

Segmented Sieve - GeeksforGeeks

Source: <https://www.geeksforgeeks.org/segmented-sieve/>

Courses Tutorials Practice Jobs DSA Tutorial Interview Questions Quizzes Must Do Advanced DSA System Design Aptitude Puzzles Interview Corner DSA Python Technical Scripter 2026 Explore DSA Fundamentals Logic Building Problems Analysis of Algorithms Data Structures Array Data Structure String in Data Structure Hashing in Data Structure Linked List Data Structure Stack Data Structure Queue Data Structure Tree Data Structure Graph Data Structure Trie Data Structure Algorithms Searching Algorithms Sorting Algorithms Introduction to Recursion Greedy Algorithms Tutorial Graph Algorithms Dynamic Programming or DP Bitwise Algorithms Advanced Segment Tree Binary Indexed Tree or Fenwick Tree Square Root (Sqrt) Decomposition Algorithm Binary Lifting Geometry Interview Preparation Interview Corner GfG160 Practice Problem GeeksforGeeks Practice - Leading Online Coding Platform Problem of The Day - Develop the Habit of Coding DSA Course 90% Refund Segmented Sieve Last Updated : 23 Jul, 2025 Given a number n, print all primes smaller than n. Input: N = 10 Output: 2, 3, 5, 7 Explanation : The output "2, 3, 5, 7" for input N = 10 represents the list of the prime numbers less than or equal to 10. Input: N = 5 Output: 2, 3, 5 Explanation : The output "2, 3, 5" for input N = 5 represents the list of the prime numbers less than or equal to 5. Try it on GfG Practice A Naive approach is to run a loop from 0 to n-1 and check each number for primeness. A Better Approach is to use Simple Sieve of Eratosthenes . C++ #include <iostream> #include <vector> void simpleSieve (int limit) { // Create a boolean array "mark[0..limit-1]" and // initialize all entries of it as true. A value // in mark[p] will finally be false if 'p' is Not // a prime, else true. std :: vector < bool > mark (limit , true); // One by one traverse all numbers so that their // multiples can be marked as composite. for (int p = 2 ; p * p < limit ; p ++) { // If p is not changed, then it is a prime if (mark [p] == true) { // Update all multiples of p for (int i = p * p ; i < limit ; i += p) mark [i] = false ; } } // Print all prime numbers and store them in prime for (int p = 2 ; p < limit ; p ++) if (mark [p] == true) std :: cout << p << " " ; } int main () { int limit = 100 ; simpleSieve (limit) ; return 0 ; } C // This functions finds all primes smaller than 'limit' // using simple sieve of eratosthenes. void simpleSieve (int limit) { // Create a boolean array "mark[0..limit-1]" and // initialize all entries of it as true. A value // in mark[p] will finally be false if 'p' is Not // a prime, else true. bool mark [limit] ; for (int i = 0 ; i < limit ; i ++) { mark [i] = true ; } // One by one traverse all numbers so that their // multiples can be marked as composite. for (int p = 2 ; p * p < limit ; p ++) { // If p is not changed, then it is a prime if (mark [p] == true) { // Update all multiples of p for (int i = p * p ; i < limit ; i += p) mark [i] = false ; } } // Print all prime numbers and store them in prime for (int p = 2 ; p < limit ; p ++) if (mark [p] == true) cout << p << " " ; } Java import java.util.Arrays ; public class Main { public static void main (String [] args) { int limit = 100 ; simpleSieve (limit) ; } // This function finds all primes smaller than 'limit' // using simple sieve of eratosthenes. static void simpleSieve (int limit) { // Create a boolean array "mark[0..limit-1]" and // initialize all entries of it as true. A value // in mark[p] will finally be false if 'p' is Not // a prime, else true. boolean [] mark = new boolean [limit] ; Arrays . fill (mark , true) ; // One by one traverse all numbers so that their // multiples can be marked as composite. for (int p = 2 ; p * p < limit ; p ++) { // If p is not changed, then it is a prime if (mark [p] == true) { // Update all multiples of p for (int i = p * p ; i < limit ; i += p) mark [i] = false ; } } // Print all prime numbers and store them in prime for (int p = 2 ; p < limit ; p ++) if (mark [p] == true) System . out . print (p + " ") ; } } Python def simple_sieve (limit) : # Create a boolean array "mark[0..limit-1]" and # initialize all entries of it as true. A value # in mark[p] will finally be false if 'p' is Not # a prime, else true. mark = [True for _ in range (limit)] # One by one traverse all numbers so that their # multiples can be marked as composite. for p in range (2 , int (limit ** 0.5) + 1) : # If p is not changed, then it is a prime if mark [p] == True : # Update all multiples of p for i in range (p * p , limit , p) : mark [i] = False # Print all prime numbers and store them in prime for p in range (2 , limit) : if mark [p] == True : print (p , end = " ") limit = 100 simple_sieve (limit) C# // This functions finds all primes smaller than 'limit' // using simple sieve of eratosthenes. static void simpleSieve (int limit) { // Create a boolean array "mark[0..limit-1]" and // initialize all entries of it as true. A value // in mark[p] will finally be false if 'p' is Not // a prime, else true. bool [] mark = new bool [limit] ; Array . Fill (mark , true) ; // One by one traverse all numbers so that their // multiples can be marked as composite. for (int p = 2 ; p * p < limit ; p ++) { // If p is not changed, then it is a prime if (mark [p] == true) { // Update all multiples of p for (int i = p * p ; i < limit ; i += p) mark [i] = false ; } } // Print all prime numbers and store them in prime for (

```

int p = 2 ; p < limit ; p ++ ) if ( mark [ p ] == true ) Console . Write ( p + " " ); } // This code is contributed
by pratham76. JavaScript function simpleSieve ( limit ) { // Create a boolean array "mark[0..limit-1]" and
// initialize all entries of it as true. A value // in mark[p] will finally be false if 'p' is Not // a prime, else true.
let mark = new Array ( limit ). fill ( true ); // One by one traverse all numbers so that their // multiples can
be marked as composite. for ( let p = 2 ; p * p < limit ; p ++ ) { // If p is not changed, then it is a prime if (
mark [ p ] === true ) { // Update all multiples of p for ( let i = p * p ; i < limit ; i += p ) mark [ i ] = false ; } }
// Print all prime numbers and store them in prime for ( let p = 2 ; p < limit ; p ++ ) if ( mark [ p ] === true
) console . log ( p + " " ); } let limit = 100 ; simpleSieve ( limit ); Problems with Simple Sieve: The Sieve
of Eratosthenes looks good, but consider the situation when n is large, the Simple Sieve faces the
following issues. An array of size  $O(n)$  may not fit in memory The simple Sieve is not cached friendly
even for slightly bigger n. The algorithm traverses the array without locality of reference Segmented
Sieve The idea of a segmented sieve is to divide the range  $[0..n-1]$  in different segments and compute
primes in all segments one by one. This algorithm first uses Simple Sieve to find primes smaller than or
equal to  $\sqrt{n}$ . Below are steps used in Segmented Sieve. Use Simple Sieve to find all primes up to the
square root of 'n' and store these primes in an array "prime[]". Store the found primes in an array
'prime[]'. We need all primes in the range  $[0..n-1]$ . We divide this range into different segments such
that the size of every segment is at-most  $\sqrt{n}$  Do following for every segment [low..high] Create an array
mark[high-low+1]. Here we need only  $O(x)$  space where x is a number of elements in a given range.
