

Introduction to Program Analysis

6. Abstract Interpretation (2): Abstract Semantics

Kihong Heo



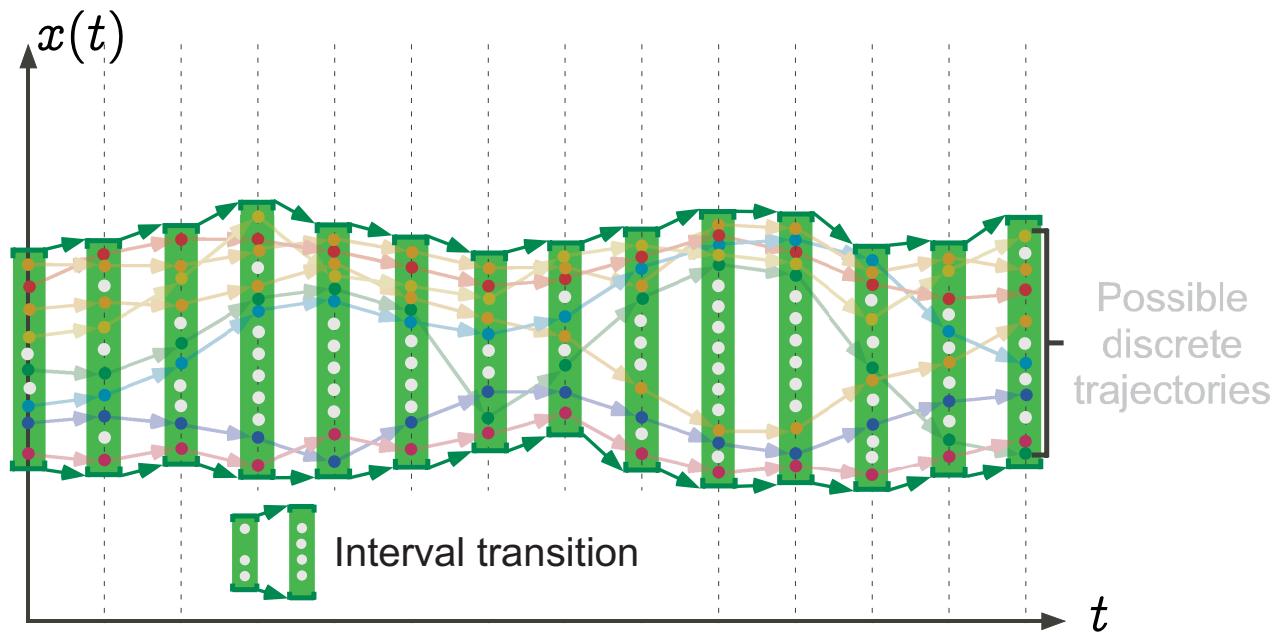
Design of Static Analysis

- Goal: **conservative** and **terminating** static analysis
- Design principles:
 - Define **concrete semantics**
 - Define **abstract semantics** (sound w.r.t the concrete semantics)
- Computation & implementation:
 - Abstract semantics of a program: **the least fixed point** of the semantic function
 - Static analyzer: **compute the least fixed point within finite time**

Step 2: Design Abstract Semantics

- Formalization of **abstract** program executions
 - Soundly subsume concrete executions
- How to subsume: different choices depending on the purposes
 - Some are more expressive than others
- Example: abstraction of {1, 3, 5, 7}
 - Integer, Positive, Odd, [1, 7], etc

Transitions of Abstract States



*from Patrick Cousot's slides

Abstract Semantics

- Define an abstract domain \mathbb{D}^\sharp (CPO)
- Define an abstract semantic function $F^\sharp : \mathbb{D}^\sharp \rightarrow \mathbb{D}^\sharp$ (monotone or extensive)

(Monotone) $\forall x^\sharp, y^\sharp \in \mathbb{D}^\sharp. x^\sharp \sqsubseteq y^\sharp \implies F^\sharp(x^\sharp) \sqsubseteq F^\sharp(y^\sharp)$

(Extensive) $\forall x^\sharp \in \mathbb{D}. x^\sharp \sqsubseteq F^\sharp(x^\sharp)$

- Static analysis is to compute an upper bound of the chain:

$$\bigsqcup_{i \geq 0} F^{\sharp i}(\perp^\sharp)$$

Q. How to ensure that the abstract semantics soundly subsume the concrete semantics?

Requirement 1: Galois Connection

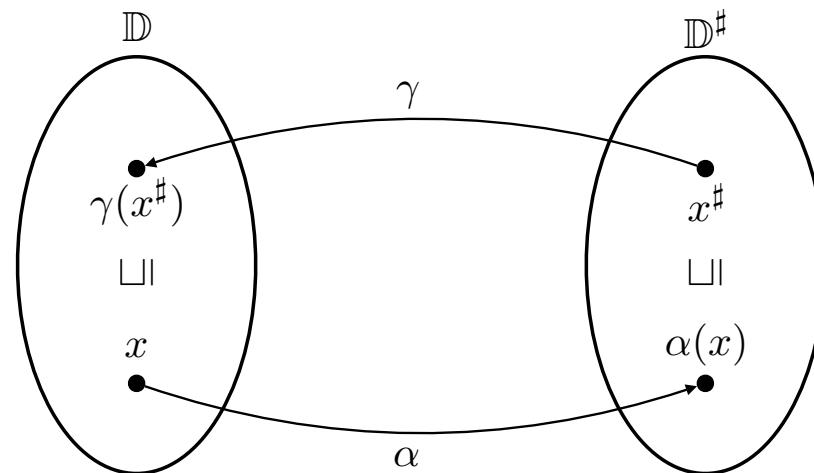
$$\mathbb{D} \begin{array}{c} \xleftarrow{\gamma} \\[-1ex] \xrightarrow{\alpha} \end{array} \mathbb{D}^\sharp$$

