

# Paths and cycles

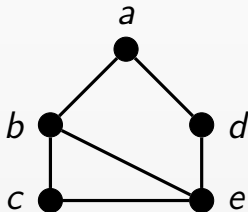
Peter Rowlett

Sheffield Hallam University

`p.rowlett@shu.ac.uk`

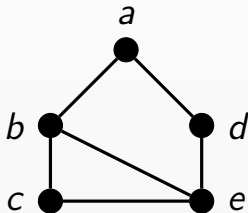
# Paths

- ▶ Some sources use more specific definitions, but we'll say a *path* is a sequence of edges which join a sequence of vertices.
- ▶ A path's *length* is the number of edges in it.
- ▶ For example, there is a path of length 3 from  $a$  to  $e$  in the graph below, it goes  $a \rightarrow b \rightarrow c \rightarrow e$ .



# Cycles

- ▶ A cycle is a path that ends at its starting vertex.
- ▶ For example, there is a cycle of length 4 from  $a$  in the graph below, it goes  $a \rightarrow b \rightarrow e \rightarrow d \rightarrow a$ .



# Adjacency matrices

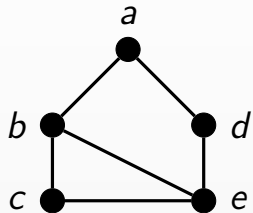
- ▶ Diagrams of dots and lines are not the only way to represent graphs. One way that can be useful is to represent the graph as a matrix.
- ▶ The *adjacency matrix* for a graph,  $G = (V, E)$ , with  $n$  vertices is an  $n \times n$  matrix **M** with entries

$$\mathbf{M} = \begin{bmatrix} m_{11} & m_{12} & \dots & m_{1n} \\ m_{21} & m_{22} & \dots & m_{2n} \\ \vdots & \vdots & \ddots & \vdots \\ m_{n1} & m_{n2} & \dots & m_{nn} \end{bmatrix}$$

$$m_{ij} = \begin{cases} 1 & \text{if } (i, j) \in E; \\ 0 & \text{otherwise.} \end{cases}$$

# Example

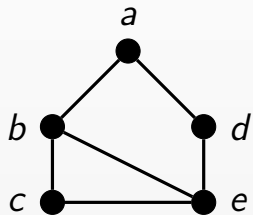
- The following are representations of the same graph.



$$\mathbf{G} = \begin{bmatrix} 0 & 1 & 0 & 1 & 0 \\ 1 & 0 & 1 & 0 & 1 \\ 0 & 1 & 0 & 0 & 1 \\ 1 & 0 & 0 & 0 & 1 \\ 0 & 1 & 1 & 1 & 0 \end{bmatrix}.$$

# Example

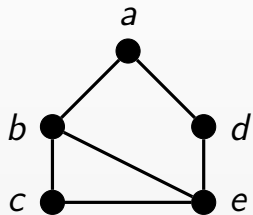
- ▶ Then entries of  $\mathbf{G}^k$  count the number of paths of length  $k$  between vertices.
- ▶ i.e. the  $(i,j)$ th entry of  $\mathbf{G}^k$  gives the number of paths of length  $k$  from vertex  $i$  to vertex  $j$ .



$$\mathbf{G}^2 = \begin{bmatrix} 2 & 0 & 1 & 0 & 2 \\ 0 & 3 & 1 & 2 & 1 \\ 1 & 1 & 2 & 1 & 1 \\ 0 & 2 & 1 & 2 & 0 \\ 2 & 1 & 1 & 0 & 3 \end{bmatrix}.$$

# Example

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$$\mathbf{G}^3 = \begin{bmatrix} 0 & 5 & 2 & 4 & 1 \\ 5 & 2 & 4 & 1 & 6 \\ 2 & 4 & 2 & 2 & 4 \\ 4 & 1 & 2 & 0 & 5 \\ 1 & 6 & 4 & 5 & 2 \end{bmatrix}.$$

