Social and Information Networks 2: Evolution

Internet Analytics (COM-308)

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Overview

- Herding and "watching thy neighbor"
 - Information cascades: why imitating your friends makes sense - and how it can lead to surprising group behavior
 - Heavy-tailed degree distributions: "the rich get richer" applied to networks
- Observing network properties
 - The importance of the observer
 - Example: your friends are more popular than you!

Watching thy Neighbor?

- Human decision-making:
 - Primary private information...
 - Heavily influenced by what decisions taken by others
- Reason:
 - Primary information: often too voluminous, noisy, not trustworthy,...
 - By imitating others, piggyback on their effort to interpret primary information
- Question:
 - Macro behavior of such systems?



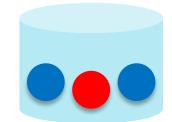
Herding and information cascades

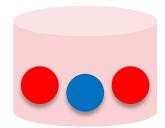
Assumptions:

- Decision: choose a restaurant, adopt new technology, political position, fashion,...
- Sequential, and each person can observe choices made earlier by others
- Each person has some private information to help guide decision: favorite food, taste,...
- Private information not observed by others (can't see what others "know"), but decisions/actions are (can see what others "do")

Herding: how it can go wrong

- Urn with 3 balls
 - A priori distribution (blue/red majority) = (0.5,0.5)
 - majority blue: 2 blue + 1 red
 - majority red: 2 red + 1 blue





- A group of people take turns:
 - Draw a ball from the urn at random
 - Check the color of the ball privately, put it back in urn
 - Announce their guess (blue/red majority) to everybody
 - Receive reward for correct guess
- Assumption:
 - DEach individual is altruistic: do what allows others to make best guess
 - Each individual is selfish = tries to make best guess for himself

Urn model: altruistic (0)

- Every person:
 - Selects a ball at random (with replacement)
 - Announces the color of the ball to everybody as their guess, even if previous information suggests a different guess
- As $n \to \infty$, majority color of urn is equal to color most frequently observed
 - Consequence of law of large numbers
- After a few "sacrifices", everybody could produce best guess
 - Sacrifice in the sense that individuals might be forced to say red (color of their ball) even if previous information suggests blue majority

Urn model: selfish (2)

- Sequential decision-making
 - Selfish guess: use previous public and new private information to maximize own reward
 - Observed color remains private
- First individual:
 - Blue ball: announce guess(1) = blue
 - Red ball: announce guess(1) = red
 - Public guess of first fully reveals private information
- Second individual:
 - If color(2) = guess(1): announce this color
 - If color(2) ≠ guess(1): does not matter (assume color(2))
 - Public guess of second fully reveals private information

Urn model: selfish (2)

- Third individual:
 - If guess(1) ≠ guess(2): announce guess(3) = color(3)
 - If guess(1) = guess(2):
 - Announce guess(3)=guess(2)=guess(1), regardless of color(3)
 - Why is this?
 - Person 3 knows that guesses 1+2 reveal perfect information
 - Therefore, regardless of color(3), guess(1)=guess(2) dominates guess
- Fourth,...,∞th individual:
 - If guess(1) = guess(2):
 - Announce guess(i) = guess(2)=guess(1), regardless of color(i)

Urn model: (2) leads to cascade

- If guess(1) = guess(2) were both wrong, then all future guesses are wrong!
- This happens with prob. 1/9
- Even though each individual is using available information in the best way to make a guess

Information cascade: suboptimal decision

- Cascade: sequential decisions
- Individual:
 - Efficiency gain by observing others' decisions
- Global behavior:
 - Primary information can "wash out"
 - Suboptimal or random decisions
- Might these be cascades:
 - Stock market gyrations, "flash crash"
 - Inexplicable shifts in popularity of {restaurants, clubs, celebrities,...}
 - Fashion, style, celebrity,...



• ...

Herding in networks

- Observation:
 - Degree distributions in networks often resemble power laws

 $A \propto B \rightarrow A$ proporcional a B

Power law:

$$P(D>d) \propto d^{-\gamma}$$
 tail = cola

Most distributions have "light tails":

•
$$P(D>d) \propto e^{-\alpha d}$$
 (or lighter/bounded)

• Exponential, Geometric, Gaussian, Poisson, ...

