

# Bridges in a graph - GeeksforGeeks

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Courses Tutorials Practice Jobs DSA Tutorial Interview Questions Quizzes Must Do Advanced DSA System Design Aptitude Puzzles Interview Corner DSA Python Technical Scripter 2026 Explore DSA Fundamentals Logic Building Problems Analysis of Algorithms Data Structures Array Data Structure String in Data Structure Hashing in Data Structure Linked List Data Structure Stack Data Structure Queue Data Structure Tree Data Structure Graph Data Structure Trie Data Structure Algorithms Searching Algorithms Sorting Algorithms Introduction to Recursion Greedy Algorithms Tutorial Graph Algorithms Dynamic Programming or DP Bitwise Algorithms Advanced Segment Tree Binary Indexed Tree or Fenwick Tree Square Root (Sqrt) Decomposition Algorithm Binary Lifting Geometry Interview Preparation Interview Corner GfG160 Practice Problem GeeksforGeeks Practice - Leading Online Coding Platform Problem of The Day - Develop the Habit of Coding DSA Course 90% Refund Bridges in a graph Last Updated : 23 Jul, 2025 Given an undirected Graph, The task is to find the Bridges in this Graph. An edge in an undirected connected graph is a bridge if removing it disconnects the graph. For a disconnected undirected graph, the definition is similar, a bridge is an edge removal that increases the number of disconnected components. Like Articulation Points , bridges represent vulnerabilities in a connected network and are useful for designing reliable networks. Try it on GfG Practice Examples: Input: Output: (0, 3) and (3, 4) Input: Output: (1, 6) Input: Output: (0, 1), (1, 2), and (2, 3) Naive Approach: Below is the idea to solve the problem: One by one remove all edges and see if the removal of an edge causes a disconnected graph. Follow the below steps to Implement the idea: For every edge  $(u, v)$ , do the following: Remove  $(u, v)$  from the graph See if the graph remains connected (either uses BFS or DFS) Add  $(u, v)$  back to the graph. Time Complexity:  $O(E^*(V+E))$  for a graph represented by an adjacency list. Auxiliary Space:  $O(V+E)$  Find Bridges in a graph using Tarjan's Algorithm. Before heading towards the approach understand which edge is termed as bridge. Suppose there exists a edge from  $u \rightarrow v$ , now after removal of this edge if  $v$  can't be reached by any other edges then  $u \rightarrow v$  edge is bridge . Our approach is based on this intuition, so take time and grasp it. ALGORITHM: - To implement this algorithm, we need the following data structures - visited[ ] = to keep track of the visited vertices to implement DFS disc[ ] = to keep track when for the first time that particular vertex is reached low[ ] = to keep track of the lowest possible time by which we can reach that vertex ' other than parent' so that if edge from parent is removed can the particular node can be reached other than parent. We will traverse the graph using DFS traversal but with slight modifications i.e. while traversing we will keep track of the parent node by which the particular node is reached because we will update the low[node] = min(low[all its adjacent node except parent]) hence we need to keep track of the parent. While traversing adjacent nodes let 'v' of a particular node let 'u', then 3 cases arise - 1. v is parent of u then, skip that iteration. 