

Data Structures and Algorithms

Lecture 19: Tree Traversals

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Outlines

- Tree traversals: *for rooted tree*
 - Preorder traversal
 - Postorder traversal

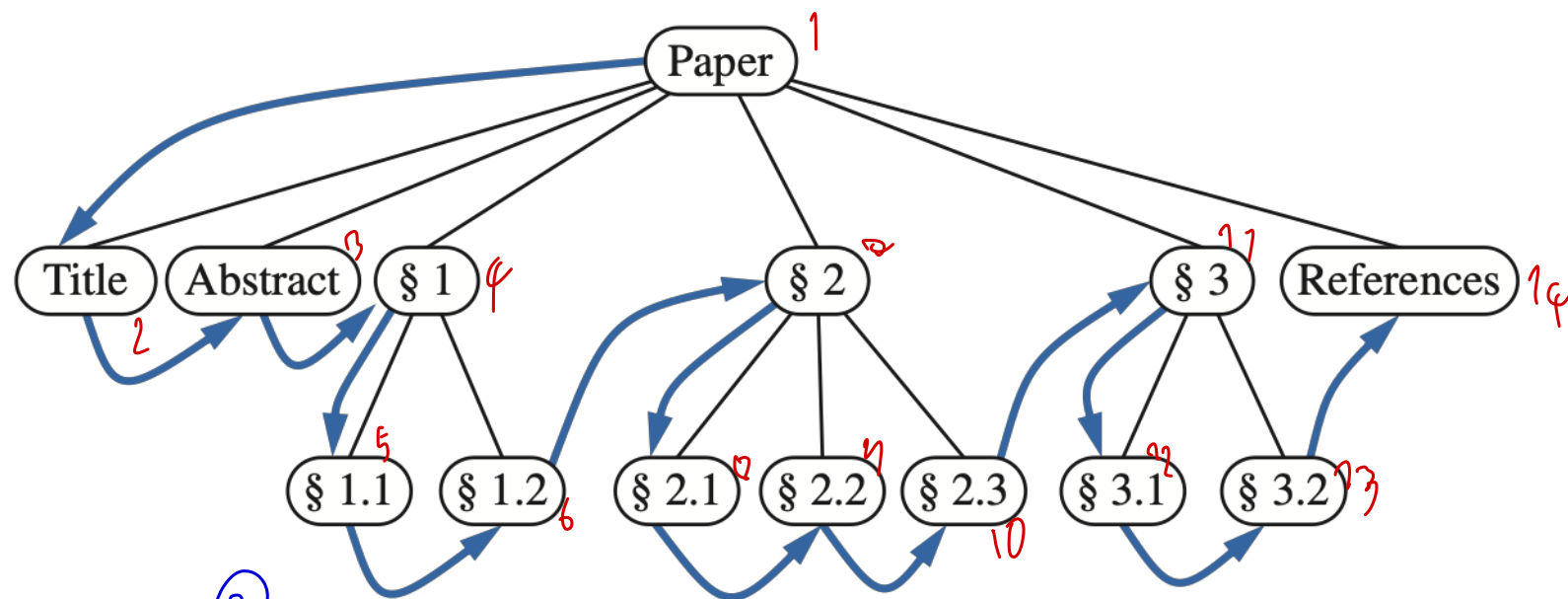
Tree Traversals

- ^{visiting all nodes in tree} **Traversal** := A systematic way of accessing (visiting, traversing) all the elements of the data structure
- **Graph traversal** := A systematic way of accessing (visiting, traversing) all the nodes and edges of a graph
 - BFS/DFS as we saw in the previous lectures
- **Tree traversal** := A systematic way of accessing all the nodes of a tree.
 - Of course, we can apply BFS/DFS to traverse a tree
 - Preorder & postorder traversals (DFS variants)