Pareto (β, γ) distribution

- Support: $d \in [\beta, \infty)$
- CCDF (Complementary Cumulative Distribution Function):

•
$$P(D > d) = 1 - F_D(d) = \begin{cases} \left(\frac{d}{\beta}\right)^{-\gamma}, & d \ge \beta \\ 1 & \text{otherwise} \end{cases}$$

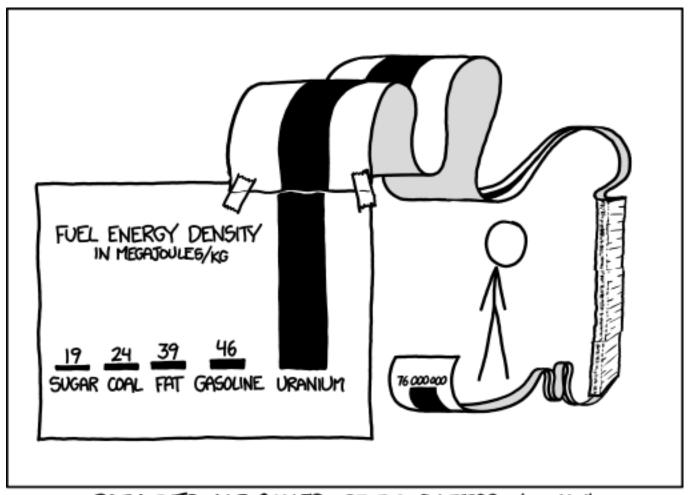
- γ : exponent, also called "Pareto index"
- Moments:

•
$$E[D^k] = \begin{cases} \frac{\beta^k \gamma}{\gamma - k}, & k < \gamma \\ \infty & \text{otherwise} \end{cases}$$

Numerical comparison exp/power

- Distribution of human height:
 - Mean = 178 cm
 - Stddev = 8 cm
- Compare tails: how tall are extremely tall people?
 - What is d^* such that $P(D > d^*) = 10^{-9}$
- Normal $N(178 \text{cm}, (8 \text{cm})^2)$:
 - $d^* = 226$ cm
- Pareto: choose β , γ s.t. first and second moments match data
 - $\gamma \cong 23, \beta \cong 170$ cm
 - $d^* = 420 \text{ cm} !!$
- Assumption very important for extremal values!

Log-log plot

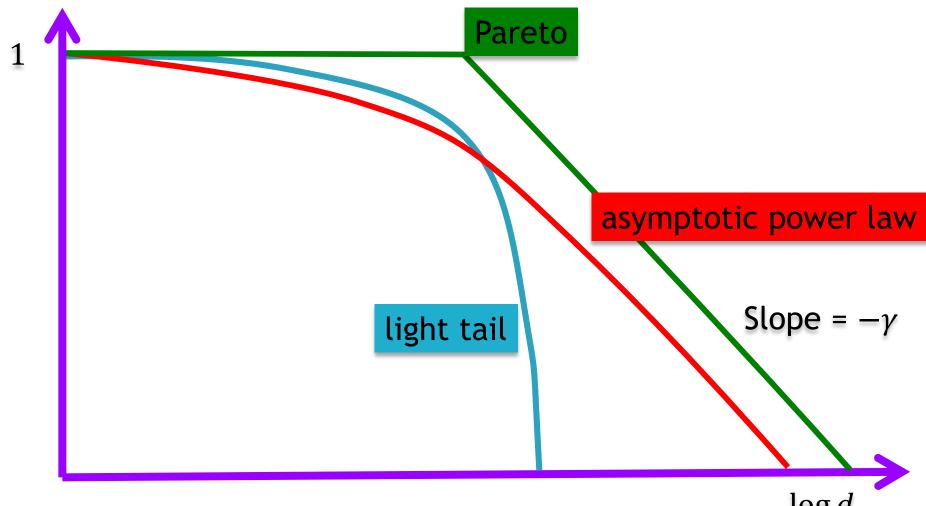


SCIENCE TIP: LOG SCALES ARE FOR QUITTERS WHO CAN'T FIND ENOUGH PAPER TO MAKE THEIR POINT PROPERLY.