- \mathbb{D} and \mathbb{D}^\sharp must be related with a Galois connection where
 - Abstraction function: $\alpha \in \mathbb{D} \rightarrow \mathbb{D}^\sharp$
 - Concretization function: $\gamma \in \mathbb{D}^\sharp \rightarrow \mathbb{D}$

$$\forall x \in \mathbb{D}, x^\sharp \in \mathbb{D}^\sharp. \alpha(x) \sqsubseteq x^\sharp \iff x \sqsubseteq \gamma(x^\sharp)$$

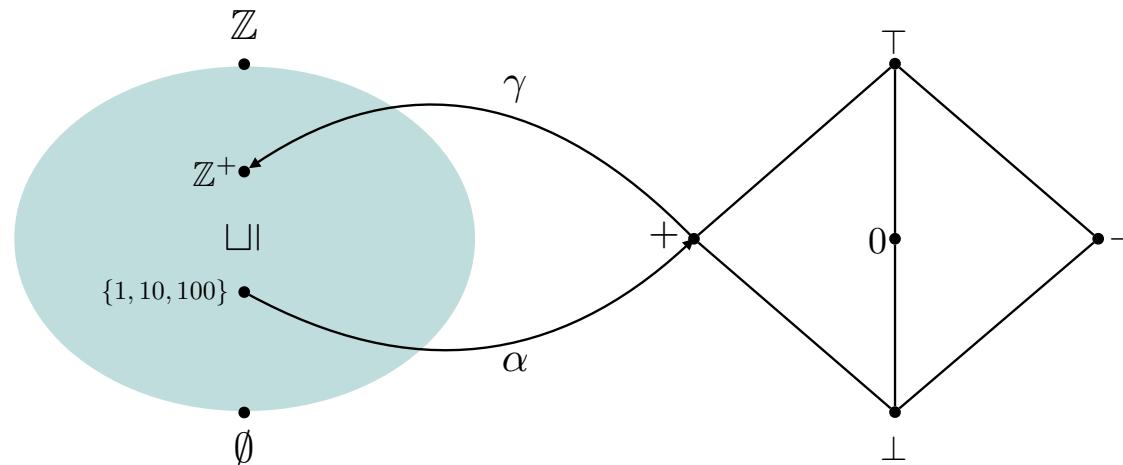
Requirement 1: Galois Connection

- Intuition: order preservation between two semantic domains



Example: Sign Abstraction

$$\wp(\mathbb{Z}) \xrightleftharpoons[\alpha]{\gamma} \{\perp, -, 0, +, \top\}$$



Example: Sign Abstraction

$$\wp(\mathbb{Z}) \xrightleftharpoons[\alpha]{\gamma} \{\perp, -, 0, +, \top\}$$

$$\alpha(Z) = \begin{cases} \perp & Z = \emptyset \\ + & \forall z \in Z. z > 0 \\ 0 & Z = \{0\} \\ - & \forall z \in Z. z < 0 \\ \top & \text{otherwise} \end{cases}$$

$$\gamma(\perp) = \emptyset$$

$$\gamma(+) = \{z \in \mathbb{Z} \mid z > 0\}$$

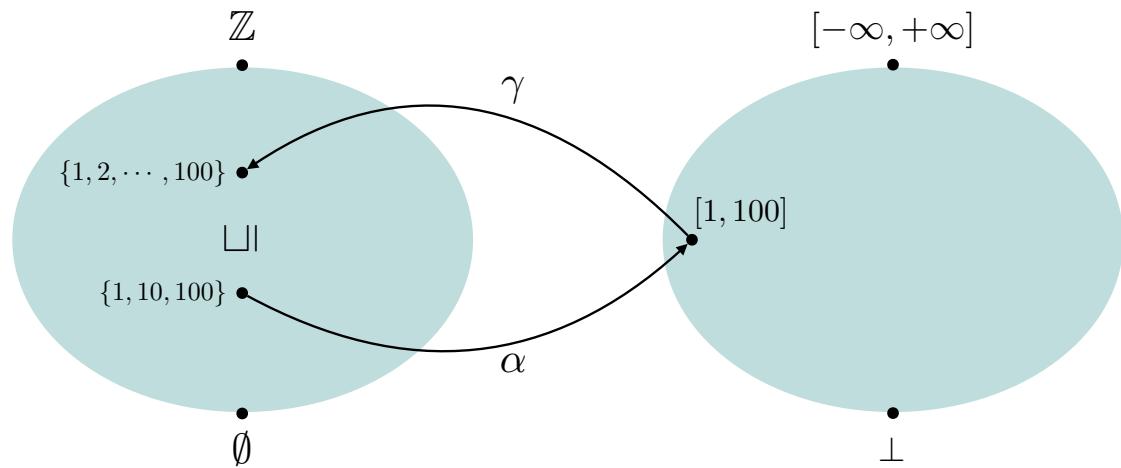
$$\gamma(0) = \{0\}$$

$$\gamma(-) = \{z \in \mathbb{Z} \mid z < 0\}$$

$$\gamma(\top) = \mathbb{Z}$$

Example: Interval Abstraction

$$\wp(\mathbb{Z}) \xrightleftharpoons[\alpha]{\gamma} \{\perp\} \cup \{[a, b] \mid a \in \mathbb{Z} \cup \{-\infty\}, b \in \mathbb{Z} \cup \{+\infty\}\}$$



