# PRIOR GEOMETRY: EXTENSIVE FORMALIZATION PART 3

## SECTION 10: ADDITIONAL FORMAL DEFINITIONS AND PROPERTIES

#### 10.1 Minimal Increment Transitions

- (10.1.1) Define a minimal increment transition between patterns x and y at (a, e) if  $d_{(a,e)}(x, y) = 1$ . That is, they differ by exactly one increment of probability distribution.
- (10.1.2) Construct a graph  $G_{(a,e)}$  with vertices as stable patterns and edges representing minimal increment transitions. This graph encodes geometric adjacency.

## 10.2 Polygonal Loops and Inside/Outside

- (10.2.1) A polygonal loop  $C = (x_1, \ldots, x_n)$  is a simple cycle in  $G_{(a,e)}$ .
- (10.2.2) Removing edges of C from  $G_{(a,e)}$  splits the graph into components. Designate one reference pattern  $x_{\text{ref}}$  known as "outside." Patterns not reachable from  $x_{\text{ref}}$  after removing C are "inside" the loop.
- (10.2.3) Thus, polygonal loops define finite analogs of planar curves and inside/outside regions, purely combinational and finite.

#### 10.3 Refinement of Numeric Scales

- (10.3.1) Consider two numeric scales:  $\{0, 1/L, \ldots, 1\}$  and  $\{0, 1/M, \ldots, 1\}$ . To combine them, find capacity a' with increments 1/R where  $R \ge \operatorname{lcm}(L, M)$ . Both scales embed into the finer scale.
- (10.3.2) This embedding shows how rational increments from different intervals can be reconciled into a common refined scale.

## 10.4 Negative Increments and Symmetry

- (10.4.1) Define negative increments by choosing a chain in the opposite direction. For any increment k/L, a corresponding pattern chain in the reverse order defines -k/L.
- (10.4.2) Thus, negative increments appear as reversible steps along stable pattern chains, ensuring that if we interpret increments as "positions," negative increments represent going backward.

## 10.5 Representing Integers and Basic Arithmetic

- (10.5.1) Integers appear as repeated increments of a chosen unit increment 1/N(a). For example, 3\*(1/N(a)) is just the third pattern after the start in a chain.
- (10.5.2) Since all arithmetic is finite, representing large integers requires longer chains or higher capacity states.

## 10.6 Algebraic Closure and Finite Domains

- (10.6.1) Within a given capacity a, the set of representable increments is finite. Addition and multiplication remain closed within this finite set of rational increments.
- (10.6.2) If a desired rational number cannot be represented exactly at capacity a, refine to a' with larger N(a') until representation is possible. If no capacity within  $\Omega$  can represent it exactly, we approximate it by the closest available increment.

## 10.7 Actions on Patterns and Dynamic Loops

- (10.7.1) Changing  $(a, e) \to (a', e')$  modifies  $\mu$  and  $p_x$  distributions, potentially altering which patterns are stable and the structure of lines and polygons.
- (10.7.2) This dynamical aspect shows that geometry and arithmetic are not fixed but depend on the joint state and can evolve over time.

## 10.8 Comparison with Classical Infinite Concepts

- (10.8.1) Classical number theory relies on infinite sets and limits. Here, no actual infinity is present. All sets are finite at each step.
- (10.8.2) Approximations to classical concepts like real numbers, continuity, or measure theory are done by capacity refinement and selection of increasingly large N(a).
- (10.8.3) Without infinite sets, no perfect realization of  $\pi, e$ , or irrational numbers occurs. Instead, we get sequences of rational increments approximating them to arbitrary precision, constrained by  $\Omega$ .

## SECTION 11: PROOF SKETCHES AND CONSISTENCY ARGUMENTS

## 11.1 Consistency with Probability Axioms

- (11.1.1) Probability measures  $\mu^{(a,e)}(x)$  and distributions  $p_x^{(a,e)}(g)$  are finite rational numbers.
- (11.1.2) Normalization conditions are easily met by assigning uniform or arbitrary rational distributions that sum to 1, requiring no infinite processes.

## 11.2 Updating without Contradictions

- (11.2.1) For each  $(a, e) \to (a', e')$ , define  $U_{\mu}$  and  $U_{p}$  as simple piecewise functions or lookup tables. Since A and E are finite, and increments are finite rationals, no contradictions arise.
- (11.2.2) Probability updates remain finite transformations, ensuring internal consistency at every step.

