

IS593: Language-based Security

7. Static Analysis for Advanced Programming Features

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Advanced Programming Features

- So far, we have learned **design and implementation** of static analysis
 - Domains, semantics, abstractions, algorithms
 - For basic programming features: assignments, conditionals, loops
- From now, we extend SmaLLVM to support more **advanced features**
 - In this lecture, pointers with dynamic allocations, and functions

The ThriLLVM Language

- Extension 1: Pointers

$L ::= l$	program label
$E ::= n$	integer
x	variable
$E \oplus E$	arithmetic operation
$E \otimes E$	comparison operation
$\&x$	location of a variable
malloc	location of a newly allocated memory
$*E$	dereference of a memory location
$\oplus ::= + - \times /$	
$\otimes ::= < \leq > \geq == !=$	
$C ::= x := E$	assignment
$\text{br } E L L$	conditional jump
$\text{goto } L$	unconditional jump
$x := \text{input}()$	input
$\text{print}(x)$	print
$*E := E$	indirect assignment



Example

x = malloc;	$\{x \mapsto a\}$
y = &x;	$\{x \mapsto a, y \mapsto x\}$
z = x;	$\{x \mapsto a, y \mapsto x, z \mapsto a\}$
*x = 5;	$\{x \mapsto a, y \mapsto x, z \mapsto a, a \mapsto 5\}$
*y = *x;	$\{x \mapsto 5, y \mapsto x, z \mapsto a, a \mapsto 5\}$

Memory addresses (variables or heap addresses)
also can be values

Concrete Domains

- Similar as before except for the followings:
 - A memory location is either a variable or a heap address
 - A value is either an integer or a memory location

$S = \mathbb{L} \times \mathbb{M}$	states
$\mathbb{M} = \mathbb{A} \rightarrow \mathbb{V}$	memories
$\mathbb{A} = \mathbb{X} \cup \mathbb{H}$	addresses (locations)
<u>$\mathbb{V} = \mathbb{Z} \cup \mathbb{A}$</u>	values
\mathbb{X}	set of variables
\mathbb{H}	set of allocated heap addresses
\mathbb{L}	set of statement labels

Domain of Heap Addresses

$$\mathbb{H} = \mathbb{N}_{site} \times \mathbb{N}$$

- A heap address is represented as a pair of an allocation-site and a counter
 - Given a program, every `malloc` expression is assumed to have a unique number μ , writing `malloc μ`
 - For example, the addresses of fresh locations from a `malloc`-site μ are $(\mu, 0), (\mu, 1), (\mu, 2), \dots$

```
while(*) {  
   $\mu$ : p = malloc(size);  
}
```

State Transition

- The semantics is specified by a transition system $(\mathbb{S}, \rightarrow)$
 - $\mathbb{S} = \mathbb{L} \times \mathbb{M}$: the set of states $\langle l, m \rangle$
 - $(\rightarrow) \subseteq \mathbb{S} \times \mathbb{S}$: the transition relation that describes computation steps
- New rule: the transition for indirect assignments

$$*E_1 := E_2 : \langle l, m \rangle \rightarrow \langle \text{next}(l), \text{update}(m, \text{eval}_{E_1}(m), \text{eval}_{E_2}(m)) \rangle$$

Semantic Operators

- The memory read/write operations

$$fetch : \mathbb{M} \times \mathbb{V} \rightarrow \mathbb{V}$$

$$fetch(m, v) = m(v)$$

$$update : \mathbb{M} \times \mathbb{V} \times \mathbb{V} \rightarrow \mathbb{M}$$

$$update(m, v_1, v_2) = m\{v_1 \mapsto v_2\}$$

- The expression-evaluation operation

$$eval_E : \mathbb{M} \rightarrow \mathbb{V}$$

$$eval_x(m) = fetch(m, x)$$

$$eval_{\&x}(m) = x$$

variable as a location &x

$$eval_{\text{malloc}_\mu}(m) = (\mu, z)$$

new number z for malloc site μ

$$eval_{*E}(m) = fetch(m, eval_E(m))$$

dereference of a location

- The memory filter operation

$$filter_E : \mathbb{M} \rightarrow \mathbb{M}$$

$$filter_E(m) = m \quad \text{if } eval_E(m) = \text{true}$$

Concrete Semantics

- Concrete domain: $\mathbb{D} = \wp(\mathbb{S})$ where $\mathbb{S} = \mathbb{L} \times \mathbb{M}$
- Concrete semantic function:

$$F : \wp(\mathbb{S}) \rightarrow \wp(\mathbb{S})$$

$$F(X) = I \cup Step(X)$$

where I is the set of initial states and $Step$ is the powerset-lifted version of \hookrightarrow

$$Step : \wp(\mathbb{S}) \rightarrow \wp(\mathbb{S})$$

$$Step(X) = \{s' \mid s \hookrightarrow s', s \in X\}$$

- Concrete semantic (i.e., reachable states): $\text{lfp } F$

Abstract Domains

- Point-wise abstractions of underlying abstractions (addresses and values)

$S^\#$	$=$	$\mathbb{L} \times M^\#$	abstract states
$M^\#$	$=$	$(\mathbb{X} \cup H^\#) \rightarrow V^\#$	abstract memories
$A^\#$	$=$	$\wp(\mathbb{X} \cup H^\#)$	abstract addresses (locations)
$V^\#$	$=$	$Z^\# \times A^\#$	abstract values
\mathbb{X}			set of variables
$H^\#$			set of abstract heap addresses
\mathbb{L}			set of statement labels

Abstraction of Values

- Kind-wise abstractions of a set of values (integers and locations)

$$\wp(\mathbb{V} = \mathbb{Z} \cup \mathbb{A}) \xrightleftharpoons[\alpha_{\mathbb{V}}]{\gamma_{\mathbb{V}}} \mathbb{V}^{\sharp} = \mathbb{Z}^{\sharp} \times \mathbb{A}^{\sharp}$$

$$\alpha_{\mathbb{V}}(V) = (\alpha_{\mathbb{Z}}(V \cap \mathbb{Z}), \alpha_{\mathbb{A}}(V \cap \mathbb{A}))$$

$$\gamma_{\mathbb{V}}(z^{\sharp}, a^{\sharp}) = \gamma_{\mathbb{Z}}(z^{\sharp}) \cup \gamma_{\mathbb{A}}(a^{\sharp})$$

- All the operators (order, join, meet, widening, narrowing, etc) are defined as kind-wise liftings

Abstraction of Locations

- Abstract all heap locations allocated at a `malloc` site μ into a single abstract address (so called *allocation-site-based abstraction*)

$$\wp(\mathbb{H} = \mathbb{N}_{site} \times \mathbb{N}) \xleftrightarrow[\alpha_{\mathbb{H}}]{\gamma_{\mathbb{H}}} \mathbb{H}^\sharp = \wp(\mathbb{N}_{site})$$

$$\alpha_{\mathbb{H}}(H) = \{\mu \mid \langle \mu, _ \rangle \in H\}$$

$$\gamma_{\mathbb{H}}(h^\sharp) = \{\langle \mu, n \rangle \mid \mu \in h^\sharp, n \in \mathbb{N}\}$$

- The abstract domain is the power set of variables and abstract heap locations
 - It is a finite set: the set of variables and `malloc` sites are finite for a program

$$\wp(\mathbb{A} = \mathbb{X} \cup \mathbb{H}) \xleftrightarrow[\alpha_{\mathbb{A}}]{\gamma_{\mathbb{A}}} \mathbb{A}^\sharp = \wp(\mathbb{X} \cup \mathbb{H}^\sharp)$$

Abstract State Transition

- The abstract semantics is defined using a transition system $(\mathbb{S}^\sharp, \hookrightarrow^\sharp)$
 - $\mathbb{S}^\sharp = \mathbb{L} \times \mathbb{M}^\sharp$: the set of states $\langle l, m^\sharp \rangle$
 - $(\hookrightarrow^\sharp) \subseteq \mathbb{S}^\sharp \times \mathbb{S}^\sharp$: the transition relation that describes computation steps
- The abstract transition for indirect assignments

$$*E_1 := E_2 : \langle l, m^\sharp \rangle \hookrightarrow^\sharp \langle \text{next}(l), \text{update}^\sharp(m^\sharp, \text{eval}_{E_1}^\sharp(m^\sharp), \text{eval}_{E_2}^\sharp(m^\sharp)) \rangle$$

Abstract Semantic Operators

- The abstract expression-evaluation operation:

$$\begin{aligned} eval_E^\sharp &: \mathbb{M}^\sharp \rightarrow \mathbb{V}^\sharp \\ eval_x^\sharp(m^\sharp) &= fetch^\sharp(m^\sharp, x) \\ eval_{\&x}^\sharp(m^\sharp) &= \{x\} \\ eval_{\text{malloc}_\mu}^\sharp(m^\sharp) &= \{\mu\} \\ eval_{*_E}^\sharp(m^\sharp) &= fetch^\sharp(m^\sharp, eval_E^\sharp(m^\sharp)) \end{aligned}$$

- The abstract memory filter operation:

$$\begin{aligned} filter_E^\sharp &: \mathbb{M}^\sharp \rightarrow \mathbb{M}^\sharp \\ filter_E^\sharp(m^\sharp) &= \alpha_{\mathbb{M}}(\{m \in \gamma_{\mathbb{M}}(m^\sharp) \mid eval_E(m) = \text{true}\}) \end{aligned}$$

Safe Memory Read Operation

- The abstract memory **read** operation $fetch^\# : \mathbb{M}^\# \times \mathbb{A}^\# \rightarrow \mathbb{V}^\#$: looks up the abstract memory entry at the abstract location
- Since the **abstract** location is a **set** of variables and malloc-sites, the result is the **join** of all the entries:

$$fetch^\#(m^\#, a^\#) = \bigsqcup_{a \in a^\#} m^\#(a)$$

Safe Memory Write Operation

- The abstract memory **write** operation: $update^\# : \mathbb{M}^\# \times \mathbb{A}^\# \times \mathbb{V}^\# \rightarrow \mathbb{M}^\#$
 - **overwrite** the memory entry (called **strong update**) when the target abstract location means a single concrete location
 - Otherwise, every entry that constitutes the target abstract location must be **joined** with the value to store (called **weak update**)

$$update^\#(m^\#, a^\#, v^\#) = \begin{cases} m^\#[x \mapsto v^\#] & \text{if } \gamma_{\mathbb{A}}(a^\#) = \{x\} \\ \bigsqcup_{a \in a^\#} m^\#[a \mapsto m^\#(a) \sqcup v^\#] & \text{otherwise} \end{cases}$$