Example: Interval Abstraction

$$\wp(\mathbb{Z}) \xrightleftharpoons[\alpha]{\gamma} \{\perp\} \cup \{[a, b] \mid a \in \mathbb{Z} \cup \{-\infty\}, b \in \mathbb{Z} \cup \{+\infty\}\}$$

$$\alpha(\emptyset) = \perp$$

$$\alpha(X) = [\min X, \max X]$$

$$\gamma(\perp) = \emptyset$$

$$\gamma([a, b]) = \{x \in \mathbb{Z} \mid a \leq x \leq b\}$$

Properties of Galois Connection

$$\forall x \in \mathbb{D}, x^\sharp \in \mathbb{D}^\sharp. \alpha(x) \sqsubseteq x^\sharp \iff x \sqsubseteq \gamma(x^\sharp)$$

- $id \sqsubseteq \gamma \circ \alpha$

$$\begin{aligned} & \alpha(x) \sqsubseteq \alpha(x) \\ \iff & x \sqsubseteq \gamma(\alpha(x)) \quad (\text{by Galois connection}) \end{aligned}$$

- $\alpha \circ \gamma \sqsubseteq id$

$$\begin{aligned} & \gamma(x^\sharp) \sqsubseteq \gamma(x^\sharp) \\ \iff & \alpha(\gamma(x^\sharp)) \sqsubseteq x^\sharp \quad (\text{by Galois connection}) \end{aligned}$$

- α is monotone

$$\begin{aligned} & x \sqsubseteq y \\ \implies & x \sqsubseteq \gamma(\alpha(y)) \quad (id \sqsubseteq \gamma \circ \alpha) \\ \iff & \alpha(x) \sqsubseteq \alpha(y) \quad (\text{by Galois connection}) \end{aligned}$$

- γ is monotone

$$\begin{aligned} & x^\sharp \sqsubseteq y^\sharp \\ \implies & \alpha(\gamma(x^\sharp)) \sqsubseteq y^\sharp \quad (\alpha \circ \gamma \sqsubseteq id) \\ \iff & \gamma(x^\sharp) \sqsubseteq \gamma(x^\sharp) \quad (\text{by Galois connection}) \end{aligned}$$

Deriving Galois Connections (1)

- Pointwise lifting:

Given a Galois connection $\mathbb{D} \xrightleftharpoons[\alpha]{\gamma} \mathbb{D}^\sharp$ and a set \mathbb{S}

$$\mathbb{S} \rightarrow \mathbb{D} \xrightleftharpoons[\alpha']{\gamma'} \mathbb{S} \rightarrow \mathbb{D}^\sharp$$

where $\alpha'(f) = \lambda x \in \mathbb{S}. \alpha(f(x))$ and $\gamma'(f^\sharp) = \lambda x \in \mathbb{S}. \gamma(f^\sharp(x))$

- Composition:

Given two galois connections $\mathbb{D}_1 \xrightleftharpoons[\alpha_1]{\gamma_1} \mathbb{D}_2 \xrightleftharpoons[\alpha_2]{\gamma_2} \mathbb{D}_3$

$$\mathbb{D}_1 \xrightleftharpoons[\alpha_2 \circ \alpha_1]{\gamma_1 \circ \gamma_2} \mathbb{D}_3$$

Deriving Galois Connections (2)

Given two Galois connections $\mathbb{D}_1 \xrightleftharpoons[\alpha_1]{\gamma_1} \mathbb{D}_1^\sharp$ and $\mathbb{D}_2 \xrightleftharpoons[\alpha_2]{\gamma_2} \mathbb{D}_2^\sharp$

- $\mathbb{D}_1 \times \mathbb{D}_2 \xrightleftharpoons[\alpha]{\gamma} \mathbb{D}_1^\sharp \times \mathbb{D}_2^\sharp$ where $\alpha = \lambda \langle x, y \rangle. \langle \alpha_1(x), \alpha_2(y) \rangle$
- $\mathbb{D}_1 + \mathbb{D}_2 \xrightleftharpoons[\alpha]{\gamma} \mathbb{D}_1^\sharp + \mathbb{D}_2^\sharp$ where $\alpha = \lambda x. \begin{cases} \alpha_1(x) & \text{if } x \in \mathbb{D}_1 \\ \alpha_2(x) & \text{if } x \in \mathbb{D}_2 \end{cases}$
- $\mathbb{D}_1 \rightarrow \mathbb{D}_2 \xrightleftharpoons[\alpha]{\gamma} \mathbb{D}_1^\sharp \rightarrow \mathbb{D}_2^\sharp$ where $\alpha = \lambda f. \alpha_2 \circ f \circ \gamma_1$

Deriving Galois Connections (3)

Given two Galois connections $\wp(\mathbb{A}) \xrightleftharpoons[\alpha_1]{\gamma_1} \mathbb{D}_1^\sharp$ and $\wp(\mathbb{B}) \xrightleftharpoons[\alpha_2]{\gamma_2} \mathbb{D}_2^\sharp$

