# Design and Implementation of Anglican Probabilistic Programming Language

David Tolpin Jan Willem van de Meent Hongseok Yang Frank Wood

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https://bitbucket.org/probprog/anglican-white-paper https://bitbucket.org/probprog/anglican http://www.robots.ox.ac.uk/~fwood/anglican/index.html



## Outline

#### Motivation

Design Outline

Implementation Highlights

Inference Algorithms

Definitions and Runtime Library

#### Intuition

## Probabilistic program:

- A program with random computations.
- Distributions are conditioned by 'observations'.
- Values of certain expressions are 'predicted' the output.

Can be written in any language (extended by sample and observe).

## Example: Model Selection

```
(let [;; Guessing a distribution
1
          dist (sample (categorical
2
                           [[normal 1] [gamma 1]
3
                            [uniform-continuous 1]
4
                            [uniform-discrete 1]]))
5
          a (sample (gamma 1 1))
6
          b (sample (gamma 1 1))
7
          d (dist a b)]
8
      ;; Observing samples from the distribution
9
      (loop [data data]
10
        (when (seq data)
11
          (let [[x & data] data]
12
            (observe d x))
13
          (recur data)))
14
      ;; Predicting a, b and the distribution
15
      (predict :a a)
16
      (predict :b b)
17
      (predict :d d))
18
```

# More examples

▶ Intruder detection — given a log of **times** and **amounts** of payments in a bank account, how likely that the baccount was compromised?

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- ➤ Counterfactual reasoning There are **two routes** from Jerusalem to Tel Aviv: 1 and 443. Based on traffic reports, I chose route 1 and was late. Would I arrive on time If I chose 443 instead?

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- Counterfactual reasoning There are two routes from Jerusalem to Tel Aviv: 1 and 443. Based on traffic reports, I chose route 1 and was late. Would I arrive on time If I chose 443 instead?
- ► (Due to Stuart Russell) If you observe that a student GPA is exactly 4.0 in a model of transcripts of students from the USA (GPA's from 0.0 to 4.0) and India (GPA's from 0.0 to 10.0) what is the probability that the student is from India?

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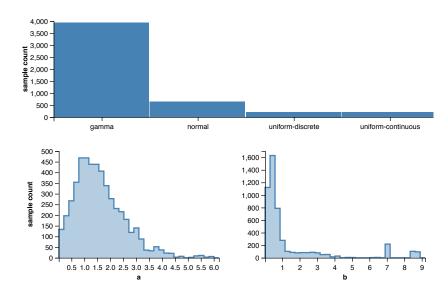
# Inference Objective

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 Continuously and infinitely generate a sequence of samples drawn from the distribution of the output expression
 — so that someone else puts it in good use (vague but common). ✓

# Example: Inference Results



## Importance Sampling

### loop

Run program, computing weight based on observations. Output result and weight.

## end loop

- ► Simple good.
- ▶ Slow convergence (unless one knows the answer) bad.

Can we do better?

# Lightweight Metropolis-Hastings (LMH)

Run program once, remembering random choices.

## loop

Uniformly select one random choice.

Propose a new value for the choice.

Re-run the program.

Accept or reject with MH probability.

Output result.

## end loop

#### Can we do better?

- Particle methods
- Variational inference
- **.**..

# Why functional?

We want a functional language because an inference algorithm controls the execution:

- ▶ A program is run many (often many hundreds of thousands) of times (with almost any algorithm).
- ▶ A program must be partially re-executed multiple times from different positions (particle methods).
- We want to reason about the distribution defined by the program.

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- Scheme (Church, Venture).
- Scala Figaro.
- Haskell Hakaru, Model-Bayes.
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As well as Python, C#, and other languages.

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- ▶ Interpeter.
- Source-to-source compiler.
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#### Anglican is

- Integrated with Clojure.
- Shares syntax.
- Alters operational semantics.

## Anglican

## A subset of Clojure, wrapped inside defquery:

- ▶ if, when, cond, case, let, and, or, fn.
- Vector destructuring in bindings of let and fn.
- Compound literals for vectors, hash maps, and sets.
- ▶ loop/recur a convenience.

#### Core library:

- ▶ All of Clojure core library, except for higher-order functions.
- ▶ map, reduce, filter, some, repeatedly, comp, partial.

Any Clojure function can be called from Anglican.

## Macro-based compilation

#### Anglican code macro-compiled into Clojure:

```
(fn loop [C23151 state data]
                                            (if (seq data)
    (loop [data data]
1
                                              (let [[x & data] data]
      (if (seq data)
                                                (->observe '023153 (gamma a b) x
        (let [[x & data] data]
                                                   (fn do23152 [_ state]
          (observe (gamma a b) x)
                                                     (loop C23151 state data))
          (recur data))
                                                  state))
        (predict :a a)))
                                              (fn []
                                                (C23151
                                                  nil
                                    10
                                                   (add-predict state :a a)))))
                                    11
```

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## Probabilistic forms

#### Two **probabilistic** forms:

- sample draws a value from a distribution;
- observe conditions the a posteriori distribution by observing a value from a distribution.

## This is where the inference algorithm steps in:

```
=> (cps-of-expression '(sample dist) 'cont)
(->sample dist cont $state)

=> (cps-of-expression '(observe dist val) 'cont)
(->observe dist val cont $state)
```

## Memoization

#### Memoization of random choice:

- A person has a random eye color.
- ▶ But the *same* person has a fixed eye color.

## Compiling a memoized function

```
(fn []
                               (cont
                                ;; every memoization gets a unique key
                                (let [M23145 (gensym "M")]
                                 (fn [C23144 $state & P23147]
                                  (if (in-mem? $state M23145 P23147)
                                   ;; previously memoized result
                                   (fn []
                                     (C23144)
                                       (get-mem $state M23145 P23147)
                         10
                                       $state))
(mem (fn [person] ...))
                                   ;; new computation
                         12
                                   (clojure.core/apply
                         13
                                    (fn [C23150 $state person]
                         14
                                      (fn □ (C23150 ... $state)))
                         15
                                    ;; memoize result in state
                         16
                                    (fn [V23146 $state]
                         17
                                     (fn []
                         18
                                      (C23144
                         19
                                        V23146
                         20
                                         (set-mem $state M23145 P23147 V23146))))
                         21
                                    $state
                         22
                                    P23147))))
                         23
                         24
                                $state))
```

## Managing stack size

- Clojure does not support tail-call optimization (TCO).
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# Managing stack size

- Clojure does not support tail-call optimization (TCO).
- Under CPS transformation the stack will explode.
- Anglican uses trampolining: every continuation call is wrapped into a thunk.

The thunk is returned ...

```
1 (fn [cont $state x y]
(fn [x y] (+ x y)) 2 (fn []
3 (cont (+ x y) $state)))
```

... and called by trampoline:

```
(defn exec
[algorithm prog value state]
(loop [step (trampoline prog value state)]
(let [next (checkpoint algorithm step)]
(if (fn? next)
(recur (trampoline next))
next))))
```

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Thank you! Questions?