2. v is visited then, update the low of u i.e. low[u] = min( low[u] , disc[v] ) this arises when a node can be visited by more than one node, but low is to keep track of the lowest possible time so we will update it. 3. v is not visited then, call the DFS to traverse ahead now update the low[u] = min( low[u] , low[v] ) as we know v can't be parent cause we have handled that case first. now check if ( low[v] > disc[u] ) i.e. the lowest possible to time to reach 'v' is greater than 'u' this means we can't reach 'v' without 'u' so the edge  $u \rightarrow v$  is a bridge. Below is the implementation of the above approach: C++ // A C++ program to find bridges in a given undirected graph #include <bits/stdc++.h> using namespace std ; // A class that represents an undirected graph class Graph { int V ; // No. of vertices list < int > \* adj ; // A dynamic array of adjacency lists void bridgeUtil ( int u , vector < bool >& visited , vector < int >& disc , vector < int >& low , int parent ) ; public : Graph ( int V ) ; // Constructor void addEdge ( int v , int w ) ; // to add an edge to graph void bridge () ; // prints all bridges } ; Graph :: Graph ( int V ) { this -> V = V ; adj = new list < int > [ V ] ; } void Graph :: addEdge ( int v , int w ) { adj [ v ]. push\_back ( w ) ; adj [ w ]. push\_back ( v ) ; // Note: the graph is undirected } // A recursive function that finds and prints bridges using // DFS traversal // u --> The vertex to be visited next // visited[] --> keeps track of visited vertices // disc[] --> Stores discovery times of visited vertices // parent[] --> Stores parent vertices in DFS tree void Graph :: bridgeUtil ( int u , vector < bool >& visited , vector < int >& disc , vector < int >& low , int parent ) { // A static variable is used for simplicity, we can // avoid use of static variable by passing a pointer. static int time = 0 ; // Mark the current node as visited visited [ u ] = true ; // Initialize discovery time and low value disc [ u ] = low [ u ] = ++ time ; // Go through all vertices adjacent to this list < int >:: iterator i ; }

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for ( i = adj [ u ]. begin (); i != adj [ u ]. end (); ++ i ) { int v = * i ; // v is current adjacent of u // 1. If v is parent if ( v == parent ) continue ; //2. If v is visited if ( visited [ v ]) { low [ u ] = min ( low [ u ], disc [ v ]); } //3. If v is not visited else { parent = u ; bridgeUtil ( v , visited , disc , low , parent ); // update the low of u as it's quite possible // a connection exists from v's descendants to u low [ u ] = min ( low [ u ], low [ v ]); // if the lowest possible time to reach vertex v // is greater than discovery time of u it means // that v can be only be reached by vertex above v // so if that edge is removed v can't be reached so it's a bridge if ( low [ v ] > disc [ u ]) cout << u << " " << v << endl ; } } } // DFS based function to find all bridges. It uses recursive // function bridgeUtil() void Graph :: bridge () { // Mark all the vertices as not visited disc and low as -1 vector < bool > visited ( V , false ); vector < int > disc ( V , -1 ); vector < int > low ( V , -1 ); // Initially there is no parent so let it be -1 int parent = -1 ; // Call the recursive helper function to find Bridges // in DFS tree rooted with vertex 'l' for ( int i = 0 ; i < V ; i ++ ) if ( visited [ i ] == false ) bridgeUtil ( i , visited , disc , low , parent ); } // Driver program to test above function int main () { // Create graphs given in above diagrams cout << " \n Bridges in first graph \n " ; Graph g1 ( 5 ); g1 . addEdge ( 1 , 0 ); g1 . addEdge ( 0 , 2 ); g1 . addEdge ( 2 , 1 ); g1 . addEdge ( 0 , 3 ); g1 . addEdge ( 3 , 4 ); g1 . bridge (); cout << " \n Bridges in second graph \n " ; Graph g2 ( 4 ); g2 . addEdge ( 0 , 1 ); g2 . addEdge ( 1 , 2 ); g2 . addEdge ( 2 , 3 ); g2 . bridge (); cout << " \n Bridges in third graph \n " ; Graph g3 ( 7 ); g3 . addEdge ( 0 , 1 ); g3 . addEdge ( 1 , 2 ); g3 . addEdge ( 2 , 0 ); g3 . addEdge ( 1 , 3 ); g3 . addEdge ( 1 , 4 ); g3 . addEdge ( 1 , 6 ); g3 . addEdge ( 3 , 5 ); g3 . addEdge ( 4 , 5 ); g3 . bridge (); return 0 ; } Java // A Java program to find bridges in a given undirected graph import java.io.