

# Storing messages with multipartite neural cliques

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## Abstract

We extend recently introduced associative memories based on clustered cliques to multipartite cliques. We propose a variant of the classic retrieving rule. We study its performance relatively to the former one for retrieving partially erased and corrupted messages. We provide both analytical and simulation results showing improvements in networks capacities.

**Index terms**— associative memory, error correcting code, cliques, multipartite cliques, neural networks

mapping doesn't matter as long as it is fixed).

Storing a message  $x$  in the network amounts to consider the network in state  $y$ , that is for all  $1 \leq j \leq l$  and  $1 \leq i \leq c$ , the neuron  $j$  in the cluster  $i$  is activated if and only if  $y_{i,j} = 1$ ; then adding all connections between activated neurons excepting for neurons sharing the same cluster. Thus we obtain a multipartite clique between activated neurons representing a message. The parameter  $a$  describing the number of activated neurons per cluster will be called number of activities.

## I Introduction

## II Networks of neural cliques

### 1 Learning messages

Let  $\mathcal{M}$  be a set of messages of length  $c$  over the alphabet  $\mathcal{A} = \{1, 2, \dots, \binom{l}{a}\}$  and  $m = |\mathcal{M}|$ . Messages in  $\mathcal{M}$  are stored in an undirected unweighted graph, or neural network, of  $c \cdot l$  vertices or neurons.

Each message  $x = (x_i)_{1 \leq i \leq c}$  is mapped to  $y = (y_i)_{1 \leq i \leq c}$  where  $y_i$  is the binary representation of the  $x_i$ -th  $a$ -combination of  $l$  elements (the

### 2 Retrieving messages

Retrieving messages occurs by message passing.

#### a The "winner takes all" rule

take only maximum score

#### b The " $a$ winners take all" rule

$a$ -rank

less biologically plausible rule

same algorithmic complexity

$w = a$  is optimal

### III Retrieval performance

#### 1 Retrieving partially erased messages

##### a Analytical result

number of neurons per cluster  $l$

number of clusters  $c$

number of erased clusters  $c_e$

number of messages :  $m$

number of activities per cluster :  $a$

density :  $d$

edge :  $(i, j)$

As the probability for  $i$  and  $j$  to be active in their respective clusters for one message is  $\frac{a}{l}$  providing messages are identically distributed, it follows that the probability for the edge  $(i, j)$  of not being added in the network for one message is  $1 - \left(\frac{a}{l}\right)^2$ . Since messages are independent, the density  $d$  can be expressed like this :

$$d = 1 - \left(1 - \left(\frac{a}{l}\right)^2\right)^m$$

Thanks to the memory effect  $\gamma$ , activated neurons in a non-erased cluster stay activated and are the only ones doing so in their cluster. For neurons in erased clusters, the maximum attainable score is  $a(c - c_e)$ . Neurons corresponding to the original message achieve this score. Besides the probability for a neuron  $v$  to achieve the maximum score in such a cluster, that is to mean  $a(c - c_e)$  is  $d^{a(c - c_e)}$ .

Since neurons representing messages are independent (as messages are themselves independent), the probability for none of the vertices of one erased cluster, excepting the correct ones, to achieve the maximum is  $(1 - d^{a(c - c_e)})^{l - a}$ .

Scores in clusters being also independent, the probability for none of the vertices in any erased

cluster, excepting the correct ones, to achieve the maximum in their respective clusters is  $(1 - d^{a(c - c_e)})^{c_e(l - a)}$ .

Whence error rate in retrieving messages is :

$$P_{err} = 1 - \left(1 - d^{a(c - c_e)}\right)^{c_e(l - a)}$$

$$\eta_m = (P_{retrieve} m)^{\frac{2(c \log_2 \left(\frac{l}{a}\right) - \log_2(m) + 1)}{c(c-1)l^2}} \times (1 - \left(\frac{l}{a}\right)^{c - c_e})^{-(m-1)}$$

$$\eta_{max} = \max_x \eta_x$$

##### b Simulations

gap of theoretical density with density in simulations : under 1 percent, see Figure 1

New rule : Improves the retrieval rate for multiple iterations and/or corrupted messages.

The simulations agree with the analytical results for density and error rate.

For one iteration (for erasures, not for errors), winner takes all is the same as  $a$ -winners take all.

For multiple iterations  $a$ -winners take all is a significant improvement.

