



Universidad de Chile

Índice

| | |
|---|----------|
| 1. Estructuras de datos | 1 |
| 1.1. Disjoint Set Union (Union-Find) | 1 |
| 1.2. Segment Tree | 1 |
| 1.3. Segment Tree Lazy | 2 |
| 1.4. Segment Tree Persistente | 2 |
| 1.5. Segment Tree Iterativo (compacto) | 3 |
| 1.6. Sparse Table | 3 |
| 1.7. Li Chao Tree (dynamic, persistent) | 4 |
| 2. Grafos | 4 |
| 2.1. LCA (Binary Lifting) | 4 |
| 2.2. LCA (Sparse Table+ETT) | 5 |
| 2.3. Heavy Light Decomposition | 5 |
| 2.4. Dinic Max Flow | 6 |
| 2.5. Centroid Decomposition | 7 |
| 3. Matemáticas | 7 |
| 3.1. Exponenciación binaria | 7 |
| 3.2. Logaritmo discreto | 7 |
| 3.3. Lema de Burnside | 8 |
| 3.4. Fórmula de inversión de Möbius | 8 |
| 4. C++ | 8 |
| 4.1. Custom set/map hash | 8 |
| 4.2. Policy Based Order Statistics Tree | 8 |

1. Estructuras de datos

1.1. Disjoint Set Union (Union-Find)

```
// find y union en  $\sim O(1)$  amortizado
struct DSU{
    vector<int> parent, sz; // sz = size
    DSU(int n){
        parent.resize(n);
        sz.resize(n);
        for(int i=0; i<n; i++){
            parent[i] = i;
            sz[i] = 1;
        }
    }
    int find_set(int v){
        if(v == parent[v]) return v;
        return parent[v] = find_set(parent[v]);
    }
    void union_set(int a, int b){
        a = find_set(a);
        b = find_set(b);
        if(a != b){
            if(sz[a] < sz[b])
                swap(a,b);
            parent[b] = a;
            sz[a] += sz[b];
        }
    }
};
```

1.2. Segment Tree

```
template<class T, T merge(T,T)>
struct segment_tree{
    int N;
    vector<T> tree;
    segment_tree(int _N){
        N = _N;
        tree.resize(4*N);
        build(0, 0, N-1);
    }

    segment_tree(vector<T> &A){
        N = int(A.size());
        tree.resize(4*N);
        build(0, 0, N-1, A);
    }

    void build(int n, int i, int j){
        if(i == j){
            tree[n] = T(); // initial value
            return;
        }
        int mid = (i+j)/2;
        build(2*n+1, i, mid);
        build(2*n+2, mid+1, j);
        tree[n] = merge(tree[2*n+1], tree[2*n+2]);
    }

    T query(int l, int r){
        return query(0, 0, N-1, l, r);
    }

    T query(int n, int i, int j, int l, int r){
        if(l <= i && j <= r) return tree[n];
        int mid = (i+j)/2;
        if(mid < l || r < i)
            return query(2*n+2, mid+1, j, l, r);
        if(j < l || r < mid+1)
            return query(2*n+1, i, mid, l, r);
        return merge(
            query(2*n+1, i, mid, l, r),
            query(2*n+2, mid+1, j, l, r));
    }

    void update(int t, T val){
        update(0, 0, N-1, t, val);
    }

    void update(int n, int i, int j, int t, T val){
        if(t < i || j < t) return;
        if(i == j){
            tree[n] = val;
            return;
        }
        int mid = (i+j)/2;
        update(2*n+1, i, mid, t, val);
        update(2*n+2, mid+1, j, t, val);
        tree[n] = merge(tree[2*n+1], tree[2*n+2]);
    }

    int search(int from, T val){
        if(!from) return search(0, 0, N-1, val);
        return search(0, 0, N-1, val+query(0, from-1));
    }

    int search(int n, int i, int j, T val){
        if(tree[n] < val) return -1;
    }
};
```

```

    if(i==j && tree[n] >= val) return i;
    int mid = (i+j)/2;
    if(tree[2*n+1] >= val) return search(2*n+1, i, mid, val);
    else return search(2*n+2, mid+1, j, val-tree[2*n+1]);
}
};

