

# Detect Cycle in a Directed Graph - GeeksforGeeks

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Courses Tutorials Practice Jobs DSA Tutorial Interview Questions Quizzes Must Do Advanced DSA System Design Aptitude Puzzles Interview Corner DSA Python Technical Scripter 2026 Explore DSA Fundamentals Logic Building Problems Analysis of Algorithms Data Structures Array Data Structure String in Data Structure Hashing in Data Structure Linked List Data Structure Stack Data Structure Queue Data Structure Tree Data Structure Graph Data Structure Trie Data Structure Algorithms Searching Algorithms Sorting Algorithms Introduction to Recursion Greedy Algorithms Tutorial Graph Algorithms Dynamic Programming or DP Bitwise Algorithms Advanced Segment Tree Binary Indexed Tree or Fenwick Tree Square Root (Sqrt) Decomposition Algorithm Binary Lifting Geometry Interview Preparation Interview Corner GfG160 Practice Problem GeeksforGeeks Practice - Leading Online Coding Platform Problem of The Day - Develop the Habit of Coding DSA Course 90% Refund Detect Cycle in a Directed Graph Last Updated : 4 Nov, 2025 Given a directed graph represented by its adjacency list `adj[][]`, determine whether the graph contains a cycle/Loop or not. A cycle is a path that starts and ends at the same vertex, following the direction of edges. Examples: Input: `adj[][] = [[1], [2], [0, 3], []]` Output: true Explanation: There is a cycle `0 -> 1 -> 2 -> 0`. Input: `adj[][] = [[2], [0], []]` Output: false Explanation: There is no cycle in the graph. Try it on GfG Practice Table of Content [Approach 1] Using DFS -  $O(V + E)$  Time and  $O(V)$  Space [Approach 2] Using Topological Sorting -  $O(V + E)$  Time and  $O(V)$  Space [Approach 1] Using DFS -  $O(V + E)$  Time and  $O(V)$  Space To detect a cycle in a directed graph, we use Depth First Search (DFS). In DFS, we go as deep as possible from a starting node. If during this process, we reach a node that we've already visited in the same DFS path, it means we've gone back to an ancestor — this shows a cycle exists. But there's a problem: When we start DFS from one node, some nodes get marked as visited. Later, when we start DFS from another node, those visited nodes may appear again, even if there's no cycle. So, using only `visited[]` isn't enough. To fix this, we use two arrays: `visited[]` - marks nodes visited at least once. `recStack[]` - marks nodes currently in the recursion (active) path. If during DFS we reach a node that's already in the `recStack`, we've found a path from the current node back to one of its ancestors, forming a cycle. As soon as we finish exploring all paths from a node, we remove it from the recursion stack by marking `recStack[node] = false`. This ensures that only the nodes in the current DFS path are tracked. Illustration: C++ //Driver Code Starts #include <iostream> #include <vector> using namespace std ; //Driver Code Ends // Utility DFS function to detect cycle in a directed graph bool isCyclicUtil ( vector < vector < int >>& adj , int u , vector < bool >& visited , vector < bool >& recStack ) { // node is already in recursion stack cycle found if ( recStack [ u ] ) return true ; // already processed no need to visit again if ( visited [ u ] ) return false ; visited [ u ] = true ; recStack [ u ] = true ; // Recur for all adjacent nodes for ( int v : adj [ u ] ) { if ( isCyclicUtil ( adj , v , visited , recStack ) ) return true ; } // remove from recursion stack before backtracking recStack [ u ] = false ; return false ; } // Function to detect cycle in a directed graph bool isCyclic ( vector < vector < int >>& adj ) { int V = adj . size () ; vector < bool > visited ( V , false ) ; vector < bool > recStack ( V , false ) ; // Run DFS from every unvisited node for ( int i = 0 ; i < V ; i ++ ) { if ( ! visited [ i ] && isCyclicUtil ( adj , i , visited , recStack ) ) return true ; } return false ; } //Driver Code Starts int main () { vector < vector < int >> adj = { { 1 } , { 2 } , { 0 , 