Source: xkcd #1162

Log-log distribution plot

• $\log P(D > d)$



Examples of observed power laws

- File sizes on a computer
- Stock market crashes
- Sizes of cities
- Phone call length
- Wealth & income distribution
- Sizes of floods
- Popularity of web pages
- Word frequencies in prose
- Degree distribution in social networks

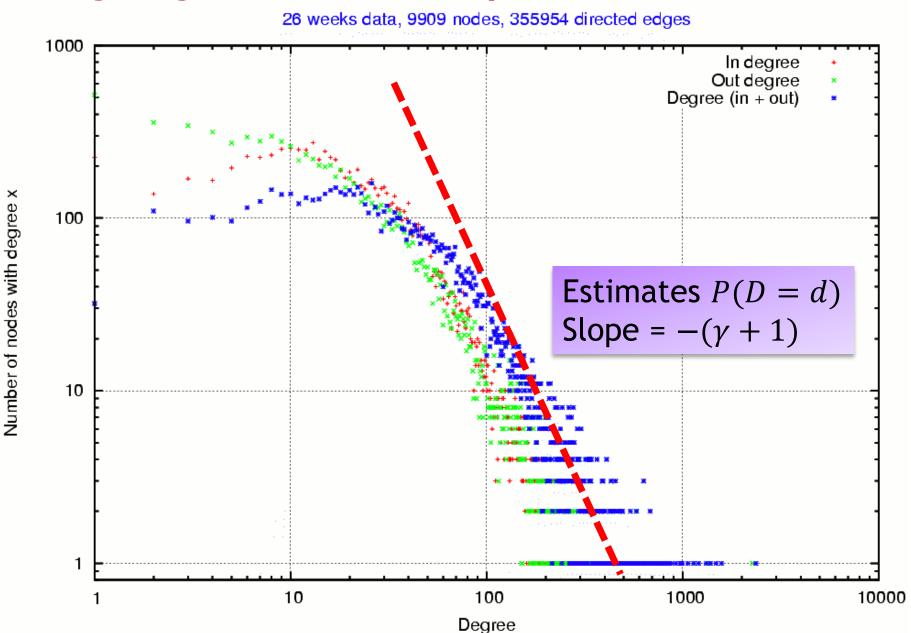
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Why worry about the tail?

 Would you like to sit on a plane engineered under a Gaussian assumption for turbulence?