- $\wp(\mathbb{A} \times \mathbb{B}) \xrightleftharpoons[\alpha]{\gamma} \mathbb{D}_1^\sharp \times \mathbb{D}_2^\sharp$ where $\alpha = \lambda X. \langle \alpha_1(\{a \mid \langle a, b \rangle \in X\}), \alpha_2(\{b \mid \langle a, b \rangle \in X\}) \rangle$
- $\wp(\mathbb{A} \times \mathbb{B}) \xrightleftharpoons[\alpha]{\gamma} \mathbb{A}' \rightarrow \mathbb{D}_2^\sharp$ where $\alpha = \lambda X. \{a \mapsto \alpha_2(S) \mid \langle a, b \rangle \in X, S = \{b \mid \langle a, b \rangle \in X\}\}$ and $\mathbb{A}' \subseteq \mathbb{A}$
- $\wp(\mathbb{A} + \mathbb{B}) \xrightleftharpoons[\alpha]{\gamma} \mathbb{D}_1^\sharp \times \mathbb{D}_2^\sharp$ where $\alpha = \lambda X. \langle \alpha_1(X \cap \mathbb{A}), \alpha_2(X \cap \mathbb{B}) \rangle$
- $\wp(\mathbb{A}) \rightarrow \wp(\mathbb{B}) \xrightleftharpoons[\alpha]{\gamma} \mathbb{D}_1^\sharp \rightarrow \mathbb{D}_2^\sharp$ where $\alpha = \lambda f. \alpha_2 \circ f \circ \gamma_1$

Example

- Concrete domain: $\mathbb{D} = \wp(\mathbb{M}) \rightarrow \wp(\mathbb{M})$ where $\mathbb{M} = \mathbb{X} \rightarrow \mathbb{Z}$
- Abstract domain: $\mathbb{D}^\sharp = \mathbb{M}^\sharp \rightarrow \mathbb{M}^\sharp$ where $\mathbb{M}^\sharp = \mathbb{X} \rightarrow \mathbb{Z}^\sharp$
- Galois connection: $\wp(\mathbb{D}) \xrightleftharpoons[\alpha]{\gamma} \mathbb{D}^\sharp$
 - Memory abstraction: $\wp(\mathbb{M}) \xrightleftharpoons[\alpha_{\mathbb{M}}]{\gamma_{\mathbb{M}}} \mathbb{M}^\sharp$ via $\wp(\mathbb{X} \rightarrow \mathbb{Z}) \xrightleftharpoons[\alpha_{\mathbb{M}}]{\gamma_{\mathbb{M}}} \mathbb{X} \rightarrow \mathbb{Z}^\sharp$
$$\gamma_{\mathbb{M}} : \mathbb{M}^\sharp \rightarrow \wp(\mathbb{M})$$
$$\gamma_{\mathbb{M}} = \lambda m^\sharp. \{m \mid \forall x. m(x) \in \gamma_{\mathbb{Z}}(m^\sharp(x))\}$$
 - Value abstraction: $\wp(\mathbb{Z}) \xrightleftharpoons[\alpha_{\mathbb{Z}}]{\gamma_{\mathbb{Z}}} \mathbb{Z}^\sharp$

Requirement 2: F and F^\sharp

- F^\sharp is a sound abstraction of F (option 1)

$$F \circ \gamma \sqsubseteq \gamma \circ F^\sharp$$

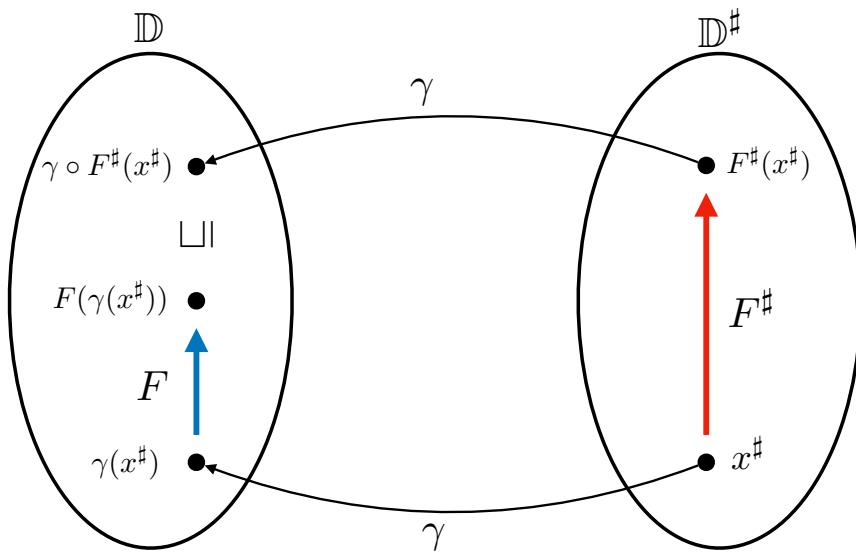
- F^\sharp is a sound abstraction of F (option 2)

$$x \sqsubseteq \gamma(x^\sharp) \implies F(x) \sqsubseteq \gamma(F^\sharp(x^\sharp))$$

Intuition: the result of one-step abstract execution (F^\sharp)
subsumes that of one-step concrete execution (F)

Sound Abstract Semantics (1)