Example

```
x = 0;  
y = 1;  
if (*) {  
    p = &x;      {x ↦ [0, 0], y ↦ [1, 1], p ↦ {x}}  
} else {  
    p = &y;      {x ↦ [0, 0], y ↦ [1, 1], p ↦ {y}}  
}  
z = *p;      {x ↦ [0, 0], y ↦ [1, 1], p ↦ {x, y}, z ↦ [0, 1]}  
*p = 2;      {x ↦ [0, 2], y ↦ [1, 2], p ↦ {x, y}, z ↦ [0, 1]}
```

Abstract Semantics

- Abstract domain: $\mathbb{D}^\sharp = \mathbb{L} \rightarrow \mathbb{M}^\sharp$

- The abstract semantic function:

$$F^\sharp : (\mathbb{L} \rightarrow \mathbb{M}^\sharp) \rightarrow (\mathbb{L} \rightarrow \mathbb{M}^\sharp)$$

$$F^\sharp(X) = \alpha(I) \sqcup Step^\sharp(X)$$

- The abstract semantic functions for transition and partition:

$$Step^\sharp : (\mathbb{L} \rightarrow \mathbb{M}^\sharp) \rightarrow (\mathbb{L} \rightarrow \mathbb{M}^\sharp)$$

$$Step^\sharp(X) = \wp(id, \sqcup) \circ \pi \circ \wp(\hookrightarrow^\sharp)(X)$$

$$\pi : \wp(\mathbb{S}^\sharp) \rightarrow (\mathbb{L} \rightarrow \wp(\mathbb{M}^\sharp))$$

$$\pi(X) = \lambda l. \{ m^\sharp \in \mathbb{M}^\sharp \mid \langle l, m^\sharp \rangle \in X \}$$

- Soundness: $\text{lfp } F \subseteq \gamma(\bigsqcup_{i \geq 0} F^\sharp(\perp))$

Soundness Proof

$\text{lfp } F \subseteq \gamma(\bigsqcup_{i \geq 0} F^\sharp(\perp))$ if $F^\sharp \circ \gamma \subseteq \gamma \circ F^\sharp$ by the fixpoint transfer theorem

It is enough to prove the following soundnesses for all semantic operator:

$$\wp(eval_E) \circ \gamma_{\mathbb{M}} \subseteq \gamma_{\mathbb{M}} \circ eval_E^\sharp$$

$$\wp(update_E) \circ \times \circ (\gamma_{\mathbb{M}}, \gamma_{\mathbb{A}}, \gamma_{\mathbb{V}}) \subseteq \gamma_{\mathbb{M}} \circ update_E^\sharp$$

$$\wp(fetch) \circ \times \circ (\gamma_{\mathbb{M}}, \gamma_{\mathbb{V}}) \subseteq \gamma_{\mathbb{V}} \circ fetch_E^\sharp$$

$$\wp(filter_E) \circ \gamma_{\mathbb{M}} \subseteq \gamma_{\mathbb{M}} \circ filter_E^\sharp$$

The ThriLLVM Language

- Extension 2: Functions

$E ::= \dots$	expression, as before
f	function name
$C ::= \dots$	statement, as before
$E(E)$	function call
return	return from call
$F ::= f(x) = C^+$	
$P ::= F^+ C^+$	program



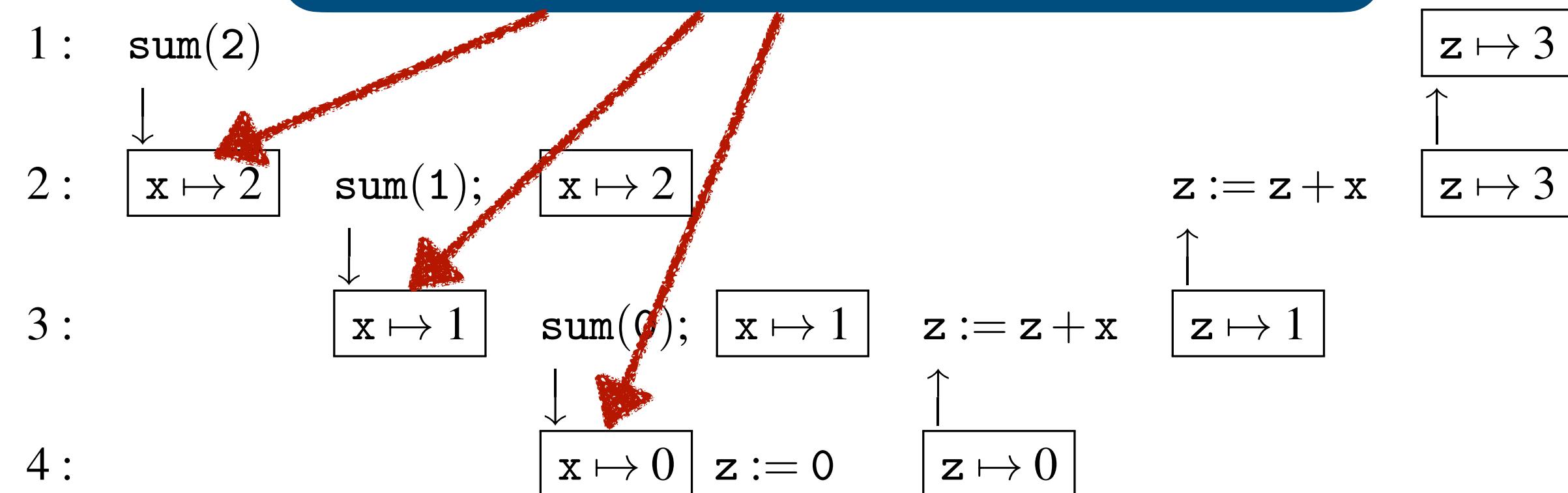
Assume:

1. Only flat function definitions (no nested function definitions)
2. Variables other than parameters are all global
3. Function names are first-class values
4. Every function has one parameter with a unique name

Example

```
void sum(int x) {  
    if (x == 0) {  
        z = 0;  
        return;  
    } else {  
        sum(x - 1);  
        z = z + x;  
        return;  
    }  
}  
  
sum(2);
```

Multiple instances of parameters
(What is the current instance of x?)



What components are needed to support functions?

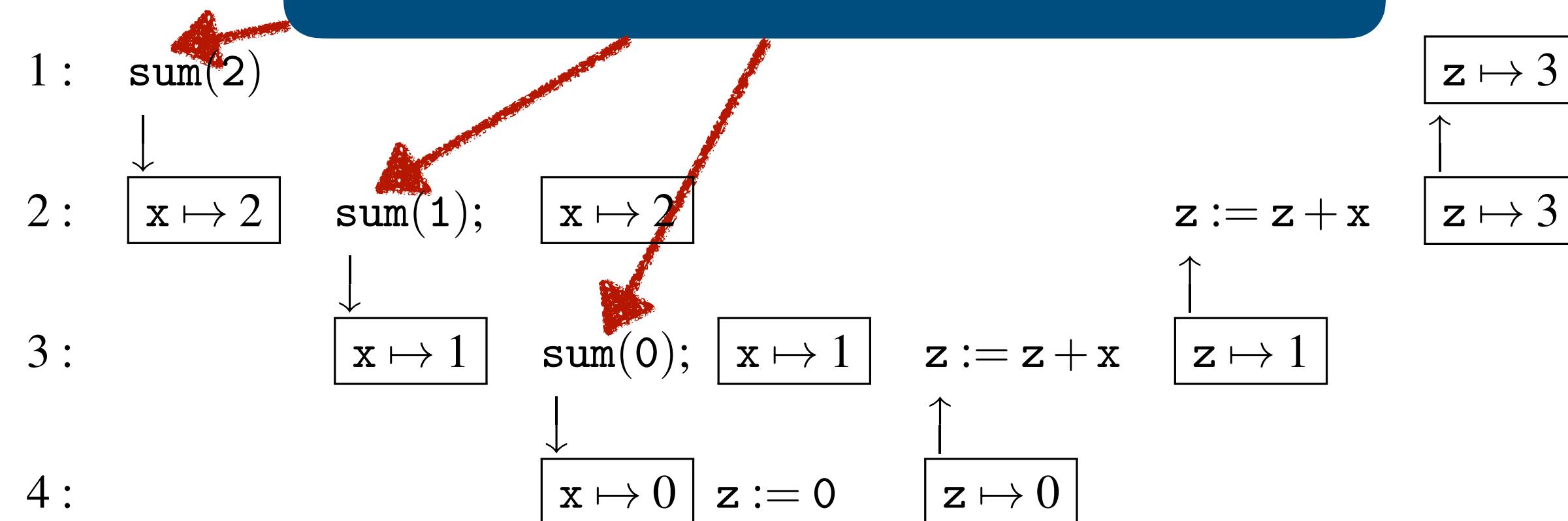
1. **Environment:** determine the current instance of the parameter

Example

```
void sum(int x) {  
    if (x == 0) {  
        z = 0;  
        return;  
    } else {  
        sum(x - 1);  
        z = z + x;  
        return;  
    }  
}
```

```
sum(2);
```

Multiple calls to the same function
(Where to return after function calls?)



What components are needed to support functions?