```

1.3. Segment Tree Lazy

```

template <class T, T merge(T, T)>
struct segment_tree{
    int N;
    vector <T> tree, lazy;
    segment_tree(int _N){
        N = _N;
        tree.resize(4*N);
        lazy.assign(4*N, T()); // modify default value
        build(0, 0, N-1);
    }

    segment_tree(vector <T> &A){
        N = A.size();
        tree.resize(4*N);
        lazy.assign(4*N, T()); // modify default value
        build(0, 0, N-1, A);
    }

    void build(int n, int i, int j){
        if(i == j){
            tree[n] = T(); // initial value
            return;
        }
        int mid = (i+j)/2;
        build(2*n+1, i, mid);
        build(2*n+2, mid+1, j);
        tree[n] = merge(tree[2*n+1], tree[2*n+2]);
    }

    void build(int n, int i, int j, vector <T> &A){
        if(i == j){
            tree[n] = A[i];
            return;
        }
        int mid = (i+j)/2;
        build(2*n+1, i, mid, A);
        build(2*n+2, mid+1, j, A);
        tree[n] = merge(tree[2*n+1], tree[2*n+2]);
    }

    void push(int n, int i, int j){
        // modify this function
        if(lazy[n]){
            tree[n] += lazy[n]*(j-i+1); // range increment
            if(i != j){

```

```

                lazy[2*n+1] += lazy[n];
                lazy[2*n+2] += lazy[n];
            }
            lazy[n] = T();
        }
    }

    T query(int l, int r){
        return query(0, 0, N-1, l, r);
    }

    T query(int n, int i, int j, int l, int r){
        push(n, i, j);
        if(l <= i && j <= r) return tree[n];
        int mid = (i+j)/2;
        if(mid < l || r < i)
            return query(2*n+2, mid+1, j, l, r);
        if(j < l || r < mid+1)
            return query(2*n+1, i, mid, l, r);
        return merge(
            query(2*n+1, i, mid, l, r),
            query(2*n+2, mid+1, j, l, r));
    }

    void update(int l, int r, T val){
        update(0, 0, N-1, l, r, val);
    }

    void update(int n, int i, int j, int l, int r, T val){
        if(l <= i && j <= r){
            lazy[n] += val; // modify this
            push(n, i, j);
            return;
        }
        push(n, i, j);
        if(r < i || j < l) return;
        int mid = (i+j)/2;
        update(2*n+1, i, mid, l, r, val);
        update(2*n+2, mid+1, j, l, r, val);
        tree[n] = merge(tree[2*n+1], tree[2*n+2]);
    }
};

```

1.4. Segment Tree Persistente

```

// Same time complexity as normal SegmentTree
// Additional O(log(n)) memory per update
template <typename T, T merge(T, T)>
struct st_node{
    st_node *left=0, *right=0;
    int i, j;
    T val;
    st_node() {}
    st_node(int _i, int _j) : i(_i), j(_j) {}

```

```

st_node(vector<T> &A){
    int N = int(A.size());
    i = 0, j = N-1;
    build(A);
}

void build(vector<ll> &A){
    if(i == j){
        val = A[i];
        return;
    }
    int mid = (i+j)/2;
    left = new st_node<T,merge>(i, mid);
    right = new st_node<T,merge>(mid+1, j);
    left->build(A);
    right->build(A);
    val = merge(left->val, right->val);
}

st_node *update(int t, ll v){
    if(t < i || j < t){
        return this;
    }
    if(i == j){
        st_node *ret = new st_node<T,merge>(*this);
        ret->val = v;
        return ret;
    }
    st_node *ret = new st_node<T,merge>(i, j);
    ret->left = left->update(t, v);
    ret->right = right->update(t, v);
    ret->val = merge(ret->left->val, ret->right->val);
    return ret;
}

ll query(int l, int r){
    if(l <= i && j <= r) return val;
    int mid = (i+j)/2;
    if(mid < l || r < i) return right->query(l, r);
    else if(j < l || r < mid+1) return left->query(l, r);
    return merge(left->query(l, r), right->query(l, r));
}
};