3 } } ; cout << ( isCyclic ( adj ) ? "true" : "false" ) << endl ; return 0 ; } //Driver Code Ends Java //Driver Code Starts import java.util.ArrayList ; public class GFG { //Driver Code Ends // Utility DFS function to detect cycle in a directed graph static boolean isCyclicUtil ( ArrayList < ArrayList < Integer >> adj , int u , boolean [] visited , boolean [] recStack ) { // Node already in recursion stack cycle found if ( recStack [ u ] ) return true ; // Already processed no need to visit again if ( visited [ u ] ) return false ; visited [ u ] = true ; recStack [ u ] = true ; // Recur for all adjacent nodes for ( int v : adj . get ( u ) ) { if ( isCyclicUtil ( adj , v , visited , recStack ) ) return true ; } // Remove from recursion stack before backtracking recStack [ u ] = false ; return false ; } // Function to detect cycle in a directed graph static boolean isCyclic ( ArrayList < ArrayList < Integer >> adj ) { int V = adj . size () ; boolean [] visited = new boolean [ V ] ; boolean [] recStack = new boolean [ V ] ; // Run DFS from every unvisited node for ( int i = 0 ; i < V ; i ++ ) { if ( ! visited [ i ] && isCyclicUtil ( adj , i , visited , recStack ) ) return true ; } return false ; } //Driver Code Starts // Function to add an edge to the adjacency list static void addEdge ( ArrayList < ArrayList < Integer >> adj , int u , int v ) { adj . get ( u ) . add ( v ) ; } public static void main ( String [] args ) { int V = 4 ; ArrayList < ArrayList < Integer >> adj = new ArrayList

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<>()); for ( int i = 0 ; i < V ; i ++ ) { adj . add ( new ArrayList <> ()); } // Add directed edges addEdge ( adj ,
0 , 1 ); addEdge ( adj , 1 , 2 ); addEdge ( adj , 2 , 0 ); addEdge ( adj , 2 , 3 ); System . out . println (
isCyclic ( adj ) ? "true" : "false" ); } } //Driver Code Ends Python # Utility DFS function to detect cycle in a
directed graph def isCyclicUtil ( adj , u , visited , recStack ): # node is already in recursion stack cycle
found if recStack [ u ]: return True # already processed no need to visit again if visited [ u ]: return False
visited [ u ] = True recStack [ u ] = True # Recur for all adjacent nodes for v in adj [ u ]: if isCyclicUtil (
adj , v , visited , recStack ): return True # remove from recursion stack before backtracking recStack [ u
] = False return False # Function to detect cycle in a directed graph def isCyclic ( adj ): V = len ( adj )
visited = [ False ] * V recStack = [ False ] * V # Run DFS from every unvisited node for i in range ( V ): if
not visited [ i ] and isCyclicUtil ( adj , i , visited , recStack ): return True return False #Driver Code Starts
if __name__ == "__main__" : adj = [[ 1 ], [ 2 ], [ 0 , 3 ]] print ( "true" if isCyclic ( adj ) else "false" ) #Driver
Code Ends C# //Driver Code Starts using System ; using System.Collections.Generic ; class GFG { //
Utility DFS function to detect cycle in a directed graph //Driver Code Ends static bool isCyclicUtil ( List <
List < int >> adj , int u , bool [] visited , bool [] recStack ) { // Node already in recursion stack cycle found
if ( recStack [ u ]) return true ; // Already processed → no need to visit again if ( visited [ u ]) return false
; visited [ u ] = true ; recStack [ u ] = true ; // Recur for all adjacent nodes foreach ( int v in adj [ u ]) { if (
isCyclicUtil ( adj , v , visited , recStack )) return true ; } // Remove from recursion stack before
backtracking recStack [ u ] = false ; return false ; } // Function to detect cycle in a directed graph static
bool isCyclic ( List < List < int >> adj ) { int V = adj . Count ; bool [] visited = new bool [ V ]; bool []
recStack = new bool [ V ]; // Run DFS from every unvisited node for ( int i = 0 ; i < V ; i ++ ) { if ( ! visited
[ i ] && isCyclicUtil ( adj , i , visited , recStack )) return true ; } return false ; } //Driver Code Starts //
Function to add an edge to the adjacency list static void addEdge ( List < List < int >> adj , int u , int v ) {