* ; import java.util.* ; import java.util.LinkedList ; // This class represents a undirected graph using adjacency list // representation class Graph { private int V ; // No. of vertices // Array of lists for Adjacency List Representation private LinkedList < Integer > adj [] ; int time = 0 ; static final int NIL = - 1 ; // Constructor @SuppressWarnings ( "unchecked" ) Graph ( int v ) { V = v ; adj = new LinkedList [ v ] ; for ( int i = 0 ; i < v ; ++ i ) adj [ i ] = new LinkedList (); } // Function to add an edge into the graph void addEdge ( int v , int w ) { adj [ v ]. add ( w ); // Add w to v's list. adj [ w ]. add ( v ); //Add v to w's list } // A recursive function that finds and prints bridges // using DFS traversal // u --> The vertex to be visited next // visited[] --> keeps track of visited vertices // disc[] --> Stores discovery times of visited vertices // parent[] --> Stores parent vertices in DFS tree void bridgeUtil ( int u , boolean visited [] , int disc [] , int low [] , int parent [] ) { // Mark the current node as visited visited [ u ] = true ; // Initialize discovery time and low value disc [ u ] = low [ u ] = ++ time ; // Go through all vertices adjacent to this Iterator < Integer > i = adj [ u ]. iterator (); while ( i . hasNext ()) { int v = i . next (); // v is current adjacent of u // If v is not visited yet, then make it a child // of u in DFS tree and recur for it. // If v is not visited yet, then recur for it if ( ! visited [ v ] ) { parent [ v ] = u ; bridgeUtil ( v , visited , disc , low , parent ); // Check if the subtree rooted with v has a // connection to one of the ancestors of u low [ u ] = Math . min ( low [ u ] , low [ v ]); // If the lowest vertex reachable from subtree // under v is below u in DFS tree, then u-v is // a bridge if ( low [ v ] > disc [ u ]) System . out . println ( u + " " + v ); } // Update low value of u for parent function calls. else if ( v != parent [ u ] ) low [ u ] = Math . min ( low [ u ] , disc [ v ]); } } // DFS based function to find all bridges. It uses recursive // function bridgeUtil() void bridge () { // Mark all the vertices as not visited boolean visited [] = new boolean [ V ] ; int disc [] = new int [ V ] ; int low [] = new int [ V ] ; int parent [] = new int [ V ] ; // Initialize parent and visited, and ap(articulation point) // arrays for ( int i = 0 ; i < V ; i ++ ) { parent [ i ] = NIL ; visited [ i ] = false ; } // Call the recursive helper function to find Bridges // in DFS tree rooted with vertex 'l' for ( int i = 0 ; i < V ; i ++ ) if ( visited [ i ] == false ) bridgeUtil ( i , visited , disc , low , parent ); } public static void main ( String args [] ) { // Create graphs given in above diagrams System . out . println ( "Bridges in first graph " ); Graph g1 = new Graph ( 5 ); g1 . addEdge ( 1 , 0 ); g1 . addEdge ( 0 , 2 ); g1 . addEdge ( 2 , 1 ); g1 . addEdge ( 0 , 3 ); g1 . addEdge ( 3 , 4 ); g1 . bridge (); System . out . println (); System . out . println ( "Bridges in Second graph " ); Graph g2 = new Graph ( 4 ); g2 . addEdge ( 0 , 1 ); g2 . addEdge ( 1 , 2 ); g2 . addEdge ( 2 , 3 ); g2 . bridge (); System . out . println (); System . out . println ( "Bridges in Third graph " ); Graph g3 = new Graph ( 7 ); g3 . addEdge ( 0 , 1 ); g3 . addEdge ( 1 , 2 ); g3 . addEdge ( 2 , 0 ); g3 . addEdge ( 1 , 3 ); g3 . addEdge ( 1 , 4 ); g3 . addEdge ( 1 , 6 ); g3 . addEdge ( 3 , 5 ); g3 . addEdge ( 4 , 5 ); g3 . bridge (); } } // This code is contributed by Aakash Hasija Python # Python program to find bridges in a given undirected graph #Complexity : O(V+E) from collections import defaultdict #This class represents an undirected graph using adjacency list representation class Graph : def __init__ ( self , vertices ): self . V = vertices #No. of vertices self . graph = defaultdict ( list ) # default dictionary to store graph self . Time = 0 # function to add an edge to graph def addEdge ( self , u , v ): self . graph [ u ]. append ( v ) self . graph [ v ]. append ( u ) "A recursive function that finds and prints bridges using DFS traversal u --> The vertex to be visited next visited[] --> keeps track of visited vertices disc[] --> Stores discovery times of visited vertices parent[] --> Stores parent vertices in DFS

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tree"" def bridgeUtil ( self , u , visited , parent , low , disc ): # Mark the current node as visited and print it
visited [ u ] = True # Initialize discovery time and low value disc [ u ] = self . Time low [ u ] = self . Time
self . Time += 1 #Recur for all the vertices adjacent to this vertex for v in self . graph [ u ]: # If v is not
visited yet, then make it a child of u # in DFS tree and recur for it if visited [ v ] == False : parent [ v ] = u
self . bridgeUtil ( v , visited , parent , low , disc ) # Check if the subtree rooted with v has a connection to
# one of the ancestors of u low [ u ] = min ( low [ u ], low [ v ]) "" If the lowest vertex reachable from
subtree under v is below u in DFS tree, then u-v is a bridge"" if low [ v ] > disc [ u ]: print ( " %d %d " % (
u , v )) elif v != parent [ u ]: # Update low value of u for parent function calls. low [ u ] = min ( low [ u ],
disc [ v ]) # DFS based function to find all bridges. It uses recursive # function bridgeUtil() def bridge (
self ): # Mark all the vertices as not visited and Initialize parent and visited, # and ap(articulation point)
arrays visited = [ False ] * ( self . V ) disc = [ float ( "Inf" )] * ( self . V ) low = [ float ( "Inf" )] * ( self . V )
parent = [ - 1 ] * ( self . V ) # Call the recursive helper function to find bridges # in DFS tree rooted with
vertex 'i' for i in range ( self . V ): if visited [ i ] == False : self . bridgeUtil ( i , visited , parent , low , disc )
# Create a graph given in the above diagram g1 = Graph ( 5 ) g1 . addEdge ( 1 , 0 ) g1 . addEdge ( 0 , 2 )
g1 . addEdge ( 2 , 1 ) g1 . addEdge ( 0 , 3 ) g1 . addEdge ( 3 , 4 ) print ( " Bridges in first graph " ) g1 .
bridge () g2 = Graph ( 4 ) g2 . addEdge ( 0 , 1 ) g2 . addEdge ( 1 , 2 ) g2 . addEdge ( 2 , 3 ) print ( " \n
Bridges in second graph " ) g2 . bridge () g3 = Graph ( 7 ) g3 . addEdge ( 0 , 1 ) g3 . addEdge ( 1 , 2 )
g3 . addEdge ( 2 , 0 ) g3 . addEdge ( 1 , 3 ) g3 . addEdge ( 1 , 4 ) g3 . addEdge ( 1 , 6 ) g3 . addEdge ( 3 ,
5 ) g3 . addEdge ( 4 , 5 ) print ( " \n Bridges in third graph " ) g3 . bridge () #This code is contributed by
Neelam Yadav C# // A C# program to find bridges // in a given undirected graph using System ; using
System.Collections.Generic ; // This class represents a undirected graph // using adjacency list
representation public class Graph { private int V ; // No. of vertices // Array of lists for Adjacency List
