

An Introduction to High Performance Computing on the Minerva Cluster

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Welcome

- ▶ Stuart Rankin — Research Computing User Services Team
- ▶ Paul Browne — Research Computing Platforms Team
- ▶ Course files can be downloaded from: www.csd3.cam.ac.uk
- ▶ Please ask questions and let us know if you need assistance.

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Plan of the Course

Part 1: Basics

Part 2: HPC Facilities

Part 2: Using HPC

10:00 WELCOME

11:30-11:45 Break

13:00-14:00 LUNCH

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16:30 CLOSE

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Part I: Basics

Basics: Why Buy a Big Computer?

What types of big problem might require a “Big Computer”?

Compute Intensive: A single problem requiring a large amount of computation.

Memory Intensive: A single problem requiring a large amount of memory.

Data Intensive: A single problem operating on a large amount of data.

High Throughput: Many unrelated problems to be executed in bulk.

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Basics: Compute Intensive Problems

- ▶ Distribute the **work** for a **single problem** across multiple CPUs to reduce the execution time as far as possible.
- ▶ Program workload must be *parallelised*:
 - Parallel programs split into copies (processes or threads).
 - Each process/thread performs a part of the work on its own CPU, concurrently with the others.
 - A well-parallelised program will fully exercise as many CPUs as there are processes/threads.
- ▶ The CPUs typically need to exchange information rapidly, requiring specialized communication hardware.
- ▶ Many use cases from Physics, Chemistry, Engineering, Astronomy, Biology...
- ▶ The traditional domain of HPC and the Supercomputer.

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Basics: Scaling & Amdahl's Law

- ▶ Using more CPUs is not necessarily faster.
- ▶ Typically parallel codes have a scaling limit.
- ▶ Partly due to the system overhead of managing more copies, but also to more basic constraints;
- ▶ Amdahl's Law (idealized):

$$S(N) = \frac{1}{(1 - p + \frac{p}{N})}$$

where

$S(N)$ is the fraction by which the program has sped up
relative to $N = 1$

p is the fraction of the program which can be parallelized

N is the number of CPUs.

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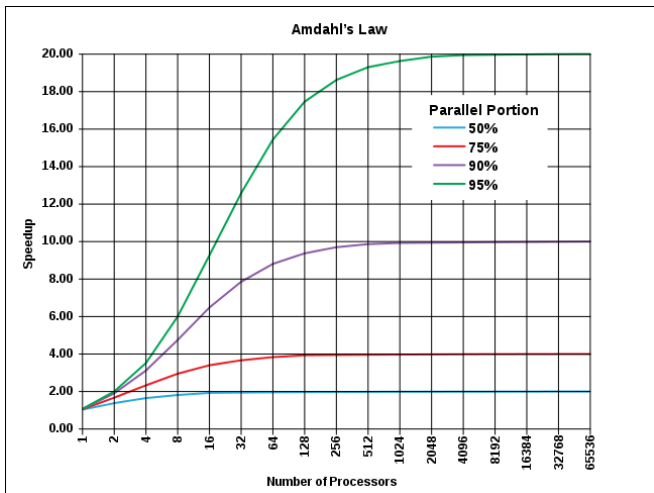
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Basics: Amdahl's Law



<http://en.wikipedia.org/wiki/File:AmdahlsLaw.svg>

The Bottom Line

- ▶ Parallelisation requires effort:
 - ▶ There are libraries to help (e.g. [OpenMP](#), [MPI](#)).
 - ▶ First optimise performance on one CPU, then make p as large as possible.
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Basics: Data Intensive Problems

- ▶ Distribute the **data** for a **single problem** across multiple CPUs to reduce the overall execution time.
- ▶ The *same* work may be done on each data segment.
- ▶ Rapid movement of data to and from disk is more important than inter-CPU communication.
- ▶ Big Data problems of great current interest -
 - ▶ Hadoop/MapReduce
 - ▶ Life Sciences (genomics) and elsewhere.