See Figure 2

" $a$ -winners take all" better than "winner takes all" with errors instead of erasures, even for one iteration.

For errors : only better than only 1 activity per cluster if multiple iterations of the  $a$ -winners take all rule. For one iteration, it is less efficient.

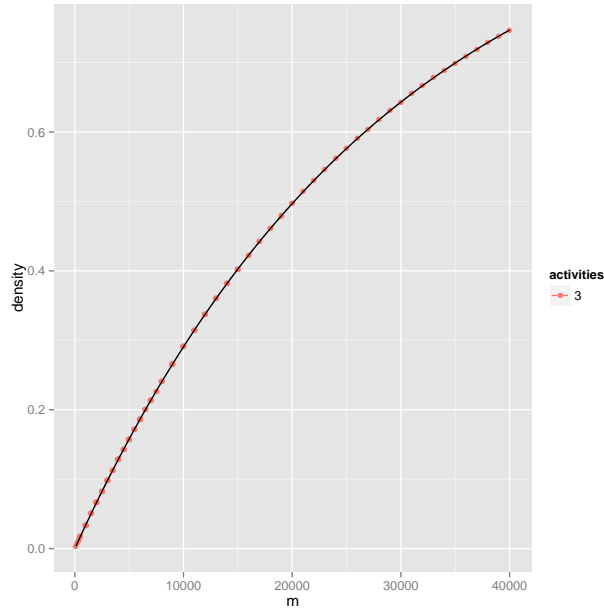


Figure 1: Theoretical and empirical densities

p

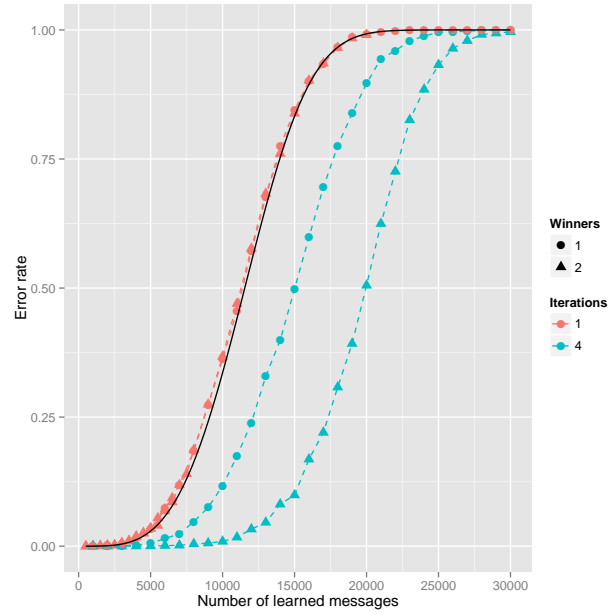


Figure 2: 2 activities per cluster, 4 clusters, 512 neurons per cluster, two erasures, each point is the mean of 5 networks with 1000 sampled messages, analytical result in continuous black

## 2 Retrieving corrupted messages

### a Analytical result

Formula for errors instead of erasures :

gc : good (à remplaceur par right/wrong ?)  
neuron in correct cluster

ge : good neuron in erroneous cluster

abe : bad activated neuron in erroneous cluster

be : others bad (à remplaceur par wrong ou incorrect etc.) neurons in erroneous cluster

bc : bad neuron in correct cluster

Correct neurons in uncorrupted clusters  
achieve at least a score of  $a(c - c_e - 1) + \gamma$ .

$$P(n_{gc} = (c - c_e - 1) + \gamma + x) = \binom{c_e}{x} d^x (1 - d)^{c_e - x}$$

Correct neurons in corrupted clusters achieve  
at least a score of  $a(c - c_e)$ .