Representation private List < int > [] adj ; int time = 0 ; static readonly int NIL = - 1 ; // Constructor Graph
( int v ) { V = v ; adj = new List < int > [ v ]; for ( int i = 0 ; i < v ; ++ i ) adj [ i ] = new List < int > (); } // Function to add an edge into the graph void addEdge ( int v , int w ) { adj [ v ]. Add ( w ); // Add w to v's
list. adj [ w ]. Add ( v ); //Add v to w's list } // A recursive function that finds and prints bridges // using
DFS traversal // u --> The vertex to be visited next // visited[] --> keeps track of visited vertices // disc[]
--> Stores discovery times of visited vertices // parent[] --> Stores parent vertices in DFS tree void
bridgeUtil ( int u , bool [] visited , int [] disc , int [] low , int [] parent ) { // Mark the current node as visited
visited [ u ] = true ; // Initialize discovery time and low value disc [ u ] = low [ u ] = ++ time ; // Go through
all vertices adjacent to this foreach ( int i in adj [ u ] ) { int v = i ; // v is current adjacent of u // If v is not
visited yet, then make it a child // of u in DFS tree and recur for it. // If v is not visited yet, then recur for it
if ( ! visited [ v ] ) { parent [ v ] = u ; bridgeUtil ( v , visited , disc , low , parent ); // Check if the subtree
rooted with v has a // connection to one of the ancestors of u low [ u ] = Math . Min ( low [ u ], low [ v ]); // If the lowest vertex reachable from subtree // under v is below u in DFS tree, then u-v is // a bridge if ( low [ v ] > disc [ u ])
Console . WriteLine ( u + " " + v ); } // Update low value of u for parent function calls. else if ( v != parent [ u ] )
low [ u ] = Math . Min ( low [ u ], disc [ v ]); } } // DFS based function to find all
bridges. It uses recursive // function bridgeUtil() void bridge () { // Mark all the vertices as not visited
bool [] visited = new bool [ V ]; int [] disc = new int [ V ]; int [] low = new int [ V ]; int [] parent = new int [ V ]
]; // Initialize parent and visited, // and ap(articulation point) arrays for ( int i = 0 ; i < V ; i ++ ) { parent [ i ]
= NIL ; visited [ i ] = false ; } // Call the recursive helper function to find Bridges // in DFS tree rooted with
vertex 'i' for ( int i = 0 ; i < V ; i ++ ) if ( visited [ i ] == false ) bridgeUtil ( i , visited , disc , low , parent );
} // Driver code public static void Main ( String [] args ) { // Create graphs given in above diagrams
Console . WriteLine ( " Bridges in first graph " ); Graph g1 = new Graph ( 5 ); g1 . addEdge ( 1 , 0 ); g1 . addEdge
( 0 , 2 ); g1 . addEdge ( 2 , 1 ); g1 . addEdge ( 0 , 3 ); g1 . addEdge ( 3 , 4 ); g1 . bridge (); Console .
WriteLine (); Console . WriteLine ( " Bridges in Second graph " ); Graph g2 = new Graph ( 4 ); g2 .
addEdge ( 0 , 1 ); g2 . addEdge ( 1 , 2 ); g2 . addEdge ( 2 , 3 ); g2 . bridge (); Console . WriteLine ();
Console . WriteLine ( " Bridges in Third graph " ); Graph g3 = new Graph ( 7 ); g3 . addEdge ( 0 , 1 ); g3 .
addEdge ( 1 , 2 ); g3 . addEdge ( 2 , 0 ); g3 . addEdge ( 1 , 3 ); g3 . addEdge ( 1 , 4 ); g3 . addEdge ( 1 ,
6 ); g3 . addEdge ( 3 , 5 ); g3 . addEdge ( 4 , 5 ); g3 . bridge (); } } // This code is contributed by
Rajput-Ji JavaScript < script > // A Javascript program to find bridges in a given undirected graph // This
class represents a directed graph using adjacency // list representation class Graph { // Constructor
constructor ( v ) { this . V = v ; this . adj = new Array ( v ); this . NIL = - 1 ; this . time = 0 ; for ( let i = 0 ;
i < v ; ++ i ) this . adj [ i ] = []; } //Function to add an edge into the graph addEdge ( v , w ) { this . adj [ v ].
