

# LeetCode: 238

## ✓ Problem Recap

We need an array `answer` such that:

$$\text{answer}[i] = \prod_{j=0, j \neq i}^{n-1} \text{nums}[j] \quad \text{answer}[i] = \prod_{\substack{j=0 \\ j \neq i}}^{n-1} \text{nums}[j]$$

Constraints:

- **$O(n)$**  time.
- **No division operation allowed.**
- **Prefer  $O(1)$  extra space** (excluding output array).

## ✓ Your Approach

### Step 1: Initialization

```
int product = 1; // product of all elements
int p2 = 1;     // product of all non-zero elements
int noz = 0;    // count zeros
```

### Step 2: First loop – Calculate products

```
for(int i=0; i<n; i++){
    if(arr[i] == 0) noz++;
    product *= arr[i]; // overall product (includes zeros)
    if(arr[i]!=0) p2 *= arr[i]; // product ignoring zeros
}
```

- `product` → full product including zeros.
- `p2` → product excluding zeros (useful when zeros exist).
- `noz` → counts how many zeros are present.

### Step 3: Handle zeros and division

```
if(noz>1) p2 = 0; // more than 1 zero → all products are zero
```

```
for(int i=0;i<n;i++){  
    if(arr[i] != 0) arr[i] = product / arr[i]; // division step  
    else arr[i] = p2;  
}
```

- If there are:
  - **0 zeros** → use `product / arr[i]`.
  - **1 zero** → only index with zero gets `p2`, others get `0`.
  - **>1 zero** → everything is `0`.

## ✓ Why is this method logically correct?

- **Mathematical logic works perfectly:**
  - For `nums[i] ≠ 0`,  
$$\text{answer}[i] = \frac{\text{total product}}{\text{nums}[i]}$$
  - For `nums[i] = 0`,  
$$\text{answer}[i] = \text{product of all non-zero elements}$$
- Handles **zero cases correctly**.
- Runs in **O(n)**.

## ✗ Why is it wrong for this problem?

### 1. Uses Division

- Problem explicitly says:

You must write an algorithm without using the division operation.

### 2. Doesn't meet the "Follow-up" constraint

- Extra variables (`product`, `p2`, `noz`) → still **O(1) extra space**, so that's fine.

- But **core violation** = usage of division.
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## ✓ Why Division is Not Allowed?

- **Risk of overflow** if the product of all numbers is huge (although 32-bit limit helps here).
  - **Division by zero** problem complicates direct approach.
  - Forcing no-division makes the solution algorithmically interesting → requires **prefix & suffix products**.
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## ✓ Correct Approach (as per constraints)

Idea:

- For each index `i`,

`answer[i]` =

$(\text{product of all elements to the left of } i) \times (\text{product of all elements to the right of } i)$   
`answer[i]`  
 $= (\text{product of all elements to the left of } i) \times (\text{product of all elements to the right of } i)$

- Build **prefix product** and **suffix product** arrays, or compute in two passes with constant extra space.
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## Steps

1. Initialize `answer` array with `1`.

2. **First pass (left to right)** → store prefix product:

`answer[i]` = product of all elements before `i`  
`answer[i] = \text{product of all elements before } i`

3. **Second pass (right to left)** → multiply by suffix product:

`answer[i]` \*= product of all elements after `i`  
`answer[i] *= \text{product of all elements after } i`

**Time:**  $O(n)$

**Space:**  $O(1)$  extra (if we ignore output array).

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## ✓ Mathematical Expression for Correct Method

$$\text{answer}[i] = (\prod_{j=0}^{i-1} \text{nums}[j]) \times (\prod_{j=i+1}^{n-1} \text{nums}[j])$$
$$\text{answer}[i] = \left( \prod_{j=0}^{i-1} \text{nums}[j] \right) \times \left( \prod_{j=i+1}^{n-1} \text{nums}[j] \right)$$

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## ✅ Explanatory Notes (Summary)

### Your method:

- ✅ Correct mathematically.
- ❌ Violates problem constraint (used division).
- ✅ Handles zeros well.

### Required method:

- Use **prefix & suffix multiplication**.
- No division allowed.

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🔥 Do you want me to **create a clean, well-commented solution for the prefix-suffix method**, along with a **PDF explanatory note comparing both approaches with diagrams and formulas**? Or should I **just give you the commented code and summary in text format here**?