

# 1046. Last Stone Weight

**Difficulty**: Easy

Topic: Heap / Priority Queue

Tags: Greedy , Heap , Priority Queue

#### Problem Description

You are given an array of integers stones, where stones[i] is the weight of the i-th stone.

We are playing a game where in each turn:

- Select the two heaviest stones.
- Smash them together:
  - If both stones have equal weight, both are destroyed.
  - If not, the smaller one is destroyed, and the bigger one becomes y-x.

Repeat this until at most one stone is left.

**Return** the weight of the last stone. If no stone is left, return **o**.

#### **Thought Process / Intuition**



- Use a **Max Heap** to always get the 2 largest stones in each round.
- Push all stones into it.
- In every iteration:
  - Pop top 2 stones.
  - If they are not equal, push their difference.
- At the end:

• Return the last stone or 0.

## **Code**

```
class Solution {
public:
  int lastStoneWeight(vector<int>& stones) {
     priority_queue<int>pq;
    for(int ele:stones){
       pq.push(ele);
     }
     while(pq.size()>1){
       int x = pq.top();
       pq.pop();
       int y = pq.top();
       pq.pop();
       if(x!=y) pq.push(x-y);
    if(pq.size()>0){
       return pq.top();
     }
     else return 0;
  }
};
```

#### Time and Space Complexity

Complexity Type	Value	Explanation
Time	O(N log N)	Each insert/pop in heap = log N, repeated N times
Space	O(N)	Priority Queue stores all stones

## 🔁 Dry Run (with 🦀 Emojis)

**Input:** stones = [2, 7, 4, 1, 8, 1]

Step	Max Heap (PQ)	Operation Performed	Remaining Stones
1	[8, 7, 4, 1, 2, 1]	Pop 8 & 7 → Push 1 (8-7)	[4, 2, 1, 1, 1]
2	[4, 2, 1, 1, 1]	Pop 4 & 2 → Push 2 (4-2)	[2, 1, 1, 1]
3	[2, 1, 1, 1]	Pop 2 & 1 → Push 1 (2-1)	[1, 1, 1]
4	[1, 1, 1]	Pop 1 & 1 → Same → Both destroyed	[1]
<b>~</b>	[1]	One stone left → Return 1	

# Final Output: 1

## ★ Takeaways

- Max Heap is best for retrieving highest values quickly.
- Greedy logic to simulate stone smashing.
- Very handy for interviews to showcase **Heap usage**.