# Example

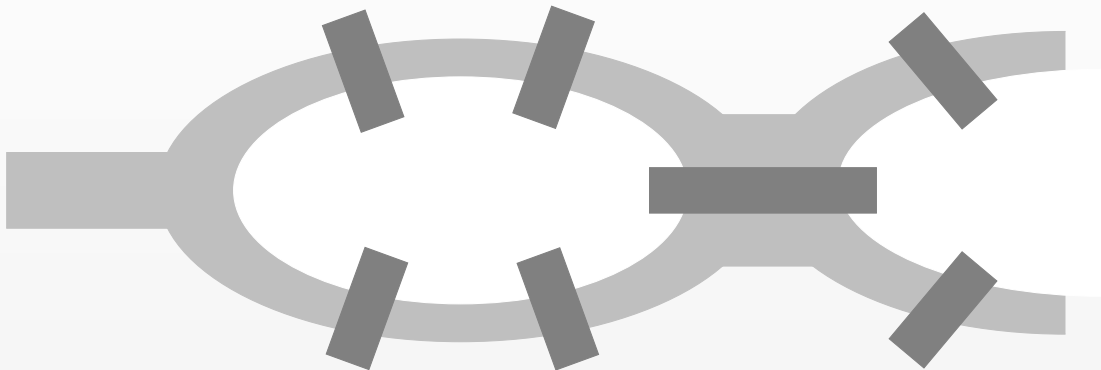
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$$\mathbf{G}^{20} = \begin{bmatrix} 10270848 & 14562688 & 11949760 & 9746560 & 15086976 \\ 14562688 & 22220608 & 17699904 & 15086976 & 21696320 \\ 11949760 & 17699904 & 14267264 & 11949760 & 17699904 \\ 9746560 & 15086976 & 11949760 & 10270848 & 14562688 \\ 15086976 & 21696320 & 17699904 & 14562688 & 22220608 \end{bmatrix}.$$



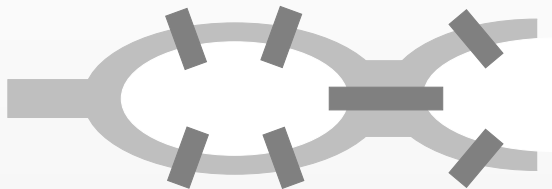
# Königsberg Bridges Problem

- In the town of Königsberg there were seven bridges across the river Pregel. Is it possible to go for a walk, crossing each bridge once, but not crossing any bridge twice?



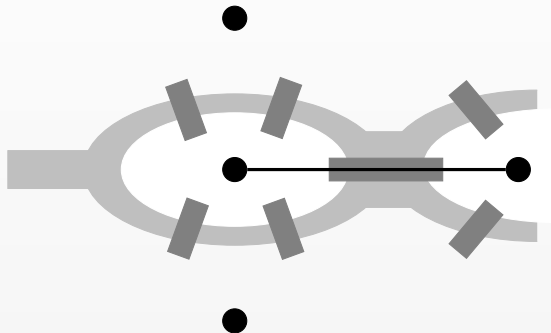
# Königsberg Bridges Problem

- ▶ This was solved by Euler in 1736 using an essentially topological argument.
- ▶ We might redraw the map as a graph.



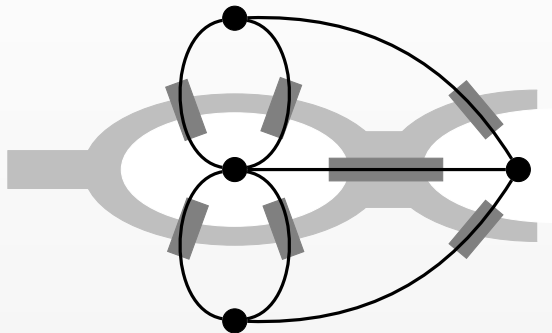
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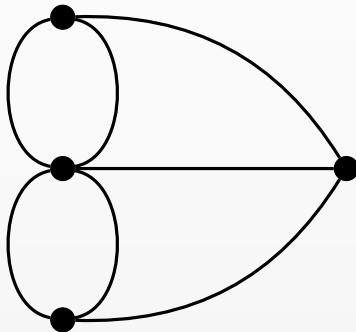
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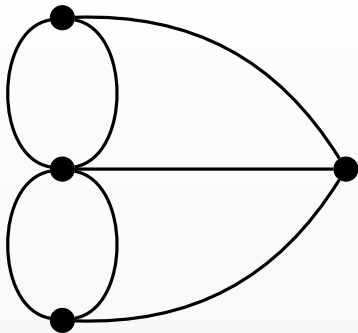
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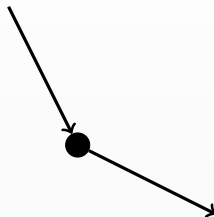
# Eulerian cycles

- Now we ask whether there is an *Eulerian cycle* – a cycle that includes every edge of the graph. (It may visit vertices more than once.)