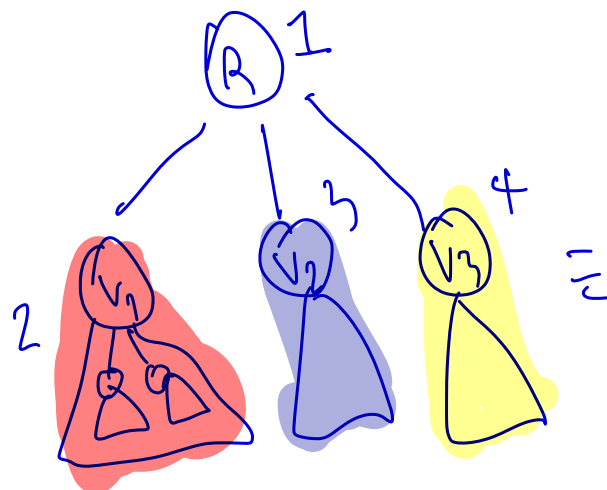
Preorder Traversal

visit root now

- **Preorder traversal:** In preorder traversal of a rooted tree T , we visit the root of T first and then the subtrees rooted at its children are traversed recursively
- If the tree is ordered, then the subtrees are traversed according to the order of the children



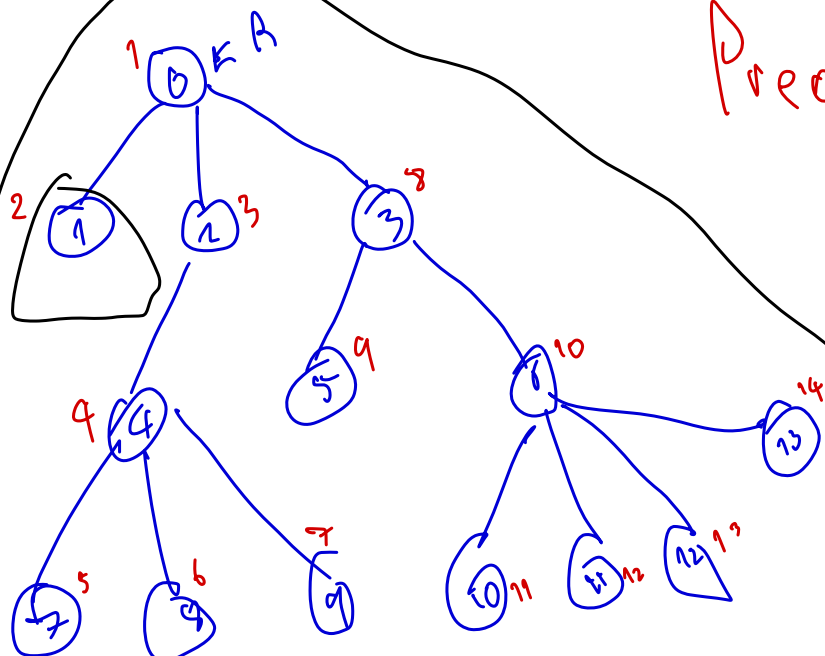
Preorder



