

## ★ Stacks and Queues - Short Notes

- Remembering stacks best example - plates in wedding,  
queue " " - ticket counter,

### ## Stacks:

- `Stack<Integer> stack = new Stack<>();`  

↓  
variable stored  
in stack memory
this new object is created in Heap
- To add items = `stack.push(34);`  
 To remove items = `stack.pop();` // Top one will be removed only
- Stacks is a class whereas Queue is an interface.

### \*\*\*

Note:- we write `Queue<Integer> queue = new LinkedList<>();` instead of " " " " " " `Queue<>();` bcoz the nature of interfaces in Java. ⚠ compilation error

- Queue is an interface & interfaces can't be instantiated directly (They only define a set of methods that class must implements)  
 e.g. Queue provides → `add()`, `remove()`, `offer()`, `peek()` & `poll()` but doesn't provide any implementation for them.
- Whereas LL is a class that implements Queue/Deque interface's methods  
 e.g. `PriorityQueue<>` is a class
- Hence when we write in this manner we're using LL object but referring it to Queue interface, this is called polymorphism
- To add items → `queue.add(3);`  
 To check top item → `queue.peek();`  
 " remove item → `queue.remove();` → // bottom one will be removed



• Deque (Deck)

[Deque<Integer> deque = new ArrayDeque<>();]

★ Similarly, here also ArrayDeque is a class that provides us to create object of methods inside deque interface & helps to implement.

∴ uses

```

deque.add(s);
deque.addLast(s);
deque.addFirst(s);
deque.removeFirst();
"    "    last();
  
```

by default

None

insert First

insert last

remove first

remove last

## \* Custom Stack Implementation

```
public class CustomStack {
```

```
    protected int[] data; → (protected access modifier is to use this in same package i.e. (Dynamic Stack))
```

Also not uses default since it is not accessible to subclasses in other packages whereas protected is accessible + (data is stored in array).

```
    private static final int DEFAULT_SIZE = 10;
```

```
    public CustomStack() {
```

```
        this(DEFAULT_SIZE); → nonparameterized constx for default size
```

```
    public CustomStack(int size) {
```

```
        this.data = new int[size] → parameterized constx for given size of array in which data is stored
```

```
    int ptr = -1; → initialised pointer
```



• To add items → `public boolean push(int item) {`  
`ptr++ ;` → ptr increased

Store that item on ptr index → `data[ptr] = item;`

`if (isFull()) {`

`out("Full");`

`return false;`

`return true;`

when ptr is at last index

i.e. `ptr == data.length - 1`

• To remove → `public int pop() throws StackException {`

`if (isEmpty()) {`

`throw new StackException("cant pop from empty");`

`return data[ptr--];`

we cannot return -1 (int) as it may be in the stack (valid no.)

when pointer is at -1

make a file of StackException which extends Exception (in built)

`super(message);`

↳ In main fu. → `psvm` throws `StackException`

\* Dynamic Stack → same as of custom stack, only diff. when insert item size functions as arraylist concept.

`public class DynamicStack extends CustomStack {`

non-parameter constructor of DynamicStack → called `super()`;

Size " " " " → " `super(size);`

@Override

`public boolean push(int item) {`

`if (this.isFull()) {`

no need to write isFull func. ∵ extended

`int[] temp = new int[data.length + 2];` ← create a temp. array (doubled if full)

Copy all prev. items → `for (int i = 0; i < data.length; i++) {`

`temp[i] = data[i];`

update the temp array → `data = temp;`

`else {`

`return super.push(item);` → directly.



## \* Custom Queue Implementation

basic starting same as of Custom Stack  
except here make end pointer & start with zero.  
(therefore, here we'll first add then increase ptr till end)

```
int end = 0;
```

add item → `public boolean insert (int item) {`  
     `if (isFull()) {` → when end ptr at ~~at~~ data.length  
         `return false;`  
     `data[end] = item;`  
     `end++;`  
     `return true;`

remove item → `public int remove () throws Exception {`  
     • `if (isEmpty()) throw exception`  
     • else firstly get the item (returning (showing) purposes) → int removed  
         = data[0];  
     • then shift all the elements to left (actual removing) → remove from first  
         `for (int i = 1; i < end; i++) {`  
             `data[i-1] = data[i]`  
         `}`  
         `end--;`  
         `return removed;`

