

Maximum sum rectangle in a 2D matrix - GeeksforGeeks

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Courses Tutorials Practice Jobs DSA Tutorial Interview Questions Quizzes Must Do Advanced DSA System Design Aptitude Puzzles Interview Corner DSA Python Technical Scripter 2026 Explore DSA Fundamentals Logic Building Problems Analysis of Algorithms Data Structures Array Data Structure String in Data Structure Hashing in Data Structure Linked List Data Structure Stack Data Structure Queue Data Structure Tree Data Structure Graph Data Structure Trie Data Structure Algorithms Searching Algorithms Sorting Algorithms Introduction to Recursion Greedy Algorithms Tutorial Graph Algorithms Dynamic Programming or DP Bitwise Algorithms Advanced Segment Tree Binary Indexed Tree or Fenwick Tree Square Root (Sqrt) Decomposition Algorithm Binary Lifting Geometry Interview Preparation Interview Corner GfG160 Practice Problem GeeksforGeeks Practice - Leading Online Coding Platform Problem of The Day - Develop the Habit of Coding DSA Course 90% Refund Maximum sum rectangle in a 2D matrix Last Updated : 28 Jul, 2025 Given a 2D matrix mat[][] of integers, find the submatrix (i.e., a rectangular section of the matrix) that has the maximum possible sum among all possible submatrices. Example: Input: mat[][] = [[1, 2, -1, -4, -20], [-8, -3, 4, 2, 1], [3, 8, 10, 1, 3], [-4, -1, 1, 7, -6]] Output: 29 Explanation: The matrix is as follows and the green rectangle denotes the maximum sum rectangle which is equal to 29. Try it on GfG Practice This problem is mainly an extension of the Maximum Subarray Sum . Table of Content [Naive approach] Iterating Over All Possible Submatrices [Better Approach] Using Prefix Sum [Expected Approach] Using Kadane's Algorithm [Naive approach] Iterating Over All Possible Submatrices We explore all possible rectangles in the given 2D array, by using four variables: two to define the left and right boundaries and two more to define the top and bottom boundaries and calculate their sums, and keep track of the maximum sum found. C++ #include <iostream> #include <vector> #include <climits> using namespace std ; int maxRectSum (vector < vector < int >> & mat) { int n = mat . size () ; int m = mat [0] . size () ; int maxSum = INT_MIN ; for (int up = 0 ; up < n ; up ++) { for (int left = 0 ; left < m ; left ++) { for (int down = up ; down < n ; down ++) { for (int right = left ; right < m ; right ++) { // Find the sum of submatrix(up, right, down, left) int sum = 0 ; for (int i = up ; i <= down ; i ++) { for (int j = left ; j <= right ; j ++) { sum += mat [i][j] ; } } // Update maxSum if sum > maxSum. if (sum > maxSum) { maxSum = sum ; } } } } } return maxSum ; } int main () { vector < vector < int >> mat = { { 1 , 2 , -1 , -4 , -20 } , { -8 , -3 , 4 , 2 , 1 } , { 3 , 8 , 10 , 1 , 3 } , { -4 , -1 , 1 , 7 , -6 } } ; cout << maxRectSum (mat) << endl ; return 0 ; } Java class GfG { static int maxRectSum (int [][] mat) { int n = mat . length ; int m = mat [0] . length ; int maxSum = Integer . MIN_VALUE ; for (int up = 0 ; up < n ; up ++) { for (int left = 0 ; left < m ; left ++) { for (int down = up ; down < n ; down ++) { for (int right = left ; right < m ; right ++) { // Find the sum of // submatrix(up, right, down, left) int sum = 0 ; for (int i = up ; i <= down ; i ++) { for (int j = left ; j <= right ; j ++) { sum += mat [i][j] ; } } // Update maxSum if sum > maxSum. if (sum > maxSum) { maxSum = sum ; } } } } } } return maxSum ; } public