

# Eulerian path and circuit for undirected graph - GeeksforGeeks

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Courses Tutorials Practice Jobs DSA Tutorial Interview Questions Quizzes Must Do Advanced DSA System Design Aptitude Puzzles Interview Corner DSA Python Technical Scripter 2026 Explore DSA Fundamentals Logic Building Problems Analysis of Algorithms Data Structures Array Data Structure String in Data Structure Hashing in Data Structure Linked List Data Structure Stack Data Structure Queue Data Structure Tree Data Structure Graph Data Structure Trie Data Structure Algorithms Searching Algorithms Sorting Algorithms Introduction to Recursion Greedy Algorithms Tutorial Graph Algorithms Dynamic Programming or DP Bitwise Algorithms Advanced Segment Tree Binary Indexed Tree or Fenwick Tree Square Root (Sqrt) Decomposition Algorithm Binary Lifting Geometry Interview Preparation Interview Corner GfG160 Practice Problem GeeksforGeeks Practice - Leading Online Coding Platform Problem of The Day - Develop the Habit of Coding DSA Course 90% Refund Eulerian path and circuit for undirected graph Last Updated : 8 Jun, 2025 Given an undirected connected graph with  $v$  nodes, and  $e$  edges, with adjacency list  $\text{adj}$ . We need to write a function that returns 2 if the graph contains an eulerian circuit or cycle , else if the graph contains an eulerian path, returns 1, otherwise, returns 0. A graph is said to be Eulerian if it contains an Eulerian Cycle, a cycle that visits every edge exactly once and starts and ends at the same vertex. If a graph contains an Eulerian Path, a path that visits every edge exactly once but starts and ends at different vertices, it is called Semi-Eulerian . Examples: Input: Output: 1 Input: Output: 2 Input: Output: 0 Try it on GfG Practice The problem can be framed as follows: "Is it possible to draw a graph without lifting your pencil from the paper and without retracing any edge?" At first glance, this may seem similar to the Hamiltonian Path problem, which is NP-complete for general graphs. However, determining whether a graph has an Eulerian Path or Cycle is much more efficient: It can be solved in  $O(v + e)$  time. Approach: The idea is to use some key properties of undirected graphs that help determine whether they are Eulerian (i.e., contain an Eulerian Path or Cycle) or not. Eulerian Cycle A graph has an Eulerian Cycle if and only if the below two conditions are true All vertices with non-zero degree are part of a single connected component. (We ignore isolated vertices— those with zero degree — as they do not affect the cycle.) Every vertex in the graph has an even degree. Eulerian Path A graph has an Eulerian Path if and only if the below two conditions are true All vertices with non-zero degree must belong to the same connected component. (Same as Eulerian Cycle) Exactly Zero or Two Vertices with Odd Degree: If zero vertices have odd degree  $\rightarrow$  Eulerian Cycle exists (which is also a path). If two vertices have odd degree  $\rightarrow$  Eulerian Path exists (but not a cycle). If one vertex has odd degree  $\rightarrow$  Not possible in an undirected graph. (Because the sum of all degrees in an undirected graph is always even.) Note: A graph with no edges is trivially Eulerian . There are no edges to traverse, so by definition, it satisfies both Eulerian Path and Cycle conditions. How Does This Work? In an Eulerian Path , whenever we enter a vertex (except start and end), we must also leave it. So all intermediate vertices must have even degree. In an Eulerian Cycle , since we start and end at the same vertex, every vertex must have even degree. This ensures that every entry into a vertex can be paired with an exit. Steps to implement the above idea: Create an adjacency list to represent the graph and initialize a visited array for DFS traversal. Perform DFS starting from the first vertex having non-zero degree to check graph connectivity. After DFS, ensure all non-zero degree vertices were visited to confirm the graph is connected. Count the number of vertices with odd degree to classify the graph as Eulerian or not. If all degrees are even, the graph has an Eulerian Circuit ; if exactly two are odd, it's a Path . If more than two vertices have odd degree or graph isn't connected, it's not Eulerian . Return 2 for Circuit, 1 for Path, and 0 when graph fails Eulerian conditions. C++ // C++ program to check whether a graph is // Eulerian Path, Eulerian Circuit, or neither  

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#include <bits/stdc++.h>
using namespace std;
// DFS to check connectivity, excluding zero-degree vertices
void dfs ( int node , vector < int > adj [] , vector < bool > & visited ) { visited [ node ] = true ;