push ( w ); //Note that the graph is undirected. this . adj [ w ]. push ( v ); } // A recursive function that
finds and prints bridges // using DFS traversal // u --> The vertex to be visited next // visited[] --> keeps
track of visited vertices // disc[] --> Stores discovery times of visited vertices // parent[] --> Stores parent

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vertices in DFS tree bridgeUtil ( u , visited , disc , low , parent ) { // Mark the current node as visited
visited [ u ] = true ; // Initialize discovery time and low value disc [ u ] = low [ u ] = ++ this . time ; // Go
through all vertices adjacent to this for ( let i of this . adj [ u ]) { let v = i ; // v is current adjacent of u // If v
is not visited yet, then make it a child // of u in DFS tree and recur for it. // If v is not visited yet, then
recur for it if ( ! visited [ v ]) { parent [ v ] = u ; this . bridgeUtil ( v , visited , disc , low , parent ); // Check if
the subtree rooted with v has a // connection to one of the ancestors of u low [ u ] = Math . min ( low [ u ]
], low [ v ]); // If the lowest vertex reachable from subtree // under v is below u in DFS tree, then u-v is //
a bridge if ( low [ v ] > disc [ u ]) document . write ( u + " " + v + "<br>" ); } // Update low value of u for
parent function calls. else if ( v != parent [ u ]) low [ u ] = Math . min ( low [ u ], disc [ v ]); } } // DFS based
function to find all bridges. It uses recursive // function bridgeUtil() bridge () { // Mark all the vertices as
not visited let visited = new Array ( this . V ); let disc = new Array ( this . V ); let low = new Array ( this . V
); let parent = new Array ( this . V ); // Initialize parent and visited, and ap(articulation point) // arrays for
( let i = 0 ; i < this . V ; i ++ ) { parent [ i ] = this . NIL ; visited [ i ] = false ; } // Call the recursive helper
function to find Bridges // in DFS tree rooted with vertex 'i' for ( let i = 0 ; i < this . V ; i ++ ) if ( visited [ i ]
== false ) this . bridgeUtil ( i , visited , disc , low , parent ); } } // Create graphs given in above diagrams
document . write ( "Bridges in first graph <br>" ); let g1 = new Graph ( 5 ); g1 . addEdge ( 1 , 0 ); g1 .
addEdge ( 0 , 2 ); g1 . addEdge ( 2 , 1 ); g1 . addEdge ( 0 , 3 ); g1 . addEdge ( 3 , 4 ); g1 . bridge (); document .
write ( "<br>" ); document . write ( "Bridges in Second graph<br>" ); let g2 = new Graph ( 4 ); g2 . addEdge ( 0 ,
1 ); g2 . addEdge ( 1 , 2 ); g2 . addEdge ( 2 , 3 ); g2 . bridge (); document . write ( "<br>" ); document .
write ( "Bridges in Third graph <br>" ); let g3 = new Graph ( 7 ); g3 . addEdge ( 0 , 1 ); g3 . addEdge ( 1 , 2 );
g3 . addEdge ( 2 , 0 ); g3 . addEdge ( 1 , 3 ); g3 . addEdge ( 1 , 4 ); g3 . addEdge ( 1 , 6 ); g3 . addEdge ( 3 , 5 );
g3 . addEdge ( 4 , 5 ); g3 . bridge (); // This code is contributed by avanitracchadiya2155 </script> Output Bridges in first graph 3 4 0 3

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Bridges in second graph 2 3 1 2 0 1

Bridges in third graph 1 6 Time Complexity: O(V+E), The above approach uses simple DFS along with Tarjan's Algorithm. So time complexity is the same as DFS which is O(V+E) for adjacency list representation of the graph. Auxiliary Space: O(V) is used for visited, disc and low arrays. Comment Article Tags: Article Tags: Graph DSA graph-connectivity