static void main (String [] args) { int [][] mat = { { 1 , 2 , -1 , -4 , -20 } , { -8 , -3 , 4 , 2 , 1 } , { 3 , 8 , 10 , 1 , 3 } , { -4 , -1 , 1 , 7 , -6 } } ; System . out . println (maxRectSum (mat)) ; } } Python def maxRectSum (mat) : n = len (mat) m = len (mat [0]) maxSum = float ('-inf') for up in range (n) : for left in range (m) : for down in range (up , n) : for right in range (left , m) : # Calculate the sum of submatrix # (up, left, down, right) subMatrixSum = 0 for i in range (up , down + 1) : for j in range (left , right + 1) : subMatrixSum += mat [i][j] # Update maxSum if a larger sum is found maxSum = max (maxSum , subMatrixSum) return maxSum if __name__ == "__main__" : mat = [[1 , 2 , -1 , -4 , -20] , [-8 , -3 , 4 , 2 , 1] , [3 , 8 , 10 , 1 , 3] , [-4 , -1 , 1 , 7 , -6]] print (maxRectSum (mat)) C# using System ; class GfG { static int maxRectSum (int [][] mat) { int n = mat . Length ; int m = mat [0] . Length ; int maxSum = int . MinValue ; for (int up = 0 ; up < n ; up ++) { for (int left = 0 ; left < m ; left ++) { for (int down = up ; down < n ; down ++) { for (int right = left ; right < m ; right ++) { // Calculate the sum of submatrix // (up, left, down, right) int subMatrixSum = 0 ; for (int i = up ; i <= down ; i ++) { for (int j = left ; j <= right ; j ++) { subMatrixSum += mat [i][j] ; } } // Update maxSum if a larger sum is found maxSum = Math . Max (maxSum , subMatrixSum) ; } } } } } return maxSum ; } static void Main (string [] args) { int [][] mat = { new int [] { 1 , 2 , -1 , -4 , -20 } , new int [] { -8 , -3 , 4 , 2 , 1 } , new int [] { 3 , 8 , 10 , 1 , 3 } , new int [] { -4 , -1 , 1 , 7 , -6 } } ; Console . WriteLine (

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maxRectSum ( mat )); } } JavaScript function maxRectSum ( mat ) { const n = mat . length ; const m =
mat [ 0 ]. length ; let maxSum = - Infinity ; for ( let up = 0 ; up < n ; up ++ ) { for ( let left = 0 ; left < m ; left
++ ) { for ( let down = up ; down < n ; down ++ ) { for ( let right = left ; right < m ; right ++ ) { // Calculate
the sum of submatrix // (up, left, down, right) let subMatrixSum = 0 ; for ( let i = up ; i <= down ; i ++ ) {
for ( let j = left ; j <= right ; j ++ ) { subMatrixSum += mat [ i ][ j ] ; } } // Update maxSum if a larger sum is //
found maxSum = Math . max ( maxSum , subMatrixSum ) ; } } } } return maxSum ; } // Driver Code const
mat = [ [ 1 , 2 , - 1 , - 4 , - 20 ], [ - 8 , - 3 , 4 , 2 , 1 ], [ 3 , 8 , 10 , 1 , 3 ], [ - 4 , - 1 , 1 , 7 , - 6 ] ]; console .
log ( maxRectSum ( mat )); Time Complexity:  $O((n*m)^3)$ , as we iterate over all the boundaries of the
rectangle in  $O((n*m)^2)$  time. For each rectangle, we find its sum in  $O(n*m)$  time. Auxiliary Space:  $O(1)$ 
[Better Approach] Using Prefix Sum In this approach, we used prefix sum matrix to efficiently compute
the sum of any submatrix. We first precompute the prefix sum such that pref[i][j] stores the sum of all
elements in the submatrix from the top-left corner (0, 0) to the cell (i, j). This preprocessing allows us to
calculate the sum of any arbitrary submatrix in constant time using the inclusion-exclusion principle,
thereby eliminating redundant computations and significantly improving performance. Step by Step
Approach: Initialize a 2D prefix sum matrix pref[][] of the same size as the input matrix. Precompute the
prefix sum for each cell : pref[i][j] = matrix[i][j] + pref[i-1][j] + pref[i][j-1] - pref[i-1][j-1] (Handle edges