Iterate through all primes found in step 1. For every prime, mark its multiples in the given range
[low..high]. In Simple Sieve, we needed  $O(n)$  space which may not be feasible for large n. Here we
need  $O(\sqrt{n})$  space and we process smaller ranges at a time (locality of reference) Below is the
implementation of the above idea. C++ // C++ program to print all primes smaller than // n using
segmented sieve #include <bits/stdc++.h> using namespace std ; // This functions finds all primes
smaller than 'limit' // using simple sieve of eratosthenes. It also stores // found primes in vector prime[]
void simpleSieve ( int limit , vector < int > & prime ) { // Create a boolean array "mark[0..n-1]" and
initialize // all entries of it as true. A value in mark[p] will // finally be false if 'p' is Not a prime, else true.
vector < bool > mark ( limit + 1 , true ); for ( int p = 2 ; p * p < limit ; p ++ ) { // If p is not changed, then it
is a prime if ( mark [ p ] == true ) { // Update all multiples of p for ( int i = p * p ; i < limit ; i += p ) mark [ i ]
= false ; } } // Print all prime numbers and store them in prime for ( int p = 2 ; p < limit ; p ++ ) { if ( mark [
p ] == true ) { prime . push_back ( p ); cout << p << " "; } } } // Prints all prime numbers smaller than 'n'
void segmentedSieve ( int n ) { // Compute all primes smaller than or equal // to square root of n using
simple sieve int limit = floor ( sqrt ( n ) ) + 1 ; vector < int > prime ; prime . reserve ( limit ); simpleSieve (
limit , prime ); // Divide the range  $[0..n-1]$  in different segments // We have chosen segment size as
 $\sqrt{n}$ . int low = limit ; int high = 2 * limit ; // While all segments of range  $[0..n-1]$  are not processed, //
process one segment at a time while ( low < n ) { if ( high >= n ) high = n ; // To mark primes in current
range. A value in mark[i] // will finally be false if 'i-low' is Not a prime, // else true. bool mark [ limit + 1 ];
memset ( mark , true , sizeof ( mark ) ); // Use the found primes by simpleSieve() to find // primes in
current range for ( int i = 0 ; i < prime . size () ; i ++ ) { // Find the minimum number in [low..high] that is //
a multiple of prime[i] (divisible by prime[i]) // For example, if low is 31 and prime[i] is 3, // we start with
33. int loLim = floor ( low / prime [ i ] ) * prime [ i ] ; if ( loLim < low ) loLim += prime [ i ] ; /* Mark multiples
of prime[i] in [low..high]: We are marking j - low for j, i.e. each number in range [low, high] is mapped to
[0, high-low] so if range is [50, 100] marking 50 corresponds to marking 0, marking 51 corresponds to 1
and so on. In this way we need to allocate space only for range */ for ( int j = loLim ; j < high ; j += prime
[ i ] ) mark [ j - low ] = false ; } // Numbers which are not marked as false are prime for ( int i = low ; i <
high ; i ++ ) if ( mark [ i - low ] == true ) cout << i << " "; // Update low and high for next segment low =
low + limit ; high = high + limit ; } } // Driver program to test above function int main () { int n = 100 ; cout
<< "Primes smaller than " << n << ": \n " ; segmentedSieve ( n ); return 0 ; } Java // Java program to
print all primes smaller than // n using segmented sieve import java.util.Vector ; import static
java.lang.Math.sqrt ; import static java.lang.Math.floor ; class Test { // This method finds all primes
smaller than 'limit' // using simple sieve of eratosthenes. It also stores // found primes in vector prime[]
static void simpleSieve ( int limit , Vector < Integer > prime ) { // Create a boolean array "mark[0..n-1]"
and initialize // all entries of it as true. A value in mark[p] will // finally be false if 'p' is Not a prime, else
true. boolean mark [] = new boolean [ limit + 1 ] ; for ( int i = 0 ; i < mark . length ; i ++ ) mark [ i ] = true ;
