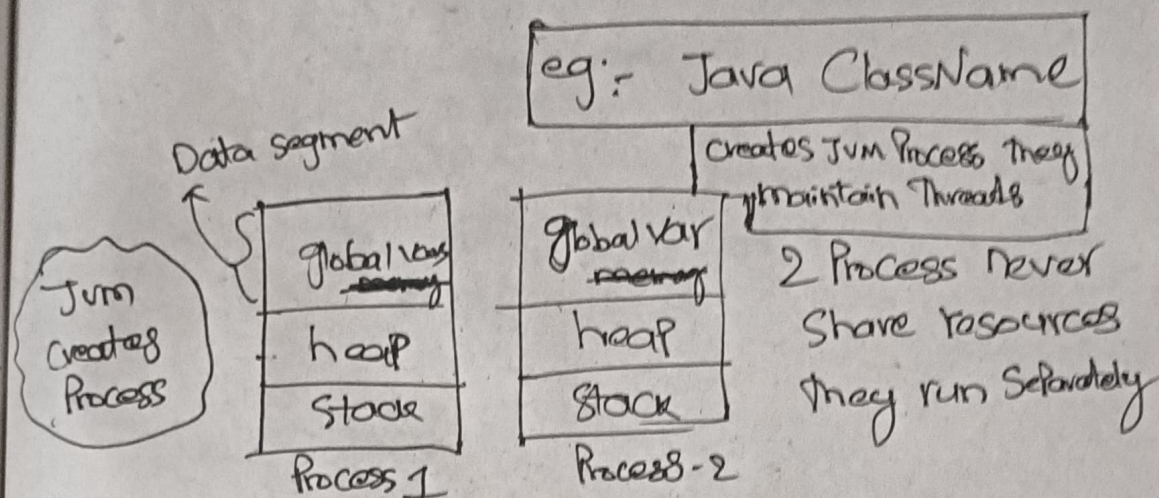


# MultiThreading

## → Process v/s Thread

Process :- It is an instance of the Program that is being executed.

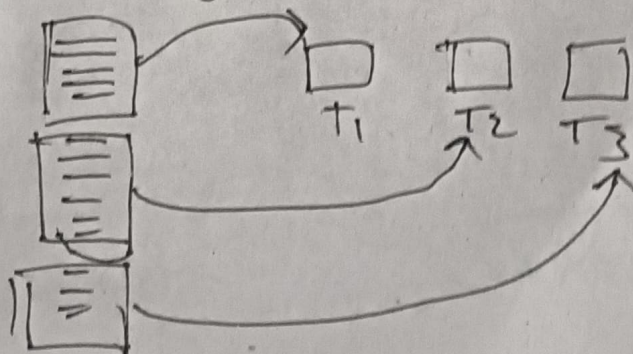


Thread :- It is a light weight Process (or)

→ Smallest sequence of instructions that are executed by CPU independently

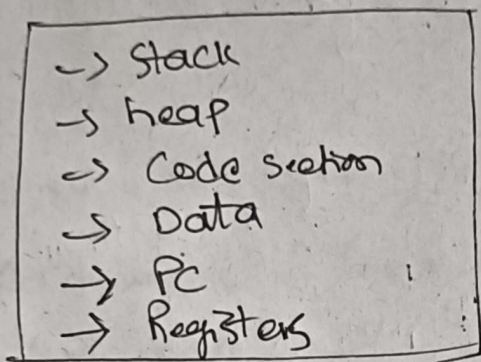
→ Process can have multiple thread.

Machine code



→ Initially we have main <sup>Thread</sup> ~~class~~ from this multiple Threads are created

eg:- Thread.currentThread().getName();

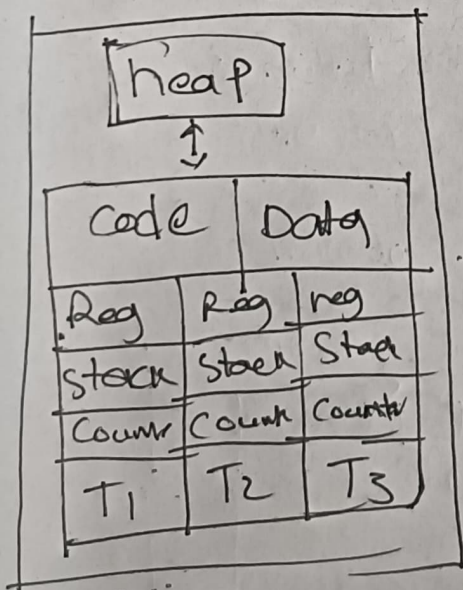


↕ JVM ↕

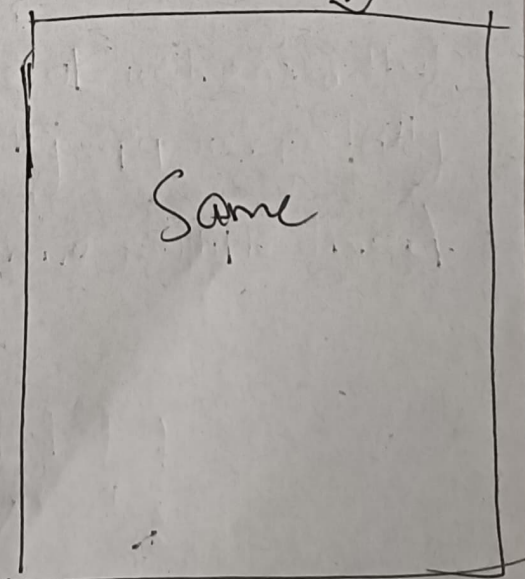
Java MyClass:



① Create <sup>new</sup> JVM Process  
x



JVM Process 1



JVM Process 2



eg:- java -Xms256m MyClass  
java -Xmx2g MyClass

Why Separate registers, Stack, Counter

### Code Segment

→ Stores ~~the~~ machine code that CPU understands

→ Data Segment (global & static data)

⊗ → read only

### Data Segment

→ Global & static vars

→ multiple threads can access data

→ So Synchronization is needed

⊗ → read/write

### Heap

→ "new", keyword creates object in heap area

→ ~~Threads~~ Threads share same heap memory

→ Synchronization is required

⊗ → read/write

## Stack:-

- Each Thread has own stack
- manages method calls, local variables, return address

## Register:-

- Store intermediate value
- JIT uses stack to reshuffle the machine code
- ⊙ → Helps in Context Switching
- Each Stack has own register

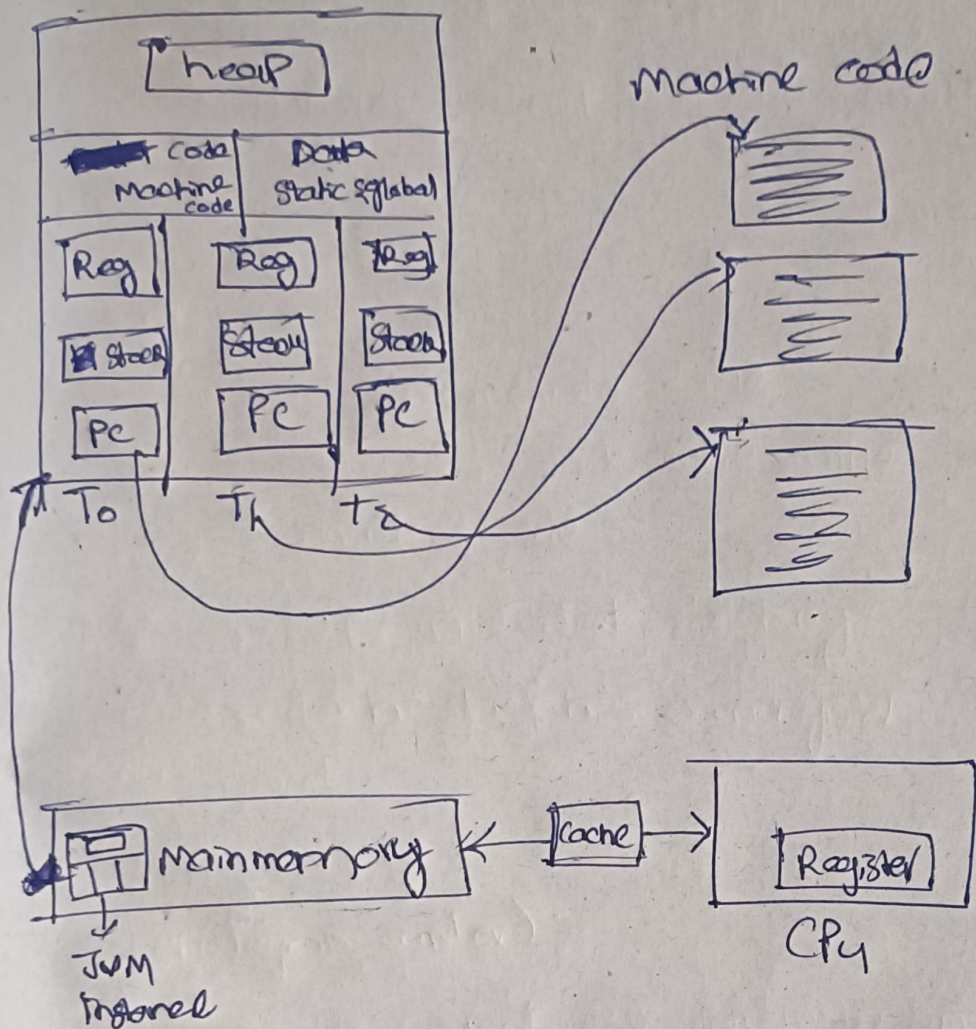
## Counter (Program Counter):-

Point to current executing instruction

after execution PC is incremented to next instruction that need to be executed



Java Myclass  $\rightarrow$  JIT compiler  $\rightarrow$  Machine code



This machine code part which  $T_0$ 's PC is pointing to is copied into  $T_0$  register and based on OS Scheduling algo's (FCFS, ~~SF~~SRF, SRTF, RR) Thread will be allocated execution time and CPU uses its registers for storing intermediate values when its time is up CPU register data copied into  $T_0$ 's individual register now its turn of  $T_1$ . This process is called Context Switching.