

C - Vacation

Problem Summary

Taro ke paas **N** din ki summer vacation hai. Har din woh teen kaam me se ek karega:

- **A**: swim (points = a_i)
- **B**: catch bugs (points = b_i)
- **C**: homework (points = c_i)

Rule: woh **do consecutive din same activity nahi kar saka**. Hume maximum total happiness score nikalna hai.

Input / Output

Input format

```
N  
a1 b1 c1  
a2 b2 c2  
...  
aN bN cN
```

Output

```
Maximum total points (single integer)
```

Constraints

- $1 \leq N \leq 10^5$
- $1 \leq a_i, b_i, c_i \leq 10^4$

- All integers
- Time limit: 2 seconds, Memory limit: 1024 MiB

Large N means solution must be O(N) time and O(N) or O(1) extra memory (but O(N) is fine here).

Key idea (Dynamic Programming)

Define $dp[i][j]$ = maximum total happiness up to day i if on day i we choose activity j .

Let activities index: 0 = A, 1 = B, 2 = C.

Transition (for $i \geq 1$):

$$\begin{aligned} dp[i][0] &= a_i + \max(dp[i-1][1], dp[i-1][2]) \\ dp[i][1] &= b_i + \max(dp[i-1][0], dp[i-1][2]) \\ dp[i][2] &= c_i + \max(dp[i-1][0], dp[i-1][1]) \end{aligned}$$

Base ($i = 0$):

$$\begin{aligned} dp[0][0] &= a_1 \\ dp[0][1] &= b_1 \\ dp[0][2] &= c_1 \end{aligned}$$

Answer = $\max(dp[N-1][0], dp[N-1][1], dp[N-1][2])$.

Why it works: on day i we pick one activity and we can add the best obtainable total from previous day that used a different activity.

Complexity

- **Time:** $O(N)$ — each day we do $O(1)$ work (3 transitions).
- **Space:** $O(N*3)$ if we store full table. Can be reduced to $O(1)$ by keeping only previous day's 3 values.

Given constraints this is perfectly fine.

Walkthrough — Sample 1

Input:

```
3
10 40 70
20 50 80
30 60 90
```

Manual DP table:

- Day 1 ($i=0$): $dp = [10, 40, 70]$
- Day 2 ($i=1$):
 - $dp[1][0] = 20 + \max(40, 70) = 20 + 70 = 90$
 - $dp[1][1] = 50 + \max(10, 70) = 50 + 70 = 120$
 - $dp[1][2] = 80 + \max(10, 40) = 80 + 40 = 120$
- Day 3 ($i=2$):
 - $dp[2][0] = 30 + \max(120, 120) = 150$
 - $dp[2][1] = 60 + \max(90, 120) = 180$
 - $dp[2][2] = 90 + \max(90, 120) = 210$

Answer = $\max(150, 180, 210) = 210$ (sequence C, B, C).

Edge cases & tips

- **N = 1:** simply $\max(a_1, b_1, c_1)$.
- **All equal values:** algorithm still works.
- **Large values:** totals can be up to $N * 10^4 \rightarrow$ for $N = 10^5$ this is 10^9 , fits in 32-bit signed int but to be safe use 64-bit (`long long`) if you want headroom.
- **Input reading:** ensure you read exactly `3*N` integers after N; otherwise the program will block waiting for input in interactive/run environments.

- **Memory optimization:** if memory tight, use three variables `prev0, prev1, prev2` and update to `cur0, cur1, cur2` per day.
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Mapping to your submitted code

- Your submitted code implements the DP exactly as explained: it reads day 1 into `dp[0]`, then loops `i=1..N-1` reading day i and computing transitions.
 - Final output prints `max(dp[N-1][0], dp[N-1][1], dp[N-1][2])` — correct.
 - Only practical improvements optional: use `long long` for dp (safer), add input validation if desired, or switch to 3 variable rolling-array to save memory.
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Pseudocode (compact)

```
read N
read a1,b1,c1
prev = [a1,b1,c1]
for i = 2 to N:
    read a,b,c
    cur0 = a + max(prev[1], prev[2])
    cur1 = b + max(prev[0], prev[2])
    cur2 = c + max(prev[0], prev[1])
    prev = [cur0, cur1, cur2]
print max(prev)
```

Quick checklist before submitting

- Make sure standard input contains N and exactly 3*N numbers.
 - No extra prompts or debug prints.
 - Check 64-bit if using extreme testcases.
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