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Basics: High Throughput

- ▶ Distribute **independent, multiple problems** across multiple CPUs to reduce the overall execution time.
- ▶ Workload is trivially (or *embarrassingly*) parallel:
 - * Workload breaks up naturally into *independent* pieces.
 - * Each piece is performed by a separate process/thread on a separate CPU (concurrently).
 - * Little or no inter-CPU communication.
- ▶ Emphasis is on throughput over a period, rather than on performance on a single problem.
- ▶ Compute intensive capable \Rightarrow high throughput capable
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Basics: Memory Intensive Problems

- ▶ Require aggregation of large memory into a **single system image** (i.e. a single computer running Linux).
- ▶ Technically more challenging to build machines (very fast, low latency interconnection between all CPUs and all memory).
- ▶ Coding/porting easier (memory appears seamless).
- ▶ Optimisation harder (memory is actually highly nonuniform).
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- ▶ Today's commodity servers already aggregate both CPUs and memory to make a single system image in a single box.
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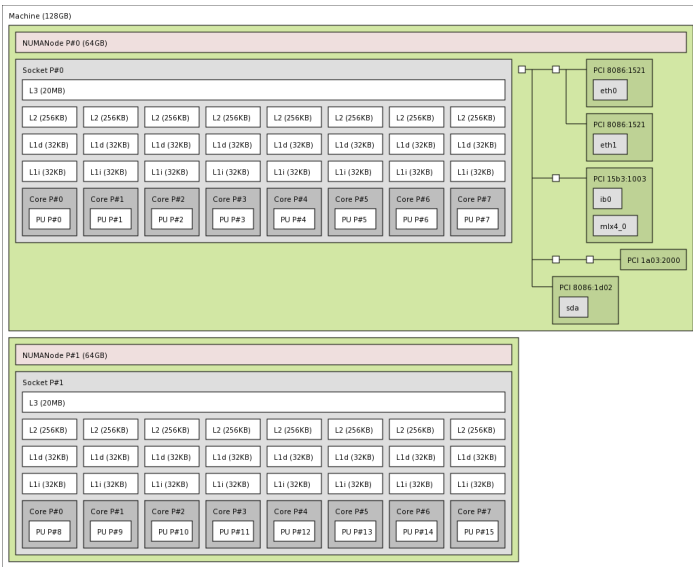
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but users still see a single computer (single system image).

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- ▶ A NUMA, shared memory, multiprocessor building block: a **node**.

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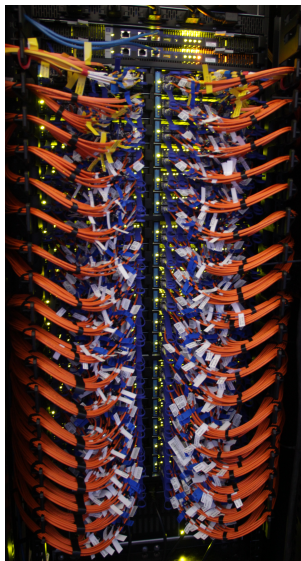
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Faster network is for inter-CPU communication across nodes.

Slower network is for management and provisioning.

Storage may use either.



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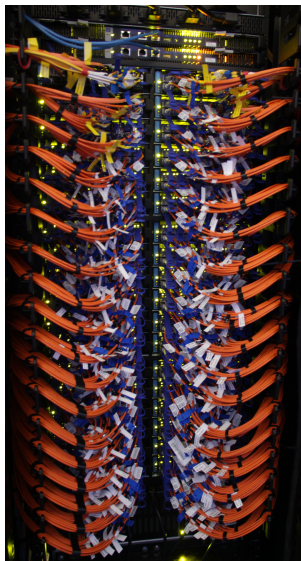
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3. Logically bind the nodes

- ▶ Clusters consist of distinct nodes (i.e. separate Linux computers) on common private network(s) and controlled centrally.
 - * Private networks allow CPUs in different nodes to communicate.
 - * Clusters are distributed memory machines:
Each process/thread sees only its local node's CPUs and memory (without help).
 - * Each process/thread must fit within a single node's memory.
- ▶ More expensive machines logically bind nodes into a single system i.e. CPUs and memory.
 - * E.g. SGI UV.
 - * Private networks allow CPUs to see CPUs and memory in other nodes.
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 - * Logically a single system - 1 big node - but very non-uniform.
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Basics: Programming a Multiprocessor Machine

► Non-parallel (serial) code

- * For a single node as for a workstation.
- * Typically run as many copies per node as CPUs, assuming node memory is sufficient.
- * Replicate across multiple nodes.

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E.g. pthreads, OpenMP.
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Basics: Summary

- ▶ **Why have a supercomputer?**

- ▶ Big single problems, many problems, Big Data.
- ▶ Most current supercomputers are clusters of separate nodes.
- ▶ Each node has multiple CPUs and non-uniform shared memory.
- ▶ Parallel code uses shared memory (pthreads/OpenMP) within a node, distributed memory (MPI) spanning multiple nodes.
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Part II: HPC Facilities

- ▶ Each compute node:
 - * 2x16 cores, Intel Skylake 2.6 GHz
 - * 192 GB RAM
 - * 100 Gb/sec Omni-Path
- ▶ 16 compute nodes
- ▶ 1 high-memory login/head node (minerva-login1.npl.co.uk), 768 GB RAM.
- ▶ 1 GPU node, 2 × NVIDIA P100 GPU, 192 GB RAM.