for ( int p = 2 ; p * p < limit ; p ++ ) { // If p is not changed, then it is a prime if ( mark [ p ] == true ) { //
Update all multiples of p for ( int i = p * p ; i < limit ; i += p ) mark [ i ] = false ; } } // Print all prime
numbers and store them in prime for ( int p = 2 ; p < limit ; p ++ ) { if ( mark [ p ] == true ) { prime . add (
p ); System . out . print ( p + " " ); } } } // Prints all prime numbers smaller than 'n' static void

```

```

segmentedSieve ( int n ) { // Compute all primes smaller than or equal // to square root of n using simple
sieve int limit = ( int ) ( floor ( sqrt ( n )) + 1 ); Vector < Integer > prime = new Vector <> (); simpleSieve (
limit , prime ); // Divide the range [0..n-1] in different segments // We have chosen segment size as
sqrt(n). int low = limit ; int high = 2 * limit ; // While all segments of range [0..n-1] are not processed, //
process one segment at a time while ( low < n ) { if ( high >= n ) high = n ; // To mark primes in current
range. A value in mark[i] // will finally be false if 'i-low' is Not a prime, // else true. boolean mark [] = new
boolean [ limit + 1 ] ; for ( int i = 0 ; i < mark . length ; i ++ ) mark [ i ] = true ; // Use the found primes by
simpleSieve() to find // primes in current range for ( int i = 0 ; i < prime . size () ; i ++ ) { // Find the
minimum number in [low..high] that is // a multiple of prime.get(i) (divisible by prime.get(i)) // For
example, if low is 31 and prime.get(i) is 3, // we start with 33. int loLim = ( int ) ( floor ( low / prime . get (
i )) * prime . get ( i )) ; if ( loLim < low ) loLim += prime . get ( i ) ; /* Mark multiples of prime.get(i) in
[low..high]: We are marking j - low for j, i.e. each number in range [low, high] is mapped to [0, high-low]
so if range is [50, 100] marking 50 corresponds to marking 0, marking 51 corresponds to 1 and so on.
In this way we need to allocate space only for range */ for ( int j = loLim ; j < high ; j += prime . get ( i ))
mark [ j - low ] = false ; } // Numbers which are not marked as false are prime for ( int i = low ; i < high ; i
++ ) if ( mark [ i - low ] == true ) System . out . print ( i + " " ) ; // Update low and high for next segment
low = low + limit ; high = high + limit ; } } // Driver method public static void main ( String args [] ) { int n =
100 ; System . out . println ( "Primes smaller than " + n + ":" ) ; segmentedSieve ( n ) ; } } Python #
Python3 program to print all primes # smaller than n, using segmented sieve import math prime = [] #
This method finds all primes # smaller than 'limit' using # simple sieve of eratosthenes. # It also stores
found primes in list prime def simpleSieve ( limit ) : # Create a boolean list "mark[0..n-1]" and # initialize
all entries of it as True. # A value in mark[p] will finally be False # if 'p' is Not a prime, else True. mark =
[ True for i in range ( limit + 1 ) ] p = 2 while ( p * p <= limit ) : # If p is not changed, then it is a prime if (
mark [ p ] == True ) : # Update all multiples of p for i in range ( p * p , limit + 1 , p ) : mark [ i ] = False p +=
1 # Print all prime numbers # and store them in prime for p in range ( 2 , limit ) : if mark [ p ] : prime .
append ( p ) print ( p , end = " " ) # Prints all prime numbers smaller than 'n' def segmentedSieve ( n ) : #
Compute all primes smaller than or equal # to square root of n using simple sieve limit = int ( math .
floor ( math . sqrt ( n )) + 1 ) simpleSieve ( limit ) # Divide the range [0..n-1] in different segments # We
have chosen segment size as sqrt(n). low = limit high = limit * 2 # While all segments of range [0..n-1]
are not processed, # process one segment at a time while low < n : if high >= n : high = n # To mark
primes in current range. A value in mark[i] # will finally be False if 'i-low' is Not a prime, # else True.