- 1. Environment:** determine the current instance of the parameter
- 2. Continuation:** remember the return context for each function call

Concrete Domains

$\langle l, m, \sigma, \kappa, \phi \rangle \in$	S	$=$	$L \times M \times E \times K \times I$	
$m \in$	M	$=$	$A \rightarrow V$	memories
$\sigma \in$	E	$=$	$X \rightarrow I$	environments
$\kappa \in$	K	$=$	$(L \times E)^*$	continuations
$\phi \in$	I			instances
$a \in$	A	$=$	$X \times I$	addresses
$v \in$	V	$=$	$Z \cup F$	values
	X			set of variables and parameters
	L			set of statement labels
	F			set of function names

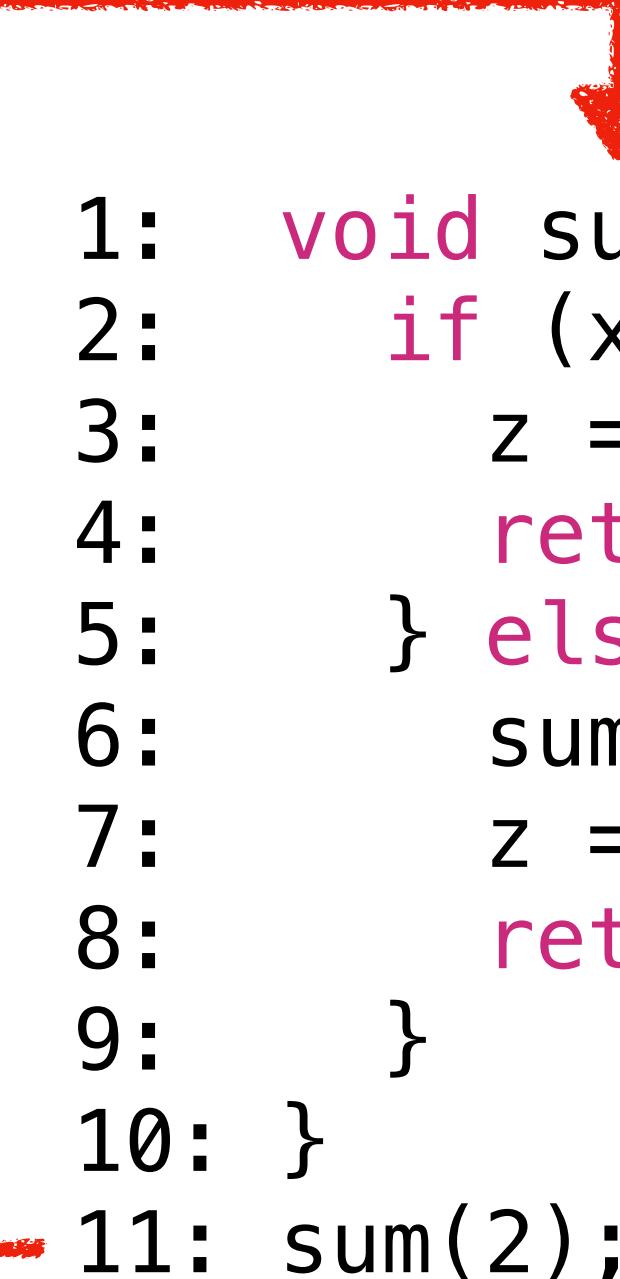
Example

```
1: void sum(int x) {  
2:     if (x == 0) {  
3:         z = 0;  
4:         return;  
5:     } else {  
6:         sum (x - 1);  
7:         z = z + x;  
8:         return;  
9:     }  
10: }  
11: sum(2);
```

Memory		Environment		Continuation	Instance
Addr	Val	Var	Instance		
		z	0		0

Example

```
1: void sum(int x) {  
2:     if (x == 0) {  
3:         z = 0;  
4:         return;  
5:     } else {  
6:         sum (x - 1);  
7:         z = z + x;  
8:         return;  
9:     }  
10: }  
11: sum(2);
```



Memory		Environment		Continuation	Instance
Addr	Val	Var	Instance		
<x, 1>	2	z	0		1
		x	1		
				11	

Example

```
1: void sum(int x) {  
2:     if (x == 0) {  
3:         z = 0;  
4:         return;  
5:     } else {  
6:         sum (x - 1); ←  
7:         z = z + x;  
8:         return;  
9:     }  
10: }  
11: sum(2);
```

Memory		Environment		Continuation	Instance
Addr	Val	Var	Instance		
<x, 1>	2	z	0		1
		x	1		
				11	

Example

```
1: void sum(int x) {  
2:     if (x == 0) {  
3:         z = 0;  
4:         return;  
5:     } else {  
6:         sum (x - 1);  
7:         z = z + x;  
8:         return;  
9:     }  
10: }  
11: sum(2);
```

Memory		Environment		Continuation	Instance
Addr	Val	Var	Instance		
<x, 1>	2	z	0		2
<x, 2>	1	x	2	6 11	

Example

```
1: void sum(int x) {  
2:     if (x == 0) {  
3:         z = 0;  
4:         return;  
5:     } else {  
6:         sum (x - 1); ←  
7:         z = z + x;  
8:         return;  
9:     }  
10: }  
11: sum(2);
```

Memory		Environment		Continuation	Instance
Addr	Val	Var	Instance		
<x, 1>	2	z	0		
<x, 2>	1	x	1		
				6 11	2

Example

```
1: void sum(int x) {  
2:     if (x == 0) {  
3:         z = 0;  
4:         return;  
5:     } else {  
6:         sum (x - 1);  
7:         z = z + x;  
8:         return;  
9:     }  
10: }  
11: sum(2);
```

Memory		Environment		Continuation	Instance
Addr	Val	Var	Instance		
<x, 1>	2	z	0		
<x, 2>	1	x	3		
<x, 3>	0			6 6 11	3

Example

```
1: void sum(int x) {  
2:     if (x == 0) {  
3:         z = 0; ←  
4:         return;  
5:     } else {  
6:         sum (x - 1);  
7:         z = z + x;  
8:         return;  
9:     }  
10: }  
11: sum(2);
```

Memory		Environment		Continuation	Instance
Addr	Val	Var	Instance		
<x, 1>	2	z	0		
<x, 2>	1	x	3		
<x, 3>	0				
<z, 0>	0				

Example

```
1: void sum(int x) {  
2:     if (x == 0) {  
3:         z = 0;  
4:         return; -----  
5:     } else {  
6:         sum (x - 1);  
7:         z = z + x; -----  
8:         return;  
9:     }  
10: }  
11: sum(2);
```

Memory		Environment		Continuation	Instance
Addr	Val	Var	Instance		
<x, 1>	2	z	0	6	
<x, 2>	1	x	2	11	
<x, 3>	0				
<z, 0>	1				

Example

```
1: void sum(int x) {  
2:     if (x == 0) {  
3:         z = 0;  
4:         return;  
5:     } else {  
6:         sum (x - 1);  
7:         z = z + x;   
8:         return;   
9:     }  
10: }  
11: sum(2);
```

Memory		Environment		Continuation	Instance
Addr	Val	Var	Instance		
<x, 1>	2	z	0		
<x, 2>	1	x	1		
<x, 3>	0				
<z, 0>	3			11	3

Example

```
1: void sum(int x) {  
2:     if (x == 0) {  
3:         z = 0;  
4:         return;  
5:     } else {  
6:         sum (x - 1);  
7:         z = z + x;  
8:         return;  
9:     }  
10: }  
11: sum(2);
```

Memory		Environment		Continuation	Instance
Addr	Val	Var	Instance		
<x, 1>	2				
<x, 2>	1				
<x, 3>	0				
<z, 0>	3				3

State Transitions

- The transition for function calls:

$$\begin{aligned} E_0(E_1) & : \langle l, m, \sigma, \kappa, \phi \rangle \\ & \hookrightarrow \langle \text{body}(\mathbf{f}), \\ & \quad \text{bind}_x(m, \phi', v), \\ & \quad \text{new-env}_x(\sigma, \phi'), \\ & \quad \text{push-context}(\kappa, l, \sigma), \\ & \quad \phi' \rangle \end{aligned}$$

where $\mathbf{f} = \text{eval}_{E_0}(m, \sigma)$, $\mathbf{x} = \text{param}(\mathbf{f})$, $v = \text{eval}_{E_1}(m, \sigma)$, $\phi' = \text{tick}(\phi)$

- The transition for returns:

$$\text{return} : \langle l, m, \sigma, \kappa, \phi \rangle \hookrightarrow \langle \text{next}(l'), m, \sigma', \kappa', \phi \rangle$$

where $\langle l', \sigma', \kappa' \rangle = \text{pop-context}(\kappa)$

Concrete Semantic Operators

- The first label of the function body: $body : \mathbb{F} \rightarrow \mathbb{L}$
- Semantic operators for function calls:

$$bind_x : \mathbb{M} \times \mathbb{I} \times \mathbb{V} \rightarrow \mathbb{M}$$
$$bind_x(m, \phi, v) = m[\langle x, \phi \rangle \mapsto v]$$

$$new-env_x : \mathbb{E} \times \mathbb{I} \rightarrow \mathbb{E}$$
$$new-env_x(\sigma, \phi) = \sigma[x \mapsto \phi]$$

$$push-context : \mathbb{K} \times \mathbb{L} \times \mathbb{E} \rightarrow \mathbb{K}$$
$$push-context(\kappa, l, \sigma) = \langle l, \sigma \rangle \cdot \kappa$$

$$tick : \mathbb{I} \rightarrow \mathbb{I}$$
$$tick(\phi) = \phi' \quad (\text{new } \phi')$$

- Semantic operators for returns:

$$pop-context : \mathbb{K} \rightarrow \mathbb{L} \times \mathbb{E} \times \mathbb{K}$$
$$pop-context(\langle l, \sigma \rangle \cdot \kappa) = \langle l, \sigma, \kappa \rangle$$