# Eulerian cycles

- For a vertex with two edges, we can enter the vertex along one and exit along the other.



# Eulerian cycles

- ▶ For a vertex with two edges, we can enter the vertex along one and exit along the other.
- ▶ However, for a vertex with one edge, we enter the vertex and cannot proceed.





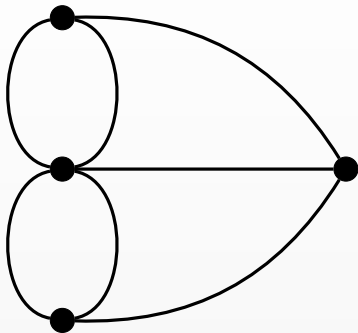
# Eulerian cycles

- ▶ For a vertex with two edges, we can enter the vertex along one and exit along the other.
- ▶ However, for a vertex with one edge, we enter the vertex and cannot proceed.
- ▶ Or if we started at that vertex, we leave and cannot return at the end of the cycle.



# Eulerian cycles

- So, for a Eulerian cycle, we require that every vertex has even degree.
- A graph which contains an Eulerian cycle is called an *Eulerian graph*.



# Eulerian path

- ▶ A related concept is an *Eulerian path*.
- ▶ This is a path that includes every edge of the graph but does not return to the starting point.
- ▶ A graph that contains an Eulerian cycle also contains an Eulerian path, since we can just leave off the edge that returns us to the starting vertex.
- ▶ A graph that contains an Eulerian path but not an Eulerian cycle is called *Semi-Eulerian*.

# Useful theorems

## Theorem

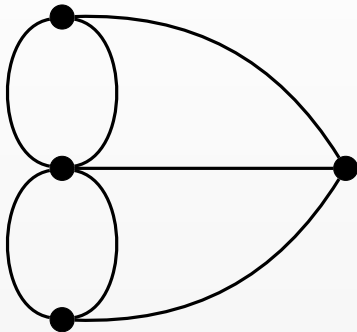
*A graph is Eulerian if and only if it is connected and every vertex has even degree.*

## Theorem

*A graph is semi-Eulerian if and only if it is connected and exactly two vertices have odd degree. An Eulerian path must start at one odd vertex and finish at the other.*

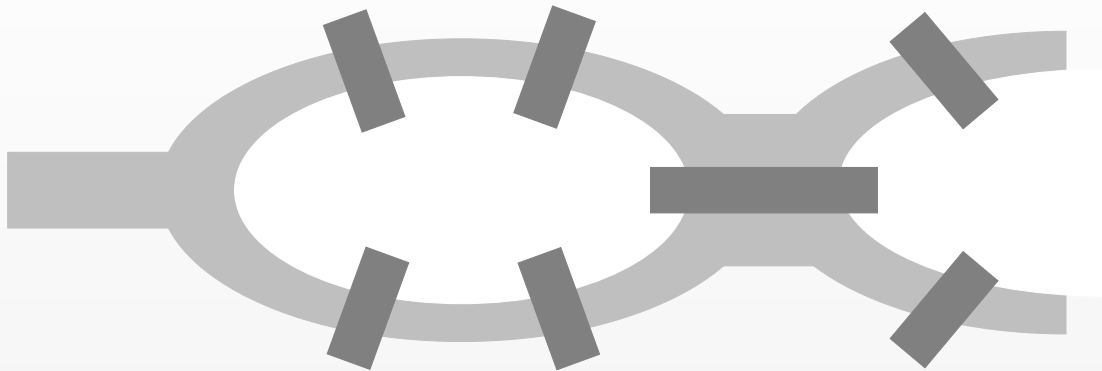
# Königsberg Bridges Problem

- ▶ In the Königsberg Bridges graph, all four vertices are odd.
- ▶ This graph is neither Eulerian nor semi-Eulerian.



