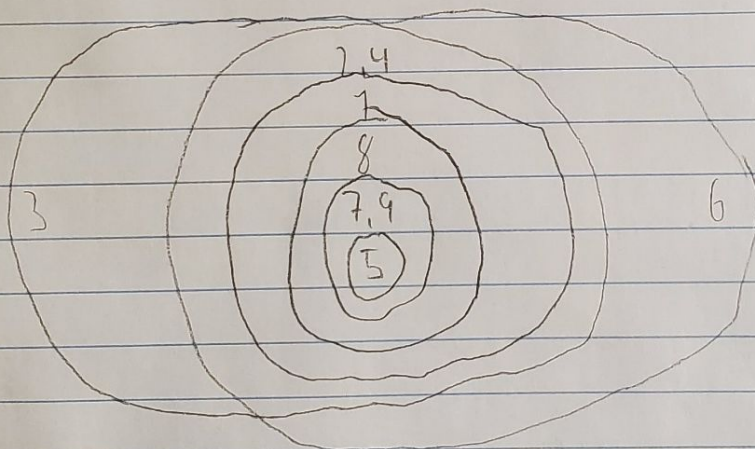


- 1)
- a) T
 - b) T
 - c) F
 - d) T
 - e) F
 - f) T
 - g) T
 - h) T
 - i) F
 - j) T

2)



(2) and (4) are equivalent
(7) and (9) are equivalent

- 3) a) Accepted: ϵ , 01
Rejected: 001, 00001

b) This language recognizes any string with the same number of 0's and 1's

Let $L = \{ w \in \{0, 1\}^* \mid w \text{ has the same number of 0's and 1's} \}$

(c) Not a regular language

4) a) $Q \Rightarrow \cdot ZV$
 $Z \rightarrow AS \mid BS \mid a \mid b \mid \epsilon$
 $S \rightarrow AS \mid BS \mid a \mid b$
 $V \rightarrow aaW \mid \epsilon$
 $W \rightarrow BW \mid b$
 $A \rightarrow a$
 $B \rightarrow b$

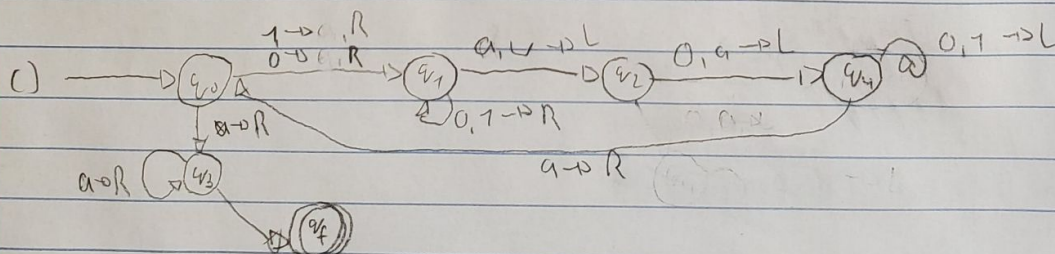
b) $Q \rightarrow ZQ \mid \epsilon$
 $Z \rightarrow AS \mid BS \mid a \mid b \mid \epsilon$
 $S \rightarrow AS \mid BS \mid a \mid b$
 $A \rightarrow a$
 $B \rightarrow b$

c) G_1 is in CNF

5) Let $A = \{w \in \{0,1\}^* \mid w0^i \text{ and } |w| = i\}$

a) 001000, 11110000

b) 1, 11



6) 1) a) Not accepted
 b) Not accepted
 c) Accepted
 d) Not accepted

6) 2) 1) Yes

- i) It is in the correct form as it contains single Turing machine encodings.
- ii) It is a property of the language of a Turing machine.
- iii) It is nontrivial

• Because some Turing Recognizable languages are of even length and some are not of even length

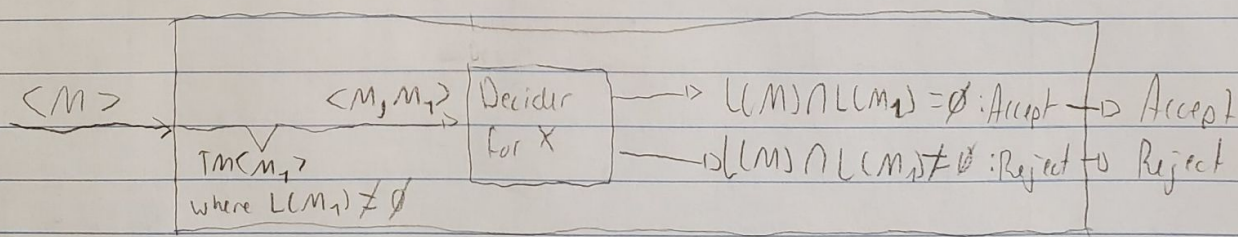
2) No.

- i) (Same reason as previous response)
- ii) (Same reason as previous response)
- iii) It is trivial

• Because all Turing Recognizable languages have $|L(M)| < \infty$

3) 1) Rice's Theorem cannot be applied

2) Assume that E_{TM} is decidable to prove that X is undecidable via contradiction.



The intersection of an empty language and non-empty language will always be empty, therefore E_{TM} is undecidable.

Since E_{TM} is undecidable, then X must also be undecidable.