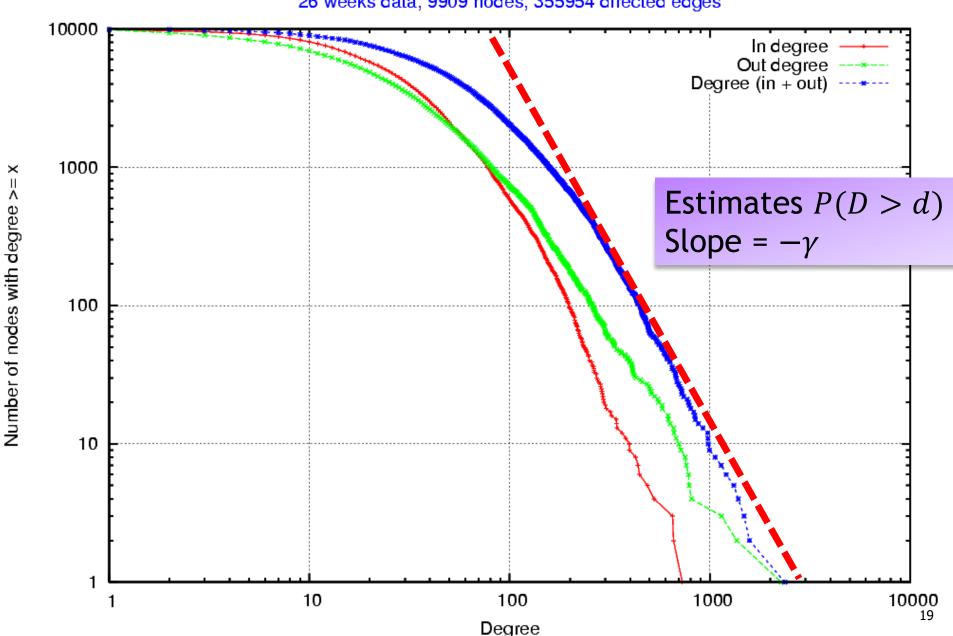
Log-log distribution plot



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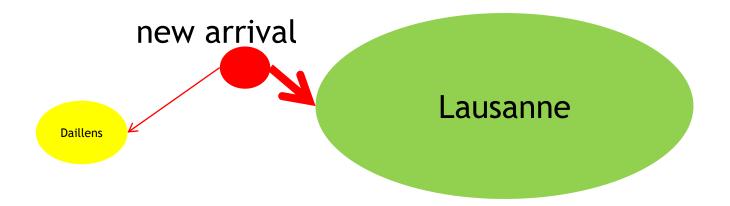
Log-log cumulative plot

26 weeks data, 9909 nodes, 355954 directed edges



One explanation: the rich get richer

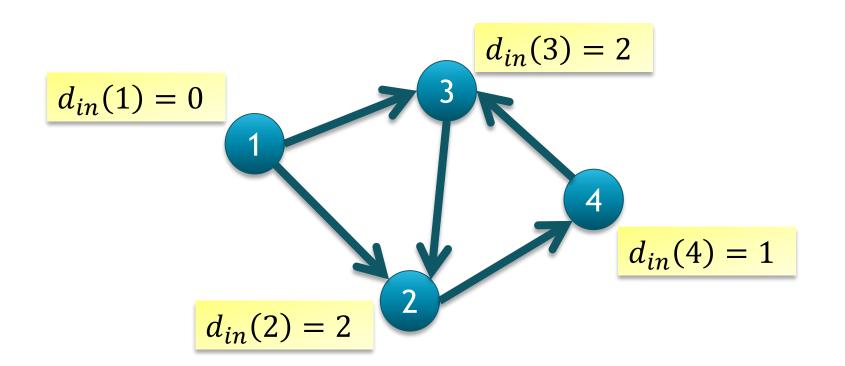
- New arrival in our region: move to Daillens or Lausanne?
 - More likely Lausanne, because more people already there
- City size distribution after many arrivals?



Also: "the first million is the hardest";-)

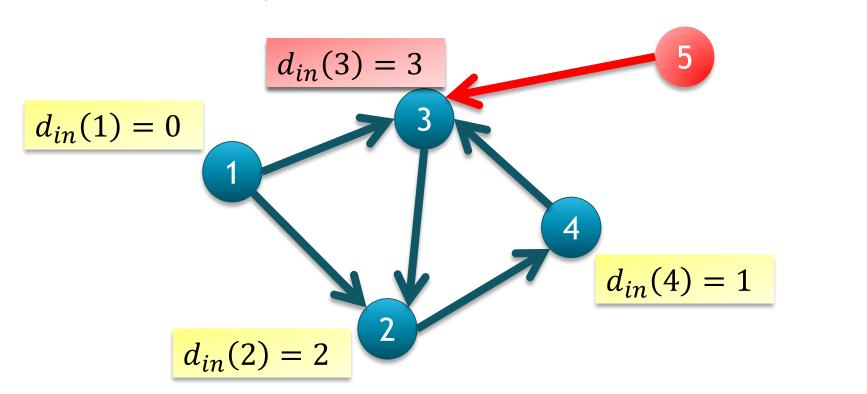
Preferential attachment in growing nets

- Growth model: nodes arrive one by one and join the existing network
 - Directed graph
 - In-degree $d_{in}(v)$ measures "popularity" and "attractiveness" of a node



Preferential attachment

- Preferential attachment: new node creates one edge
- Prob. of connecting to v is $\propto d_{in}(v)$
 - Intuition: high-degree easier to meet; more popular; more useful;...