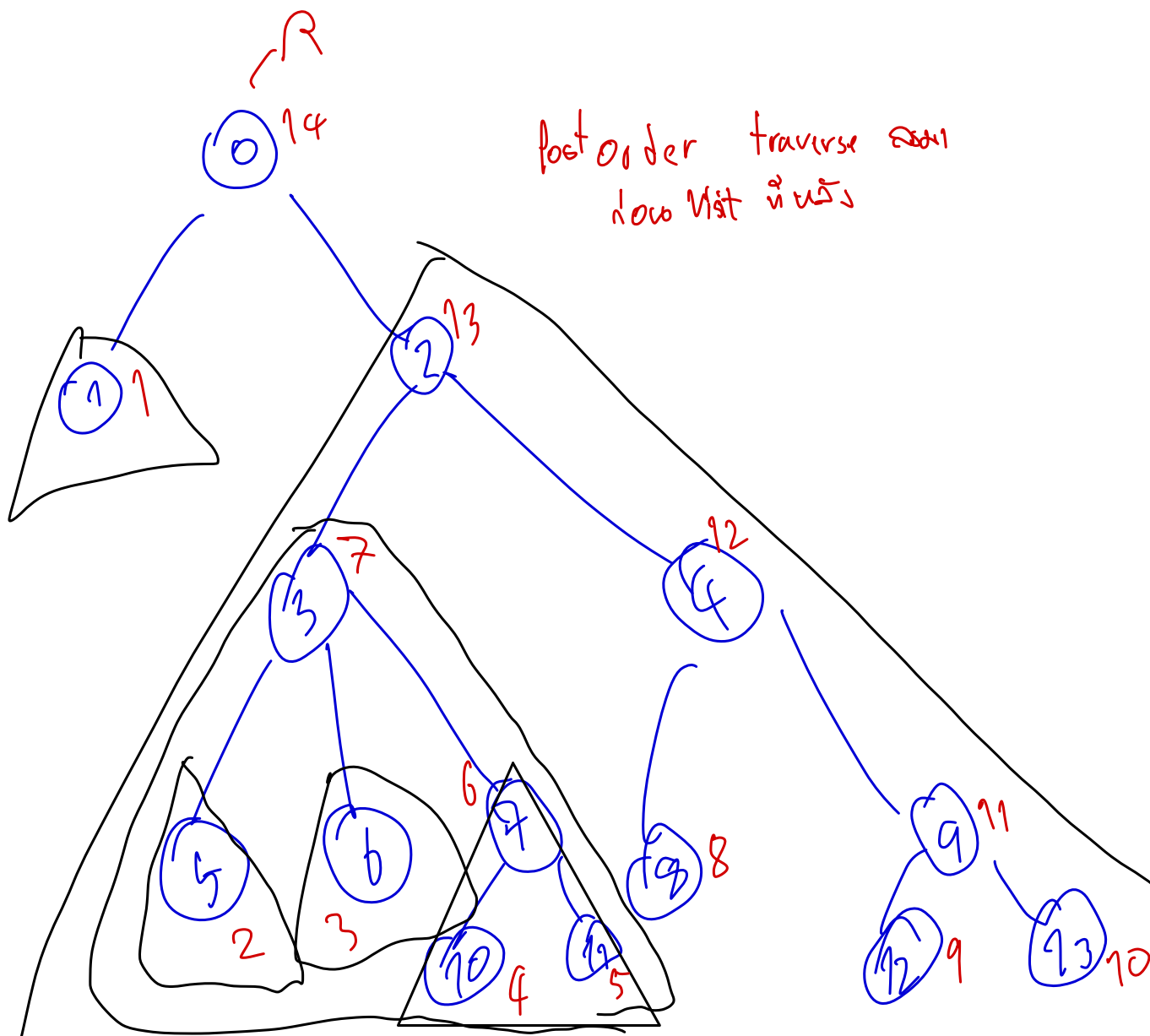
sequence of actions:

1. visit root
2. traverse subtree rooted at v_1 , a child of r
3. v_2

Preorder



post order traverse
now visit root



Preorder Traversal: Pseudocode (Root, Left, Right)

- In preorder traversal of a rooted tree T , we visit the root of T first and then the subtrees rooted at its children are traversed recursively

preorder(r, T):