Concrete Semantics

- Concrete semantic function:

$$F : \wp(\mathbb{S}) \rightarrow \wp(\mathbb{S})$$

$$F(X) = I \cup Step(X)$$

where I is the set of initial states and $Step$ is the powerset-lifted version of \hookrightarrow

$$Step : \wp(\mathbb{S}) \rightarrow \wp(\mathbb{S})$$

$$Step(X) = \{s' \mid s \hookrightarrow s', s \in X\}$$

- Concrete semantic (i.e., reachable states): $\text{lfp } F$

Abstract Domains

- The abstract domain $\mathbb{D}^\sharp = \mathbb{L} \rightarrow \mathbb{M}^\sharp \times \mathbb{E}^\sharp \times \mathbb{K}^\sharp \times \mathbb{I}^\sharp$
- Element-wise Galois connection

$$\wp(\mathbb{L} \times \mathbb{M} \times \mathbb{E} \times \mathbb{K} \times \mathbb{I}) \xrightleftharpoons[\alpha]{\gamma} \mathbb{L} \rightarrow \mathbb{M}^\sharp \times \mathbb{E}^\sharp \times \mathbb{K}^\sharp \times \mathbb{I}^\sharp$$

where

$$\wp(\mathbb{M}) \xrightleftharpoons[\alpha_{\mathbb{M}}]{\gamma_{\mathbb{M}}} \mathbb{M}^\sharp \qquad \wp(\mathbb{E}) \xrightleftharpoons[\alpha_{\mathbb{E}}]{\gamma_{\mathbb{E}}} \mathbb{E}^\sharp$$

$$\wp(\mathbb{K}) \xrightleftharpoons[\alpha_{\mathbb{K}}]{\gamma_{\mathbb{K}}} \mathbb{K}^\sharp \qquad \wp(\mathbb{I}) \xrightleftharpoons[\alpha_{\mathbb{I}}]{\gamma_{\mathbb{I}}} \mathbb{I}^\sharp$$

Abstract Domains (Cont'd)

- An example of a memory abstract domain:

$$M^\# = (\mathbb{X} \times \mathbb{I}^\#) \rightarrow V^\# \quad \text{where} \quad \wp(M) \xrightleftharpoons[\alpha_M]{\gamma_M} M^\# \quad \text{and} \quad M = (\mathbb{X} \times \mathbb{I}) \rightarrow V$$

- An example of an environment abstract domain:

$$E^\# = \mathbb{X} \rightarrow \mathbb{I}^\# \quad \text{where} \quad \wp(E) \xrightleftharpoons[\alpha_E]{\gamma_E} E^\# \quad \text{and} \quad E = \mathbb{X} \rightarrow \mathbb{I}$$

- An example of a value abstract domain:

$$V^\# = Z^\# \times F^\# \quad \text{where} \quad \wp(V) \xrightleftharpoons[\alpha_V]{\gamma_V} V^\# \quad \wp(Z) \xrightleftharpoons[\alpha_Z]{\gamma_Z} Z^\# \quad \wp(F) \xrightleftharpoons[\alpha_F]{\gamma_F} F^\#$$

- An example of a location abstract domain:

$$L^\# \quad \text{where} \quad \wp(L) \xrightleftharpoons[\alpha_L]{\gamma_L} L^\#$$

Abstractions of instances and
continuations are parameters

Abstract Semantics

- The abstract semantic function:

$$F^\sharp : (\mathbb{L} \rightarrow \mathbb{M}^\sharp \times \mathbb{E}^\sharp \times \mathbb{K}^\sharp \times \mathbb{I}^\sharp) \rightarrow (\mathbb{L} \rightarrow \mathbb{M}^\sharp \times \mathbb{E}^\sharp \times \mathbb{K}^\sharp \times \mathbb{I}^\sharp)$$

$$F^\sharp(X) = \alpha(I) \sqcup Step^\sharp(X)$$

- The abstract semantic functions for transition and partition:

$$Step^\sharp : (\mathbb{L} \rightarrow \mathbb{M}^\sharp \times \mathbb{E}^\sharp \times \mathbb{K}^\sharp \times \mathbb{I}^\sharp) \rightarrow (\mathbb{L} \rightarrow \mathbb{M}^\sharp \times \mathbb{E}^\sharp \times \mathbb{K}^\sharp \times \mathbb{I}^\sharp)$$

$$Step^\sharp(X) = \wp(id, \sqcup) \circ \pi \circ \check{\wp}(\hookrightarrow^\sharp)(X)$$

$$\pi : \wp(\mathbb{S}^\sharp) \rightarrow (\mathbb{L} \rightarrow \wp(\mathbb{M}^\sharp \times \mathbb{E}^\sharp \times \mathbb{K}^\sharp \times \mathbb{I}^\sharp))$$

$$\pi(X) = \lambda l. \{ \langle m^\sharp, \sigma^\sharp, \kappa^\sharp, \phi^\sharp \rangle \mid \langle l, m^\sharp, \sigma^\sharp, \kappa^\sharp, \phi^\sharp \rangle \in X \}$$

- Soundness: $\text{lfp } F \subseteq \gamma(\bigsqcup_{i \geq 0} F^\sharp(\perp))$

Abstract State Transition

- The transition for function calls:

$$E_0(E_1) : \langle l, m^\#, \sigma^\#, \kappa^\#, \phi^\# \rangle \xrightarrow{\#} \langle body(f),$$

How to abstract different machine states at each call?

How to abstract different continuations at each call?

$$\begin{aligned} & bind_x^\#(m^\#, \phi'^\#, v^\#), \\ & new-env_x^\#(\sigma^\#, \phi'^\#), \\ & push-context^\#(\kappa^\#, l, \sigma^\#), \\ & \phi'^\# \rangle \end{aligned}$$

where $f \in eval_{E_0}^\#(m^\#, \sigma^\#)$, $x = param(f)$, $v^\# = eval_{E_1}^\#(m^\#, \sigma^\#)$, $\phi'^\# = tick^\#(\phi^\#)$

- The transition for returns:

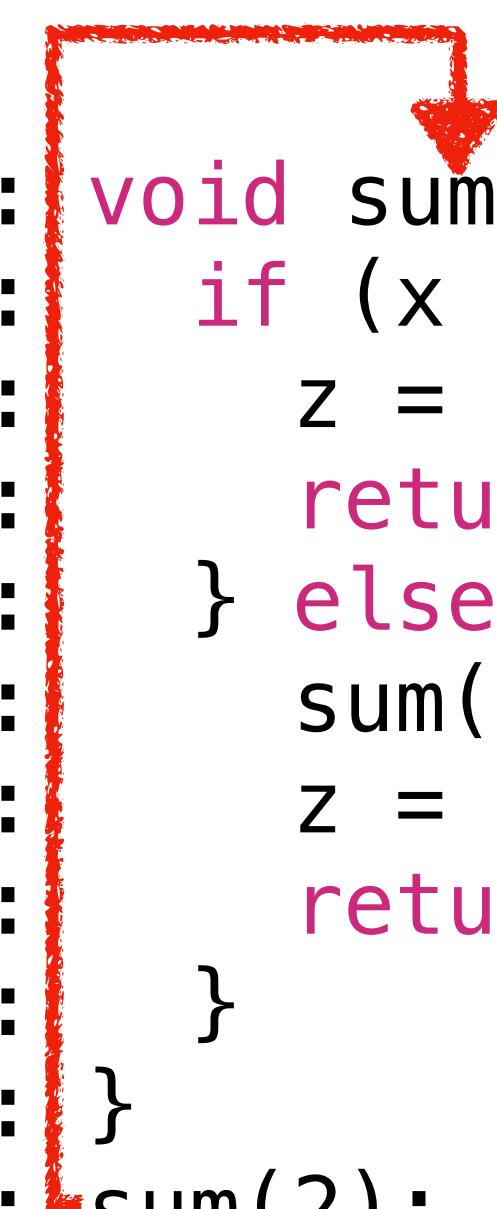
$$\text{return} : \langle l, m^\#, \sigma^\#, \kappa^\#, \phi^\# \rangle \xrightarrow{\#} \langle \text{next}(l'), m^\#, \sigma'^\#, \kappa'^\#, \phi^\# \rangle$$

where $\langle l'^\#, \sigma'^\#, \kappa'^\# \rangle = pop-context^\#(\kappa)$ and $l' \in l^\#$

How to abstract different continuations at each call?

Context-sensitivity

- How to abstract different instances of machine states at each call?
- How to abstract different continuations at each call?



Memory			Continuation
Label	x	z	
1			
3			
4			
6			
7			
8			
12			

A large red question mark is positioned below the continuation table.

```
1: void sum(int x) {
2:     if (x == 0) {
3:         z = 0;
4:         return;
5:     } else {
6:         sum(x - 1);
7:         z = z + x;
8:         return;
9:     }
10: }
11: sum(2);
```

Abstract Domains for Context-insensitive Analysis

- Context-insensitive: abstract all possible contexts into a single abstract one
- An example of an abstract instance domain:

$$\wp(\mathbb{I}) \xrightleftharpoons[\alpha_{\mathbb{I}}]{\gamma_{\mathbb{I}}} \mathbb{I}^{\sharp} \quad \text{where} \quad \mathbb{I}^{\sharp} = \{\cdot\}$$

Aggregate all instances from different contexts into a single abstract element

- An example of an abstract continuation domain:

$$\wp(\mathbb{K}) \xrightleftharpoons[\alpha_{\mathbb{K}}]{\gamma_{\mathbb{K}}} \mathbb{K}^{\sharp} \quad \text{where} \quad \mathbb{K} = (\mathbb{L} \times \mathbb{E})^* \quad \text{and} \quad \mathbb{K}^{\sharp} = \{\cdot\}$$

Aggregate all continuations from different contexts into a single abstract element

Context-insensitive Analysis

```
1: void sum(int x) {  
2:     if (x == 0) {  
3:         z = 0;  
4:         return;  
5:     } else {  
6:         sum(x - 1);  
7:         z = z + x;  
8:         return;  
9:     }  
10: }  
11: sum(2);
```

Label	x	z	Continuation
1	[2,2]		
3			•
4			
6			
7			
8			
12			

Context-insensitive Analysis

```
1: void sum(int x) {  
2:   if (x == 0) {  
3:     z = 0;  
4:     return;  
5:   } else {  
6:     sum(x - 1); ←  
7:     z = z + x;  
8:     return;  
9:   }  
10: }  
11: sum(2);
```

Label	x	z	Continuation
1	[2.2]		
3			
4			
6	[2.2]		•
7			
8			
12			

Context-insensitive Analysis

```
1: void sum(int x) {  
2:     if (x == 0) {  
3:         z = 0;  
4:         return;  
5:     } else {  
6:         sum(x - 1);  
7:         z = z + x;  
8:         return;  
9:     }  
10: }  
11: sum(2);
```