```

1.5. Segment Tree Iterativo (compacto)

```

template<class T, T m(T, T)> struct iter_seg_tree{
    int n; vector<T> ST;
    iter_seg_tree(vector<T> &a){
        n = a.size(); ST.resize(n << 1);
        for (int i=n;i<(n<<1);i++)ST[i]=a[i-n];
        for (int i=n-1;i>0;i--)ST[i]=m(ST[i<<1],ST[i<<1|1]);
    }
    void update(int pos, T val){ // replace with val

```

```

        ST[pos += n] = val;
        for (pos >= 1; pos > 0; pos >= 1)
            ST[pos] = m(ST[pos<<1], ST[pos<<1|1]);
    }
    T query(int l, int r){ // [l, r]
        T ansL, ansR; bool hasL = 0, hasR = 0;
        for (l += n, r += n + 1; l < r; l >= 1, r >= 1) {
            if (l & 1)
                ansL=(hasL?m(ansL,ST[l++]):ST[l++]),hasL=1;
            if (r & 1)
                ansR=(hasR?m(ST[--r],ansR):ST[--r]),hasR=1;
        }
        if (!hasL) return ansR; if (!hasR) return ansL;
        return m(ansL, ansR);
    }
};
// Example:
iter_seg_tree<int, my_merge_function> st;

```

1.6. Sparse Table

```

// O(nlogn) preprocesamiento, O(1) query en rango
// para función idempotente (como min, max, gcd, etc)
struct sparse_table{
    int n;
    vector<int> logs;
    vector<vector<ll>> table;
    sparse_table(vector<ll> &A){
        n = A.size();
        logs.resize(n+1);
        logs[1] = 0;
        for(int i=2; i<=n; i++){
            logs[i] = logs[i/2] + 1;
        }
        table.assign(logs[n]+1, vector<ll>(n,0));
        for(int i=0; i<=logs[n]; i++){
            int cur_len = 1 << i;
            for(int j=0; j+cur_len-1<n; j++){
                if(cur_len == 1){
                    table[i][j] = A[j];
                }
                else{
                    table[i][j] = min(table[i-1][j], table[i-1][j+cur_len
                        ↪ /2]);
                }
            }
        }
    }
    ll query(int i, int j){
        int p = logs[j-i+1];
        int len = 1 << p;
        return min(table[p][i], table[p][j-len+1]);
    }
};

```

1.7. Li Chao Tree (dynamic, persistent)

```

struct line{
    ll a,b;
    line(){}
    line(ll _a, ll _b) : a(_a), b(_b) {}
    ll eval(ll x) { return a*x+b; }
};

// Dynamic/persistent min Li Chao tree
// Tested on: https://codeforces.com/contest/319/problem/C (add_line,
// ↪ query)
// Tested on: https://www.acmicpc.net/problem/3319 (padd_line, query)
struct lc_node{
    lc_node *left=0, *right=0;
    ll i, j;
    line val;
    lc_node(ll _i, ll _j, line _val) : i(_i), j(_j), val(_val) {}

    // Non-persistent line add
    void add_line(ll a, ll b){
        line v(a,b);
        add_line(v);
    }

    void add_line(line &v){
        ll cur_left=val.eval(i), cur_right=val.eval(j);
        ll new_left=v.eval(i), new_right=v.eval(j);
        if(cur_left <= new_left && cur_right <= new_right) return;
        if(cur_left > new_left && cur_right > new_right){
            val=v;
            return;
        }
        ll mid = (i+j)>>1;
        if(cur_left > new_left) swap(val, v);
        if(val.eval(mid) < v.eval(mid)){
            if(!right) right = new lc_node(mid+1, j, v);
            else right->add_line(v);
        }
        else{
            swap(val, v);
            if(!left) left = new lc_node(i, mid, v);
            else left->add_line(v);
        }
    }

    // Persistent line add
    lc_node *padd_line(ll a, ll b){
        line v(a,b);
        return padd_line(v);
    }

    lc_node *padd_line(line &v){
        ll cur_left=val.eval(i), cur_right=val.eval(j);
        ll new_left=v.eval(i), new_right=v.eval(j);