visit node r

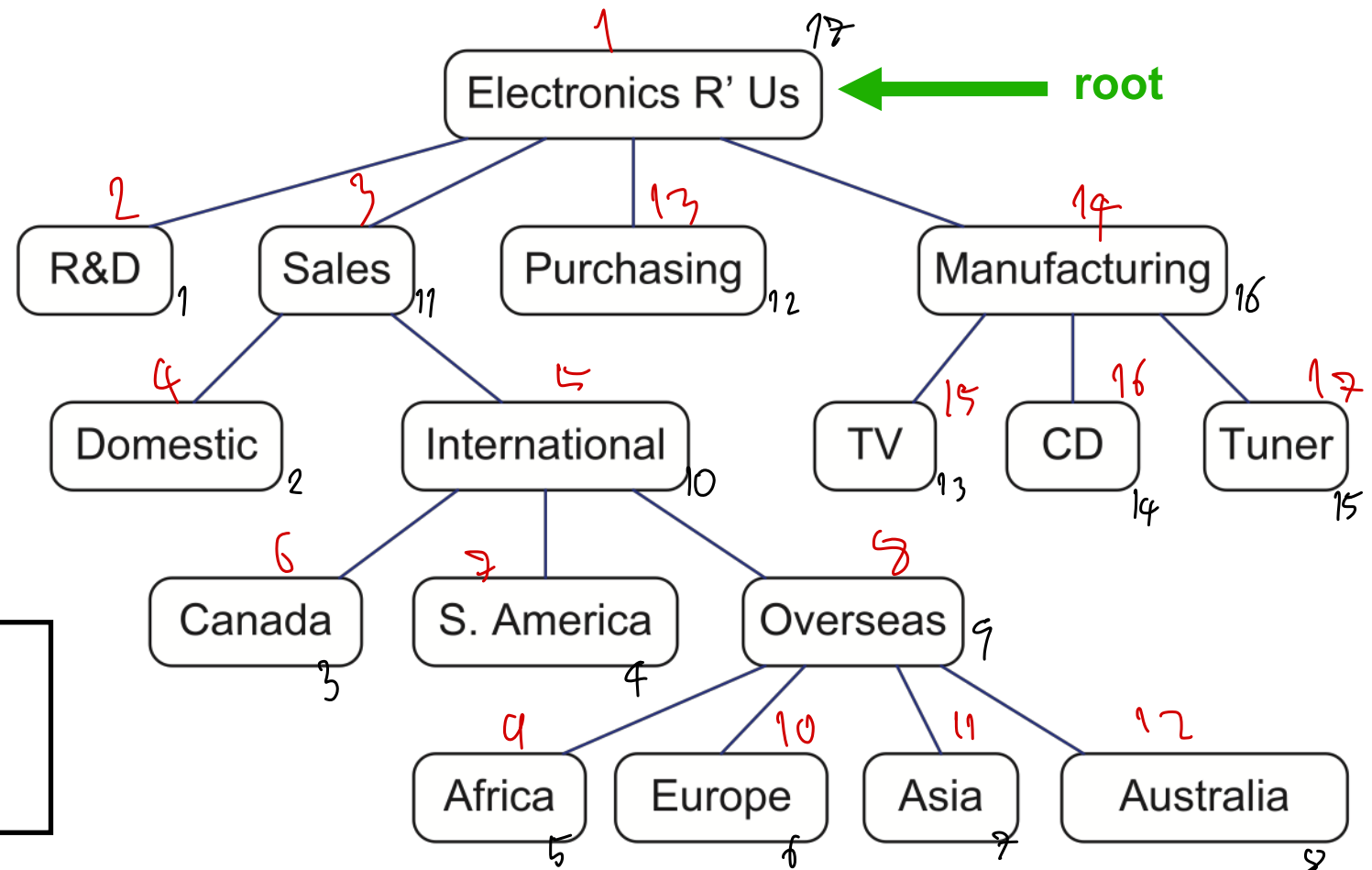
for each child q of r :

preorder(q, T)

- In other words, we visit the *root* first, then recursively visit the *left* child, and its *right* siblings