$$P(n_{ge} = (c - c_e) + x) = \binom{c_e - 1}{x} d^x (1 - d)^{c_e - 1 - x}$$

$$P(n_{abe} = \gamma + x) = \binom{c - 1}{x} d^x (1 - d)^{c - 1 - x}$$

$$P(n_{be} = x) = P(n_{bc} = x) = \binom{c - 1}{x} d^x (1 - d)^{c - 1 - x}$$

$$P_{retrieve} = \left[ \sum_{n=1}^{c-1+\gamma} P(n_{gc} = n) \left[ \sum_{x=0}^{n-1} P(n_{bc} = x) \right]^{l-1} \right]^{c-c_e} \quad (1)$$

$$\times \left[ \sum_{n=1}^{c-1+\gamma} P(n_{ge} = n) \left[ \sum_{x=0}^{n-1} P(n_{bc} = x) \right]^{l-2} \left[ \sum_{x=0}^{n-1} P(n_{abc} = x) \right]^{c_e} \right] \quad (2)$$

generalisation to multiple activities per cluster

:

$$P(n_{gc} = a(c - c_e - 1) + \gamma + x) = \binom{ac_e}{x} d^x (1 - d)^{ace - x}$$

$$P(n_{ge} = a(c - c_e) + x) = \binom{a(c_e - 1)}{x} d^x (1 - d)^{a(c_e - 1) - x}$$

$$P(n_{abe} = \gamma + x) = \binom{a(c - 1)}{x} d^x (1 - d)^{a(c - 1) - x}$$

$$P(n_{be} = x) = P(n_{bc} = x) = \binom{a(c - 1)}{x} d^x (1 - d)^{a(c - 1) - x}$$

$$P_{retrieve} = \left[ \sum_{n=1}^{a(c-1)+\gamma} \left[ \sum_{k=1}^a \binom{a}{k} P(n_{gc} = n)^k P(n_{gc} > n)^{a-k} \right] P(n_{bc} < n)^{l-a} \right]^{c-c_e} \\ \times \left[ \sum_{n=1}^{a(c-1)+\gamma} \left[ \sum_{k=1}^a \binom{a}{k} P(n_{ge} = n)^k P(n_{ge} > n)^{a-k} \right] P(n_{bc} < n)^{l-2a} P(n_{abc} < n)^a \right]^{c_e}$$

To be really precise : (ne pas oublier de tenir compte changement des lois)

$$P_{retrieve} = \left[ \sum_{n=1}^{a(c-1)+\gamma} \left[ \sum_{k=1}^a \binom{a}{k} P(n_{gc} = n)^k P(n_{gc} > n)^{a-k} \right] P(n_{bc} < n)^{l-a} \right]^{c-c_e} \\ \times \left[ \sum_{y=0}^a \left( \frac{\binom{l-y}{a} - \sum_{z=y+1}^a \binom{l-z}{a}}{\binom{l}{a}} \right) \sum_{n=1}^{a(c-1)+\gamma} \left[ \sum_{k=1}^a \binom{a}{k} P(n_{ge} = n)^k P(n_{ge} > n)^{a-k} \right] P(n_{bc} < n)^{l-a-y} P(n_{abc} < n)^{a-y} \right]$$

$$P_{err} = 1 - P_{retrieve}$$

See 3

## b Simulations

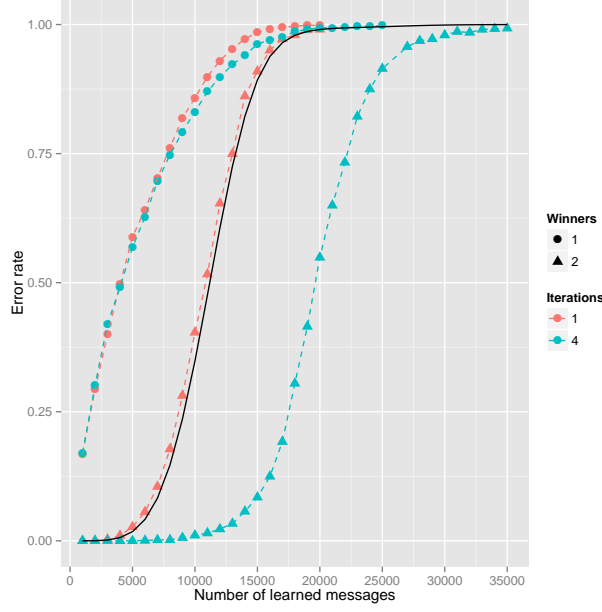


Figure 3: 2 activities per cluster, 4 clusters, 512 neurons per cluster, one erroneous cluster, each point is the mean of 5 networks with 1000 sampled messages, analytical result in continuous black