for ( int neighbor : adj [ node ] ) { if ( ! visited [ neighbor ] ) { dfs ( neighbor , adj , visited ); } } }
// Function to check Eulerian Path or Circuit
int isEulerCircuit ( int v , vector < int > adj [] ) { vector < bool > visited ( v , false );
// Find first vertex with non-zero degree
int start = -1 ;
for ( int i = 0 ; i < v ; i ++ ) { if ( adj [ i ]. size () > 0 ) { start = i ; break ; } }
// No edges: graph is trivially Eulerian if ( start == -1 ) { return 2 ; }
// DFS
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from the first non-zero degree vertex dfs ( start , adj , visited ); // Check if all non-zero degree vertices
are connected for ( int i = 0 ; i < v ; i ++ ) { if ( adj [ i ]. size () > 0 && ! visited [ i ] ) { return 0 ; // Not
connected } } // Count vertices with odd degree int odd = 0 ; for ( int i = 0 ; i < v ; i ++ ) { if ( adj [ i ]. size
() % 2 != 0 ) { odd ++ ; } } // Apply Eulerian rules if ( odd == 0 ) { return 2 ; } else if ( odd == 2 ) { return 1 ;
} else { return 0 ; } } // Driver code int main () { int v = 5 ; vector < int > adj [ 5 ] = { { 1 , 2 , 3 } , { 0 , 2 } , { 1 ,
0 } , { 0 , 4 } , { 3 } } ; cout << isEulerCircuit ( v , adj ); return 0 ; } Java // Java program to check whether a
graph is // Eulerian Path, Eulerian Circuit, or neither import java.util.* ; class GfG { // DFS to check
connectivity, excluding zero-degree vertices static void dfs ( int node , List < Integer >[] adj , boolean []
visited ) { visited [ node ] = true ; for ( int neighbor : adj [ node ] ) { if ( ! visited [ neighbor ] ) { dfs (
neighbor , adj , visited ); } } } // Function to check Eulerian Path or Circuit static int isEulerCircuit ( int v ,
List < Integer >[] adj ) { boolean [] visited = new boolean [ v ] ; // Find first vertex with non-zero degree
int start = - 1 ; for ( int i = 0 ; i < v ; i ++ ) { if ( adj [ i ]. size () > 0 ) { start = i ; break ; } } // No edges:
graph is trivially Eulerian if ( start == - 1 ) { return 2 ; } // DFS from the first non-zero degree vertex dfs (
start , adj , visited ); // Check if all non-zero degree vertices are connected for ( int i = 0 ; i < v ; i ++ ) { if ( adj [ i ]. size ()
> 0 && ! visited [ i ] ) { return 0 ; // Not connected } } // Count vertices with odd degree int odd = 0 ; for ( int i = 0 ;
i < v ; i ++ ) { if ( adj [ i ]. size () % 2 != 0 ) { odd ++ ; } } // Apply Eulerian rules if ( odd == 0 ) { return 2 ;
} else if ( odd == 2 ) { return 1 ; } else { return 0 ; } } public static void main ( String []
args ) { int v = 5 ; List < Integer >[] adj = new ArrayList [ v ] ; for ( int i = 0 ; i < v ; i ++ ) { adj [ i ] = new
ArrayList <> (); } adj [ 0 ]. addAll ( Arrays . asList ( 1 , 2 , 3 ) ); adj [ 1 ]. addAll ( Arrays . asList ( 0 , 2 ) );
adj [ 2 ]. addAll ( Arrays . asList ( 1 , 0 ) ); adj [ 3 ]. addAll ( Arrays . asList ( 0 , 4 ) ); adj [ 4 ]. addAll (
Arrays . asList ( 3 ) ); System . out . println ( isEulerCircuit ( v , adj )); } } Python # Python program to
check whether a graph is # Eulerian Path, Eulerian Circuit, or neither # DFS to check connectivity,
excluding zero-degree vertices def dfs ( node , adj , visited ): visited [ node ] = True for neighbor in adj [
node ]: if not visited [ neighbor ]: dfs ( neighbor , adj , visited ) # Function to check Eulerian Path or
Circuit def isEulerCircuit ( v , adj ): visited = [ False ] * v # Find first vertex with non-zero degree start = -
1 for i in range ( v ): if len ( adj [ i ]) > 0 : start = i break # No edges: graph is trivially Eulerian if start == -
1 : return 2 # DFS from the first non-zero degree vertex dfs ( start , adj , visited ) # Check if all non-zero
degree vertices are connected for i in range ( v ): if len ( adj [ i ]) > 0 and not visited [ i ]: return 0 # Not
connected # Count vertices with odd degree odd = 0 for i in range ( v ): if len ( adj [ i ]) % 2 != 0 : odd += 1 # Apply
Eulerian rules if odd == 0 : return 2 elif odd == 2 : return 1 else : return 0 if __name__ ==
"__main__" : v = 5 adj = [ [ 1 , 2 , 3 ] , [ 0 , 2 ] , [ 1 , 0 ] , [ 0 , 4 ] , [ 3 ] ] print ( isEulerCircuit ( v , adj )) C# // C#
program to check whether a graph is // Eulerian Path, Eulerian Circuit, or neither using System ;
using System.Collections.Generic ; class GfG { // DFS to check connectivity, excluding zero-degree
vertices static void dfs ( int node , List < int > [] adj , bool [] visited ) { visited [ node ] = true ; foreach ( int
neighbor in adj [ node ] ) { if ( ! visited [ neighbor ] ) { dfs ( neighbor , adj , visited ); } } } // Function to
check Eulerian Path or Circuit static int isEulerCircuit ( int v , List < int > [] adj ) { bool [] visited = new
bool [ v ] ; // Find first vertex with non-zero degree int start = - 1 ; for ( int i = 0 ; i < v ; i ++ ) { if ( adj [ i ].