mark = [ True for i in range ( limit + 1 ) ] # Use the found primes by simpleSieve() # to find primes in
current range for i in range ( len ( prime )) : # Find the minimum number in [low..high] # that is a multiple
of prime[i] # (divisible by prime[i]) # For example, if low is 31 and prime[i] is 3, # we start with 33. loLim
= int ( math . floor ( low / prime [ i ] ) * prime [ i ] ) if loLim < low : loLim += prime [ i ] # Mark multiples of
prime[i] in [low..high]: # We are marking j - low for j, i.e. each number # in range [low, high] is mapped
to [0, high-low] # so if range is [50, 100] marking 50 corresponds # to marking 0, marking 51
corresponds to 1 and # so on. In this way we need to allocate space # only for range for j in range (
loLim , high , prime [ i ] ) : mark [ j - low ] = False # Numbers which are not marked as False are prime for
i in range ( low , high ) : if mark [ i - low ] : print ( i , end = " " ) # Update low and high for next segment low
= low + limit high = high + limit # Driver Code n = 100 print ( "Primes smaller than" , n , ":" )
segmentedSieve ( 100 ) # This code is contributed by bhavyadeep C# // C# program to print // all
primes smaller than // n using segmented sieve using System ; using System.Collections ; class GFG {
// This method finds all primes // smaller than 'limit' using simple // sieve of eratosthenes. It also stores //
found primes in vector prime[] static void simpleSieve ( int limit , ArrayList prime ) { // Create a boolean
array "mark[0..n-1]" // and initialize all entries of it as // true. A value in mark[p] will finally be // false if 'p'
is Not a prime, else true. bool [] mark = new bool [ limit + 1 ] ; for ( int i = 0 ; i < mark . Length ; i ++ )
mark [ i ] = true ; for ( int p = 2 ; p * p < limit ; p ++ ) { // If p is not changed, then it is a prime if ( mark [ p ]
== true ) { // Update all multiples of p for ( int i = p * p ; i < limit ; i += p ) mark [ i ] = false ; } } // Print all
prime numbers and store them in prime for ( int p = 2 ; p < limit ; p ++ ) { if ( mark [ p ] == true ) { prime .
Add ( p ) ; Console . Write ( p + " " ) ; } } } // Prints all prime numbers smaller than 'n' static void
segmentedSieve ( int n ) { // Compute all primes smaller than or equal // to square root of n using simple
sieve int limit = ( int ) ( Math . Floor ( Math . Sqrt ( n )) + 1 ) ; ArrayList prime = new ArrayList () ;
simpleSieve ( limit , prime ) ; // Divide the range [0..n-1] in // different segments We have chosen //
segment size as sqrt(n). int low = limit ; int high = 2 * limit ; // While all segments of range // [0..n-1] are
not processed, // process one segment at a time while ( low < n ) { if ( high >= n ) high = n ; // To mark
primes in current range. // A value in mark[i] will finally // be false if 'i-low' is Not a prime, // else true.

```

```

bool [] mark = new bool [ limit + 1 ]; for ( int i = 0 ; i < mark . Length ; i ++ ) mark [ i ] = true ; // Use the
found primes by // simpleSieve() to find // primes in current range for ( int i = 0 ; i < prime . Count ; i ++ )
{ // Find the minimum number in // [low..high] that is a multiple // of prime.get(i) (divisible by //
prime.get(i)) For example, // if low is 31 and prime.get(i) // is 3, we start with 33. int loLim = (( int ) Math .