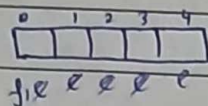
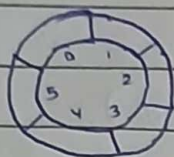
peek item (front) → `public int front () throws Exception`  
     `if (empty) → ✓`  
     `return data[0];` → before  
     → when end ptr at 0;  
     " is full & end ptr out of bound pe rehita hai

display → `for i = 0 till end {`  
     `cout (data[i] + " < ");`  
     `}`



lly, as of Dynamic stack we can do for Dynamic Queue (only insert changes) but it is extended from circular queue.

## \* Circular queue implementation



Concept: consider items only in b/w first (f) and end (e).

\* Note:- after inserting till last size of array order use  $[end \% size]$   $\rightarrow$  adding again.   
  $\because$  we can add  $\infty$  items in circular queue.  $\because$  only (f, e) will be consider

e.g. size of array/queue is 5 + we are adding 6th item it will be stored at index 1  $\because (6 \% 5)$ .

implementation:- basics, same as of custom queue +   
  $\uparrow$  is full/empty   
 initialise  $\rightarrow$  protected int end = 0   
 " " first = 0

add items  $\rightarrow$  public boolean insert (int item) {

if full  $\rightarrow$  return false

else  $\rightarrow$  data[end] = item;   
 end++;

end = 0   
  $\rightarrow$  1

\* end = end % data.length;   
 end\_ptr++;

1 % 5 = 1  $\checkmark \rightarrow$  check

return true;

remove  $\rightarrow$  same as queue +  $\checkmark$  else {

int removed = data[front];   
 front++;

update front also  $\rightarrow$  front = front % data.length;   
 end\_ptr--;

return removed.

peek (front) → same as Queue from

```

public int front() throws Exception {
    if (isEmpty()) {
        // ...
    }
    return ...
}
    
```

display → if empty → print empty → return .

else → int i = front : (starting from the front)

```

do {
    Sout (data[i] + "→");
    i++;
    i %= data.length → "circular condition" → socho
    } while (i != end);
    Sout ("End");
}
    
```

don't memorize



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travelling

1. Implement Queue using stacks

g. using two stacks  $\rightarrow$  functions used (push, peek, pop and empty)  
make FIFO

Concept  $\Rightarrow$  ~~add~~ items in stack first

• put those in second stack till empty :  $\left[ \frac{1}{2} \right]$  : using pop of first  
then remove from this stack.

Peek → again put that back in first & return peak  
 first follow above same steps  
 use bead restore bhi to karoge na. :-  
 get that element i.e. removed → using second.pop  
 then till second is not empty put it in first.

Steps  $\Rightarrow$  • firstly declare both stacks

- In myQueue method.  $\rightarrow$  initialise both stacks.

- In Push " → ~~add~~ push in first stack

- In POP  $\rightarrow$  under condition that second is empty  
now while first is not empty put items from first  $\rightarrow$  second  
return second.pop

ditto ← "In peek" → " " " " " "

of pop " " " " " " " " " "

- In empty  $\rightarrow$  return first empty & second empty.

## 2. Game of Two Stocks

Q. There will be two stacks A & B from both the elements you've to take the max. no. of items till max limit.

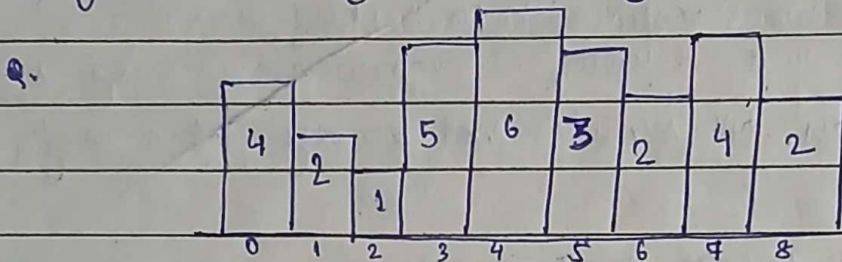
Steps  $\Rightarrow$  • Create two stack method which takes max. no of both stocks in which ~~all~~<sup>we'll</sup> call the actual stocks - 1.



