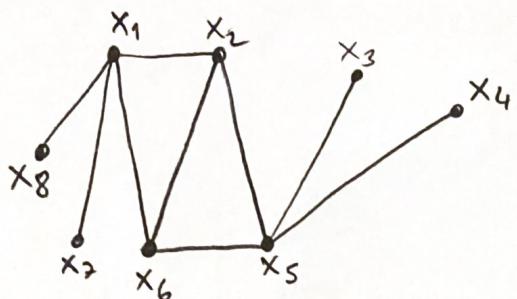


6.042 - Lecture 7 - Matching Problems

Def Given graph $G = (V, E)$ a matching is a subgraph of G where every node has degree 1.

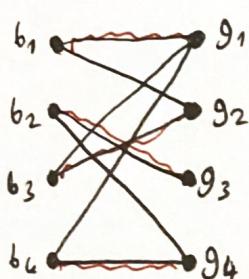
Ex



$\{x_1-x_6, x_2-x_5\}$ is a matching of size 2

Def A matching is perfect if it has size $\frac{|V|}{2}$

Ex

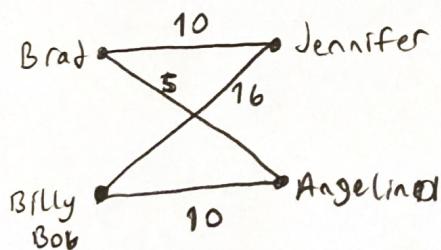


$\{b_1-g_1, b_2-g_2, b_3-g_3, b_4-g_4\}$ is a perfect matching

Def The weight of a matching M is the sum of the weights on the edges in M

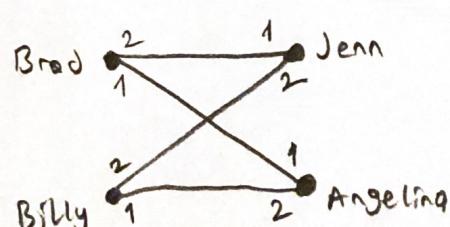
Def A min-weight matching for G is a perfect matching for G with the min weight.

Ex



Brad - Jenn
Billy - Angelina } 20

Def Given a matching M , $x \& y$ form a rogue couple if they each other to their mates in M



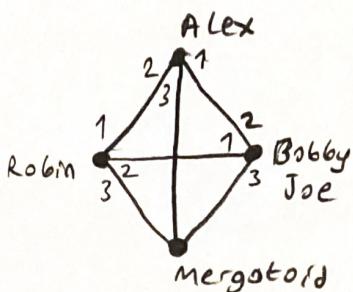
If we married Brad to Jenn and Billy to Angelina, Brad and Angelina form a rogue couple because they like each other better than who they were hooked up with.

Def A matching is stable if there are no rogue couples.

Goal: Find a perfect matching that is stable.

Thm \rightarrow stable matching for this graph

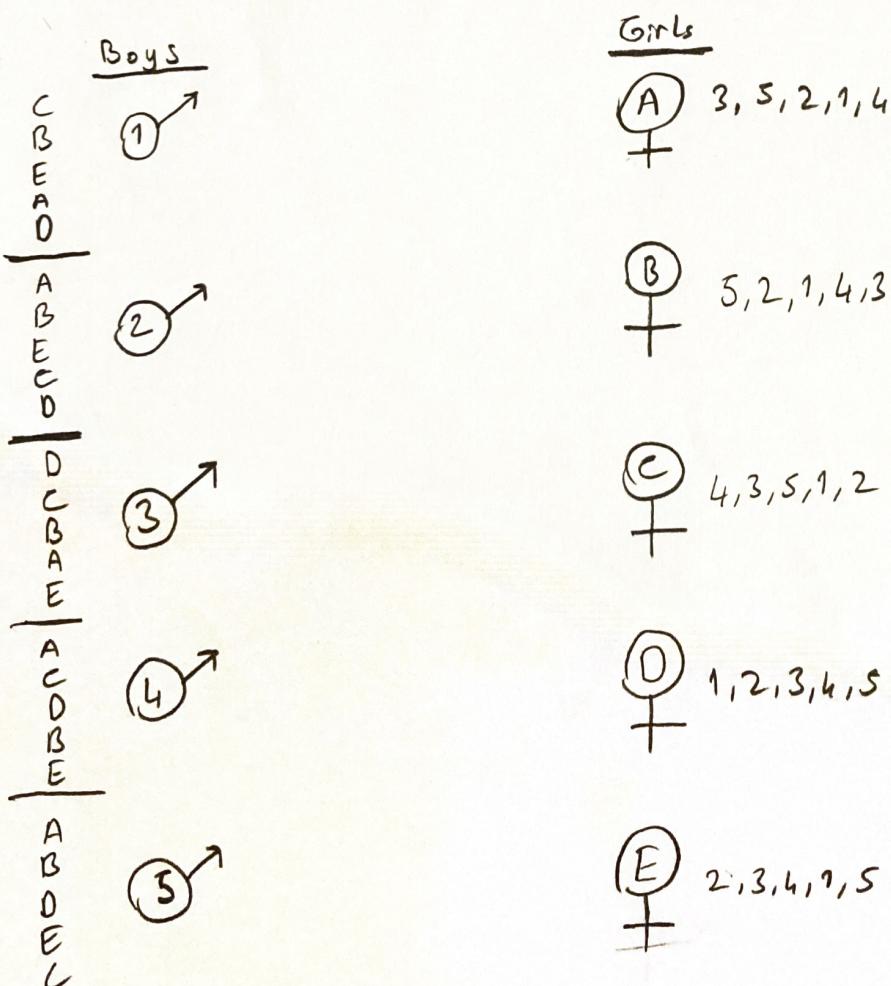
PF By contradiction. Assume \exists stable matching M
 Then Mergatroid match w/ someone in M.
 WLOG (by symmetry), assume Mergatroid matched
 to Alex. Alex & Robin form a rogue couple for M.