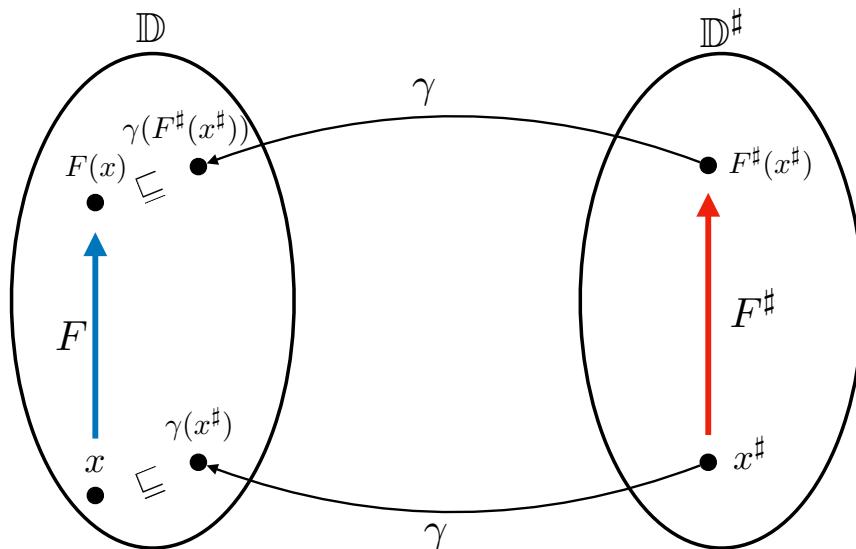
$$F \circ \gamma \sqsubseteq \gamma \circ F^\sharp$$



Intuition: the result of one-step abstract execution (F^\sharp)
subsumes that of one-step concrete execution (F)

Sound Abstract Semantics (2)

$$x \sqsubseteq \gamma(x^\sharp) \implies F(x) \sqsubseteq \gamma(F^\sharp(x^\sharp))$$



Intuition: the result of one-step abstract execution (F^\sharp)
subsumes that of one-step concrete execution (F)

Soundness

- Static analysis result $\bigsqcup_{i \geq 0} F^{\sharp i}(\perp)$ soundly subsumes all possible executions

$$\text{lfp } F \sqsubseteq \gamma\left(\bigsqcup_{i \geq 0} F^{\sharp i}(\perp)\right)$$

- How to guarantee the soundness?
- How to compute the sound result within finite time?

Fixpoint Transfer Theorems

- With option 1

Theorem (Fixpoint Transfer 1). *Let \mathbb{D} and \mathbb{D}^\sharp be related by Galois connection $\mathbb{D} \xleftarrow[\alpha]{\gamma} \mathbb{D}^\sharp$. Let $F : \mathbb{D} \rightarrow \mathbb{D}$ be a continuous function and $F^\sharp : \mathbb{D}^\sharp \rightarrow \mathbb{D}^\sharp$ be a monotone or extensive function such that $F \circ \gamma \sqsubseteq \gamma \circ F^\sharp$. Then,*

$$\text{lfp } F \sqsubseteq \gamma\left(\bigsqcup_{i \geq 0} F^{\sharp i}(\perp^\sharp)\right).$$

- With option 2

Theorem (Fixpoint Transfer 2). *Let \mathbb{D} and \mathbb{D}^\sharp be related by Galois connection $\mathbb{D} \xleftarrow[\alpha]{\gamma} \mathbb{D}^\sharp$. Let $F : \mathbb{D} \rightarrow \mathbb{D}$ be a continuous function and $F^\sharp : \mathbb{D}^\sharp \rightarrow \mathbb{D}^\sharp$ be a monotone or extensive function such that $x \sqsubseteq \gamma(x^\sharp) \implies F(x) \sqsubseteq \gamma(F^\sharp(x^\sharp))$. Then,*

$$\text{lfp } F \sqsubseteq \gamma\left(\bigsqcup_{i \geq 0} F^{\sharp i}(\perp^\sharp)\right).$$

Proof

Theorem (Fixpoint Transfer 1). Let \mathbb{D} and \mathbb{D}^\sharp be related by Galois connection $\mathbb{D} \xrightleftharpoons[\alpha]{\gamma} \mathbb{D}^\sharp$. Let $F : \mathbb{D} \rightarrow \mathbb{D}$ be a continuous function and $F^\sharp : \mathbb{D}^\sharp \rightarrow \mathbb{D}^\sharp$ be a monotone or extensive function such that $F \circ \gamma \sqsubseteq \gamma \circ F^\sharp$. Then,

$$\text{lfp } F \sqsubseteq \gamma\left(\bigcup_{i \geq 0} F^{\sharp i}(\perp^\sharp)\right).$$

Proof.

1. First we prove $\forall n \in \mathbb{N}. F^n(\perp) \sqsubseteq \gamma(F^{\sharp n}(\perp^\sharp))$ by induction. The base case is trivial. The inductive case is as follows:

$$\begin{aligned} F^{n+1}(\perp) &= F \circ F^n(\perp) \\ &\sqsubseteq F \circ \gamma(F^{\sharp n}(\perp^\sharp)) \quad (\text{by induction hypothesis and monotonicity of } F) \\ &\sqsubseteq \gamma \circ F^\sharp \circ F^{\sharp n}(\perp^\sharp) \quad (\text{by assumption } F \circ \gamma \sqsubseteq \gamma \circ F^\sharp) \\ &= \gamma(F^{\sharp n+1}(\perp^\sharp)) \end{aligned}$$

2. $\{F^i(\perp)\}_i$ is a chain because F is continuous (so monotone). Then, the least upper bound of the chain $\bigcup_{i \geq 0} F^i(\perp)$ exists because \mathbb{D} is a CPO.

