

✓ LeetCode 1497 — Check If Array Pairs Are Divisible by k

📌 Problem

Given an array `arr` of even length and integer `k`, check if you can pair elements such that:

```
[
(arr[i] + arr[j]) % k == 0
]
```

Return `true` if possible, otherwise `false`.

🎯 Key Insight

For each number `x`, use remainder:

```
[
r = (x % k + k) % k
]
```

A pair is valid only if:

```
[

$$r + r_2 = k \quad \Rightarrow \quad r_2 = k - r$$

]
```

🧠 Remainder Pairing Rules

Case	Condition
<code>r == 0</code>	frequency must be even
<code>k</code> is even & <code>r == k/2</code>	frequency must be even

Case	Condition
Otherwise	<code>freq[r] == freq[k-r]</code>

Algorithm

1. Convert all elements to safe remainders
2. Count frequency of each remainder
3. Validate pairing rules

Time / Space Complexity

Complexity	Value
Time	O(n)
Space	O(k)

Code

```
class Solution {
public:
    bool canArrange(vector<int>& arr, int k) {
        vector<int> modified(arr.size(), 0);

        // Convert to positive remainders
        for(int i = 0; i < arr.size(); i++) {
            modified[i] = ((arr[i] % k) + k) % k;
        }

        unordered_map<int, int> mp;

        // Count frequency of each remainder
        for(int r : modified) {
            mp[r]++;
        }
    }
};
```

```

// Validate pairs
for(auto &ele : mp) {
    int r = ele.first;
    int freq = ele.second;

    // Case 1: remainder 0 must be even
    if(r == 0 && freq % 2 != 0) return false;

    // Case 2: if k even & r = k/2
    if(k % 2 == 0 && r == k/2 && freq % 2 != 0) return false;

    // Case 3: freq[r] == freq[k-r]
    if(r != 0 && r != k-r) {
        if(mp[r] != mp[k-r]) return false;
    }
}

return true;
}
};

```

Example

Input

arr = [1,2,3,4,5,10,6,7,8,9], k = 5

Remainders

1→1, 2→2, 3→3, 4→4, 0→5

Valid pairs

(1,4), (2,3), (0,0) ...

✓ Output: `true`



Notes

- Works for negative numbers
- Greedy pairing using remainders avoids $O(n^2)$

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📌 Problem Summary

Given an array `arr` of even length `n` and an integer `k`, check if it's possible to divide the array into `n/2` pairs such that **sum of each pair is divisible by `k`**.

✨ Key Idea

For every number `x`, calculate its remainder `r = x % k`.

To form valid pairs:

- `(r + r2) % k == 0` → `r2 = k - r`

🧠 Important Rules

Remainder	What to Check
<code>r = 0</code>	Count of such numbers must be even
<code>r = k/2</code> (only when k is even)	Count must be even
Other <code>r</code>	Count of remainder <code>r</code> must equal count of <code>k-r</code>

Also handle negative numbers properly using:

$$((x \% k) + k) \% k$$



Approach

1. Convert all numbers to positive remainders
2. Count frequency of each remainder
3. Validate rules above



Your Code Logic Breakdown

✓ Convert to safe remainders

```
modified[i] = ((arr[i] % k) + k) % k;
```

✓ Count frequency using map

```
mp[rem]++;
```

✓ Check remainder pairing rules

- Remainder `0` → must be even
- For remainder `r` → check `freq[r] == freq[k-r]`



Time & Space Complexity

Complexity	Value
Time	O(n)
Space	O(k) (map for remainder counts)



Example

Input:

```
arr = [1,2,3,4,5,10,6,7,8,9], k = 5
```

Remainders count:

1→1, 2→2, 3→3, 4→4, 0→5

Pairs:

1 & 4, 2 & 3, 0 & 0...

✓ Valid

Edge Cases

- Negative numbers
- k even → check $k/2$ remainder
- All elements divisible by k

Final Conclusion

We match remainders such that the **sum becomes multiple of k** .

Remainder pairing technique solves cleanly in **$O(n)$** .