- The next method must contain  $\text{max no.}$ , both stacks,  $\text{sum till}$   $(x)$  here & count.  $(\text{sum})$
- If  $\text{sum} > x$  then return count.
- else make two recursion call in  $a + b$  in which we'll copy the range and increase count & sum by  $\text{stack}[0]$  in both  $\text{ans1}$  &  $\text{ans2}$ .  
value at first

3.   
 Leetcode Hard

### Largest rectangle in Histogram



- Concept  $\rightarrow$
- get the height of each indices  
 $\Rightarrow a = \{4, 2, 1, 5, 6, 3, 2, 4, 2\}$
  - get the prev. smaller (any <sup>nearest</sup> item is less than that particular index then mention its index  
 $\Rightarrow ps = \{-1, -1, -1, 2, 3, 2, 2, 6, 2\}$   
 if not found
  - lly get for next smaller  
 $\Rightarrow ns = \{3, 3, 3, 4, 9, 4, \}$   
 if not found  
 $\Rightarrow ns = \{1, 2, 3, 5, 5, 6, 9, 8, 9\}$

• Then area of the ques. will be:-

$$A = ps - ns = (ns[i] - ps[i] - 1) * a[i]$$

Steps  $\Rightarrow$  • create an array of  $ps$  & call its items from a fn.

lly for  $ns$ .

- get  $\text{maxArea}$  using formula for individual indices



- In prevSmaller method  $\rightarrow$ 
  - create an array  $ps$  & stack of integer
  - for 0 to  $a.length$  till while stack is not empty and  $\bullet$  top of stack ( $a[s.peek()]$ ) is greater than  $a[i]$ 
    - $s.pop()$ ;
  - If  $s.isEmpty() \rightarrow ps[i] = -1$ , else  $ps[i] = s.peek()$ ;
  - then push ~~at~~  $i$  after checking (always)
  - & return  $ps$ .
- illy for nextSmaller method  $\rightarrow$  everything except go from last to first and for if stack is empty  $ms[i] = b.length$ ; // obviously not coz array size might change.

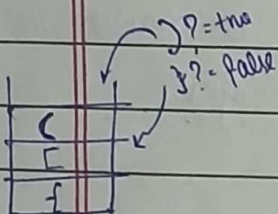
#### 4. Valid Parenthesis

e.  $\{ [ ( ) ] \} \rightarrow true$ ,  $\{ [ ( ] [ ] \} \rightarrow false$ ,  $\{ [ ( ] ) \} \rightarrow false$

Concept  $\rightarrow$  Create a character type array.

for each element in string  $s$  ( $s.toCharArray()$ ) character

- if  $ch$  is opening brackets then push them in.
- for closing brackets if only if stack is empty or just prev character (using pop) is not <sup>its</sup> opening bracket then false



- In the end return  $stack.isEmpty()$ ;  $\rightarrow$  not true directly becoz if  $s = '[$   $\therefore$  wrong Ans.  $\therefore$  <sup>\*</sup>if we're popping out we'll make sure all opening & closing are together successfully.

#### 5. Minimum insertions to balance a Parentheses String

e. for every opening bracket it needs two closing brackets to make it balanced. eg.  $s = "))(())(" \therefore O/P = 3 \rightarrow$  1 "(" for first "))( and 2 ")" for last "(".

e.g.  $s = "(())"$   $\therefore op = 1$ .



Steps  $\Rightarrow$

6. Minimum Add to make Parentheses Valid

eg. C)CC  $\rightarrow$  2

Steps  $\Rightarrow$  • make character stack

• for each character in string

if  $ch = )$   $\rightarrow$  closing  $\left( \begin{matrix} 4f \\ \star \end{matrix} \right)$  top of stack is "C" opening equal

•  $\therefore$  pop no use to count

• else push

$\rightarrow$  • ~~if~~ ~~else~~ else  $\rightarrow$  push that ch first then check

• return `stack.size()`;  $\rightarrow$   $\because$  we've pushed only those items in stacks to make valid

$\therefore$  returning integer.