3. $\{F^{\sharp i}(\perp^\sharp)\}_i$ is a chain because F^\sharp is monotone or extensive. Then, $\{\gamma(F^{\sharp i}(\perp^\sharp))\}_i$ is also a chain because γ is monotone. Therefore, the least upper bound of the chain $\bigcup_{i \geq 0} \{\gamma(F^{\sharp i}(\perp^\sharp))\}_i$ exists.

4. Finally,

$$\begin{aligned} \text{lfp } F &= \bigcup_{i \geq 0} F^i(\perp) \sqsubseteq \bigcup_{i \geq 0} \gamma(F^{\sharp i}(\perp^\sharp)) \\ &\sqsubseteq \gamma\left(\bigcup_{i \geq 0} (F^{\sharp i}(\perp^\sharp))\right) \quad (\text{by monotonicity of } \gamma) \end{aligned}$$

Abstraction of Compositions

- Composition of sound semantic functions is also sound

$$\begin{array}{ccc} \mathbb{D}_1 & \xrightleftharpoons[\alpha_1]{\gamma_1} & \mathbb{D}_1^\sharp \\ f \downarrow & & \downarrow f^\sharp \\ \mathbb{D}_2 & \xrightleftharpoons[\alpha_2]{\gamma_2} & \mathbb{D}_2^\sharp \\ g \downarrow & & \downarrow g^\sharp \\ \mathbb{D}_3 & \xrightleftharpoons[\alpha_3]{\gamma_3} & \mathbb{D}_3^\sharp \end{array} \quad \longrightarrow \quad g \circ f \circ \gamma_1 \sqsubseteq \gamma_3 \circ g^\sharp \circ f^\sharp$$

- Implication: soundness proof can be compositional

Abstract Semantics of Expressions

$$\llbracket E \rrbracket^\sharp : \mathbb{M}^\sharp \rightarrow \mathbb{Z}^\sharp$$

$$\llbracket n \rrbracket^\sharp = \lambda m^\sharp. \alpha(\{n\})$$

$$\llbracket x \rrbracket^\sharp = \lambda m^\sharp. m^\sharp(x)$$

$$\llbracket E_1 \odot E_2 \rrbracket^\sharp = \lambda m^\sharp. \llbracket E_1 \rrbracket^\sharp(m^\sharp) \odot^\sharp \llbracket E_2 \rrbracket^\sharp(m^\sharp)$$

- Example
 - Sign domain $\mathbb{Z}^\sharp = \{\perp, -, 0, +, \top\}$
 - Interval domain $\mathbb{Z}^\sharp = \{\perp\} \cup \{[a, b] \mid a \in \mathbb{Z} \cup \{-\infty\}, b \in \mathbb{Z} \cup \{+\infty\}\}$
 - Soundness: $\llbracket E \rrbracket_\wp \circ \gamma_{\mathbb{M}} \subseteq \gamma_{\mathbb{Z}} \circ \llbracket E \rrbracket^\sharp$

Example: Binary Operator

$$\wp(\mathbb{Z}) \xrightleftharpoons[\alpha]{\gamma} \{\perp, -, 0, +, \top\}$$

- Define an abstract addition operator $+^\sharp : \mathbb{Z}^\sharp \rightarrow \mathbb{Z}^\sharp \rightarrow \mathbb{Z}^\sharp$ (as precise as possible)

	\perp	$-$	0	$+$	\top
\perp	\perp	\perp	\perp	\perp	\perp
$-$	\perp	$-$	$-$	\top	\top
0	\perp	$-$	0	$+$	\top
$+$	\perp	\top	$+$	$+$	\top
\top	\perp	\top	\top	\top	\top

Soundness Proof

$$\llbracket E \rrbracket_{\wp} \circ \gamma_{\mathbb{M}} \subseteq \gamma_{\mathbb{Z}} \circ \llbracket E \rrbracket^{\sharp}$$

- $E : n$

$$\begin{aligned}\llbracket n \rrbracket_{\wp} \circ \gamma_{\mathbb{M}}(m^{\sharp}) &= \{n\} \\ &= \gamma_{\mathbb{Z}} \circ \llbracket n \rrbracket^{\sharp}(m^{\sharp})\end{aligned}$$

- $E : x$

- $E : E_1 \odot E_2$

- Exercise!

- For each operator, prove using induction

$$\begin{aligned}\llbracket x \rrbracket_{\wp} \circ \gamma_{\mathbb{M}}(m^{\sharp}) &= \llbracket x \rrbracket_{\wp} \{m \mid \forall x. m(x) \in \gamma_{\mathbb{Z}}(m^{\sharp}(x))\} \\ &= \{m(x) \mid \forall x. m(x) \in \gamma_{\mathbb{Z}}(m^{\sharp}(x))\} \\ &\subseteq \gamma_{\mathbb{Z}}(m^{\sharp}(x)) \\ &= \gamma_{\mathbb{Z}} \circ \llbracket x \rrbracket^{\sharp}(m^{\sharp})\end{aligned}$$