* Without Loss of Generality (WLOG) means all of the choices will end up with same result, so choose and proof one choice will be enough.

stable Marriage Problem

Goal: Find perfect matching w/o rogue couples



Serenaders

Girls	Day 1	Day 2	Day 3	Day 4
A	2, 4, 5	5	5	5
B	-	2	2, 1	2
C	1	1, 4	4	4
D	3	3	3	3
E	-	-	-	1

Crossouts

Boys

Boys
1
2
3
4
5

C B

A

C

Need to show:

- TMA terminates (quickly)
- Everyone get married
- No rogue couples
- Fairness

Thm 1] TMA terminates in $\leq N^2 + 1$ days.

Pf] By contradiction. Suppose TMA does not terminate in $N^2 + 1$ days.

Claim: If we don't terminate on a day, then some boy crosses a girl off his list that night

N list with N names $\Rightarrow \leq N^2$ crossouts $\neq N^2 + 1$ crossouts \square

Let $P ::=$ "If a girl G ever rejected a boy B , then G has a suitor who she prefers to B ."

Lemma 1] P is an invariant for TMA

Pf] By induction on # days

Base Case: Day 0. No one is rejected yet \Rightarrow vacuously true

Ind. Step: Assume P holds at the end of day d

Case 1: G rejects B on day $d+1$. Then there was SMA, a better boy
 $\Rightarrow P$ true at $d+1$

Case 2: G rejected B before $d+1$.

$P \Rightarrow G$ had a better suitor on day d

$\Rightarrow G$ has same or better suitor on $d+1 \Rightarrow P$ true on $d+1 \square$

(3)

Thm 2) Everyone is married in TMA

Pf) By contradiction. Assume that some boy B is not married at the end.

$\Rightarrow B$ rejected by every girl

\Rightarrow Every girl has better suitor (Lemma 1)

\Rightarrow Every girl married

\Rightarrow Every boy married \rightarrow Contradiction \square

Thm 3) TMA produces a stable matching

Pf) Let Bob & Gail be any pair that are not married

Case 1: Gail rejected Bob

\Rightarrow Gail marries smn. better than Bob (Lemma 1)

\Rightarrow Gail & Bob are not rogue

Case 2: Gail did not reject Bob

\Rightarrow Bob never serenaded Gail

\Rightarrow Gail is lower on Bob's list than Bob's wife

\Rightarrow Bob & Gail are not rogue

\Rightarrow No rogue couples \rightarrow contradiction \square

Let S = set of all stable matchings $S \neq \emptyset$

For each person P , we define the realm of possibility for P to be

$$\{Q \mid \exists M \in S, \{P, Q\} \subseteq M\}$$

Def) A person's optimal mate is his/her favorite from the realm of possibility

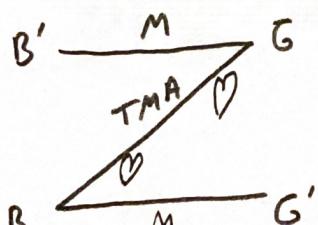
Def) " " " pessimal " " " least favorite " " " " "

Thm 4) TMA marries every boy w/ his optimal mate

Thm 5) " " " girl " " her pessimal "

Pf of Thm 5) By contradiction

Suppose that \exists stable matching $M \ni$ girl G who fares worse than TMA



$B \& G$ rogue in $M \rightarrow$ contradiction
 $\Rightarrow M$ is not stable \square