```

```

        if(cur_left <= new_left && cur_right <= new_right) return
            ↪ this;
        lc_node *ret = new lc_node(*this);
        if(cur_left > new_left && cur_right > new_right){
            ret->val = v;
            return ret;
        }
        ll mid = (i+j)>>1;
        if(cur_left > new_left) swap(ret->val, v);
        if(ret->val.eval(mid) < v.eval(mid)){
            if(!ret->right) ret->right = new lc_node(mid+1, j,
                ↪ v);
            else ret->right = ret->right->padd_line(v);
        }
        else{
            swap(ret->val, v);
            if(!ret->left) ret->left = new lc_node(i, mid, v);
            else ret->left = ret->left->padd_line(v);
        }
        return ret;
    }

    ll query(ll x){
        if(i == j) return val.eval(x);
        ll mid = (i+j)>>1;
        if(x <= mid && left) return min(val.eval(x), left->query(x)
            ↪ );
        else if(x >= mid+1 && right) return min(val.eval(x), right
            ↪ ->query(x));
        return val.eval(x);
    }
};

/* Example:
lc_node *root = new lc_node(min_val, max_val, line(b[0],0));
for(int i=1; i<n; i++){
    dp[i] = root->query(a[i]);
    root->add_line(b[i],dp[i]);
}
*/

```

2. Grafos

2.1. LCA (Binary Lifting)

```

struct LCA{ // Uses Binary Lifting. O(nlogn) preprocessing, O(logn) query.
    int n, l, timer=0;
    vector <vector<int>> up;
    vector <int> enter, exit;
    LCA(vector <vector<int>> &adj, int root=0){
        n = adj.size();
        l = ceil(log2(n));
        enter.resize(n);
        exit.resize(n);

```

```

        up.resize(n, vector<int>(l+1));
        dfs(root, root, adj);
    }

    void dfs(int u, int p, vector<vector<int>> &adj){
        enter[u] = timer++;
        up[u][0] = p;
        for(int j=1; j<=l; j++){
            up[u][j] = up[up[u][j-1]][j-1];
        }
        for(int v : adj[u]){
            if(v != p) dfs(v, u, adj);
        }
        exit[u] = timer++;
    }

    bool is_ancestor(int u, int v){ // v is ancestor of u
        return enter[u] <= enter[v] && exit[u] >= exit[v];
    }

    int query(int u, int v){
        if(is_ancestor(u,v)) return u;
        if(is_ancestor(v,u)) return v;
        for(int i=l; i>=0; i--){
            if(!is_ancestor(up[u][i], v)){
                u = up[u][i];
            }
        }
        return up[u][0];
    }
};

```

2.2. LCA (Sparse Table+ETT)

```

// LCA with SparseTable and Euler Tour
// Requires sparse table of pair<int,int> with min operation
// O(n log n) preprocessing, O(1) query
struct LCA{
    SparseTable st;
    int time=0;
    vector<pair<int,int>> euler;
    vector<int> left, right;
    vector<bool> vis;
    LCA(vector<vector<int>> &adj, int root=0){
        int n = int(adj.size());
        left.resize(n);
        right.resize(n);
        vis.assign(n, false);
        dfs(root, adj);
        st = SparseTable(euler);
    }

    void dfs(int u, vector<vector<int>> &adj, int depth=0){
        vis[u] = 1;

