

Concurrency Theory

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Sheet 6
Due: Monday, 2025-12-08

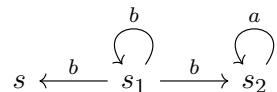
Exercise 6.1

Use two separate formulas to describe the properties in Hennessy-Milner Logic with recursion.

- There is an infinite a -labelled computation path in which all states have an outgoing b -labelled transition.
- Every b -labelled computation path leads to a state from which an a -labelled transition is not possible.

Exercise 6.2

Consider the LTS



Compute all fixpoints of the functions

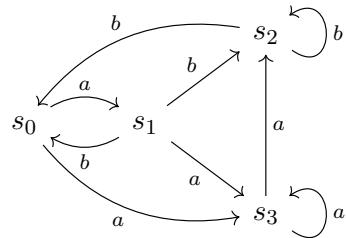
- $\llbracket \langle a \rangle tt \vee [b]X \rrbracket$
- $\llbracket \langle a \rangle tt \vee ([b]X \wedge \langle b \rangle tt) \rrbracket$

Exercise 6.3

Given a LTS $L = (S, Act, \rightarrow)$, show that $\llbracket F \rrbracket : 2^S \rightarrow 2^S$ is a monotonic function over the complete lattice $(2^S, \subseteq)$, for all formulas F , expressible in HML with recursion.

Exercise 6.4

Consider the LTS



- a) Compute $\llbracket \langle b \rangle [a] tt \wedge \langle b \rangle [b] X \rrbracket (\{s_0, s_2\})$
- b) Compute the set of processes satisfying the property

$$X \stackrel{\text{min}}{=} \langle b \rangle \langle a \rangle tt \vee \langle b \rangle [b] X$$

- c) Compute the sets of processes satisfying the mutual recursive equational system

$$\begin{aligned} A &\stackrel{\text{max}}{=} [a] B \\ B &\stackrel{\text{max}}{=} \langle a \rangle C \wedge [b] B \\ C &\stackrel{\text{max}}{=} [b] B \end{aligned}$$