Label	x	z	Continuation
1	[1,2]		
3			•
4			
6	[2,2]		
7			
8			
12			

Context-insensitive Analysis

```
1: void sum(int x) {  
2:   if (x == 0) {  
3:     z = 0;  
4:     return;  
5:   } else {  
6:     sum(x - 1); ←  
7:     z = z + x;  
8:     return;  
9:   }  
10: }  
11: sum(2);
```

Label	x	z	Continuation
1	[1,2]		
3			•
4			
6	[1,2]		
7			
8			
12			

Context-insensitive Analysis

```
1: void sum(int x) {  
2:     if (x == 0) {  
3:         z = 0;  
4:         return;  
5:     } else {  
6:         sum(x - 1);  
7:         z = z + x;  
8:         return;  
9:     }  
10: }  
11: sum(2);
```

Label	x	z	Continuation
1	[0,2]		
3			•
4			
6	[1,2]		
7			
8			
12			

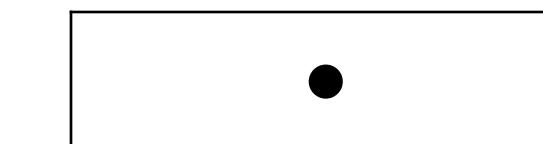
Context-insensitive Analysis

```
1: void sum(int x) {  
2:   if (x == 0) {  
3:     z = 0; ← [red box]  
4:     return;  
5:   } else {  
6:     sum(x - 1);  
7:     z = z + x;  
8:     return;  
9:   }  
10: }  
11: sum(2);
```

Memory

Label	x	z
1	[0,2]	
3	[0,2]	
4	[0,2]	[0,0]
6	[1,2]	
7		
8		
12		

Continuation



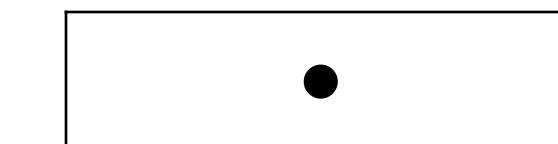
Context-insensitive Analysis

```
1: void sum(int x) {  
2:     if (x == 0) {  
3:         z = 0;  
4:         return;   
5:     } else {  
6:         sum(x - 1);  
7:         z = z + x;   
8:         return;  
9:     }  
10:  
11: sum(2); 
```

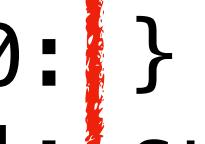
Memory

Label	x	z
1	[0,2]	
3	[0,2]	
4	[0,2]	[0,0]
6	[1,2]	
7	[0,2]	[0,0]
8	[0,2]	[0,2]
12	[0,0]	[0,0]

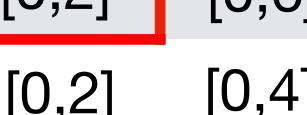
Continuation



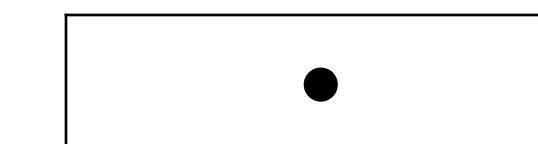
Context-insensitive Analysis

```
1: void sum(int x) {  
2:     if (x == 0) {  
3:         z = 0;  
4:         return;  
5:     } else {  
6:         sum(x - 1);  
7:         z = z + x;   
8:         return;   
9:     }   
10: }   
11: sum(2); 
```

Memory

Label	x	z
1	[0,2]	
3	[0,2]	
4	[0,2]	[0,0]
6	[1,2]	
7	[0,2]	[0,4] 
8	[0,2]	[0,6]
12	[0,2]	[0,4] 

Continuation

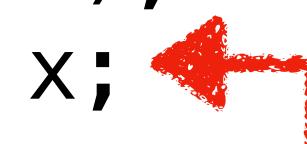
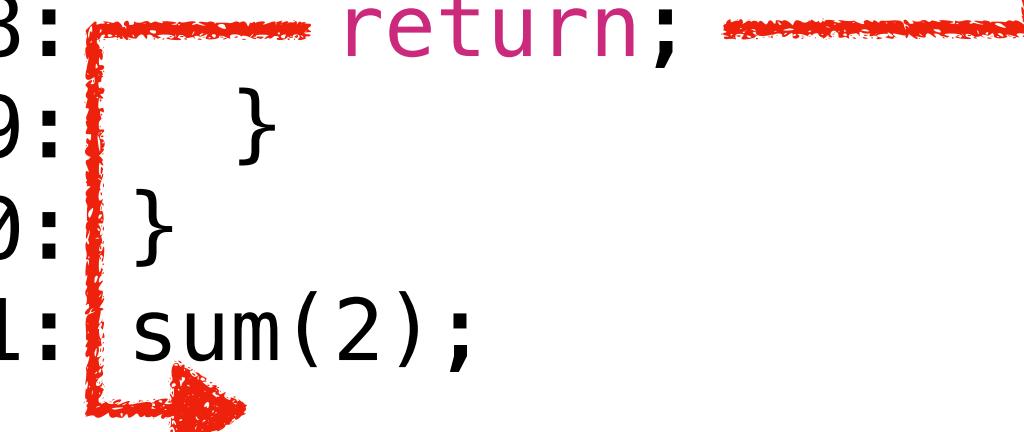
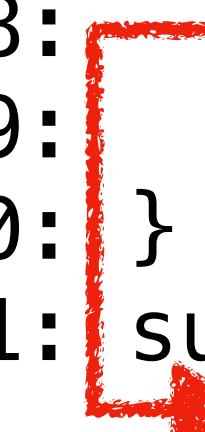


Context-insensitive Analysis

```
1: void sum(int x) {  
2:     if (x == 0) {  
3:         z = 0;  
4:         return;  
5:     } else {  
6:         sum(x - 1);  
7:         z = z + x;   
8:         return;   
9:     }  
10:  
11: sum(2); 
```

Label	x	z	Continuation
1	[0,2]		
3	[0,2]		
4	[0,2]	[0,0]	
6	[1,2]		
7	[0,2]	[0,6] 	
8	[0,2]	[0,8]	
12	[0,2]	[0,6] 	

Context-insensitive Analysis

```
1: void sum(int x) {  
2:   if (x == 0) {  
3:     z = 0;  
4:     return;  
5:   } else {  
6:     sum(x - 1);  
7:     z = z + x;   
8:   return;   
9: }   
10:  
11: sum(2); 
```

Label	x	z	Continuation
1	[0,2]		
3	[0,2]		
4	[0,2]	[0,0]	
6	[1,2]		
7	[0,2]	[0,+\infty]	
8	[0,2]	[0,+\infty]	
12	[0,2]	[0,+\infty]	

Fixed point!

Abstract Semantics for Context-insensitive Analysis

- The abstract semantic operators for function calls:

$$\begin{aligned} bind_x^\# : \mathbb{M}^\# \times \mathbb{I}^\# \times \mathbb{V}^\# &\rightarrow \mathbb{M}^\# \\ bind_x^\#(m^\#, \phi^\#, v^\#) &= m^\# \{ \langle x, \bullet \rangle \mapsto v^\# \} \end{aligned}$$

$$\begin{aligned} new-env_x^\# : \mathbb{E}^\# \times \mathbb{I}^\# &\rightarrow \mathbb{E}^\# \\ new-env_x^\#(\sigma^\#, \phi^\#) &= \sigma^\# \{ x \mapsto \bullet \} \end{aligned}$$

$$\begin{aligned} push-context^\# : \mathbb{K}^\# \times \mathbb{L} \times \mathbb{E}^\# &\rightarrow \mathbb{K}^\# \\ push-context^\#(\kappa^\#, l, \sigma^\#) &= \bullet \end{aligned}$$

Abstract all continuations

$$\begin{aligned} tick^\# : \mathbb{I}^\# &\rightarrow \mathbb{I}^\# \\ tick^\#(\phi^\#) &= \bullet \end{aligned}$$

Abstract all instances

- The abstract semantic operator for returns:

$$pop-context^\# : \mathbb{K}^\# \rightarrow \mathbb{L}^\# \times \mathbb{E}^\# \times \mathbb{K}^\#$$

$$pop-context^\#(\kappa^\#) = \langle l^\#, \sigma^\#, \bullet \rangle$$

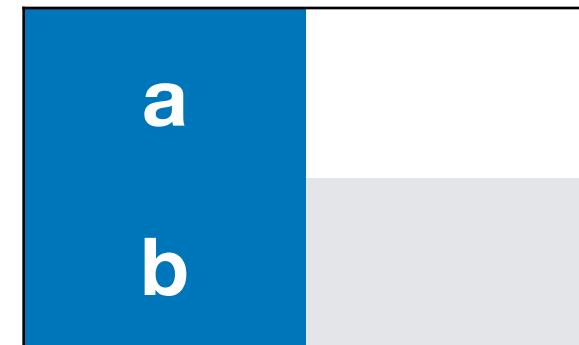
where $l^\# = \{ l \mid \langle l, -, -, -, - \rangle \hookrightarrow^\# \langle body(f), -, -, -, - \rangle \}$

and $\sigma^\# = \lambda x. \bullet$

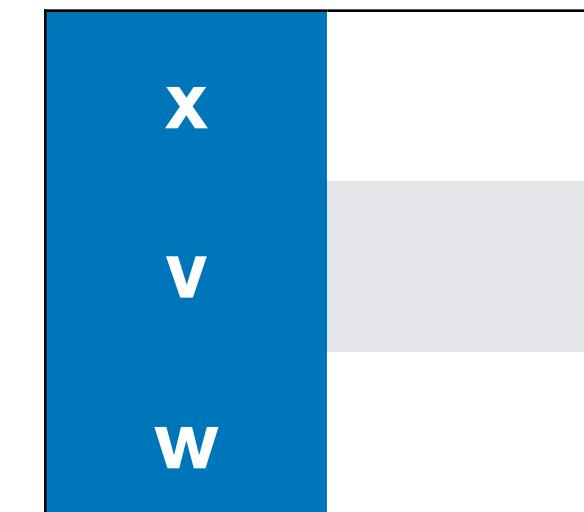
All possible call-sites for f

Example

```
void main() {  
    a = f(1);  
    b = f(3);  
}
```



```
int f(int x) {  
    v = g(x);  
    w = g(x + 1);  
    return v + w;  
}
```