Floor (( double )( low / ( int ) prime [ i ] ) * ( int ) prime [ i ] ); if ( loLim < low ) loLim += ( int ) prime [ i ]; /*
Mark multiples of prime.get(i) in [low..high]: We are marking j - low for j, i.e. each number in range [low,
high] is mapped to [0, high-low] so if range is [50, 100] marking 50 corresponds to marking 0, marking
51 corresponds to 1 and so on. In this way we need to allocate space only for range */ for ( int j = loLim ;
j < high ; j += ( int ) prime [ i ] ) mark [ j - low ] = false ; } // Numbers which are not marked as false are
prime for ( int i = low ; i < high ; i ++ ) if ( mark [ i - low ] == true ) Console . Write ( i + " " ); // Update low
and high for next segment low = low + limit ; high = high + limit ; } } // Driver code static void Main () { int
n = 100 ; Console . WriteLine ( "Primes smaller than " + n + ":" ); segmentedSieve ( n ); } } // This code
is contributed by mits JavaScript // JavaScript program to print all primes smaller than // n using
segmented sieve // This functions finds all primes smaller than 'limit' // using simple sieve of
eratosthenes. It also stores // found primes in vector prime[] let res = "" ; function simpleSieve ( limit ,
prime ) { // Create a boolean array "mark[0..n-1]" and initialize // all entries of it as true. A value in
mark[p] will // finally be false if 'p' is Not a prime, else true. let mark = new Array ( limit + 1 ). fill ( true );
for ( let p = 2 ; p * p < limit ; p ++ ) { // If p is not changed, then it is a prime if ( mark [ p ] === true ) { //
Update all multiples of p for ( let i = p * p ; i < limit ; i += p ) { mark [ i ] = false ; } } // Print all prime
numbers and store them in prime for ( let p = 2 ; p < limit ; p ++ ) { if ( mark [ p ] === true ) { prime . push
( p ); res = res + p + " " ; } } } // Prints all prime numbers smaller than 'n' function segmentedSieve ( n ) {
// Compute all primes smaller than or equal // to square root of n using simple sieve let limit = Math .
floor ( Math . sqrt ( n ) ) + 1 ; let prime = new Array ( limit ); simpleSieve ( limit , prime ); // Divide the
range [0..n-1] in different segments // We have chosen segment size as sqrt(n). let low = limit ; let high
= 2 * limit ; // While all segments of range [0..n-1] are not processed, // process one segment at a time
while ( low < n ) { if ( high >= n ) { high = n ; } // To mark primes in current range. A value in mark[i] // will
finally be false if 'i-low' is Not a prime, // else true. let mark = new Array ( limit + 1 ). fill ( true ); // Use the
found primes by simpleSieve() to find // primes in current range for ( let i = 0 ; i < prime . length ; i ++ ) {
// Find the minimum number in [low..high] that is // a multiple of prime[i] (divisible by prime[i]) // For
example, if low is 31 and prime[i] is 3, // we start with 33. let loLim = Math . floor ( low / prime [ i ] ) *
prime [ i ]; if ( loLim < low ) loLim += prime [ i ]; } /* Mark multiples of prime[i] in [low..high]: We are
marking j - low for j, i.e. each number in range [low, high] is mapped to [0, high-low] so if range is [50,
100] marking 50 corresponds to marking 0, marking 51 corresponds to 1 and so on. In this way we
need to allocate space only for range */ for ( let j = loLim ; j < high ; j += prime [ i ] ) { mark [ j - low ] =
false ; } } // Numbers which are not marked as false are prime for ( let i = low ; i < high ; i ++ ) { if ( mark [
i - low ] == true ) { res = res + i + " " ; } } // Update low and high for next segment low = low + limit ; high =
high + limit ; } console . log ( res ); } // Driver program to test above function let n = 100 ; console . log (
"Primes smaller than" , n ); segmentedSieve ( n ); // The code is contributed by Gautam goel
(gautamgoel962) Output Primes smaller than 100: 2 3 5 7 11 13 17 19 23 29 31 37 41 43 47 53 59 61
67 71 73 79 83 89 97 Time Complexity : O(n * ln(sqrt(n))) Auxiliary Space: O(sqrt(n)) Note that time
complexity (or a number of operations) by Segmented Sieve is the same as Simple Sieve . It has
advantages for large 'n' as it has better locality of reference thus allowing better caching by the CPU
and also requires less memory space. Comment Article Tags: Article Tags: Mathematical DSA sieve
Prime Number

```