separately to avoid index out-of-bound errors.) Iterate over all possible top-left corners ( up, left ) and
bottom-right corners ( down, right ) to define submatrices. For each such submatrix, compute the sum
using the prefix sum matrix in constant time: sum = pref[down][right] - pref[up - 1][right] -
pref[down][left-1] + pref[up-1][left - 1] Track the maximum sum obtained across all submatrices. C++
#include <iostream> #include <vector> #include <climits> using namespace std ; int findSum ( int up ,
int left , int down , int right , vector < vector < int >> & pref ) { // Start with the sum of the entire submatrix
(0, 0) to (down, right) int sum = pref [ down ][ right ] ; // Subtract the area to the left of the submatrix, if it
exists if ( left - 1 >= 0 ) { sum -= pref [ down ][ left - 1 ] ; } // Subtract the area above the submatrix, if it
exists if ( up - 1 >= 0 ) { sum -= pref [ up - 1 ][ right ] ; } // Add back the overlapping area that was
subtracted twice if ( up - 1 >= 0 && left - 1 >= 0 ) { sum += pref [ up - 1 ][ left - 1 ] ; } return sum ; } //
function to find the maximum sum rectangle in a 2D matrix int maxRectSum ( vector < vector < int >> &
mat ) { int n = mat . size () ; int m = mat [ 0 ]. size () ; // Initialize the prefix sum matrix vector < vector < int
>> pref ( n , vector < int > ( m , 0 )); // Row-wise sum for ( int i = 0 ; i < n ; i ++ ) { for ( int j = 0 ; j < m ; j
++ ) { pref [ i ][ j ] = mat [ i ][ j ] ; if ( j - 1 >= 0 ) { pref [ i ][ j ] += pref [ i ][ j - 1 ] ; } } } // Column-wise sum for (
int j = 0 ; j < m ; j ++ ) { for ( int i = 0 ; i < n ; i ++ ) { if ( i - 1 >= 0 ) { pref [ i ][ j ] += pref [ i - 1 ][ j ] ; } } } int
maxSum = INT_MIN ; for ( int up = 0 ; up < n ; up ++ ) { for ( int left = 0 ; left < m ; left ++ ) { for ( int down
= up ; down < n ; down ++ ) { for ( int right = left ; right < m ; right ++ ) { // Find the sum of submatrix(up,
right, down, left) int sum = findSum ( up , left , down , right , pref ) ; // Update maxSum if sum > maxSum.
if ( sum > maxSum ) { maxSum = sum ; } } } } } return maxSum ; } int main () { vector < vector < int >>
mat = { { 1 , 2 , -1 , -4 , -20 }, { -8 , -3 , 4 , 2 , 1 }, { 3 , 8 , 10 , 1 , 3 }, { -4 , -1 , 1 , 7 , -6 } }; cout <<
maxRectSum ( mat ) << endl ; } Java import java.util.* ; class GfG { static int findSum ( int up , int left ,
int down , int right , int [][] pref ) { // Start with the sum of the entire submatrix // (0, 0) to (down, right) int
sum = pref [ down ][ right ] ; // Subtract the area to the left of the submatrix, if it exists if ( left - 1 >= 0 ) {
sum -= pref [ down ][ left - 1 ] ; } // Subtract the area above the submatrix, if it exists if ( up - 1 >= 0 ) {
sum -= pref [ up - 1 ][ right ] ; } // Add back the overlapping area that was subtracted twice if ( up - 1 >= 0
&& left - 1 >= 0 ) { sum += pref [ up - 1 ][ left - 1 ] ; } return sum ; } // function to find the maximum sum
rectangle in a 2D matrix public static int maxRectSum ( int [][] mat ) { int n = mat . length ; int m = mat [ 0
]. length ; // Initialize the prefix sum matrix int [][] pref = new int [ n ][ m ] ; // Row-wise sum for ( int i = 0 ;
i < n ; i ++ ) { for ( int j = 0 ; j < m ; j ++ ) { pref [ i ][ j ] = mat [ i ][ j ] ; if ( j - 1 >= 0 ) { pref [ i ][ j ] += pref [ i ][
j - 1 ] ; } } } // Column-wise sum for ( int j = 0 ; j < m ; j ++ ) { for ( int i = 0 ; i < n ; i ++ ) { if ( i - 1 >= 0 ) {
pref [ i ][ j ] += pref [ i - 1 ][ j ] ; } } } // Find the maximum sum rectangle int maxSum = Integer .