```

```

        left[u] = int(euler.size());
        euler.push_back({depth,u});
        for(int v : adj[u]){
            if(!vis[v]) dfs(v, adj, depth+1);
            right[u] = int(euler.size());
            euler.push_back({depth,u});
        }

        int query(int u, int v){
            if(left[u] > right[v]) swap(u,v);
            return st.query(left[u], right[v]);
        }
};

```

2.3. Heavy Light Decomposition

```

// Heavy Light decomposition of a tree
// Queries in O(log²(n))
// requires: segment_tree
// querying on edges: store edge value in child, change enter[u] in query
template <class T, T merge(T, T)>
struct heavy_light{
    // depth: node depth;
    // sz: subtree size
    // enter: discovery time (index in euler tour)
    // par: parent node
    // head: head of node's chain
    vector<int> depth, sz, enter, par, head;
    segment_tree<T, merge> st;
    vector<T> euler;
    vector<vector<int>> &adj;
    vector<T> &val;
    int time=0;
    /* adj: adjacency list
     * val: value associated with each node
     * merge: merge function for queries
     */
    heavy_light(vector<vector<int>> &adj, vector<T> &val, int
        ↪ root=0) : adj(_adj), val(_val) {
        int n = int(adj.size());
        depth.resize(n); sz.resize(n);
        enter.resize(n); par.resize(n);
        euler.resize(n); head.resize(n);
        par[root] = -1;
        depth[root] = 0;
        dfs1(root);
        dfs2(root, root);
        st = segment_tree<T, merge>(euler);
    }

    void dfs1(int u){ // first dfs, computes depth and sz
        sz[u]=1;
        for(int v : adj[u]){
            if(v != par[u]){

```

```

        par[v] = u;
        depth[v] = depth[u]+1;
        dfs1(v);
        sz[u] += sz[v];
    }
}

void dfs2(int u, int h){ // second dfs, computes hld
    head[u] = h;
    enter[u] = time++;
    euler[enter[u]] = val[u];
    int mx=-1;
    for(int v : adj[u]){
        if(par[u] != v && (mx==-1 || sz[v]>sz[mx])) mx=v;
    }
    if(mx != -1) dfs2(mx, h);
    for(int v : adj[u]){
        if(v != par[u] && v != mx)
            dfs2(v, v);
    }
}

T query(int u, int v){
    T ans = T(); // identity element
    while(head[u] != head[v]){ // find LCA
        if(depth[head[u]] > depth[head[v]]) swap(u, v);
        ans = merge(ans, st.query(enter[head[v]], enter[v]
        ↪ ));
        v = par[head[v]];
    }
    if(depth[u] > depth[v]) swap(u, v); // make sure "u" is
    ↪ LCA
    ans = merge(ans, st.query(enter[u], enter[v])); // enter[u]
    ↪ +1 for edge queriesk
    return ans;
}

void update(int u, T x){
    st.update(enter[u], x);
}

};

```

2.4. Dinic Max Flow

```

// Dinic Max Flow  $O(V^2 E)$ 
struct FlowEdge{
    int u,v;
    ll cap, flow = 0;
    FlowEdge(int u,int v,ll cap):u(u),v(v),cap(cap){}
};

struct Dinic{
    const ll flow_inf = 1e18;
    vector<FlowEdge> edges;
    vector< vector<int> > gr;
    int s,t,n,m=0;
    vector<int> lvl,idx;