Let's Do Preorder Traversal

● Pre
● Post

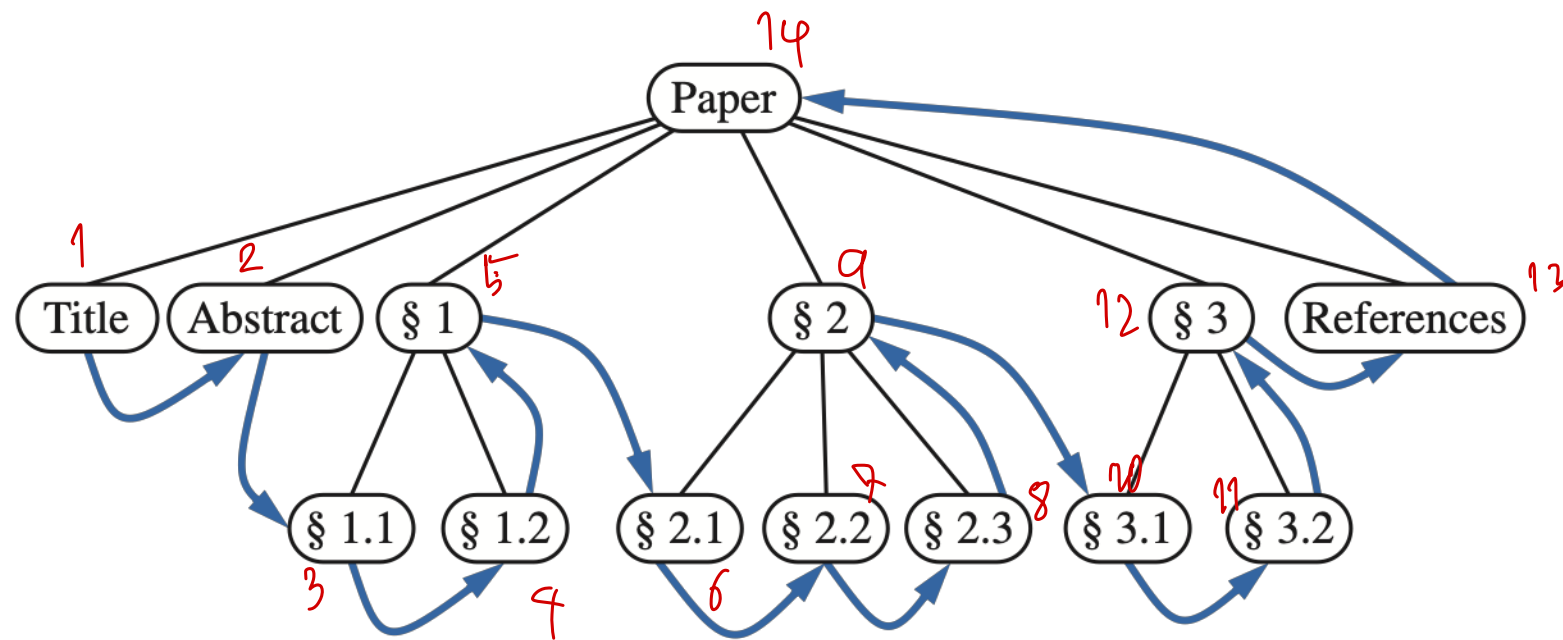


```
preorder(r,T):  
  visit node r  
  for each child q of r:  
    preorder(q,T)
```

- Let's try to perform preorder traversal of the above rooted tree

Postorder Traversal

- **Postorder traversal** : As opposed to preorder traversal, in postorder traversal of a rooted tree T , we recursively traverse the subtrees rooted at the children of the root first and then visits the root
- If the tree is ordered, then the subtrees are traversed according to the order of the children



Postorder Traversal: Pseudocode (Left, Right, Root)

- In post traversal of a rooted tree T , we recursively traverse the subtrees rooted at the children of the root first and then visits the root