MIN_VALUE ; for ( int up = 0 ; up < n ; up ++ ) { for ( int left = 0 ; left < m ; left ++ ) { for ( int down = up ;
down < n ; down ++ ) { for ( int right = left ; right < m ; right ++ ) { // Find the sum of the submatrix // (up,
left) to // (down, right) int sum = findSum ( up , left , down , right , pref ) ; // Update maxSum if sum >
maxSum if ( sum > maxSum ) { maxSum = sum ; } } } } } return maxSum ; } public static void main (
String [] args ) { int [][] mat = { { 1 , 2 , -1 , -4 , -20 }, { -8 , -3 , 4 , 2 , 1 }, { 3 , 8 , 10 , 1 , 3 }, { -4 , -1 , 1
, 7 , -6 } }; System . out . println ( maxRectSum ( mat )); } } Python def findSum ( up , left , down , right ,
pref ): # Start with the sum of the entire submatrix (0, 0) to (down, right) totalSum = pref [ down ][ right ]
# Subtract the area to the left of the submatrix, if it exists if left - 1 >= 0 : totalSum -= pref [ down ][ left -
1 ] # Subtract the area above the submatrix, if it exists if up - 1 >= 0 : totalSum -= pref [ up - 1 ][ right ] #
Add back the overlapping area that was subtracted twice if up - 1 >= 0 and left - 1 >= 0 : totalSum +=

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pref [ up - 1 ][ left - 1 ] return totalSum # Function to find the maximum sum rectangle in a 2D matrix
def maxRectSum ( mat ): n = len ( mat ) m = len ( mat [ 0 ] ) # Initialize the prefix sum matrix pref = [[ 0 ] * m
for _ in range ( n )] # Row-wise sum for i in range ( n ): for j in range ( m ): pref [ i ][ j ] = mat [ i ][ j ] if j - 1
>= 0 : pref [ i ][ j ] += pref [ i ][ j - 1 ] # Column-wise sum for j in range ( m ): for i in range ( n ): if i - 1 >= 0
: pref [ i ][ j ] += pref [ i - 1 ][ j ] # Find the maximum sum rectangle maxSum = float ( '-inf' ) for up in
range ( n ): for left in range ( m ): for down in range ( up , n ): for right in range ( left , m ): totalSum =
findSum ( up , left , down , right , pref ) maxSum = max ( maxSum , totalSum ) return maxSum if
__name__ == "__main__" : mat = [ [ 1 , 2 , - 1 , - 4 , - 20 ], [ - 8 , - 3 , 4 , 2 , 1 ], [ 3 , 8 , 10 , 1 , 3 ], [ - 4 , -
1 , 1 , 7 , - 6 ] ] print ( maxRectSum ( mat ))
C# using System ; class GfG { static int findSum ( int up , int
left , int down , int right , int[,] pref ) { // Start with the sum of the entire submatrix from (0,0) to
(down,right) int sum = pref [ down , right ]; // Subtract the area to the left of the submatrix, if it exists if (
left - 1 >= 0 ) sum -= pref [ down , left - 1 ]; // Subtract the area above the submatrix, if it exists if ( up - 1
>= 0 ) sum -= pref [ up - 1 , right ]; // Add back the overlapping area that was subtracted twice if ( up - 1
>= 0 && left - 1 >= 0 ) sum += pref [ up - 1 , left - 1 ]; return sum ; } static int maxRectSum ( int[,] mat ) {
int n = mat . GetLength ( 0 ) ; int m = mat . GetLength ( 1 ) ; // Initialize the prefix sum matrix int[,] pref =
new int [ n , m ] ; // Compute row-wise prefix sum for ( int i = 0 ; i < n ; i ++ ) { for ( int j = 0 ; j < m ; j ++ ) {
pref [ i , j ] = mat [ i , j ] ; if ( j - 1 >= 0 ) pref [ i , j ] += pref [ i , j - 1 ] ; } } // Compute column-wise prefix sum
for ( int j = 0 ; j < m ; j ++ ) { for ( int i = 0 ; i < n ; i ++ ) { if ( i - 1 >= 0 ) pref [ i , j ] += pref [ i - 1 , j ] ; } } int
maxSum = int . MinValue ; // Check all possible submatrices for ( int up = 0 ; up < n ; up ++ ) { for ( int
left = 0 ; left < m ; left ++ ) { for ( int down = up ; down < n ; down ++ ) { for ( int right = left ; right < m ;
right ++ ) { int totalSum = findSum ( up , left , down , right , pref ) ; if ( totalSum > maxSum ) maxSum =
