



746. Min Cost Climbing Stairs

Problem Statement

You are given an integer array `cost` where `cost[i]` is the cost of the i -th step on a staircase.

Once you pay the cost, you can either climb **one or two steps**.

You can start from **index 0 or index 1**.

Return the **minimum cost** to reach the **top of the floor**.

Example 1

Input: `cost = [10,15,20]`

Output: `15`

Explanation:

- Start from index `1`.
- Pay `15` and climb two steps to reach the top.

✓ Total cost = `15`

Example 2

Input: `cost = [1,100,1,1,1,100,1,1,100,1]`

Output: `6`

Explanation:

Optimal path $\rightarrow 0 \rightarrow 2 \rightarrow 4 \rightarrow 6 \rightarrow 7 \rightarrow 9 \rightarrow \text{top}$

✓ Total cost = `1 + 1 + 1 + 1 + 1 + 1 = 6`

 **Approach: Dynamic Programming (Top-Down with Memoization)**

 **Intuition**

At each step i , we can either come from:

- step $i-1$ (pay $\text{cost}[i-1]$)
- or step $i-2$ (pay $\text{cost}[i-2]$)

So,

```
dp[i] = cost[i] + min(dp[i-1], dp[i-2])
```

We finally take the **minimum of the last two steps** to reach the top.

Recurrence Relation

```
helper(i) = cost[i] + min(helper(i-1), helper(i-2))
```

Base cases:

```
helper(0) = cost[0]  
helper(1) = cost[1]
```

Final answer:

```
min(helper(n-1), helper(n-2))
```

Code (C++)

```
class Solution {  
public:  
    int helper(vector<int>& cost, int i, vector<int>& dp) {  
        if (i == 0 || i == 1) return cost[i];  
        if (dp[i] != -1) return dp[i];  
        return dp[i] = cost[i] + min(helper(cost, i - 1, dp), helper(cost, i - 2, dp));  
    }  
  
    int minCostClimbingStairs(vector<int>& cost) {  
        int n = cost.size();  
        vector<int> dp(n, -1);  
    }  
};
```

```
        return min(helper(cost, n - 1, dp), helper(cost, n - 2, dp));
    }
};
```

Time Complexity


- **O(n)** → Each state (step) is computed once.

Space Complexity

- **O(n)** → For the recursion + memoization array.

Alternate (Bottom-Up DP)

```
int minCostClimbingStairs(vector<int>& cost) {
    int n = cost.size();
    for (int i = 2; i < n; i++) {
        cost[i] += min(cost[i-1], cost[i-2]);
    }
    return min(cost[n-1], cost[n-2]);
}
```

 **Optimized:** No extra space used.

Key Takeaways

- Can start from step 0 or 1.
- Either climb one or two steps each time.
- Final cost = `min(dp[n-1], dp[n-2])`.
- Classic **Dynamic Programming** pattern (like Fibonacci).