Count > 0 ) { start = i ; break ; } } // No edges: graph is trivially Eulerian if ( start == - 1 ) { return 2 ; } // DFS
from the first non-zero degree vertex dfs ( start , adj , visited ); // Check if all non-zero degree
vertices are connected for ( int i = 0 ; i < v ; i ++ ) { if ( adj [ i ]. Count > 0 && ! visited [ i ] ) { return 0 ; // Not
connected } } // Count vertices with odd degree int odd = 0 ; for ( int i = 0 ; i < v ; i ++ ) { if ( adj [ i ].
Count % 2 != 0 ) { odd ++ ; } } // Apply Eulerian rules if ( odd == 0 ) { return 2 ; } else if ( odd == 2 ) { return 1 ;
} else { return 0 ; } } static void Main () { int v = 5 ; List < int > [] adj = new List < int > [ v ]; for ( int i = 0 ;
i < v ; i ++ ) { adj [ i ] = new List < int > (); } adj [ 0 ]. AddRange ( new int [] { 1 , 2 , 3 } ); adj [ 1 ].
AddRange ( new int [] { 0 , 2 } ); adj [ 2 ]. AddRange ( new int [] { 1 , 0 } ); adj [ 3 ]. AddRange ( new int [] { 0 ,
4 } ); adj [ 4 ]. AddRange ( new int [] { 3 } ); Console . WriteLine ( isEulerCircuit ( v , adj )); } } JavaScript // JavaScript
program to check whether a graph is // Eulerian Path, Eulerian Circuit, or neither // DFS to check connectivity,
excluding zero-degree vertices function dfs ( node , adj , visited ) { visited [ node ] = true ; for ( let
neighbor of adj [ node ] ) { if ( ! visited [ neighbor ] ) { dfs ( neighbor , adj , visited ); } } } // Function to
check Eulerian Path or Circuit function isEulerCircuit ( v , adj ) { let visited = new Array ( v ). fill ( false );
// Find first vertex with non-zero degree let start = - 1 ; for ( let i = 0 ; i < v ; i ++ ) { if ( adj [ i ]. length > 0 ) {
start = i ; break ; } } // No edges: graph is trivially Eulerian if ( start === - 1 ) { return 2 ; } // DFS from the
first non-zero degree vertex dfs ( start , adj , visited ); // Check if all non-zero degree vertices are connected
for ( let i = 0 ; i < v ; i ++ ) { if ( adj [ i ]. length > 0 && ! visited [ i ] ) { return 0 ; // Not connected } } // Count
vertices with odd degree let odd = 0 ; for ( let i = 0 ; i < v ; i ++ ) { if ( adj [ i ]. length % 2 !== 0 ) { odd ++ ;
} } // Apply Eulerian rules if ( odd === 0 ) { return 2 ; } else if ( odd === 2 ) { return 1 ; } else { return 0 ; } } // Driver
Code let v = 5 ; let adj = [ [ 1 , 2 , 3 ] , [ 0 , 2 ] , [ 1 , 0 ] , [ 0 , 4 ] , [ 3 ] ]

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], [ 0 , 4 ], [ 3 ] ]; console . log ( isEulerCircuit ( v , adj )); Output 1 Time Complexity:  $O(v + e)$ , DFS traverses all vertices and edges to check connectivity and degree. Space Complexity:  $O(v + e)$ , Space used for visited array and adjacency list to represent graph. Comment Article Tags: Article Tags: Graph DSA Euler-Circuit