totalSum ; } } } } return maxSum ; } static void Main ( string [] args ) { int[,] mat = new int [,] { { 1 , 2 , - 1 ,
- 4 , - 20 }, { - 8 , - 3 , 4 , 2 , 1 }, { 3 , 8 , 10 , 1 , 3 }, { - 4 , - 1 , 1 , 7 , - 6 } } ; Console . WriteLine (
maxRectSum ( mat )) ; } }
JavaScript // Function to get the sum of submatrix (r1, c1) to (r2, c2) function
findSum ( up , left , down , right , pref ) { // Start with the sum of the entire submatrix (0, 0) to (down,
right) let sum = pref [ down ][ right ]; // Subtract the area to the left of the submatrix, if it exists if ( left - 1
>= 0 ) { sum -= pref [ down ][ left - 1 ] ; } // Subtract the area above the submatrix, if it exists if ( up - 1 >=
0 ) { sum -= pref [ up - 1 ][ right ] ; } // Add back the overlapping area that was subtracted twice if ( up - 1
>= 0 && left - 1 >= 0 ) { sum += pref [ up - 1 ][ left - 1 ] ; } return sum ; } // function to find the maximum
sum rectangle in a 2D matrix function maxRectSum ( mat ) { const n = mat . length ; const m = mat [ 0 ].
length ; // Create and initialize prefix sum matrix const pref = Array . from ( { length : n } , () => Array ( m
). fill ( 0 )) ; // Row-wise prefix sum for ( let i = 0 ; i < n ; i ++ ) { for ( let j = 0 ; j < m ; j ++ ) { pref [ i ][ j ] =
mat [ i ][ j ] ; if ( j - 1 >= 0 ) { pref [ i ][ j ] += pref [ i ][ j - 1 ] ; } } } // Column-wise prefix sum for ( let j = 0 ; j <
m ; j ++ ) { for ( let i = 0 ; i < n ; i ++ ) { if ( i - 1 >= 0 ) { pref [ i ][ j ] += pref [ i - 1 ][ j ] ; } } } let maxSum = -
Infinity ; // Try all submatrices for ( let up = 0 ; up < n ; up ++ ) { for ( let left = 0 ; left < m ; left ++ ) { for (
let down = up ; down < n ; down ++ ) { for ( let right = left ; right < m ; right ++ ) { let sum = findSum ( up ,
left , down , right , pref ) ; if ( sum > maxSum ) { maxSum = sum ; } } } } } return maxSum ; } // Driver
code const mat = [ [ 1 , 2 , - 1 , - 4 , - 20 ], [ - 8 , - 3 , 4 , 2 , 1 ], [ 3 , 8 , 10 , 1 , 3 ], [ - 4 , - 1 , 1 , 7 , - 6 ] ] ;
console . log ( maxRectSum ( mat )) ; // Output: 29 Time Complexity : O((n*m) ^ 2 ) Auxiliary Space:
O(n*m), due to the prefix sum matrix. [Expected Approach] Using Kadane's Algorithm Instead of
checking every possible submatrix, we fix two column boundaries : left and right. For each pair of
columns, we compress the 2D matrix into a 1D array , where each element represents the sum of rows
between the current left and right columns. This 1D array (temp[]) represents the column-wise sum for
a fixed column window. Now, we can apply Kadane's algorithm on this temp[] array to find the
maximum subarray sum , which corresponds to a rectangular submatrix with the maximum sum
between columns left and right. C++ #include <iostream> #include <vector> #include <limits> using
namespace std ; // Kadane's algorithm to find the maximum sum // subarray in a 1D array int kadane (
vector < int > & temp ) { int rows = temp . size () ; int currSum = 0 ; int maxSum = INT_MIN ; for ( int i = 0
; i < rows ; i ++ ) { currSum += temp [ i ] ; // Update maxSum if the current sum is greater if ( maxSum <
currSum ) { maxSum = currSum ; } // If the current sum becomes negative, reset it to 0 if ( currSum < 0 )
{ currSum = 0 ; } } return maxSum ; } // Function to find the maximum sum rectangle in a 2D matrix int
maxRectSum ( vector < vector < int >> & mat ) { int rows = mat . size () ; int cols = mat [ 0 ]. size () ; int
maxSum = INT_MIN ; // Initialize a temporary array to store row-wise // sums between left and right
boundaries vector < int > temp ( rows ) ; // Check for all possible left and right boundaries for ( int left = 0
