description
<b>Answer</b>	<code>min(arr[n-1], arr[n-2])</code>
<b>Complexity</b>	Time: O(n), Space: O(1)

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