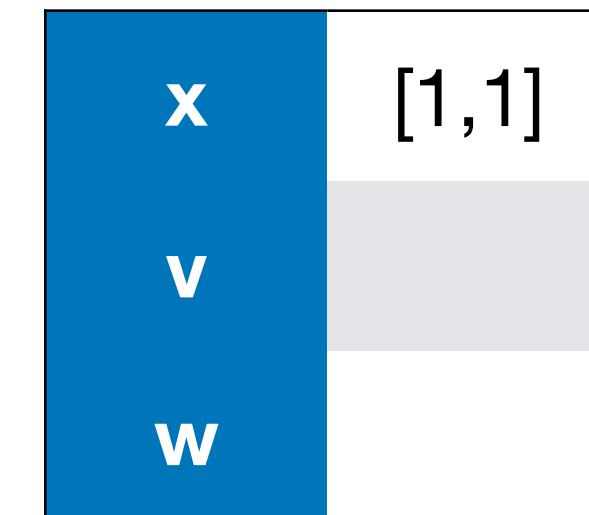
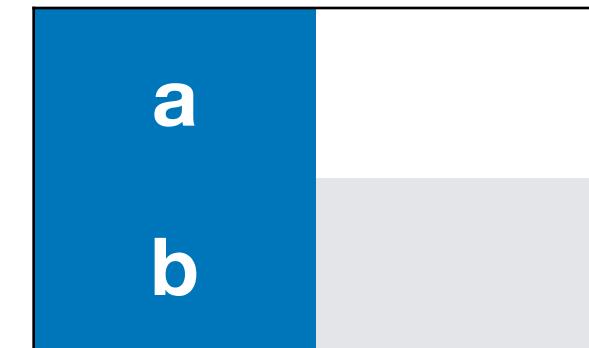


```
int g(int y) {  
    return y;  
}
```



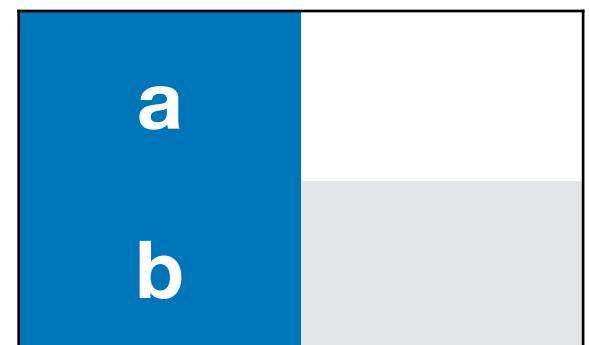
Example

```
void main() { → int f(int x) { → int g(int y) {  
    a = f(1); ← v = g(x); ← return y;  
    b = f(3); ← w = g(x + 1); }  
} } return v + w;
```

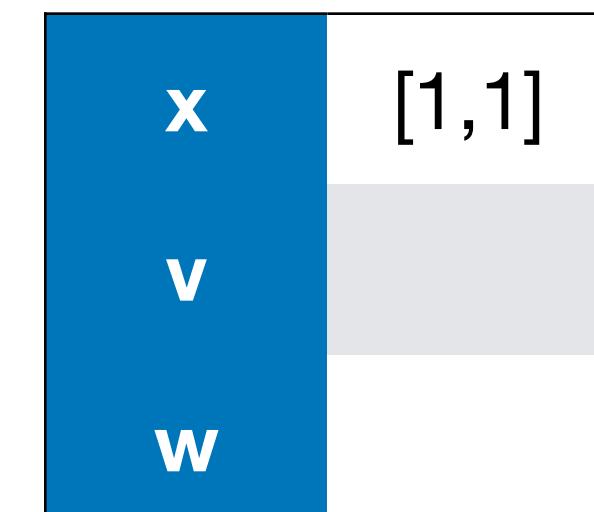


Example

```
void main() {  
    a = f(1);  
    b = f(3);  
}
```



```
int f(int x) {  
    v = g(x);  
    w = g(x + 1);  
    return v + w;  
}  
  
int g(int y) {  
    return y;  
}
```



Example

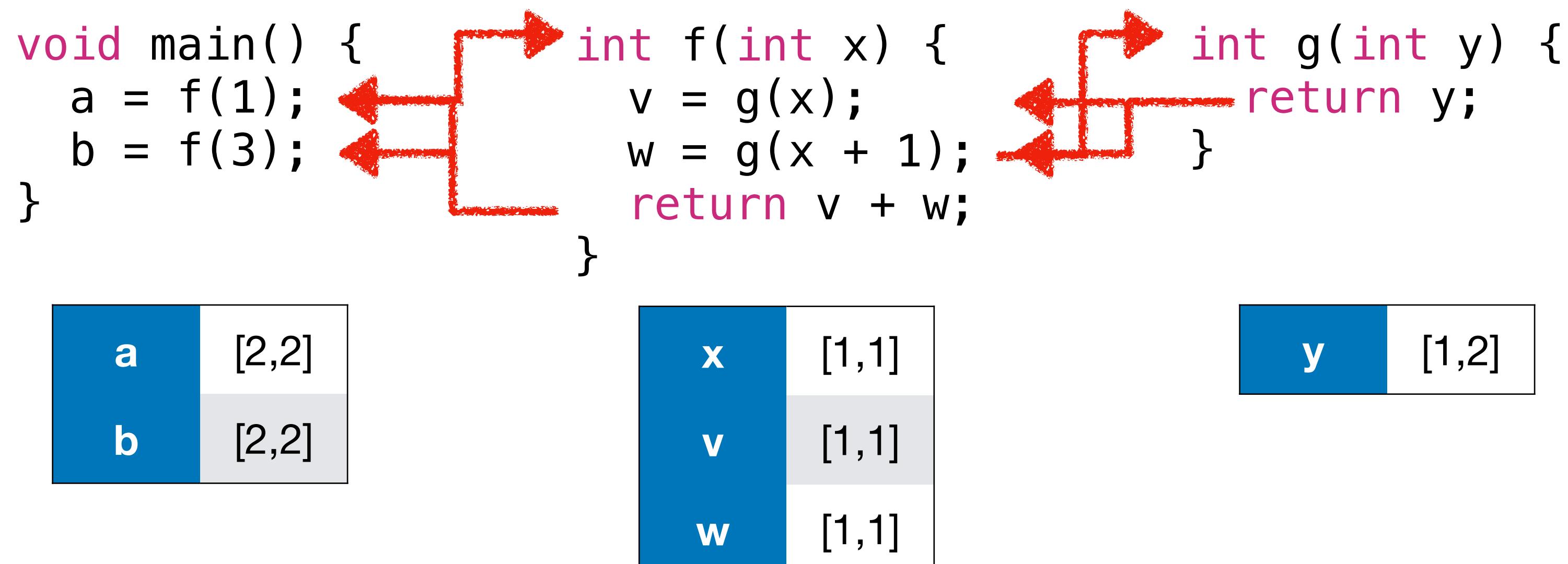
```
void main() {  
    a = f(1); ←  
    b = f(3); ←  
}  
  
int f(int x) {  
    v = g(x); →  
    w = g(x + 1); →  
    return v + w;  
}
```

a	
b	

x	[1,1]
v	[1,1]
w	[1,1]

y	[1,1]
---	-------

Example



Example

```
void main() {  
    a = f(1);  
    b = f(3);  
}  
  
int f(int x) {  
    v = g(x);  
    w = g(x + 1);  
    return v + w;  
}  
  
int g(int y) {  
    return y;  
}
```

a	[2,2]
b	[2,2]

x	[1,3]
v	[1,2]
w	[1,2]

y	[1,2]
---	-------

...

Example

```
void main() {  
    a = f(1);  
    b = f(3);  
}
```

a	[2,8]
b	[2,8]

```
int f(int x) {  
    v = g(x);  
    w = g(x + 1);  
    return v + w;  
}
```

x	[1,3]
v	[1,4]
w	[1,4]

```
int g(int y) {  
    return y;  
}
```

y	[1,4]
---	-------

Fixed point!

Context-sensitivity

- How to abstract different instances of machine states at each call?
- How to abstract different continuations at each call?

```
1: void sum(int x) {  
2:   if (x == 0) {  
3:     z = 0;  
4:     return;  
5:   } else {  
6:     sum(x - 1);  
7:     z = z + x;  
8:     return;  
9:   }  
10: }  
11: sum(2);
```

Label	Memory		Continuation
	x	z	
1	[0,2]		
3	[0,2]		
4	[0,2]	[0,0]	
6	[1,2]		
7	[0,2]	[0,4]	
8	[0,2]	[0,2]	
12	[0,2]	[0,4]	



Context-sensitive Analysis

- Distinguish different contexts (instances and continuations)
- In principle, many possible designs for context-sensitive analysis
- In this lecture, we use a well-known approach based on *call-string*
 - Call-string: an ordered sequence of call-sites

Context-sensitive Analysis

```
1: void sum(int x) {  
2:     if (x == 0) {  
3:         z = 0;  
4:         return;  
5:     } else {  
6:         sum(x - 1);  
7:         z = z + x;  
8:         return;  
9:     }  
10: }  
11: sum(2);
```

Memory

Label	x<11>	Continuation
1	[2,2]	
3		
4		
6		
7		
8		
12		11

Context-sensitive Analysis

Memory

Continuation

```
1: void sum(int x) { [ ]  
2:   if (x == 0) { [ ]  
3:     z = 0;  
4:     return;  
5:   } else { [ ]  
6:     sum(x - 1); [ ]  
7:     z = z + x;  
8:     return;  
9:   }  
10: }  
11: sum(2);
```