Abstract Semantics of Conditions

$$\llbracket B \rrbracket^\sharp : \mathbb{M}^\sharp \rightarrow \mathbb{M}^\sharp$$

$$\llbracket \text{true} \rrbracket^\sharp = \lambda m^\sharp. m^\sharp$$

$$\llbracket \text{false} \rrbracket^\sharp = \lambda m^\sharp. \perp$$

- Example

- Sign domain $\llbracket x < 0 \rrbracket^\sharp = \lambda m^\sharp. \begin{cases} \perp & \text{if } m^\sharp(x) \in \{+, 0, \perp\} \\ m^\sharp\{x \mapsto -\} & \text{o.w.} \end{cases}$
- Interval domain $\llbracket x < n \rrbracket^\sharp = \lambda m^\sharp. \begin{cases} \perp & \text{if } m^\sharp(x) = [a, b] \wedge a > n \\ m^\sharp\{x \mapsto [a, n-1]\} & \text{if } a \leq n \leq b \\ m^\sharp & \text{if } b < n \end{cases}$
- Soundness: $\llbracket B \rrbracket_\wp \circ \gamma_{\mathbb{M}} \subseteq \gamma_{\mathbb{M}} \circ \llbracket B \rrbracket^\sharp$

Soundness Proof

$$\llbracket B \rrbracket_{\wp} \circ \gamma_{\mathbb{M}} \subseteq \gamma_{\mathbb{M}} \circ \llbracket B \rrbracket^{\sharp}$$

- $B : x < 0$ and $\llbracket x < 0 \rrbracket^{\sharp} = \lambda m^{\sharp}. \begin{cases} \perp & \text{if } m^{\sharp}(x) \in \{+, 0, \perp\} \\ m^{\sharp}\{x \mapsto -\} & \text{o.w.} \end{cases}$

$$\begin{aligned} \llbracket x < 0 \rrbracket_{\wp} \circ \gamma_{\mathbb{M}}(m^{\sharp}) &= \llbracket x < 0 \rrbracket_{\wp}\{m \mid \forall i. m(i) \in \gamma_{\mathbb{Z}}(m^{\sharp}(i))\} \\ &= \{m \mid \forall i. m(i) \in \gamma_{\mathbb{Z}}(m^{\sharp}(i)), m(x) < 0\} \end{aligned}$$

- Case 1: $m^{\sharp}(x) \in \{+, 0, \perp\}$

$$LHS = \emptyset = RHS$$

- Case 2: $m^{\sharp}(x) \in \{-, \top\}$

$$\begin{aligned} \gamma_{\mathbb{M}} \circ \llbracket x < 0 \rrbracket^{\sharp}(m^{\sharp}) &= \gamma_{\mathbb{M}}(m^{\sharp}\{x \mapsto -\}) \\ &= \{m \mid \forall i. m(i) \in \gamma_{\mathbb{Z}}(m^{\sharp}\{x \mapsto -\}(i)), m(x) < 0\} \\ &= LHS \end{aligned}$$

Abstract Semantics of Commands

$$\llbracket C \rrbracket^\# : \mathbb{M}^\# \rightarrow \mathbb{M}^\#$$

$$\llbracket \text{skip} \rrbracket^\# = \lambda m^\#.m^\#$$

$$\llbracket C_0 ; C_1 \rrbracket^\# = \lambda m^\#. \llbracket C_1 \rrbracket^\# \circ \llbracket C_0 \rrbracket^\#(m^\#)$$

$$\llbracket x := E \rrbracket^\# = \lambda m^\#.m^\#\{x \mapsto \llbracket E \rrbracket^\#(m^\#)\}$$

$$\llbracket \text{input}(x) \rrbracket^\# = \lambda m^\#.m^\#\{x \mapsto \alpha(\mathbb{Z})\}$$

$$\llbracket \text{if } B \text{ then } C_1 \text{ else } C_2 \rrbracket^\# = \lambda m^\#. \llbracket C_1 \rrbracket^\# \circ \llbracket B \rrbracket^\#(m^\#) \sqcup \llbracket C_2 \rrbracket^\# \circ \llbracket \neg B \rrbracket^\#(m^\#)$$

$$\llbracket \text{while } B \text{ } C \rrbracket^\# = \lambda m^\#. \llbracket \neg B \rrbracket^\# (\text{lfp} \lambda X. m^\# \sqcup \llbracket C \rrbracket^\# \circ \llbracket B \rrbracket^\#(X))$$

- Soundness: $\llbracket C \rrbracket_\wp \circ \gamma_{\mathbb{M}} \subseteq \gamma_{\mathbb{M}} \circ \llbracket C \rrbracket^\#$ and $\cup \circ (\gamma_{\mathbb{M}}, \gamma_{\mathbb{M}}) \subseteq \gamma_{\mathbb{M}} \circ \sqcup$

Soundness Proof

$$\llbracket C \rrbracket_\wp \circ \gamma_{\mathbb{M}} \subseteq \gamma_{\mathbb{M}} \circ \llbracket C \rrbracket^\sharp$$

- $C_1 ; C_2$

$$\begin{aligned} \llbracket C_0 ; C_1 \rrbracket_\wp \circ \gamma_{\mathbb{M}}(m^\sharp) &= \llbracket C_0 \rrbracket_\wp \circ \llbracket C_1 \rrbracket_\wp \circ \gamma_{\mathbb{M}}(m^\sharp) \\ &\subseteq \llbracket C_0 \rrbracket_\wp \circ \gamma_{\mathbb{M}} \circ \llbracket C_1 \rrbracket^\sharp(m^\sharp) \\ &\subseteq \gamma_{\mathbb{M}} \circ \llbracket C_0 \rrbracket^\sharp \circ \llbracket C_1 \rrbracket^\sharp(m^\sharp) \\ &= \gamma_{\mathbb{M}} \circ \llbracket C_0 ; C_1 \rrbracket^\sharp(m^\sharp) \end{aligned}$$

