

## Binary Tree Postorder Traversal - Explanation & Dry Run

### Problem Statement

Given the root of a binary tree, return the postorder traversal of its nodes' values.

### Postorder Traversal Rule:

- Traverse the **left** subtree.
  - Traverse the **right** subtree.
  - Visit the **root** node.
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### Code Explanation

#### C++ Implementation

```
/**
 * Definition for a binary tree node.
 * struct TreeNode {
 *     int val;
 *     TreeNode *left;
 *     TreeNode *right;
 *     TreeNode() : val(0), left(nullptr), right(nullptr) {}
 *     TreeNode(int x) : val(x), left(nullptr), right(nullptr) {}
 *     TreeNode(int x, TreeNode *left, TreeNode *right) : val(x), left(left), right(right) {}
 * };
 */

class Solution {
public:
    void helper(TreeNode* root, vector<int> &ans) {
        if (root == NULL) return; // Base case
        helper(root->left, ans); // Traverse left subtree
        helper(root->right, ans); // Traverse right subtree
        ans.push_back(root->val); // Visit the root
    }

    vector<int> postorderTraversal(TreeNode* root) {
```

```

        vector<int> ans;

        helper(root, ans);

        return ans;
    }
};

```

### Breakdown of Code:

#### 1. helper Function:

- Recursively visits each node in **postorder sequence**.
- Base case: If root == NULL, return immediately.
- First calls itself recursively for **left** subtree.
- Then calls itself recursively for **right** subtree.
- Finally, adds root->val to the result list (ans).

#### 2. postorderTraversal Function:

- Creates an empty vector ans to store the traversal.
- Calls the helper function with the root node.
- Returns the final result.

### Dry Run with Examples

#### Example 1:

Input: root = [1, null, 2, 3]

#### Tree Structure:

```

1
 \
  2
 /
3

```

#### Recursive Calls:

1. helper(1) → helper(NULL) → ans = []
2. helper(2) → helper(3) → ans = []
3. ans.push\_back(3) → ans = [3]
4. ans.push\_back(2) → ans = [3, 2]

5. `ans.push_back(1) → ans = [3, 2, 1]`

**Output: [3, 2, 1]**

---

### Example 2:

**Input: root = [1,2,3,4,5,null,8,null,null,6,7,9]**

**Tree Structure:**

```

  1
 / \
2   3
/\  \
4 5  8
/\  \
6 7  9
```

**Recursive Calls:**

1. `helper(4) → ans = [4]`
2. `helper(6) → ans = [4, 6]`
3. `helper(7) → ans = [4, 6, 7]`
4. `helper(5) → ans = [4, 6, 7, 5]`
5. `helper(2) → ans = [4, 6, 7, 5, 2]`
6. `helper(9) → ans = [4, 6, 7, 5, 2, 9]`
7. `helper(8) → ans = [4, 6, 7, 5, 2, 9, 8]`
8. `helper(3) → ans = [4, 6, 7, 5, 2, 9, 8, 3]`
9. `helper(1) → ans = [4, 6, 7, 5, 2, 9, 8, 3, 1]`

**Output: [4, 6, 7, 5, 2, 9, 8, 3, 1]**

---

### Edge Cases

#### 1. Empty Tree:

- **Input:** `root = []`
- **Output:** `[]`

#### 2. Single Node:

- **Input:** `root = [1]`

- **Output:** [1]

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### Time & Space Complexity

- **Time Complexity:**  $O(N)$ , as every node is visited once.
  - **Space Complexity:**  $O(N)$  (worst case for recursion stack if the tree is skewed).
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### Follow-up: Iterative Solution (Using Stack)

To solve the problem iteratively, we can use a stack:

```
vector<int> postorderTraversal(TreeNode* root) {
```

```
    vector<int> ans;
```

```
    if (!root) return ans;
```

```
    stack<TreeNode*> st;
```

```
    TreeNode* last = NULL;
```

```
    TreeNode* curr = root;
```

```
    while (!st.empty() || curr) {
```

```
        if (curr) {
```

```
            st.push(curr);
```

```
            curr = curr->left;
```

```
        } else {
```

```
            TreeNode* node = st.top();
```

```
            if (node->right && node->right != last) {
```

```
                curr = node->right;
```

```
            } else {
```

```
                ans.push_back(node->val);
```

```
                last = node;
```

```
                st.pop();
```

```
            }
```

```
        }
```

```
    }
```

```
    return ans;
```

}

#### Advantages of Iterative Approach:

- **Avoids recursion depth issues** (useful for deep trees).
  - **More memory-efficient in some cases.**
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#### Conclusion

- We explored a recursive solution for **Postorder Traversal** of a binary tree.
- Dry-ran multiple test cases to understand recursive calls.
- Discussed an **iterative stack-based** approach as an alternative.
- **Postorder traversal follows the sequence: Left → Right → Root.**

This concludes our explanation of **Binary Tree Postorder Traversal!**