```

```

    Dinic (int n,int s,int t):n(n),s(s),t(t){
        gr.resize(n);
        lvl.resize(n);
        idx.resize(n);
    }

    void add_edge(int u,int v,ll cap){
        edges.emplace_back(u,v,cap);
        edges.emplace_back(v,u,0); //cap si bidireccional
        gr[u].push_back(m++);
        gr[v].push_back(m++);
    }

    bool run_bfs(){
        queue<int> bfs;
        bfs.push(s);
        fill(lvl.begin(),lvl.end(),-1);
        lvl[s] = 0;
        while (!bfs.empty()){
            int no = bfs.front();
            bfs.pop();
            for (int ne:gr[no]){
                if (lvl[edges[ne].v] == -1 && edges[ne].cap - edges[ne].
                ↪ flow > 0){
                    lvl[edges[ne].v] = lvl[no] + 1;
                    bfs.push(edges[ne].v);
                }
            }
        }
        return lvl[t] != -1;
    }

    ll dfs(int u,ll cflow){
        if (cflow == 0 || u == t) return cflow;
        while (idx[u] < gr[u].size()){
            int edg = gr[u][idx[u]++];
            int v = edges[edg].v;
            if (lvl[u]+1 == lvl[v] && edges[edg].cap - edges[edg].flow >
            ↪ 0){
                ll rflow = dfs(v,min(cflow,edges[edg].cap - edges[edg].
                ↪ flow));
                if (rflow){
                    edges[edg].flow += rflow;
                    edges[edg^1].flow -= rflow;
                    return rflow;
                }
            }
        }
        return 0;
    }

    ll flow(){
        ll f = 0;
        while (true){
            if (run_bfs()){
                fill(idx.begin(),idx.end(),0);
                ll cf;
                while (cf = dfs(s,flow_inf)) f += cf;
            } else break;
        }
        return f;
    }

```

```

    }
};

```

2.5. Centroid Decomposition

```

struct cenDec{
    // cen_p[i]: el padre de i en el grafo de centroides
    // cen_d[i][j]: distancia entre j y su correspondiente centroide en el
    //             ↪ nivel i
    // cen_h[i] es la lista de hijos de i en el grafo de centroides
    vector<int> cen_p, path;
    vector< vector<int> > &gr, cen_d, cen_h;
    vector<bool> on;
    int n, cen_r;
    int pathm(int no, int p){
        path[no] = 1;
        for (int ne:gr[no]){
            if (on[ne] && ne != p){
                path[no] += pathm(ne, no);
            }
        }
        return path[no];
    }
    void dec(int c, int p, int l){
        pathm(c, -1);
        int cen = c, las = -1;
        while (cen != las){
            las = cen;
            for (int ne:gr[cen]){
                if (on[ne] && path[ne] > path[cen]/2){
                    cen = ne;
                    break;
                }
            }
            path[las] -= path[cen];
            path[cen] += path[las];
        }
        cen_p[cen] = p;
        if (p != -1) cen_h[p].push_back(cen);
        if (l > cen_d.size()) cen_d.push_back(vector<int>(n, -1));
        cen_d[l-1][cen] = 0;
        queue<int> bfs;
        bfs.push(cen);
        while (!bfs.empty()){
            int cno = bfs.front();
            bfs.pop();
            for (int ne:gr[cno]){
                if (on[ne] && cen_d[l-1][ne] == -1){
                    cen_d[l-1][ne] = cen_d[l-1][cno] + 1;
                    bfs.push(ne);
                }
            }
        }
    }
    on[cen] = false;

```

```

        for (int ne:gr[cen]){
            if (on[ne]){
                dec(ne, cen, l+1);
            }
        }
    }
    cenDec(vector< vector<int> > &gr):gr(_gr){
        n = gr.size();
        path.resize(n);
        cen_p.resize(n);
        cen_h.resize(n);
        on.assign(n, true);
        dec(0, -1, 1);
        for (int i=0; i<n; i++){
            if (cen_p[i] == -1){
                cen_r = i;
                break;
            }
        }
    }
};

```

3. Matemáticas

3.1. Exponenciación binaria

```

const ll MOD; // MOD variable global
ll binpow(ll a, ll b){
    a %= MOD;
    ll ans=1;
    while(b > 0){
        if(b & 1)
            ans = ans * a % MOD;
        a = a * a % MOD;
        b >>= 1;
    }
    return ans;
}

```

3.2. Logaritmo discreto

```

// Si (a,b) coprimos, retorna mínimo x tal que a^x = b (mod m) o -1 si no
//             ↪ existe.
// Asume 0^0 = 1
// fuente: cp-algorithms
int dlog(int a, int b, int m) {
    a %= m, b %= m;
    int n = sqrt(m) + 1;

    int an = 1;
    for (int i = 0; i < n; ++i)