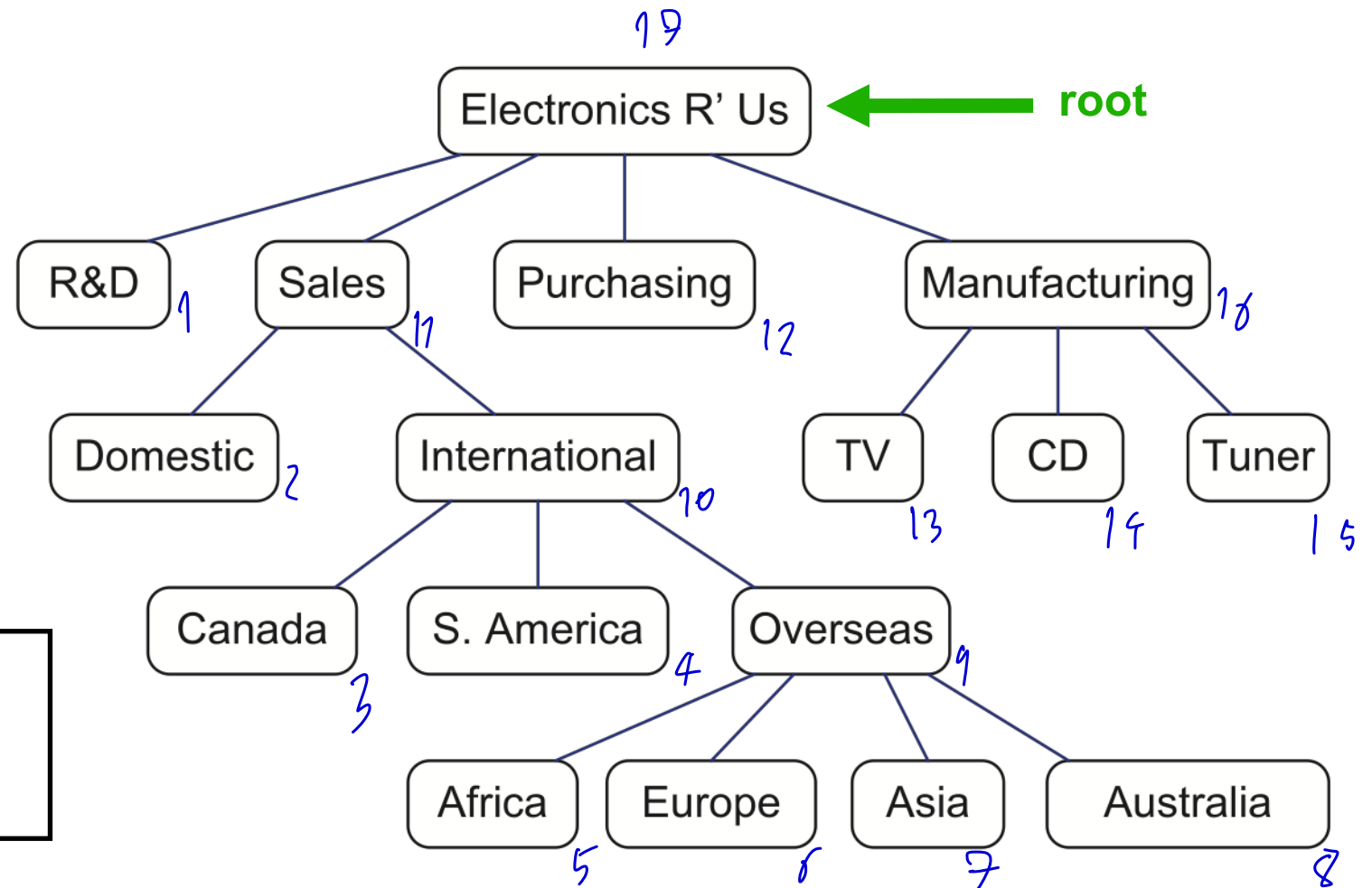
```
postorder(r,T):  
    for each child q of r:  
        postorder(q,T)  
    visit node r
```

- In other words, we visit recursively visit the *left* child, and its *right* siblings, then visit the *root*

Let's Do Postorder Traversal

Node with leaf
does 1 parent

```
postorder(r,T):
  for each child q of r:
    postorder(q,T)
  visit node r
```