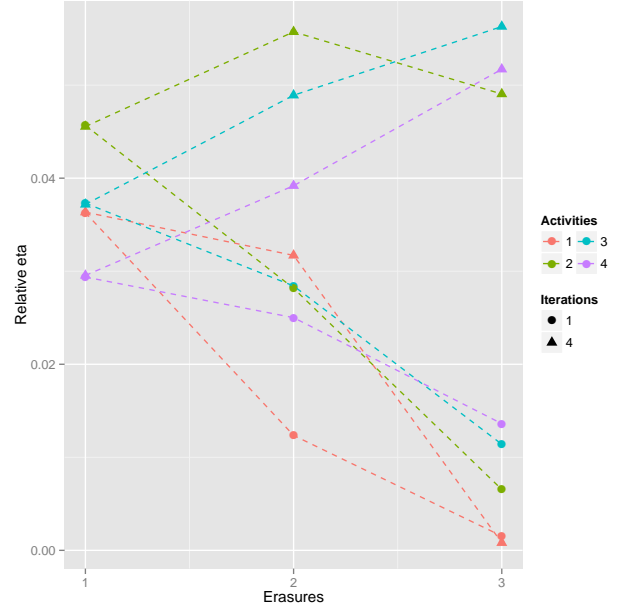


Figure 4: Comparison between different number of activities per cluster : 4 clusters, 512 neurons per cluster, 1000 sampled messages

## IV Conclusion

## V Resilience

Spéculations en Français :

Peut-être un intérêt pour cluter based associative memories build from unreliable storage : si une arête porte moins d'info, on peut peut-être en supprimer plus (mais rajout ?!)

Marche mieux pour des bruits importants

$$P(n_{v_c} = n_0 + \gamma) = \binom{a(c_k-1)}{n_0} (1 -$$

$$\psi)^{n_0} \psi^{a(c_k-1)-n_0}$$

$$P(n_{v_c} = n_0) = \binom{ac_k}{n_0} (1 - \psi)^{n_0} \psi^{ac_k - n_0}$$

$$P_+ = \psi(1 - d) + (1 - \psi)d$$

$$P(n_v = x) = \binom{ac_k}{x} P_+^x (1 - P_+)^{ac_k - x}$$

Sans tenir compte du choix au hasard si plus  
sont activés.

a priori for noise on edges :

$$P(\text{no other vertex activated in this cluster}) =$$

$$\sum_{n=1}^{a(c-c_e)} \left[ \sum_{k=1}^a \binom{a}{k} P(n_{v_c} = n)^k P(n_{v_c} > n)^{a-k} \right] \left[ \sum_{x=0}^{n-1} P(n_v = x) \right]^{l-a}$$

$$P(\text{no other vertex activated in any cluster}) =$$

$$\left( \sum_{n=1}^{a(c-c_e)} \left[ \sum_{k=1}^a \binom{a}{k} P(n_{v_c} = n)^k P(n_{v_c} > n)^{a-k} \right] \left[ \sum_{x=0}^{n-1} P(n_v = x) \right]^{l-a} \right)^{c_e}$$

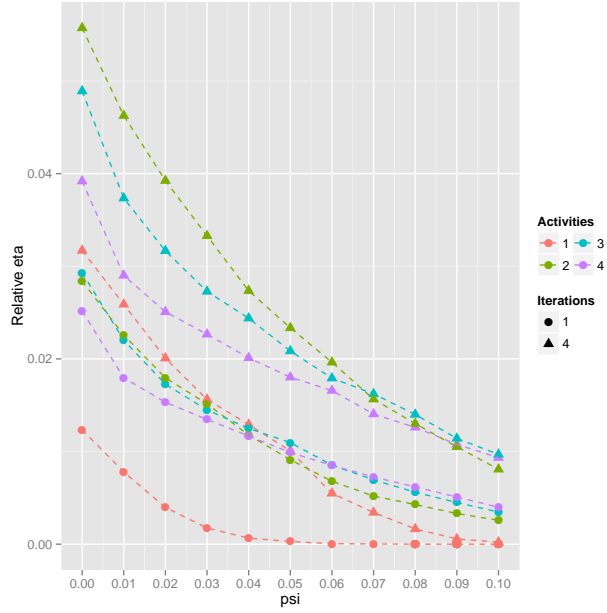


Figure 5: Comparison between different number of activities per cluster : 4 clusters, 512 neurons per cluster, 1000 sampled messages

## References

- [LPGRG14] François Leduc-Primeau, Vincent Gripon, Michael Rabbat, and Warren Gross. Cluster-based associative memories built from unreliable storage. In *ICASSP*, May 2014. To appear.