- $\text{if } B \text{ then } C_1 \text{ else } C_2$

$$\begin{aligned} \llbracket \text{if } B \text{ then } C_1 \text{ else } C_2 \rrbracket_\wp \circ \gamma_{\mathbb{M}}(m^\sharp) &= \llbracket C_1 \rrbracket_\wp \circ \llbracket B \rrbracket_\wp \circ \gamma_{\mathbb{M}}(m^\sharp) \cup \llbracket C_2 \rrbracket_\wp \circ \llbracket \neg B \rrbracket_\wp \circ \gamma_{\mathbb{M}}(m^\sharp) \\ &\subseteq \llbracket C_1 \rrbracket_\wp \circ \gamma_{\mathbb{M}} \circ \llbracket B \rrbracket^\sharp(m^\sharp) \cup \llbracket C_2 \rrbracket_\wp \circ \gamma_{\mathbb{M}} \circ \llbracket B \rrbracket^\sharp(m^\sharp) \\ &\subseteq \gamma_{\mathbb{M}} \circ \llbracket C_1 \rrbracket^\sharp \circ \llbracket B \rrbracket^\sharp(m^\sharp) \cup \gamma_{\mathbb{M}} \circ \llbracket C_2 \rrbracket^\sharp \circ \llbracket B \rrbracket^\sharp(m^\sharp) \\ &\subseteq \gamma_{\mathbb{M}}(\llbracket C_1 \rrbracket^\sharp \circ \llbracket B \rrbracket^\sharp(m^\sharp) \sqcup \llbracket C_2 \rrbracket^\sharp \circ \llbracket B \rrbracket^\sharp(m^\sharp)) \\ &= \gamma_{\mathbb{M}} \circ \llbracket \text{if } B \text{ then } C_1 \text{ else } C_2 \rrbracket^\sharp(m^\sharp) \end{aligned}$$

- $\text{while } B \text{ } C$

$$\begin{aligned} \llbracket \text{while } B \text{ } C \rrbracket_\wp \circ \gamma_{\mathbb{M}}(m^\sharp) &= \llbracket \neg B \rrbracket_\wp(\text{lfp} \lambda X. \gamma_{\mathbb{M}}(m^\sharp) \cup \llbracket C \rrbracket_\wp \circ \llbracket B \rrbracket_\wp(X)) \\ &\subseteq \llbracket \neg B \rrbracket_\wp \circ \gamma_{\mathbb{M}}(\text{lfp} \lambda X. m^\sharp \sqcup \llbracket C \rrbracket^\sharp \circ \llbracket B \rrbracket^\sharp(X)) \\ &\subseteq \gamma_{\mathbb{M}} \circ \llbracket \neg B \rrbracket^\sharp(\text{lfp} \lambda X. m^\sharp \sqcup \llbracket C \rrbracket^\sharp \circ \llbracket B \rrbracket^\sharp(X)) \end{aligned}$$

- $(\text{lfp} \lambda X. \gamma_{\mathbb{M}}(m^\sharp) \cup \llbracket C \rrbracket_\wp \circ \llbracket B \rrbracket_\wp(X)) \subseteq \gamma_{\mathbb{M}}(\text{lfp} \lambda X. m^\sharp \sqcup \llbracket C \rrbracket^\sharp \circ \llbracket B \rrbracket^\sharp(X))$

by the fixed point transfer theorem if

$$\begin{aligned} &(\lambda X. \gamma_{\mathbb{M}}(m^\sharp) \cup \llbracket C \rrbracket_\wp \circ \llbracket B \rrbracket_\wp(X)) \circ \gamma_{\mathbb{M}}(m^\sharp) \\ &= \gamma_{\mathbb{M}}(m^\sharp) \cup \llbracket C \rrbracket_\wp \circ \llbracket B \rrbracket_\wp(\gamma_{\mathbb{M}}(m^\sharp)) \\ &\subseteq \gamma_{\mathbb{M}}(m^\sharp) \cup \llbracket C \rrbracket_\wp \circ \gamma_{\mathbb{M}} \circ \llbracket B \rrbracket^\sharp(m^\sharp) \\ &\subseteq \gamma_{\mathbb{M}}(m^\sharp) \cup \gamma_{\mathbb{M}} \circ \llbracket C \rrbracket^\sharp \circ \llbracket B \rrbracket^\sharp(m^\sharp) \\ &\subseteq \gamma_{\mathbb{M}}(m^\sharp \sqcup \llbracket C \rrbracket^\sharp \circ \llbracket B \rrbracket^\sharp(m^\sharp)) \\ &= \gamma_{\mathbb{M}} \circ (\lambda X. m^\sharp \sqcup \llbracket C \rrbracket^\sharp \circ \llbracket B \rrbracket^\sharp(X))(m^\sharp) \end{aligned}$$

Summary

- Abstract semantics: a **sound approximation** of the concrete semantics
- **Soundness guarantee** by Fixpoint transfer theorem
 - Galois connection, sound abstract semantic function
- Static analysis design: design of **precise and efficient** abstract semantics
- Plan: compute the abstract semantics within a **finite time**