; left < cols ; left ++ ) { // Reset the temporary array for each new left boundary for ( int i = 0 ; i < rows ; i
++ ) temp [ i ] = 0 ; for ( int right = left ; right < cols ; right ++ ) { // Update the temporary array with the
current // column's values for ( int row = 0 ; row < rows ; row ++ ) { temp [ row ] += mat [ row ][ right ] ; } //

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Find the maximum sum of the subarray for the // current column boundaries int sum = kadane ( temp );
// Update the maximum sum found so far maxSum = max ( maxSum , sum ); } } return maxSum ; } int
main () { vector < vector < int >> mat = { { 1 , 2 , -1 , -4 , -20 }, { -8 , -3 , 4 , 2 , 1 }, { 3 , 8 , 10 , 1 , 3 }, { -4 ,
-1 , 1 , 7 , -6 } }; int res = maxRectSum ( mat ); cout << res << endl ; return 0 ; } Java class GfG { //
Kadane's algorithm to find the maximum sum subarray // in a 1D array static int kadane ( int [] temp ) {
int rows = temp . length ; int currSum = 0 ; int maxSum = Integer . MIN_VALUE ; for ( int i = 0 ; i < rows ;
i ++ ) { currSum += temp [ i ] ; // Update maxSum if the current sum is greater if ( maxSum < currSum ) {
maxSum = currSum ; } // If the current sum becomes negative, reset it // to 0 if ( currSum < 0 ) {
currSum = 0 ; } } return maxSum ; } // Function to find the maximum sum rectangle in a 2D // matrix
static int maxRectSum ( int [][] mat ) { int rows = mat . length ; int cols = mat [ 0 ] . length ; int maxSum =
Integer . MIN_VALUE ; // Initialize a temporary array to store row-wise // sums between left and right
boundaries int [] temp = new int [ rows ] ; // Check for all possible left and right boundaries for ( int left =
0 ; left < cols ; left ++ ) { // Reset the temporary array for each new left // boundary for ( int i = 0 ; i < rows
; i ++ ) { temp [ i ] = 0 ; } for ( int right = left ; right < cols ; right ++ ) { // Update the temporary array with
the // current column's values for ( int row = 0 ; row < rows ; row ++ ) { temp [ row ] += mat [ row ][ right ]
; } // Find the maximum sum of the subarray for // the current column boundaries int sum = kadane (
temp ); // Update the maximum sum found so far maxSum = Math . max ( maxSum , sum ); } } return
maxSum ; } public static void main ( String [] args ) { int [][] mat = { { 1 , 2 , -1 , -4 , -20 }, { -8 , -3 , 4 ,
2 , 1 }, { 3 , 8 , 10 , 1 , 3 }, { -4 , -1 , 1 , 7 , -6 } }; int res = maxRectSum ( mat ); System . out . println (
res ); } } Python def kadane ( temp ): rows = len ( temp ) currSum = 0 maxSum = float ( '-inf' ) for i in
range ( rows ): currSum += temp [ i ] # Update maxSum if the current sum is greater if maxSum <
currSum : maxSum = currSum # If the current sum becomes negative, reset it to 0 if currSum < 0 :
currSum = 0 return maxSum def maxRectSum ( mat ): rows = len ( mat ) cols = len ( mat [ 0 ] ) maxSum =
float ( '-inf' ) # Initialize a temporary array to store row-wise # sums between left and right boundaries
temp = [ 0 ] * rows # Check for all possible left and right boundaries for left in range ( cols ): # Reset the
temporary array for each new left # boundary temp = [ 0 ] * rows for right in range ( left , cols ): # Update
the temporary array with the current # column's values for row in range ( rows ): temp [ row ] += mat [
row ][ right ] # Find the maximum sum of the subarray for the # current column boundaries sumValue =
kadane ( temp ) # Update the maximum sum found so far maxSum = max ( maxSum , sumValue )
return maxSum if __name__ == "__main__": mat = [ [ 1 , 2 , -1 , -4 , -20 ], [ -8 , -3 , 4 , 2 , 1 ], [ 3 , 8 ,