```



```

    an = (an * 111 * a) % m;

unordered_map<int, int> vals;
for (int q = 0, cur = b; q <= n; ++q) {
    vals[cur] = q;
    cur = (cur * 111 * a) % m;
}

for (int p = 1, cur = 1; p <= n; ++p) {
    cur = (cur * 111 * an) % m;
    if (vals.count(cur)) {
        int ans = n * p - vals[cur];
        return ans;
    }
}
return -1;
}

```

3.3. Lema de Burnside

El lema de Burnside sirve para problemas de conteo donde hay que contar solo una vez cada simetría.

Sea G un grupo finito actuando sobre un conjunto finito X . Para $g \in G$, denotamos como X^g los elementos de X que están fijos por g . El lema da una fórmula para el número de órbitas $|X/G|$:

$$|X/G| = \frac{1}{|G|} \sum_{g \in G} |X^g|$$

Ejemplo: Contemos collares de n perlas, donde cada perla tiene m posibles colores. Dos collares son simétricos si son idénticos bajo alguna rotación. Así, cada órbita representa un collar, y el grupo G se compone de las n rotaciones posibles: $0, 1, \dots, n-1$ pasos en algún sentido.

Entonces, contamos cuántos collares permanecen invariantes luego de aplicar una rotación de k pasos. Con cero pasos, todos los m^n collares permanecen fijos, y con 1 paso, los m collares donde todas las perlas tienen el mismo color permanecen fijos. En general, un total de $m^{\gcd(k,n)}$ collares están fijos con k pasos, porque bloques de tamaño $\gcd(k,n)$ se reemplazan unos a los otros. Finalmente, por el Lema de Burnside, la cantidad de collares distintos es:

$$\frac{1}{n} \sum_{k=0}^{n-1} m^{\gcd(k,n)}.$$

3.4. Fórmula de inversión de Möbius

Sea (P, \leq) un poset, la función de Möbius μ de P se define recursivamente para elementos de P como:

$$\begin{aligned} \mu(s, s) &= 1 \\ \mu(s, u) &= - \sum_{s \leq t < u} \mu(s, t). \end{aligned}$$

Si cada ideal principal de P es finito, $f: P \rightarrow R$ es una función, y existe una función g que cumple

$$g(y) = \sum_{x \leq y} f(x),$$

luego, se tiene

$$f(y) = \sum_{x \leq y} g(x) \mu(x, y).$$

4. C++

4.1. Custom set/map hash

```

struct custom_hash {
    static uint64_t splitmix64(uint64_t x) {
        // http://xorshift.di.unimi.it/splitmix64.c
        x += 0x9e3779b97f4a7c15;
        x = (x ^ (x >> 30)) * 0xbf58476d1ce4e5b9;
        x = (x ^ (x >> 27)) * 0x94d049bb133111eb;
        return x ^ (x >> 31);
    }
    size_t operator()(uint64_t x) const {
        static const uint64_t FIXED_RANDOM = chrono::steady_clock
            ::now().time_since_epoch().count();
        return splitmix64(x + FIXED_RANDOM);
    }
};
unordered_map<ll, int, custom_hash> safe_map;
unordered_set<ll, custom_hash> safe_set;

```

4.2. Policy Based Order Statistics Tree

```

#include <ext/pb_ds/assoc_container.hpp>
using namespace __gnu_pbds;
typedef tree<int, null_type, less<int>, rb_tree_tag,
            tree_order_statistics_node_update> indexed_set;
/*
usage:
indexed_set s;

```

```
s.insert(2);  
s.insert(3);  
s.insert(7);  
s.insert(9);  
auto x = s.find_by_order(2);  
cout << *x << "\n"; // 7  
cout << s.order_of_key(7) << "\n"; // 2  
*/
```