## 11.3 Constructing a Model

- (11.3.1) Start with minimal sets:  $A = \{a1\}, E = \{e1\}, X = \{x1\}, G_{x1} = \{g11, g12\}, \text{ etc.}$
- (11.3.2) Assign  $\mu^{(a1,e1)}(x1) = 1$  and  $p_{x1}^{(a1,e1)}(g11) = 1$ , trivial stable pattern.
- (11.3.3) Add more states, patterns, and features step by step, always ensuring finite rational increments. This yields at least one non-empty model, proving the system is consistent.

## SECTION 12: APPLICATION TO INTEGER SCALES AND RATIONALS

## 12.1 Integer Approximations

- (12.1.1) Fix capacity a. The smallest increment is 1/N(a). Counting k increments from a baseline pattern defines k/N(a).
- (12.1.2) If N(a) = 10, for example, increments are multiples of 0.1. Integers like 1,2,3 appear as 10/10, 20/10, 30/10 but must remain within representable sets.

## 12.2 Rational Density under Refinement

- (12.2.1) Increasing capacity state  $a \to a'$  with N(a') = N(a) \* Q introduces finer increments of size 1/(N(a) \* Q), making the rationals denser.
- (12.2.2) Thus, by repeated refinements (bounded by  $\Omega$ ), we simulate rational density up to the complexity limit.

## 12.3 Operations with Minimal Error

- (12.3.1) For any target rational number r, we can approximate r by k/N(a') at a sufficiently high capacity a'.
- (12.3.2) The error  $|r k/N(a')| < \epsilon$  if we choose a' so  $N(a') > k/(\epsilon)$ . This process requires no infinite limit, just a finite large N(a').

## SECTION 13: STRUCTURED COMPLEXES AND MAPPINGS

## 13.1 Complexes of Patterns

- (13.1.1) Define a finite set of patterns arranged so that minimal increment edges form higher structures like cubes, tetrahedra, etc.
- (13.1.2) Each complex remains finite, and each topological property (connectedness, cycles) is computed combinationally.

## 13.2 Maps Between Pattern Sets

- (13.2.1) A map  $F: X \to X$  between pattern sets can be defined increment-wise. If increments correspond to rational numbers, F can represent linear transformations over finite vector spaces.
- (13.2.2) Changing capacity refines these maps, allowing finer approximations of linear maps and eigen decompositions.

## SECTION 14: NO SOLIPSISM PRINCIPLE

## 14.1 Joint Dependence on (a,e)

- (14.1.1) Probability structures  $p_x^{(a,e)}$  depend on both observer state a and environment e. This prevents solipsism: distributions are not purely subjective but constrained by external signals.
- (14.1.2) As  $(a, e) \rightarrow (a', e')$ , environment changes force recalculation of probabilities, ensuring external grounding.

## 14.2 Capacity and External Constraints

- (14.2.1) Higher capacity states reveal finer distinctions in signals  $S_e$ . If  $S_e$  changes drastically, patterns compatible with  $S_e$  must also change.
- (14.2.2) This coupling ensures that the finite increments are always tied to external conditions, not abstractly chosen.

## SECTION 15: STABILITY AND EMERGENT REG-ULARITIES

#### 15.1 Attractors as Stable Patterns

- (15.1.1) If repeated environmental interactions lead to certain x always having high  $\mu^{(a,e)}(x)$  and stable  $p_x^{(a,e)}$ , x is an emergent stable pattern.
- (15.1.2) Such attractors represent geometric objects: points, lines, etc., formed by consistent probability assignments.

#### 15.2 Parameter Variation

- (15.2.1) Varying capacity systematically reveals more subtle patterns and increments.
- (15.2.2) At low capacity, multiple stable patterns might appear merged as one pattern. At high capacity, they differentiate into several distinct stable patterns.

## SECTION 16: SYNTHESIS OF THE FORMAL SYSTEM

## 16.1 Summary

- Finite sets  $A, E, X, G_x$
- Probability increments quantized by N(a)
- Capacity V(a) controls granularity
- Stability defines geometric-like entities
- Refined increments yield rational arithmetic
- Complex algebraic structures built from finite increments
- Approximation of classical concepts via capacity refinement
- Finite complexity  $\Omega$  forbids true infinity

#### 16.2 Concluding Formal Remarks

- (16.2.1) The entire framework is self-contained, finite, and relies on no infinite sets, no limits, and no continuity arguments.
- (16.2.2) By careful selection of increments, chains, and distributions, all classical-like structures (numbers, geometry, algebra) appear as finite approximations.
- (16.2.3) This ensures a consistent, finite, rational-based mathematical universe where complexity and capacity guide the emergence of structured concepts.