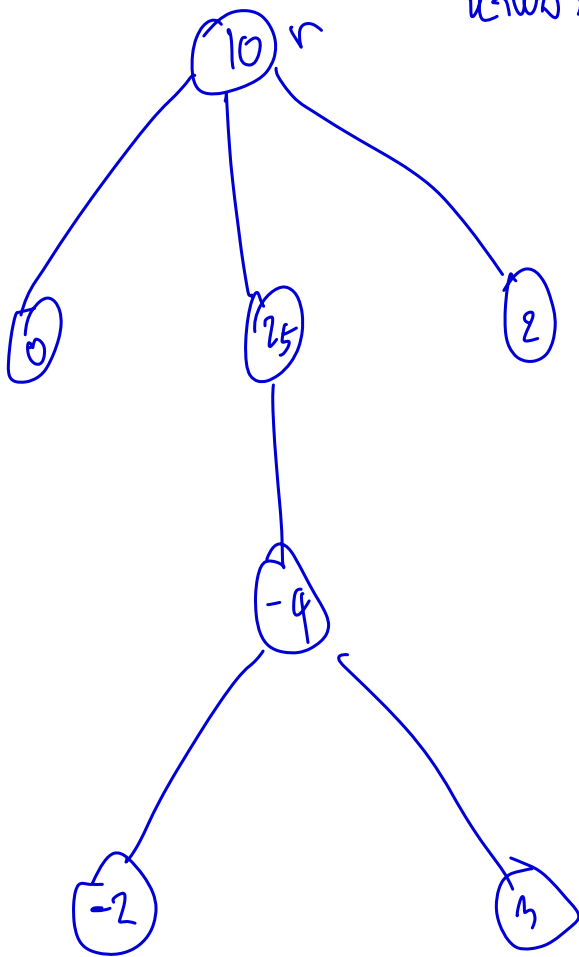


- Let's try to perform a postorder traversal of the rooted tree above

Complexity of Traversals on Ordered Trees

Tree traversal algorithms	Complexity
preorder	$O(n)$
postorder	$O(n)$
	Remarks: a tree of n nodes can only have $n-1$ edges. A traversal need to visit every node of the tree

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$i = 0$
 Preorder(r, T)

$s++ = r.key$

for each child q of r :

Preorder(q, T)