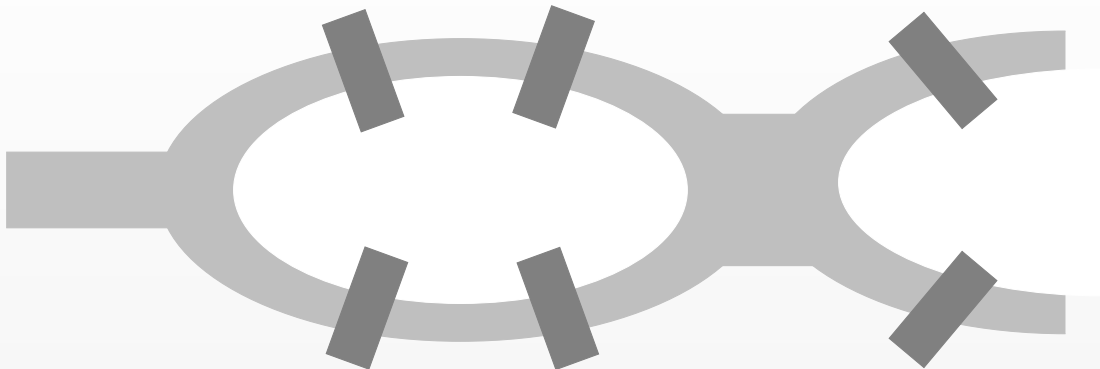
# Königsberg Bridges Problem

- This means we cannot cross every bridge exactly once, we must cross at least one bridge more than once.



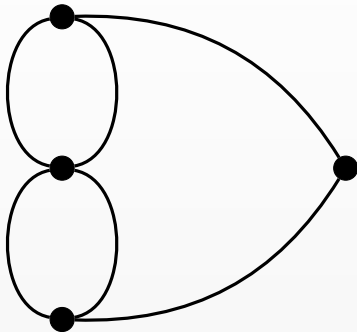
# Different bridges

- If we remove the central bridge, we obtain a different graph.



# Different bridges

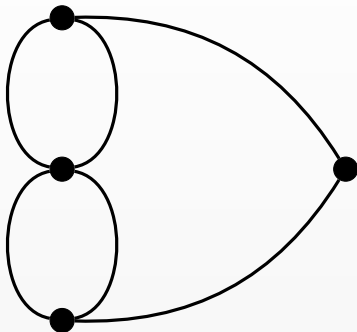
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# Different bridges

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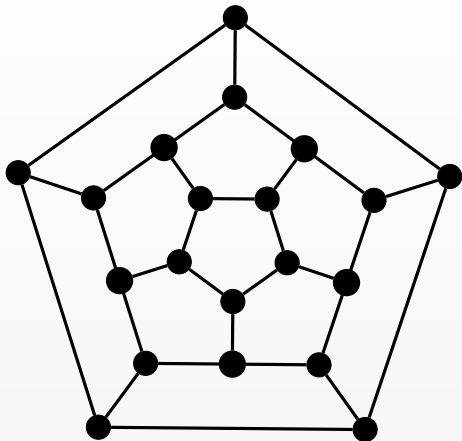
- Since we now have exactly two vertices of odd degree, we can make an Eulerian path on this graph.

# Hamiltonian graphs

- ▶ A path that visits every vertex exactly once is called a *Hamiltonian path*.
- ▶ A cycle that visits every vertex once and then returns to the starting vertex is called a *Hamiltonian cycle*.
- ▶ A graph that includes a Hamiltonian cycle is called a *Hamiltonian graph*, and a graph that includes a Hamiltonian path is called a *Hamiltonian path*.
- ▶ A graph that contains an Hamiltonian cycle also contains an Hamiltonian path, since we can just leave off the edge that returns us to the starting vertex.
- ▶ (Note this is about visiting vertices, rather than edges.)

# The Icosian Game

- ▶ The name comes from this game invented by mathematician William Rowan Hamilton in 1857.
- ▶ The game is based on finding cycles on this graph.



# Example

- ▶ This graph contains a Hamiltonian cycle (e.g.  $a \rightarrow b \rightarrow c \rightarrow e \rightarrow d \rightarrow a$ ).
- ▶ And therefore also a Hamiltonian path (e.g.  $a \rightarrow b \rightarrow c \rightarrow e \rightarrow d$ ).

