

while $\text{sizes}(L) < 0$: $L = -\text{sizes}(L)$

Write L to the LABEL lattice, sizes(L) \rightarrow 1
- both are occupied :

find both root labels, take the smaller
and unify them

$$(L_1 \sqsubset L_2 : \quad \text{SIZES}(L_1) \pm = \text{SIZES}(L_2) + 1 \\ \text{SIZES}(L_2) = -L_1)$$

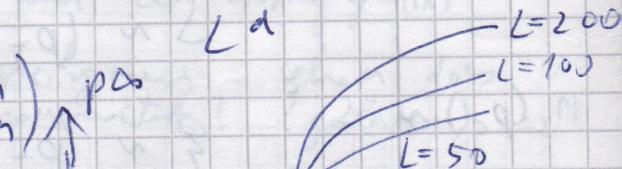
at the end the `sizes` array has what we need.

~~fill table rows~~: TRICK: do this alg as occupations
 are generated.
 we can forget old rows. (we need L^{d-1} storage)

$$\left\{ \begin{array}{l} \text{in } d=2 \quad p = \frac{1}{2} \text{ for triangular bath} \\ p = 0.5927 \text{ square bath} \end{array} \right\} \quad \overbrace{\quad}^{\alpha = 4/3} \quad \overbrace{\quad}^{\beta = 5/36}$$

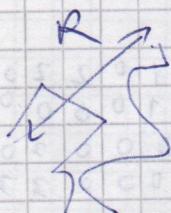
for finite lattice e.g. $P_{\infty} = \frac{\text{points in largest cluster}}{V^d}$

(Generalizations bond percolation) \uparrow
directed percolation



Random walks

Conformation of a polymer chain



$$R \sim N^{\nu}$$

Simplest model: Random walk

$\rightarrow R \sim t^{1/2}$ disagrees with experiment.

No back track RW : do not intend turn back

SAR, self avoiding random walk: do not go to a site already visited (monomers repulse)

What should be the weight of a Random walk?

Re: all walks of length N are equally probable

Saws = — || — as long as they are Saws

\Rightarrow generate RD's, throw away if they intersect
 works for small N , otherwise most RD's are thrown away (attrition problem)

Reptation algorithm:

1. Assume we have a walk.
2. delete one point from a random end
3. add one point to a random end with random direction. ~~accept the new walk if it is SAW
otherwise use old as new config.~~

Pivot algorithm Use lattice symmetries

$$V = \frac{3}{4} \quad d=2$$

$$V = 0.59 \dots \quad d=3$$

$$V = \frac{1}{2} \quad d=4$$

Check how many steps are available
 N (no backtrace)

$$N = \prod_{i=1}^d \frac{L_i}{(2d-1)} \Leftarrow \text{number of steps possible at step } i$$

Generalization: potential between monomers.

now different scales have different weights.

Importance Sampling

Usual question: $\langle F \rangle = \frac{1}{Z} \int F(p, q) e^{-\beta H(p, q)} dp dq$

$e^{-\beta H(p, q)} \approx 0$ for most configurations

$$\langle F \rangle = \int F(E) \cdot \rho(E) \cdot e^{-\beta E} dE$$

$\rho(E) \cdot e^{-\beta E}$ is sharply peaked

uniform random configurations have too little weight ($\rho \approx 0$) or too large Energy $\rho(E) e^{-\beta E} \approx 0$

if we sample with prob. $e^{-\beta H(p, q)}$

$$\langle F \rangle = \frac{1}{N} \sum_i F(p_i, q_i)$$

Metropolis algorithm

Markov chain

i - computer time

$$\phi_i \rightarrow \phi_{i+1} \rightarrow \phi_{i+2} \dots$$

$$\text{prob. to go from } \phi \text{ to } \phi' = T(\phi'| \phi)$$

T does not depend on i , only on ϕ not older $\phi_s \Rightarrow$

\Rightarrow Markov process

$$P(Q) = \text{prob. of state } Q$$

Master eq.:

discretised

$$\frac{\partial P(Q, t)}{\partial t} = \sum_{Q'} w(Q \leftarrow Q') P(Q') - \sum_{Q'} w(Q' \leftarrow Q) P_Q$$

$$P_j^{(n+1)} = T_{jk} P_k^{(n)}$$

$$P_j^{(n+1)} - P_j^{(n)} = T_{jk} P_k^{(n)} - P_j^{(n)} =$$

in equilibrium

$$P_{i+1}(\phi) = \sum_{\phi'} T(\phi | \phi') P_i(\phi) \quad \begin{matrix} \cancel{T(\phi' | \phi)} \\ \cancel{P_i(\phi')} \end{matrix}$$

$$\sum_{\phi'} T_{ik} P_k^{(n)} - \sum_{k'} T_{ki} P_j^{(n)}$$

Balance equation
no source or sink

$$T(\phi) = \sum_{\phi'} T(\phi | \phi') \pi(\phi') = \underbrace{\sum_{\phi'} T(\phi | \phi') \pi(\phi')}_{\phi'} \underbrace{\pi(\phi)}_{\phi}$$