Preferential attachment

- Node with in-degree 0 never gets "started"
- Need another assumption:
 - With prob. α , new node connects uniformly at random
 - With prob. (1α) , preferential attachment

$$d_{in}(3) = 2$$

$$d_{in}(1) = 0$$

$$P[\text{new edge} = (t+1, v)] = \alpha \frac{1}{t} + (1-\alpha) \frac{d_v}{\sum d_v} = \alpha \frac{1}{t} + (1-\alpha) \frac{d_v}{t}$$

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Pref attachment: analysis

- Evolution of this system:
 - Graph structure only matters through in-degrees
- Markov chain {X_j(t)}: # nodes with in-degree = j at time t
 - Total # of nodes and edges at time t = t
 - Notation: $X_j := X_j(t)$
- Drift:

•
$$P(X_j(t+1)=X_j(t)+1) = \alpha \frac{X_{j-1}}{t} + (1-\alpha)(j-1)\frac{X_{j-1}}{t}$$

prob. of selecting a node uniformly prob. of selecting a node \propto degree

•
$$P(X_j(t+1)=X_j(t)-1) = \alpha \frac{X_j}{t} + (1-\alpha)j \frac{X_j}{t}$$

Pref attachment: analysis

• Combined drift (pretend X_i , $t \in \mathbb{R}$)

- Assume as $t \to \infty$, degree sequence converges $\frac{X_j}{t} \to c_j$, then solve for c_j :
 - c_i : fraction of nodes with degree j

•
$$c_j = \alpha(c_{j-1} - c_j) + (1 - \alpha)((j-1)c_{j-1} - jc_j)$$

$$\frac{c_j}{c_{j-1}} = \frac{\alpha + (1-\alpha)(j-1)}{1+\alpha + (1-\alpha)j} = 1 - \frac{2-\alpha}{1+\alpha + (1-\alpha)j}$$

• Asymptotically for large j, this is $\cong 1 - \frac{2-\alpha}{1-\alpha}j^{-1}$

Pref attachment: analysis

• Note that
$$\left(\frac{j}{j-1}\right)^{-(\gamma+1)} = \left(1 - \frac{1}{j}\right)^{\gamma+1} \sim 1 - \frac{\gamma+1}{j}$$

$$So \gamma = \frac{2-\alpha}{1-\alpha} - 1 = \frac{1}{1-\alpha}$$

Putting together:

- $c_j \propto j^{-(\gamma+1)}$: asymptotic power law
- The stronger the preferential attachment (α smaller), the "heavier" the tail of the degree distribution (γ smaller)
- Arguments can be made rigorous

Network effects and "winner-takes-all"

- Other examples of "rich-get-richer" phenomena:
 - Facebook vs {friendster, sixdegrees, xing,...}
 - Android vs iPhone
 - Technology standards: BluRay,...
- Metcalfe's Law:
 - The value of a network is proportional to n^2
 - Because the value to an individual is proportional to n
- Lock-in
 - Being early is very important

Observer: Friendship Paradox

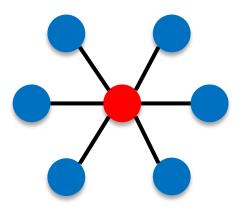
- "Your friends have more friends than you"
- Experiment:
 - Get on facebook and compute the average # friends of your friends
 - How does this compare to your own # friends?

Friendship Paradox

- Formally:
 - Social network = G(V, E)
 - d_v : degree of node v
 - n = |V|: number of nodes, m = |E|: number of edges
- Average number of friends: $\mu = \frac{\sum d_v}{n}$
- How to talk about average number of friends' friends?
 - Natural measure: $\frac{1}{n}\sum_{u\in V}\frac{1}{d_u}\sum_{v\in N_u}d_v$
 - Easier to analyze: degree "seen" by random edge

Friendship Paradox

• Star network (|V| = n):



- Avg degree: $\frac{1}{n}((n-1)+(n-1))=\frac{2(n-1)}{n}\to 2$
- Avg degree of neighbors:

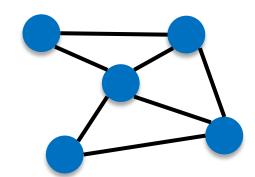
$$\frac{1}{n}\big((n-1)^2+1\big)\to n$$

• Degree of random edge: $\frac{1}{2}(n-1) + \frac{1}{2} = \frac{n}{2}$

Sampling nodes vs sampling edges

Average degree over nodes:

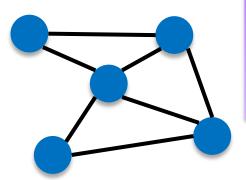
$$\mu = \frac{\sum d_v}{n} = 2\frac{m}{n}$$



"look at each person"

Average degree over edges:

$$\frac{\sum_{(u,v)\in E} d_v}{2m}$$



"look at each person's list of friends"

Friendship Paradox

Lemma:

$$\frac{\sum_{(u,v)\in E} d_v}{2m} = \mu \left(1 + \frac{\sigma^2}{\mu^2}\right)$$

Degree (empirical) variance:

$$\sigma^{2} = \frac{1}{n} \sum_{v \in V} d_{v}^{2} - \left(\frac{1}{n} \sum_{v \in V} d_{v}\right)^{2}$$
$$= \widehat{Var}[d_{v}]$$

Friendship Paradox

Proof:

$$\frac{\sum_{(u,v)\in E} d_v}{2m} =$$

- $= \frac{\sum_{v \in V} d_v^2}{2m} = \text{(because } v \text{ appears } d_v \text{ times in sum over } E)$
- $= \frac{\sum_{v \in V} d_v^2}{\mu n} = \text{(because avg degree is } \mu = 2m/n\text{)}$

$$= \frac{\sigma^2 + \left(\frac{1}{n}\sum_v d_v\right)^2}{\mu}$$

Why is it important?

- Epidemiology:
 - Best protection for a population with a given budget?
 - Assume social network is not knowable globally
- Two strategies:
 - (a) immunize a random set of people
 - (b) immunize random friends of a random set of people
- Friendship Paradox:
 - (b) better than (a)!
 - Bias towards "higher-degree friends"
- Other applications:
 - Finding good monitors, trend-setters, etc.

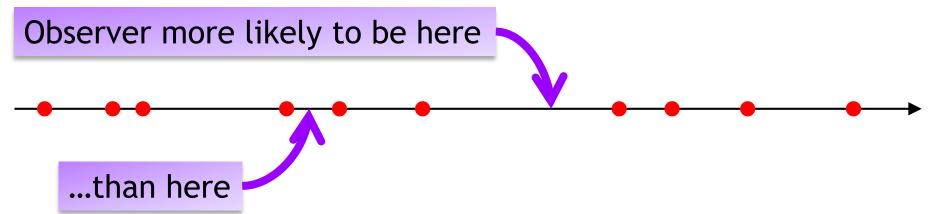
The observer matters

- Other examples:
 - Occupancy distribution:
 - Suppose a train is full 50% of the time, and empty 50% of the time
 - Observer: train is full 100% of the time



The observer matters

- Waiting time:
 - Suppose buses arrive a Poisson(λ) point process
 - Average interarrival interval: 1/λ
 - Observer point of view:
 - Residual time (until next bus): mean = $1/\lambda$
 - Since last bus: mean = $1/\lambda$
 - Mean observed interval length: 2/λ!



Summary and lessons

Herding

- Following others' decisions: natural social mechanism, can lead to suboptimal global behavior
- Information cascades: watching others can wash out primary information
- Rich-get-richer: huge differences in {wealth/degree/influence/membership/...}, winnertakes-all markets

Observing

- Choice of observer sampling bias
- Paradox: average friend is more popular than average individual
- Next week:
 - Processes on networks: epidemics, sampling

References

- [D. Easley and J. Kleinberg, Networks, Crowds, and Markets (chapter 16), 2010]
- [Grossglauser & Thiran, COM-512: Models and Methods for Random Networks (class notes)]