Равенство в Теории Типов

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Propositions as Types

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\begin{array}{ccc} \mathsf{Proposition} & \leftrightarrow & \mathsf{Type} \\ \mathsf{Proof} & \leftrightarrow & \mathsf{Program} \\ \mathsf{Proof simplification} & \leftrightarrow & \mathsf{Computation} \end{array}
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- ► Syntactic =. $\frac{A}{A \wedge B}$
- = on propositions.
- = as proposition.

Intentional Type Theory

Martin-Löf 1975

Definitional equality $=_{def}$

Definition,
$$\alpha\beta$$
-equivalence, SUB $\frac{a =_{\text{def}} b}{x[a] =_{\text{def}} x[b]}$, $\frac{a : A \qquad A =_{\text{def}} B}{a : B}$.

Propositional equality (a.k.a. Identity) I(A, a, b)

$$\frac{c[x]:C[x,x,r(x)]}{f(x,x,r(x)) =_{\mathsf{def}} c[x]}$$

$$J: ((x:A) \to C(x,x,r(x)))$$

$$\to (xy:A)(p:I(A,x,y))$$

$$\to C(x,y,p)$$

Definitional (syntactic, intentional) equality Definition, α -equivalence, SUB. Judgemental equality A=B, a=b:A Defined for every type, $\beta\eta$, [x:A] $\xi \frac{f(x)=g(x):B(x)}{\lambda x.\ f(x)=\lambda x.\ g(x):\Pi(x:A)B(x)},$ $\frac{a:A}{a:B}, \frac{A=B}{a:B}, \frac{a=b:A}{a=b:B}.$

$$\mathsf{ER}\,\frac{c:I(A,a,b)}{a=b:A}\;,\;\;\mathsf{UIP}\,\frac{c:I(A,a,b)}{c=r:I(A,a,b)}\;.$$

Observational Type Theory

Definitional equality $\alpha\beta\eta$.

Propositional equality

$$(x:A) = (y:B)$$

$$refl: (x:A) \rightarrow (x:A) = (x:A)$$

$$Coercion \frac{p:S_0 = S_1 \quad s_0:S_0}{s_0[p\rangle:S_1}$$

$$Coherence \frac{p:S_0 = S_1 \quad s_0:S_0}{s_0[p\rangle:(s_0:S_0) = (s_0[p\rangle:S_1)}$$

Result: propositional equality is extensional, definitional equality is decidable.

Computational Type Theory

Based on ETT. But ETT is formal type theory.

Definitional equality

Define the meaning of x = y : A for the types you want. x : A is then a shortcut for x = x : A.

Propositional equality

The same as judgemental.

Quotient types

 $(x,y):A/\!/E$. Change an equivalence relation to E.

Homotopy Type Theory

Definitional equality

ITT $+ \eta$.

Propositional equality

Proof relevance to 11. Equality proof is a path between values. No UIP but funext.

Unified Type Theory

ITT + impredicative Prop. Equality is Leibniz: $\forall P: A \rightarrow Prop. \ P(a) \implies P(b)$.

ZOMBIE

Forget automatic β . Get equality relations from the context and use them.