adj [ u ]. Add ( v ); } static void Main () { int V = 4 ; List < List < int >> adj = new List < List < int >> ( ); for (
int i = 0 ; i < V ; i ++ ) { adj . Add ( new List < int > ( )); } // Add directed edges addEdge ( adj , 0 , 1 );
addEdge ( adj , 1 , 2 ); addEdge ( adj , 2 , 0 ); addEdge ( adj , 2 , 3 ); Console . WriteLine ( isCyclic ( adj
) ? "true" : "false" ); } } //Driver Code Ends JavaScript // Utility DFS function to detect cycle in a directed
graph function isCyclicUtil ( adj , u , visited , recStack ) { // node is already in recursion stack cycle
found if ( recStack [ u ]) return true ; // already processed no need to visit again if ( visited [ u ]) return
false ; visited [ u ] = true ; recStack [ u ] = true ; // Recur for all adjacent nodes for ( let v of adj [ u ]) { if (
isCyclicUtil ( adj , v , visited , recStack )) return true ; } // remove from recursion stack before
backtracking recStack [ u ] = false ; return false ; } // Function to detect cycle in a directed graph
function isCyclic ( adj ) { const V = adj . length ; const visited = Array ( V ). fill ( false ); const recStack =
Array ( V ). fill ( false ); // Run DFS from every unvisited node for ( let i = 0 ; i < V ; i ++ ) { if ( ! visited [ i ]
&& isCyclicUtil ( adj , i , visited , recStack )) return true ; } return false ; } // Driver Code //Driver Code
Starts let V = 4 ; const adj = [[ 1 ],[ 2 ],[ 0 , 3 ]]; console . log ( isCyclic ( adj ) ? "true" : "false" ); //Driver
Code Ends Output true [Approach 2] Using Topological Sorting - O(V + E) Time and O(V) Space The
idea is to use Kahn's algorithm because it works only for Directed Acyclic Graphs (DAGs). So, while
performing topological sorting using Kahn's algorithm, if we are able to include all the vertices in the
topological order, it means the graph has no cycle and is a DAG. However, if at the end there are still
some vertices left (i.e., their in-degree never becomes 0), it means those vertices are part of a cycle.
Hence, if we cannot get all the vertices in the topological sort, the graph must contain at least one cycle.
C++ //Driver Code Starts #include <iostream> #include <vector> #include <queue> using namespace
std ; //Driver Code Ends bool isCyclic ( vector < vector < int >> & adj ) { int V = adj . size (); // Array to
store in-degree of each vertex vector < int > inDegree ( V , 0 ); queue < int > q ; // Count of visited
(processed) nodes int visited = 0 ; //Compute in-degrees of all vertices for ( int u = 0 ; u < V ; ++ u ) { for
( int v : adj [ u ]) { inDegree [ v ] ++ ; } } // Add all vertices with in-degree 0 to the queue for ( int u = 0 ; u
< V ; ++ u ) { if ( inDegree [ u ] == 0 ) { q . push ( u ); } } // Perform BFS (Topological Sort) while ( ! q .
empty () ) { int u = q . front (); q . pop (); visited ++ ; // Reduce in-degree of neighbors for ( int v : adj [ u ])
{ inDegree [ v ] -- ; if ( inDegree [ v ] == 0 ) { // Add to queue when in-degree becomes 0 q . push ( v ); } }
} // If visited nodes != total nodes, a cycle exists return visited != V ; } //Driver Code Starts int main () {
vector < vector < int >> adj = { { 1 }, { 2 }, { 0 , 3 } }; cout << ( isCyclic ( adj ) ? "true" : "false" ) << endl ;
return 0 ; } //Driver Code Ends Java //Driver Code Starts import java.util.Queue ; import
java.util.LinkedList ; import java.util.ArrayList ; public class GFG { //Driver Code Ends static boolean
isCyclic ( ArrayList < ArrayList < Integer >> adj ) { int V = adj . size (); // Array to store in-degree of each
vertex int [] inDegree = new int [ V ]; Queue < Integer > q = new LinkedList <> ( ); // Count of visited
(processed) nodes int visited = 0 ; // Compute in-degrees of all vertices for ( int u = 0 ; u < V ; ++ u ) { for
( int v : adj . get ( u )) { inDegree [ v ] ++ ; } } // Add all vertices with in-degree 0 to the queue for ( int u = 0

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; u < V ; ++ u ) { if ( inDegree [ u ] == 0 ) { q . add ( u ); } } // Perform BFS (Topological Sort) while ( ! q .
