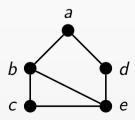
# Paths and cycles

#### Peter Rowlett

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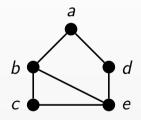
#### **Paths**

- ➤ Some sources use more specific definitions, but we'll say a *path* is a sequence of edges which join a sequence of vertices.
- ► A path's *length* is the number of edges in it.
- ▶ For example, there is a path of length 3 from a to e in the graph below, it goes  $a \rightarrow b \rightarrow c \rightarrow e$ .



### Cycles

- ► A cycle is a path that ends at its starting vertex.
- ▶ For example, there is a cycle of length 4 from a in the graph below, it goes  $a \rightarrow b \rightarrow e \rightarrow d \rightarrow a$ .



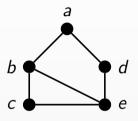
# Adjacency matrices

- ▶ Diagrams of dots and lines are not the only way to represent graphs. One way that can be useful is to represent the graph as a matrix.
- ▶ The *adjacency matrix* for a graph, G = (V, E), with n vertices is an  $n \times n$  matrix M with entries

$$\mathbf{M} = egin{bmatrix} m_{11} & m_{12} & \dots & m_{1n} \ m_{21} & m_{22} & \dots & m_{2n} \ dots & dots & \ddots & dots \ m_{n1} & m_{n2} & \dots & m_{nn} \end{bmatrix}$$

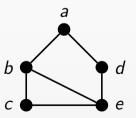
$$m_{ij} = \begin{cases} 1 & \text{if } (i,j) \in E; \\ 0 & \text{otherwise.} \end{cases}$$

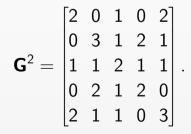
▶ The following are representations of the same graph.



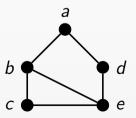
$$\mathbf{G} = egin{bmatrix} 0 & 1 & 0 & 1 & 0 \ 1 & 0 & 1 & 0 & 1 \ 0 & 1 & 0 & 0 & 1 \ 1 & 0 & 0 & 0 & 1 \ 0 & 1 & 1 & 1 & 0 \ \end{pmatrix}$$

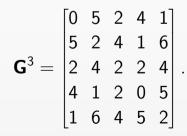
- ▶ We can perform matrix operations on **G**.
- For example, entries of  $\mathbf{G}^k$  count the number of paths of length k between vertices.
- ▶ i.e. the (i, j)th entry of  $\mathbf{G}^k$  gives the number of paths of length k from vertex i to vertex j.





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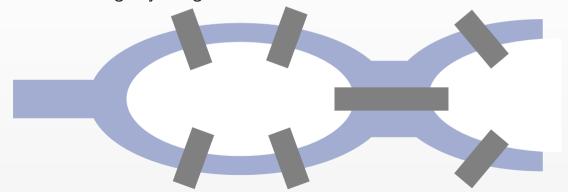


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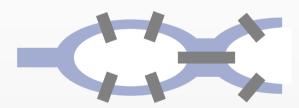
$$\mathbf{G}^{20} = \begin{bmatrix} 10270848 & 14562688 & 11949760 & 9746560 & 15086976 \\ 14562688 & 22220608 & 17699904 & 15086976 & 21696320 \\ 11949760 & 17699904 & 14267264 & 11949760 & 17699904 \\ 9746560 & 15086976 & 11949760 & 10270848 & 14562688 \\ 15086976 & 21696320 & 17699904 & 14562688 & 22220608 \end{bmatrix}$$

Peter Rowlett SHU Paths and cycles 5 / 16

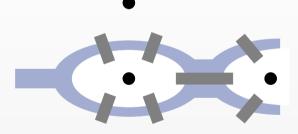
▶ In the town of Königsberg there were seven bridges across the river Pregel. Is it possible to go for a walk, crossing each bridge once, but not crossing any bridge twice?



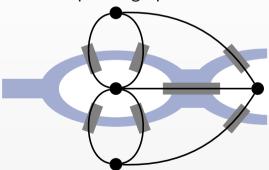
- ► This was solved by Euler in 1736 using an essentially topological argument.
- ▶ We might redraw the map as a graph.