11

Label	X<11>
1	[2,2]
3	
4	
6	[2,2]
7	
8	
12	

Context-sensitive Analysis

```
1: void sum(int x) {  
2:     if (x == 0) {  
3:         z = 0;  
4:         return;  
5:     } else {  
6:         sum(x - 1);  
7:         z = z + x;  
8:         return;  
9:     }  
10: }  
11: sum(2);
```

Label	X<11>	X<6,11>
1	[2,2]	[1,1]
3		
4		
6	[2,2]	
7		
8		
12		

6
11

Context-sensitive Analysis

```
1: void sum(int x) {  
2:   if (x == 0) {  
3:     z = 0;  
4:     return;  
5:   } else {  
6:     sum(x - 1); ←  
7:     z = z + x;  
8:     return;  
9:   }  
10: }  
11: sum(2);
```

Memory			Continuation
Label	$x_{<11>}$	$x_{<6,11>}$	
1	[2,2]	[1,1]	
3			
4			
6	[2,2]	[1,1]	
7			
8			
12			

Context-sensitive Analysis

```
1: void sum(int x) {  
2:     if (x == 0) {  
3:         z = 0;  
4:         return;  
5:     } else {  
6:         sum(x - 1);  
7:         z = z + x;  
8:         return;  
9:     }  
10: }  
11: sum(2);
```

Memory

Label	X<11>	X<6,11>	X<6,6,11>
1	[2,2]	[1,1]	[0,0]
3			
4			
6	[2,2]	[1,1]	
7			
8			
12			

Continuation

6
6
11

Context-sensitive Analysis

```
1: void sum(int x) {  
2:   if (x == 0) {  
3:     z = 0; ← [red box]  
4:     return;  
5:   } else {  
6:     sum(x - 1);  
7:     z = z + x;  
8:     return;  
9:   }  
10: }  
11: sum(2);
```

Label	X<11>	X<6,11>	X<6,6,11>
1	[2,2]	[1,1]	[0,0]
3			[0,0]
4			[0,0]
6	[2,2]	[1,1]	
7			
8			
12			

Continuation

6
6
11

Context-sensitive Analysis

```
1: void sum(int x) {  
2:     if (x == 0) {  
3:         z = 0;  
4:         return;     
5:     } else {  
6:         sum(x - 1);  
7:         z = z + x;     
8:         return;  
9:     }  
10: }  
11: sum(2);
```

Memory					Continuation
Label	X<11>	X<6,11>	X<6,6,11>	z	
1	[2,2]	[1,1]	[0,0]		
3			[0,0]		
4			[0,0]	[0,0]	
6	[2,2]	[1,1]			
7			[0,0]	[0,0]	
8					
12					

Context-sensitive Analysis

```
1: void sum(int x) {  
2:     if (x == 0) {  
3:         z = 0;  
4:         return;  
5:     } else {  
6:         sum(x - 1);  
7:         z = z + x;   
8:         return;   
9:     }  
10: }  
11: sum(2);
```

Label	Memory				Continuation
	X<11>	X<6,11>	X<6,6,11>	z	
1	[2,2]	[1,1]	[0,0]		
3			[0,0]		
4			[0,0]	[0,0]	
6	[2,2]	[1,1]			11
7			[0,0]	[0,1]	
8					
12					

Context-sensitive Analysis

```
1: void sum(int x) {  
2:   if (x == 0) {  
3:     z = 0;  
4:     return;  
5:   } else {  
6:     sum(x - 1);  
7:     z = z + x;  
8:   } return;  
9: }  
10:} sum(2);
```

Label	Memory				Continuation
	X<11>	X<6,11>	X<6,6,11>	z	
1	[2,2]	[1,1]		[0,0]	
3				[0,0]	
4			[0,0]	[0,0]	
6	[2,2]	[1,1]			
7			[0,0]	[0,1]	
8				[0,3]	
12				[0,3]	

Problems of Full Context-sensitivity

- Precision vs Cost

Label	x	z
1	[0,2]	
3	[0,0]	
4	[0,0]	[0,0]
6	[1,2]	
7	[0,2]	[0,4]
8	[0,2]	[0,2]
12	[0,2]	[0,4]

Context-insensitive

Label	$x_{<11>}$	$x_{<6,11>}$	$x_{<6,6,11>}$	z
1	[2,2]	[1,1]	[0,0]	
3			[0,0]	
4			[0,0]	[0,0]
6	[2,2]	[1,1]		
7			[0,0]	[0,1]
8				[0,3]
12				[0,3]

Fully Context-sensitive

- May not terminate because of infinitely long call-strings

```
1: void sum(int x) {
2:   if (???) {
3:     z = 0;
4:     return;
5:   } else {
6:     sum(x - 1);
7:     z = z + x;
8:     return;
9:   }
10: }
11: sum(2);
```

Label	$x_{<11>}$	$x_{<6,11>}$	$x_{<6,6,11>}$	$x_{<6,6,6,11>}$...
1	[2,2]	[1,1]	[0,0]	[-1.-1]	
3					
4					
6	[2,2]	[1,1]	[0,0]	[-1.-1]	
7					
8					
12					

Partial Context-sensitivity

- The most common way: keep only the top-most k continuations (so-called k -CFA)
 - $k = 0$: ignore all contexts, i.e., context-insensitive
 - $k = \infty$: keep all contexts, i.e., fully context-sensitive

$$\mathbb{I}^\sharp = \bigcup_{0 \leq i \leq k} \wp(\mathbb{L}^i) \quad \mathbb{K}^\sharp = \bigcup_{0 \leq i \leq k} \wp((\mathbb{L} \times \mathbb{E}^\sharp)^i)$$

$$\begin{aligned}\alpha_{\mathbb{K}}(K) &= \{\kappa^\sharp \mid \kappa^\sharp \text{ is a prefix of } \kappa \in K \text{ and } |\kappa^\sharp| \leq k\} \\ \gamma_{\mathbb{K}}(\kappa^\sharp) &= \{\kappa \in \mathbb{K} \mid \kappa^\sharp \text{ is a prefix of } \kappa\}\end{aligned}$$

Abstract Semantics for Partially Context-sensitive Analysis

- The abstract semantic operators for function calls:

$$\begin{aligned} bind_x^\# : \mathbb{M}^\# \times \mathbb{I}^\# \times \mathbb{V}^\# &\rightarrow \mathbb{M}^\# \\ bind_x^\#(m^\#, \phi^\#, v^\#) &= \bigsqcup_{\iota \in \phi^\#} m^\#[\langle x, \iota \rangle \mapsto v^\#] \end{aligned}$$

$$\begin{aligned} new-env_x^\# : \mathbb{E}^\# \times \mathbb{I}^\# &\rightarrow \mathbb{E}^\# \\ new-env_x^\#(\sigma^\#, \phi^\#) &= \sigma^\#[x \mapsto \phi^\#] \end{aligned}$$

$$\begin{aligned} push-context^\# : \mathbb{K}^\# \times \mathbb{L}^\# \times \mathbb{E}^\# &\rightarrow \mathbb{K}^\# \\ push-context^\#(\kappa^\#, l, \sigma^\#) &= \{\lceil \langle l, \sigma^\# \rangle \cdot K^\# \rceil_k \mid K^\# \in \kappa^\#\} \end{aligned}$$

Top-k contexts

$$\begin{aligned} tick^\# : \mathbb{I}^\# &\rightarrow \mathbb{I}^\# \\ tick^\#(\phi^\#) &= \{\lceil l \cdot K^\# \rceil_k \mid K^\# \in \kappa^\#\} \end{aligned}$$

Top-k contexts

- The abstract semantic operator for returns:

$$\begin{aligned} pop-context^\# : \mathbb{K}^\# &\rightarrow \mathbb{L}^\# \times \mathbb{E}^\# \times \mathbb{K}^\# \\ pop-context^\#(\kappa^\#) &= \bigsqcup \{pop^\#(K^\#) \mid K^\# \in \kappa^\#\} \end{aligned}$$

$$\text{where } pop^\#(K^\#) = \begin{cases} \langle l^\#, \sigma^\#, K^{\#'} \rangle & \text{if } K^\# = \langle l^\#, \sigma^\# \rangle \cdot K^{\#'} \\ \langle \{l \mid \langle l, -, -, -, - \rangle \hookrightarrow^\# \langle body(f), -, -, -, - \rangle\}, \lambda x. \{e\}, e \rangle & \text{if } K^\# = \epsilon \end{cases}$$

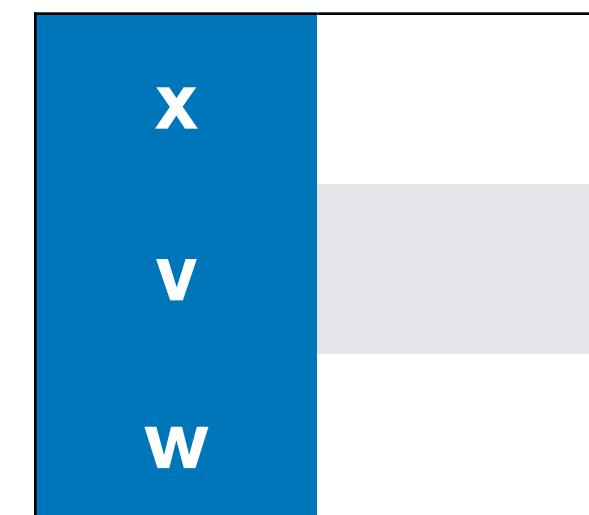
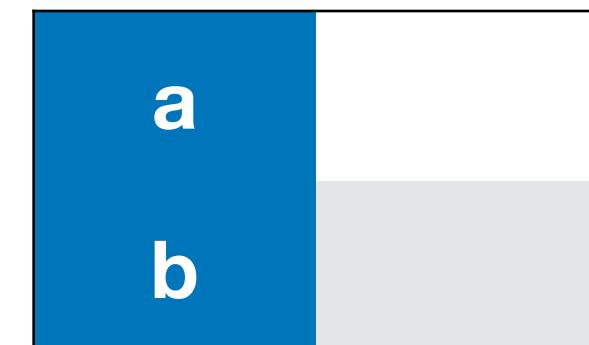
All possible call-sites for f