isEmpty () ) { int u = q . poll (); visited ++ ; // Reduce in-degree of neighbors for ( int v : adj . get ( u ) ) {
inDegree [ v ] -- ; if ( inDegree [ v ] == 0 ) { // Add to queue when in-degree becomes 0 q . add ( v ); } } } //
If visited nodes != total nodes, a cycle exists return visited != V ; } //Driver Code Starts // Function to add
an edge to the adjacency list static void addEdge ( ArrayList < ArrayList < Integer >> adj , int u , int v ) {
adj . get ( u ) . add ( v ); } public static void main ( String [] args ) { int V = 4 ; ArrayList < ArrayList <
Integer >> adj = new ArrayList <> (); for ( int i = 0 ; i < V ; i ++ ) { adj . add ( new ArrayList <> () ); } // Add
edges addEdge ( adj , 0 , 1 ); addEdge ( adj , 1 , 2 ); addEdge ( adj , 2 , 0 ); addEdge ( adj , 2 , 3 );
System . out . println ( isCyclic ( adj ) ? "true" : "false" ); } } //Driver Code Ends Python #Driver Code
Starts from collections import deque #Driver Code Ends def isCyclic ( adj ): V = max ( max ( sub ) if sub
else 0 for sub in adj ) + 1 # Array to store in-degree of each vertex inDegree = [ 0 ] * V q = deque () #
Count of visited (processed) nodes visited = 0 # Compute in-degrees of all vertices for u in range ( V ):
for v in adj [ u ]: inDegree [ v ] += 1 # Add all vertices with in-degree 0 to the queue for u in range ( V ): if
inDegree [ u ] == 0 : q . append ( u ) # Perform BFS (Topological Sort) while q : u = q . popleft () visited
+= 1 # Reduce in-degree of neighbors for v in adj [ u ]: inDegree [ v ] -= 1 if inDegree [ v ] == 0 : # Add
to queue when in-degree becomes 0 q . append ( v ) # If visited nodes != total nodes, a cycle exists
return visited != V #Driver Code Starts if __name__ == "__main__" : adj = [[ 1 ],[ 2 ],[ 0 , 3 ], []] print (
"true" if isCyclic ( adj ) else "false" ) #Driver Code Ends C# //Driver Code Starts using System ; using
System.Collections.Generic ; class GFG { //Driver Code Ends static bool isCyclic ( List < List < int >>
adj ) { int V = adj . Count ; // Array to store in-degree of each vertex int [] inDegree = new int [ V ]; Queue
< int > q = new Queue < int > (); // Count of visited (processed) nodes int visited = 0 ; // Compute
in-degrees of all vertices for ( int u = 0 ; u < V ; ++ u ) { foreach ( int v in adj [ u ] ) { inDegree [ v ] ++ ; } }
// Add all vertices with in-degree 0 to the queue for ( int u = 0 ; u < V ; ++ u ) { if ( inDegree [ u ] == 0 ) { q
. Enqueue ( u ); } } // Perform BFS (Topological Sort) while ( q . Count > 0 ) { int u = q . Dequeue ();
visited ++ ; // Reduce in-degree of neighbors foreach ( int v in adj [ u ] ) { inDegree [ v ] -- ; if ( inDegree [
v ] == 0 ) { // Add to queue when in-degree becomes 0 q . Enqueue ( v ); } } } // If visited nodes != total
nodes, a cycle exists return visited != V ; } //Driver Code Starts // Function to add an edge to the
adjacency list static void addEdge ( List < List < int >> adj , int u , int v ) { adj [ u ] . Add ( v ); } static void
Main () { int V = 4 ; List < List < int >> adj = new List < List < int >> (); for ( int i = 0 ; i < V ; i ++ ) { adj .
Add ( new List < int > () ); } // Add edges addEdge ( adj , 0 , 1 ); addEdge ( adj , 1 , 2 ); addEdge ( adj , 2
, 0 ); addEdge ( adj , 2 , 3 ); Console . WriteLine ( isCyclic ( adj ) ? "true" : "false" ); } } //Driver Code
Ends JavaScript //Driver Code Starts const Denque = require ( "denque" ); //Driver Code Ends function
isCyclic ( adj ) { const V = adj . length ; // Array to store in-degree of each vertex const inDegree = new
Array ( V ) . fill ( 0 ); const q = new Denque (); // Count of visited (processed) nodes let visited = 0 ;
//Compute in-degrees of all vertices for ( let u = 0 ; u < V ; ++ u ) { for ( let v of adj [ u ] ) { inDegree [ v ]
++ ; } } // Add all vertices with in-degree 0 to the queue for ( let u = 0 ; u < V ; ++ u ) { if ( inDegree [ u ]
=== 0 ) { q . push ( u ); } } // Perform BFS (Topological Sort) while ( ! q . isEmpty () ) { const u = q . shift
(); visited ++ ; // Reduce in-degree of neighbors for ( let v of adj [ u ] ) { inDegree [ v ] -- ; if ( inDegree [ v ]
=== 0 ) { // Add to queue when in-degree becomes 0 q . push ( v ); } } } // If visited nodes != total
nodes, a cycle exists return visited !== V ; } //Driver Code Starts // Driver Code const adj = [ [ 1 ],[ 2 ],[ 0
, 3 ],[]]; console . log ( isCyclic ( adj ) ? "true" : "false" ); //Driver Code Ends Output true Similar Article:
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