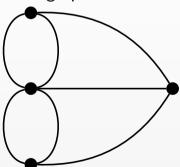
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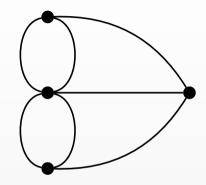
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Now we ask whether there is an Eulerian cycle − a cycle that includes every edge of the graph exactly once. (It may visit vertices more than once.)



► For a vertex with two edges, we can enter the vertex along one and exit along the other.



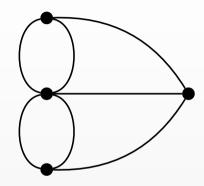
- ► For a vertex with two edges, we can enter the vertex along one and exit along the other.
- However, for a vertex with one edge, we enter the vertex and cannot proceed.



- ► For a vertex with two edges, we can enter the vertex along one and exit along the other.
- However, for a vertex with one edge, we enter the vertex and cannot proceed.
- Or if we started at that vertex, we leave and cannot return at the end of the cycle.



- ➤ So, for a Eulerian cycle, we require that every vertex has even degree.
- A graph which contains an Eulerian cycle is called an Eulerian graph.



# Eulerian path

- ► A related concept is an *Eulerian path*.
- ► This is a path that includes every edge of the graph but does not return to the starting point.
- ► A graph that contains an Eulerian cycle also contains an Eulerian path, since a cycle is a type of path.
- ► A graph that contains an Eulerian path but not an Eulerian cycle is called *Semi-Eulerian*.

#### Useful theorems

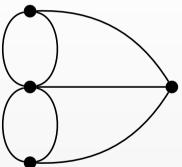
#### **Theorem**

A graph is Eulerian if and only if it is connected and every vertex has even degree.

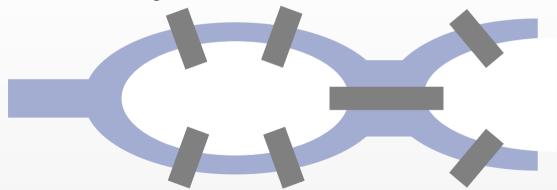
#### **Theorem**

A graph is semi-Eulerian if and only if it is connected and exactly two vertices have odd degree. An Eulerian path must start at one odd vertex and finish at the other.

- ▶ In the Königsberg Bridges graph, all four vertices are odd.
- ► This graph is neither Eulerian nor semi-Eulerian.

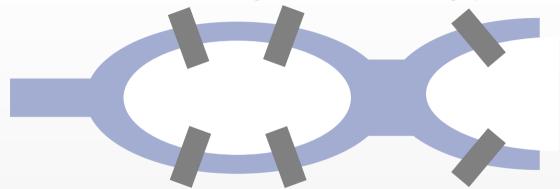


► This means we cannot cross every bridge exactly once, we must cross at least one bridge more than once.



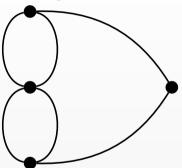
# Different bridges

▶ If we remove the central bridge, we obtain a different graph.



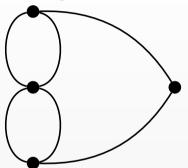
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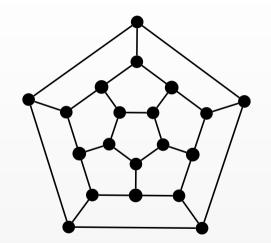
► Since we now have exactly two vertices of odd degree, we can make an Eulerian path on this graph.

### Hamiltonian graphs

- ▶ A path that visits every vertex exactly once is called a *Hamiltonian* path.
- A cycle that visits every vertex once and then returns to the starting vertex is called a *Hamiltonian cycle*.
- A graph that includes a Hamiltonian cycle is called a *Hamiltonian* graph, and a graph that includes a Hamiltonian path is called a semi-Hamiltonian graph.
- ► A graph that contains an Hamiltonian cycle also contains an Hamiltonian path, since a cycle is a type of path.
- ► (Note this is about visiting vertices, rather than edges.)

#### The Icosian Game

- The name comes from this game invented by mathematician William Rowan Hamilton in 1857.
- ► The game is based on finding cycles on this graph.



- This graph contains a Hamiltonian cycle (e.g.  $a \rightarrow b \rightarrow c \rightarrow e \rightarrow d \rightarrow a$ ).
- ▶ And therefore also a Hamiltonian path (e.g.  $a \rightarrow b \rightarrow c \rightarrow e \rightarrow d$ ).