10 , 1 , 3 ], [ -4 , -1 , 1 , 7 , -6 ] ] res = maxRectSum ( mat ) print ( res ) C# using System ; class GfG {
// Function to apply Kadane's algorithm to find the // maximum sum subarray static int kadane ( int []
temp ) { int rows = temp . Length ; int currSum = 0 ; int maxSum = int . MinValue ; for ( int i = 0 ; i < rows
; i ++ ) { currSum += temp [ i ] ; // Update maxSum if the current sum is greater if ( maxSum < currSum )
{ maxSum = currSum ; } // If the current sum becomes negative, reset it // to 0 if ( currSum < 0 ) {
currSum = 0 ; } } return maxSum ; } // Function to find the maximum sum of submatrix static int
maxRectSum ( int [, ] mat ) { int rows = mat . GetLength ( 0 ) ; int cols = mat . GetLength ( 1 ) ; int
maxSum = int . MinValue ; // Initialize a temporary array to store row-wise // sums between left and right
boundaries int [] temp = new int [ rows ] ; // Check for all possible left and right boundaries for ( int left =
0 ; left < cols ; left ++ ) { // Reset the temporary array for each new left // boundary Array . Clear ( temp ,
0 , rows ); for ( int right = left ; right < cols ; right ++ ) { // Update the temporary array with the // current
column's values for ( int row = 0 ; row < rows ; row ++ ) { temp [ row ] += mat [ row , right ] ; } // Find the
maximum sum of the subarray for // the current column boundaries int sumValue = kadane ( temp ); //
Update the maximum sum found so far maxSum = Math . Max ( maxSum , sumValue ); } } return
maxSum ; } static void Main () { int [, ] mat = { { 1 , 2 , -1 , -4 , -20 }, { -8 , -3 , 4 , 2 , 1 }, { 3 , 8 , 10 , 1 ,
3 }, { -4 , -1 , 1 , 7 , -6 } }; int res = maxRectSum ( mat ); Console . WriteLine ( res ); } } JavaScript //
Function to apply Kadane's algorithm to find the maximum // sum subarray function kadane ( temp ) {
let rows = temp . length ; let currSum = 0 ; let maxSum = - Infinity ; for ( let i = 0 ; i < rows ; i ++ ) {
currSum += temp [ i ] ; // Update maxSum if the current sum is greater if ( maxSum < currSum ) {
maxSum = currSum ; } // If the current sum becomes negative, reset it to // 0 if ( currSum < 0 ) {
currSum = 0 ; } } return maxSum ; } // Function to find the maximum sum of submatrix function
maxRectSum ( mat ) { const rows = mat . length ; const cols = mat [ 0 ] . length ; let maxSum = - Infinity ;
// Initialize a temporary array to store row-wise sums // between left and right boundaries let temp = new
Array ( rows ). fill ( 0 ); // Check for all possible left and right boundaries for ( let left = 0 ; left < cols ; left
++ ) { // Reset the temporary array for each new left // boundary temp . fill ( 0 ); for ( let right = left ; right
< cols ; right ++ ) { // Update the temporary array with the current // column's values for ( let row = 0 ;
row < rows ; row ++ ) { temp [ row ] += mat [ row ][ right ] ; } // Find the maximum sum of the subarray for

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the // current column boundaries let sumValue = kadane ( temp , rows ); // Update the maximum sum
found so far maxSum = Math . max ( maxSum , sumValue ); } } return maxSum ; } // Driver Code const
mat = [ [ 1 , 2 , - 1 , - 4 , - 20 ], [ - 8 , - 3 , 4 , 2 , 1 ], [ 3 , 8 , 10 , 1 , 3 ], [ - 4 , - 1 , 1 , 7 , - 6 ] ]; const res =
maxRectSum ( mat ); console . log ( res ); Time complexity :  $O(n*m^2)$  Auxiliary Space:  $O(n)$  Comment
Article Tags: Article Tags: Dynamic Programming Matrix DSA Arrays Amazon Samsung Accolite
FactSet + 4 More
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