In Matrix notation

$$T_{ik} \phi_k = \phi_i$$

eigenvector with e.v. = 1

Let's choose $T(\phi | \phi)$ such that we have
detailed balance

$$T(\phi' | \phi) \pi(\phi) = T(\phi | \phi') \pi(\phi')$$

current is the same in both directions

$$T_{ij} \phi'_j = T_{ji} \phi'_i$$

consider the evolution of $P_0 \xrightarrow{T} P_1 \xrightarrow{T} P_2 \xrightarrow{T} \dots P^{(n)} \xrightarrow{T} P^{(n+1)} \xrightarrow{T} P^{(n+2)} \xrightarrow{T} \dots$

does it converge to π ?

$$\bar{\sigma}_N = \sum_i |P_i^{(n)} - \pi_i|$$

$$\bar{\sigma}_{N+1} = \sum_i |T_{ij} P_j^{(n)} - \pi_i| = \sum_i |T_{ij} P_j^{(n)} - T_{ji} \pi_j| =$$

$$= \sum_i \left| \sum_j (T_{ij} P_j^{(n)} - T_{ji} \pi_j) \right| \leq \sum_{i,j} |T_{ij}| (P_j^{(n)} - \pi_j) = \bar{\sigma}_N$$

$$\left. \begin{aligned} \sum_i P_i^{(n)} - \pi_i &= 0 \\ T_{ij} &> 0 \end{aligned} \right\} \Rightarrow \text{equality only if } \bar{\sigma}_N = 0$$

\Rightarrow convergence to unique eq. π

$\forall i, j \in N \quad T_{ij} > 0 \text{ irreducible}$
 $\forall i, j \in N \quad T_{ij} > 0 \text{ aperiodic}$
 $\forall i \in N \quad (T^N)_{ii} > 0$

Metropolis alg.

starting from state ϕ_0 . After $n-1$ steps, do

1. proposal $T_0(\phi' | \phi_{n-1})$
2. Acceptance probability $T_A(\phi' | \phi_n)$
if accepted $\phi_n = \phi'$
else $\phi_n = \phi_{n-1}$

To satisfy detailed balance, we must have

$$T_0(\phi | \phi') \cdot T_A(\phi' | \phi) \cdot \pi(\phi') = T_0(\phi' | \phi) \cdot T_A(\phi | \phi') \cdot \pi(\phi)$$

$$\frac{T_0(\phi | \phi') T_A(\phi | \phi')}{T_0(\phi' | \phi) T_A(\phi' | \phi)} = \frac{\pi(\phi)}{\pi(\phi')}$$

$$\frac{T_A(\phi | \phi')}{T_A(\phi' | \phi)} = \frac{\pi(\phi)}{\pi(\phi')} \cdot \frac{T_0(\phi' | \phi)}{T_0(\phi | \phi')} = F(\phi | \phi')$$

we can choose e.g. $T_A(\phi | \phi') = \begin{cases} F(\phi | \phi') & \text{if } F(\phi | \phi') < 1 \\ 1 & \text{otherwise} \end{cases}$

$$\text{or } T_A(\phi | \phi') = \frac{F}{1 + F}$$

Usually $T_0(\phi | \phi') = T_0(\phi' | \phi)$ reversibility of proposal

$$\pi(\phi) = e^{-S[\phi]}$$

$$\frac{\pi(\phi)}{\pi(\phi')} = e^{-(S(\phi) - S(\phi'))} = e^{-\Delta S}$$

In this case

$$\text{acceptance prob} = \begin{cases} e^{-\Delta S} & \text{if } \Delta S > 0 \\ 1 & \text{if } \Delta S < 0 \end{cases}$$

$T_0(\phi | \phi')$ should be optimized

change too big \Rightarrow acceptance small
change too small \Rightarrow conf. space exploration is slow

e.g. for Ising model

local change $s_i \rightarrow -s_i$ OK for most cases
cluster alg: change many spins at once except near critical phase transition