struct node {

int key

struct node * firstChild

struct node * nextSibling

}

Postorder (Node *root) {

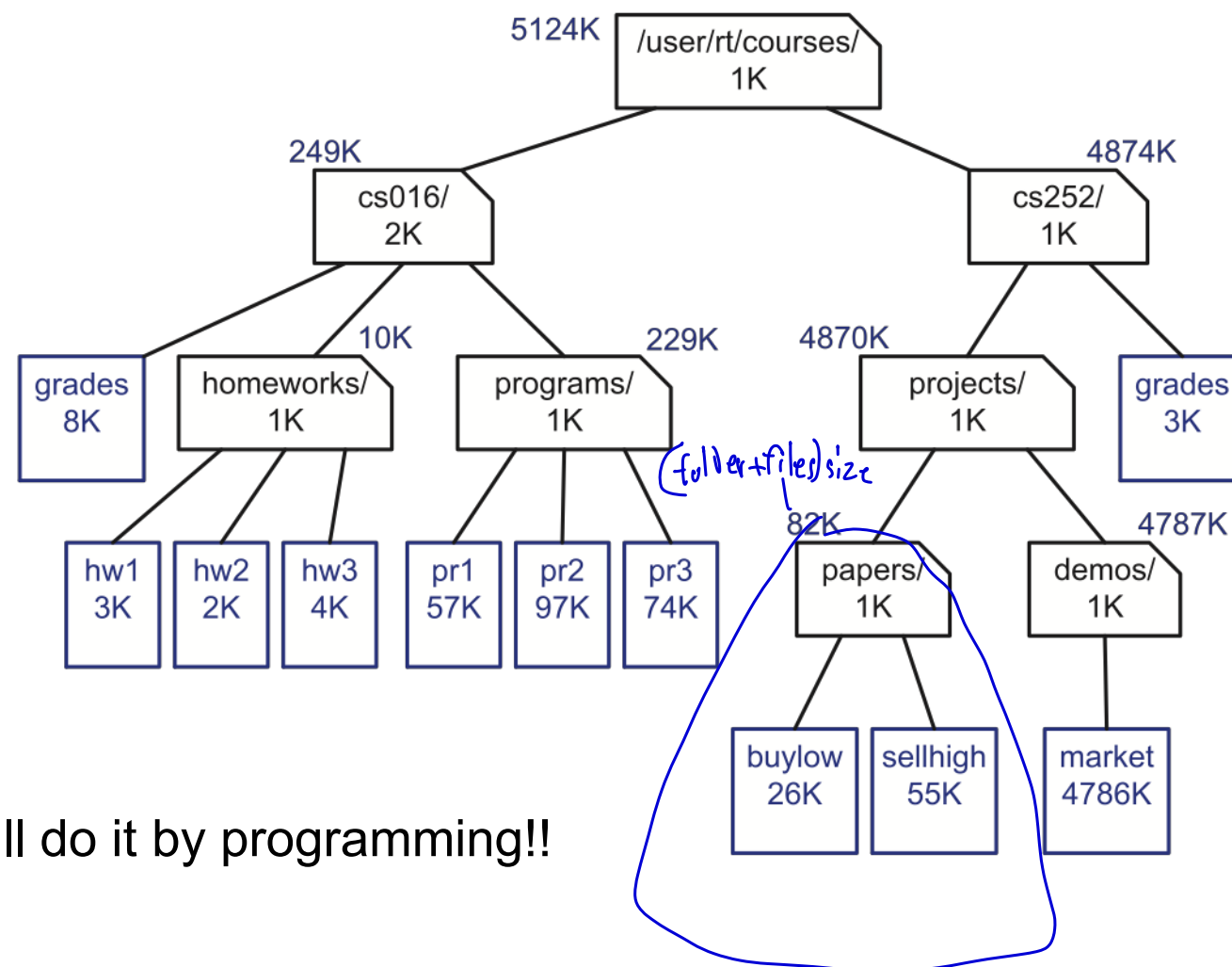
// visit root

for (Node* child = root->firstchild; child != Null; child = child->nextsibling)

Postorder (child)

Programming Exercise

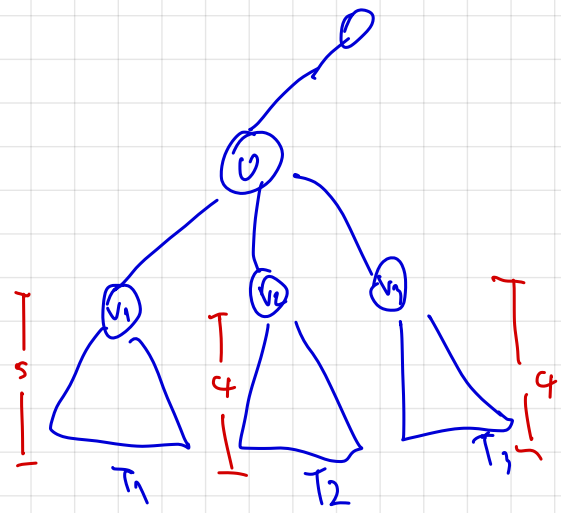
- Now that you know everything about preorder/postorder traversal. Your next assignments will be to construct the tree below which represents the directory structure of the root folder /user/rt/courses/ as well as to compute the file sizes of each directory that resides in the root folder.



- Of course, you will do it by programming!!

- જાણ $\text{Height} = 0$

- u માટે $\text{height}(u) = \max_{v \in \text{child}(u)} \{ \text{height}(v) \} + 1$



બિન
Postorder*

