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Congestion Control Using FEC for Conversational Media
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Abstract

This document describes a new mechanism for conversational multimedia flows. The proposed mechanism uses Forward Error Correction (FEC) encoded RTP packets (redundant packets) along side media packets to probe for available network capacity. A straightforward interpretation is, the sending endpoint increases the sending rate by keeping the media rate constant but increases the amount of FEC. If no losses and discards occur, the endpoint can then increase the media rate. If losses occur, the redundant FEC packets help in recovering the lost packets. Consequently, the endpoint can vary the FEC bit rate to conservatively (by a small amount) or aggressively (by a large amount) probe for available network capacity.

Status of This Memo

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1. Introduction

The Real-time Transport Protocol (RTP) [RFC3550] is widely used in voice telephony and video conferencing systems. Many of these systems run over best-effort UDP/IP networks, and are required to implement congestion to adapt the transmission rate of the multimedia streams to match the available network capacity, while maintaing the user-experience [I-D.ietf-rmcat-cc-requirements]. The circuit breakers [I-D.ietf-avtcore-rtp-circuit-breakers] describe a minimal set of conditions when an RTP media stream is causing severe congestion and should cease transmission. Consequently, the congestion control algorithm are expected to avoid triggering these conditions.

Conversational multimedia systems use Negative Acknowlegment (NACK), Forward Error Correction (FEC), and Reference Picture Selection (RPS) to protect against packet loss. These are used in addition to the codec-dependent resilience methods (for e.g., full intra-refresh and error-concealment). In this way, the multimedia system is anyway trading off part of the sending rate for redundancy or retransmissions to reduce the effects of packet loss. An endpoint often prefers using FEC in high latency networks where retransmissions may arrive later than the playout time of the packet (due to the size of the dejitter buffer) [Holmer13]. Therefore, the endpoint needs to adapt the transmission rate to best fit the changing network capacity and the amount of redundancy based on the observed/expected loss rate and network latency. Figure 1 shows the applicatbility of different error-resilience schemes based on the end-to-end latency and the observed packet loss [Devadoss08].

```
  ^
  |            .__________.
  |            |          |
  |            |  UEP/FEC |
l |____________|____.     |
a |            |    |     |
t |       RPS  |    |     |
e |_______.    |    |     |
n |       |    |    |     |
c |       |    |____|_____|
y | NACK  |         |
  |       |         |
  +-------------------------------->
         Packet loss
```

Figure 1: Applicability of Error Resilience Schemes based on the network delay and observed packet loss

In this document, we describe the use of FEC packets not only for error-resilience but also as a probing mechanism for congestion control (ramping up the transmission rate).

2. Terminology

The terminology defined in RTP [RFC3550], RTP Profile for Audio and Video Conferences with Minimal Control [RFC3551], RTCP Extended Report (XR) [RFC3611], Extended RTP Profile for RTCP-based Feedback (RTP/AVPF) [RFC4585], RTP Retransmission Payload Format [RFC4588], Forward Error Correction (FEC) Framework [RFC6363], and Support for Reduced-Size RTCP [RFC5506] apply.

3. Concept: FEC for Congestion Control

FEC is one method for providing error-resilience, it improves reliability by adding redundant data to the primary media flow, which is used by received to recover packets that have been lost due to congestion or bit-errors. The congestion control algorithm on the other hand aims at maximizing the network path utilization, but risks over-estimating the avaiable end-to-end network capacity leading to congestion (and therefore losses).

The main idea behind using FEC for congestion control is as follows: the sending endpoint chooses a high FEC rate to aggressively probe for available capacity and conversely chooses a low FEC rate to conservatively probe for available capacity. During the ramp up, if a packet is lost and the FEC packet arrives in time for decoding, the receiver is be able to recover the lost packet; if no packet is lost, the sender is able to increase the media encoding rate by swapping out a part of the FEC rate. This method can be especially useful when the sending rate is close to the bottleneck link rate: by choosing an appropriate FEC rate, the endpoint is able to probe for available capacity without changing the target media rate and therefore not affecting the user-experience. Therefore, the congestion control algorithm is always able to probe for available capacity, as improved reliability compensates for possible errors resulting from overuse (i.e., increase in observed latency and/or losses).

```
+------------+          Good conditions          +-----------+
|            |---------------------------------->|           |
|    STAY    |                                   |   PROBE   |
|            |<----------------------------------|           |
+------------+     Probed, but Loss recovered    +-----------+
     |                                              |  |  |
     |   ___________________________________________|  |  |
     |  |                                              |  |
     |  |                                              |  |
+------------+                                   +------------+
|            |          Unstable conditions      |            |
|   REDUCE   |<----------------------------------|  INCREASE  |
|            |                                   |            |
+------------+                                   +------------+
```

Figure 2: State machine of a Congestion Control enabling FEC.

3.1. State Machine

Figure 2 illustrates the the state machine of a congestion control algorithm incorporating FEC for probing. The state-machine includes 4 states: STAY, PROBE, INCREASE, REDUCE.

3.2. FEC Scheme

[RFC6363] describes a framework for using Forward Error Correction (FEC) codes with RTP and allows any FEC code to be used with the framework. For this proposal, the FEC packets are created by XORing

RTP media packets, the resulting redundant RTP packets are encoded using the scheme defined in [RFC5109].

The endpoint MAY use a single-frame FEC or a multi-frame FEC for protecting the media stream. A single-frame FEC protects against a single packet loss and fails when burst loss occurs. Using multi-frame FEC helps mitigate these issues at the cost of higher overhead and latency in recovering lost packets. [Holmer13] shows examples of using a single- and multi-frame FEC.

3.3. Applicability to RMCAT

[Open issue: Apply to Google’s Congestion Control algorithm [I-D.alvestrand-rmcat-congestion] and/or other proposals (e.g., NADA).]

4. Security Considerations

TBD

5. IANA Considerations

There are no IANA impacts in this memo.

6. Acknowledgements

This document is based on the results published in [Nagy14], it has benefited from several rounds of peer-review and hallway discussions.

7. References

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Appendix A. Simulations

This document is based on the results published in [Nagy14]. See the paper for ns-2 and testbed results; more results based on the scenarios listed in [I-D.sarker-rmcat-eval-test] will be published shorty.

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