Example

- Suppose $k > 2$

```
void main() {  
1:   a = f(1);  
2:   b = f(3);  
}  
  
int f(int x) {  
3:   v = g(x);  
4:   w = g(x + 1);  
5:   return v + w;  
}
```

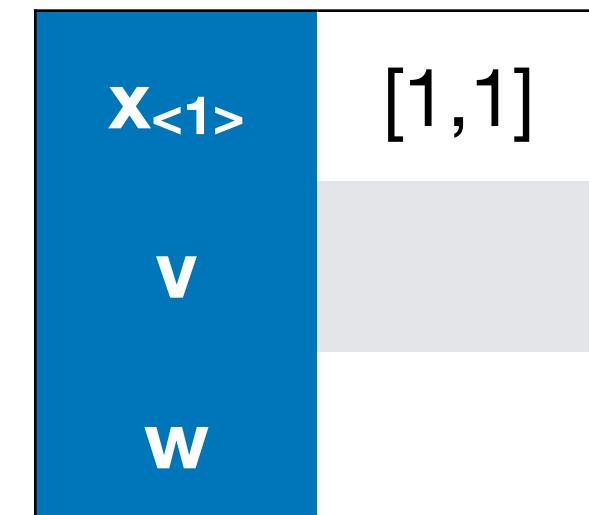
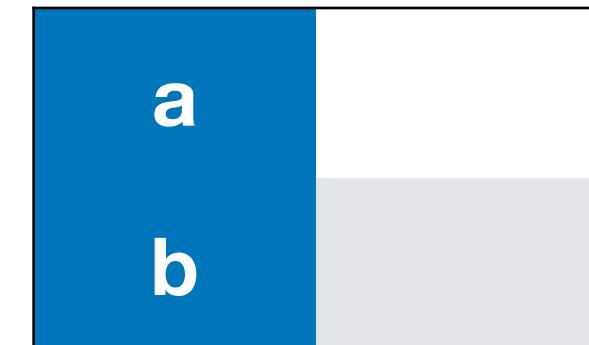
```
int g(int y) {  
6:   return y;  
}
```



Example

- Suppose $k > 2$

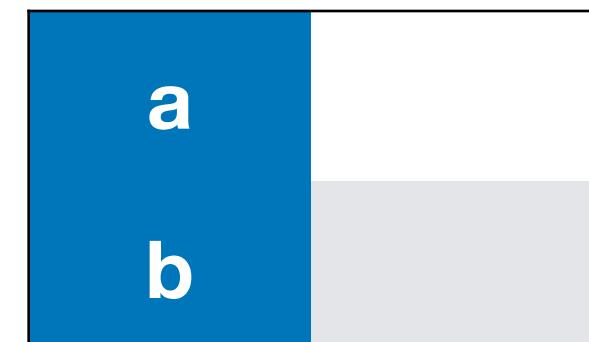
```
void main() {  
1:   a = f(1);  ↗ int f(int x) {  
2:   b = f(3);  ↗ 3:   v = g(x);  ↗ int g(int y) {  
}          }      4:   w = g(x + 1);  
                  5:   return v + w;  
                  }      6:   return y;  
}
```



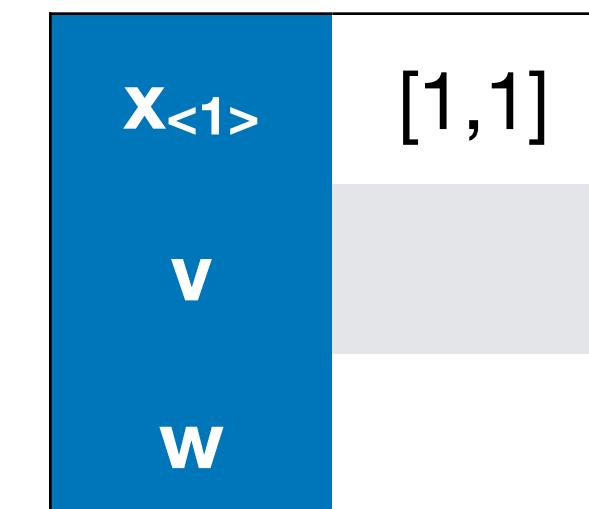
Example

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int f(int x) {  
3:   v = g(x);  
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}
```



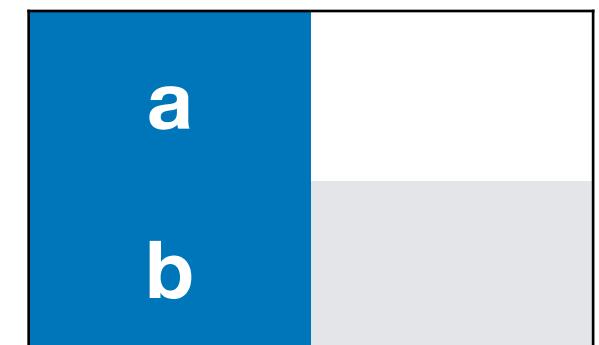
```
int g(int y) {  
6:   return y;  
}
```



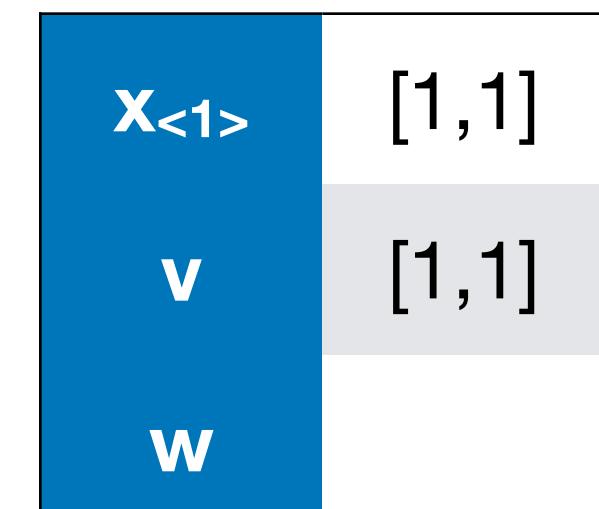
Example

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void main() {  
1:   a = f(1);  
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}
```



```
int f(int x) {  
3:   v = g(x);  
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}
```



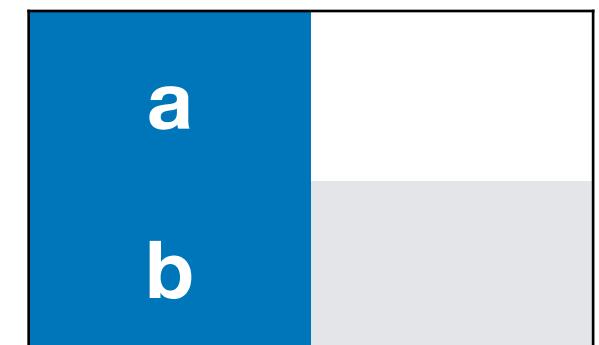
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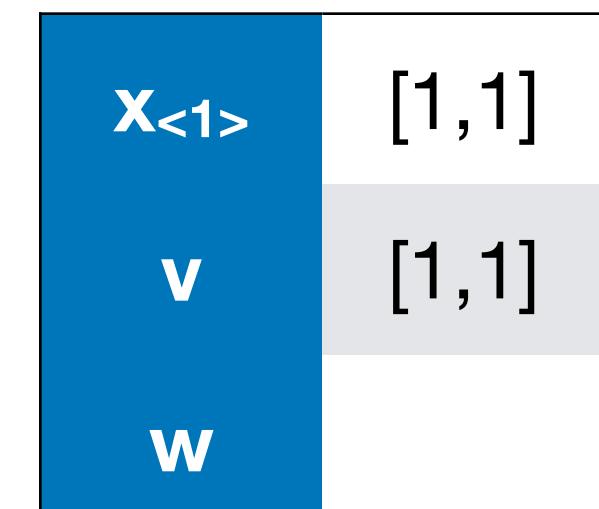
Example

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```
int f(int x) {  
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4:   w = g(x + 1);  
5:   return v + w;  
}
```



```
int g(int y) {  
6:   return y;  
}
```



Example

- Suppose $k > 2$

```
void main() {  
1:   a = f(1);  
2:   b = f(3);  
}
```

a	
b	

```
int f(int x) {  
3:   v = g(x);  
4:   w = g(x + 1);  
5:   return v + w;  
}
```

x<1>	[1,1]
v	[1,1]
w	

```
int g(int y) {  
6:   return y;  
}
```

y<3,1>	[1,1]
y<4,1>	[2,2]

Example

- Suppose $k > 2$

```
void main() {  
1:   a = f(1);  
2:   b = f(3);  
}
```

a	
b	

```
int f(int x) {  
3:   v = g(x);  
4:   w = g(x + 1);  
5:   return v + w;  
}
```

x<1>	[1,1]
v	[1,1]
w	[2,2]

```
int g(int y) {  
6:   return y;  
}
```

y<3,1>	[1,1]
y<4,1>	[2,2]

Example

- Suppose $k > 2$

```
void main() {  
1:   a = f(1); ←  
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}  
  
int f(int x) {  
3:   v = g(x);  
4:   w = g(x + 1);  
5:   return v + w;  
}  
  
int g(int y) {  
6:   return y;  
}
```

a	[3,3]
b	

x<1>	[1,1]
v	[1,1]
w	[2,2]

y<3,1>	[1,1]
y<4,1>	[2,2]

...

Example

- Suppose $k > 2$

```
void main() {  
1:   a = f(1);  
2:   b = f(3);  
}
```

a	[3,3]
b	[7,7]

```
int f(int x) {  
3:   v = g(x);  
4:   w = g(x + 1);  
5:   return v + w;  
}
```

x<1>	[1,1]
x<2>	[3,3]
v	[1,1]
w	[2,2]

```
int g(int y) {  
6:   return y;  
}
```

y<3,1>	[1,1]
y<4,1>	[2,2]
y<3,2>	[3,3]
y<3,2>	[4,4]

Fixed point!

Summary

- Handing more **advanced programming features** (pointers and functions)
 - Crucial parts for precision and scalability
- Many possible **choices of abstractions** for such features
 - E.g., allocation-site-based heap abstraction, call-string-based context abstraction, etc