

Revision Notes: Level Order Traversal of a Binary Tree

Problem Statement

Given the root of a binary tree, return the level order traversal of its nodes' values. (i.e., from left to right, level by level).

Example 1:

Input: root = [3,9,20,null,null,15,7]

Output: [[3],[9,20],[15,7]]

Example 2:

Input: root = [1]

Output: [[1]]

Example 3:

Input: root = []

Output: []

Constraints:

- The number of nodes in the tree is in the range [0, 2000].
 - $-1000 \leq \text{Node.val} \leq 1000$
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Code Implementation:

```
class Solution {
public:
    int level(TreeNode* root){
        if(root == NULL) return 0;
        return 1 + max(level(root->right), level(root->left));
    }

    void nth_level(TreeNode* root, int lvl, int t_lvl, vector<int> v){
        if(root == NULL) return;
        if(lvl == t_lvl){
            v.push_back(root->val);
            return;
        }
        nth_level(root->left, lvl+1, t_lvl, v);
        nth_level(root->right, lvl+1, t_lvl, v);
    }

    void lOrder(TreeNode* root, vector<vector<int>>& ans){
        int n = level(root);
        for(int i = 0; i < n; i++){
            vector<int> v;
            nth_level(root, 0, i, v);
            ans.push_back(v);
        }
    }

    vector<vector<int>> levelOrder(TreeNode* root){
        vector<vector<int>> ans;
        lOrder(root, ans);
        return ans;
    }
};
```

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Explanation & Logic

1. Finding the Tree Depth (level function)

- This function calculates the depth of the binary tree recursively.
- If the root is NULL, return 0.
- Otherwise, return $1 + \max(\text{depth of left subtree}, \text{depth of right subtree})$.

2. Traversing Nodes at a Given Level (nth_level function)

- This function is used to collect all node values at a specific level (t_lvl).
- If the current level (lvl) matches t_lvl , add the node's value to the vector.
- Recursively call this function for left and right subtrees with $lvl + 1$.

3. Collecting Nodes for Each Level (IOrder function)

- First, find the depth of the tree (n).
- Iterate from level 0 to $n-1$, calling `nth_level` to collect node values for each level.
- Store the collected values in `ans`.

4. Final Level Order Traversal (levelOrder function)

- Calls `IOrder` to populate the answer vector and returns it.

Time Complexity Analysis

- `level()` function runs in $O(N)$ time (worst-case scenario for skewed trees).
- `nth_level()` is called for every level, leading to a **worst-case complexity of $O(N^2)$** for skewed trees.
- **Overall Complexity: $O(N^2)$** (not the most optimal approach but works well for small constraints).

Alternative Approach (Using Queue - BFS)

- A more efficient approach is to use **Breadth-First Search (BFS)** with a queue.
- This approach processes nodes level by level in $O(N)$ time complexity.

Key Takeaways

- The given approach follows a recursive method to perform level order traversal.
 - It uses a depth-first strategy to collect node values at each level.
 - While it works correctly, a BFS-based approach using a queue is more efficient. And that too is mentioned in after this approach...
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