- ▶ Each compute node:
 - * 32 CPUs
 - * 6 GB per CPU
 - * 10 GB/sec (for MPI and storage)
- ▶ 16 compute nodes
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- ▶ Cambridge Service for Data Driven Discovery
 - ▶ Peta4 — Intel CPU cluster
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- ▶ Each compute node:
 - * 2x16 cores, Intel Skylake 2.6 GHz
 - * 192 GB or 384 GB RAM
 - * 100 Gb/sec Omni-Path
- ▶ 768 compute nodes
- ▶ 8 login nodes (login-cpu.hpc.cam.ac.uk)

- ▶ Each compute node:
 - * 32 CPUs
 - * 6 GB or 12 GB per CPU
 - * 10 GB/sec (for MPI and storage)
- ▶ 768 compute nodes
- ▶ 8 login nodes (login-cpu.hpc.cam.ac.uk)

- ▶ CPUs are **general purpose**
- ▶ Some types of parallel workload fit vector processing well:
 - ▶ Single Instruction, Multiple Data (SIMD)
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- ▶ Each compute node:
 - * 4 × NVIDIA P100 GPU
 - * 1x12 cores, Intel Broadwell 2.2 GHz
 - * 96 GB RAM
 - * 100 Gb/sec (4X EDR) Infiniband.
- ▶ 90 compute nodes.
- ▶ 8 login nodes (login-gpu.hpc.cam.ac.uk).

- ▶ Each compute node:
 - * 4 GPUs
 - * 12 CPUs
 - * 96 GB RAM
 - * 10 GB/sec (for MPI and storage)
- ▶ 90 compute nodes.
- ▶ 8 login nodes (login-gpu.hpc.cam.ac.uk).

Peta4-KNL (Intel Phi)

- ▶ Each compute node:
 - * 64 cores, Intel Phi 7210
 - * 96 GB RAM
 - * 100 Gb/sec Omni-Path
- ▶ 342 compute nodes
- ▶ Shared login nodes with Peta4-Skylake

Peta4-KNL (Intel Phi)

- ▶ Each compute node:
 - * 256 CPUs
 - * 96 GB RAM
 - * 10 GB/sec (for MPI and storage)
- ▶ 342 compute nodes
- ▶ Shared login nodes with Peta4-Skylake

- ▶ Minerva uses NFS to share user directories between all cluster nodes (150 TB).
- ▶ CSD3 uses the Lustre cluster filesystem:
 - * Very scalable, high bandwidth.
 - * Multiple RAID6 back-end disk volumes.
 - * Multiple object storage servers.
 - * Single metadata server.
 - * Tape-backed HSM on newest filesystems.
 - * 12 GB/sec overall read or write.
 - * Prefers big read/writes over small.

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Obtaining an Account and Support

- ▶ Please contact the NPL IT Service Desk:
 - ▶ itservicedesk@npl.co.uk
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Part III: Using HPC

Using HPC: Connecting

- ▶ SSH secure protocol only.

Using HPC: Connecting

- ▶ SSH secure protocol only.
Supports login, file transfer, remote desktop. . .

Connecting: Windows Clients

- ▶ putty, pscp, psftp
<http://www.chiark.greenend.org.uk/~sgtatham/putty/download.html>
- ▶ WinSCP
<http://winscp.net/eng/download.php>
- ▶ TurboVNC (remote desktop, 3D optional)
<http://sourceforge.net/projects/turbovnc/files/>
- ▶ Cygwin
<http://cygwin.com/install.html>
- ▶ MobaXterm
<http://mobaxterm.mobatek.net/>

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<http://sourceforge.net/projects/turbovnc/files/>
- ▶ Cygwin (provides an application environment similar to Linux)
<http://cygwin.com/install.html>
Includes X server for displaying graphical applications running remotely.
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Connecting: Linux/MacOSX/UNIX Clients

- ▶ ssh, scp, sftp, rsync
Installed (or installable).
- ▶ TurboVNC (remote desktop, 3D optional)
<http://sourceforge.net/projects/turbovnc/files/>
- ▶ On MacOSX, install XQuartz to display remote graphical applications.
<http://xquartz.macosforge.org/landing/>

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Connecting: Login

- ▶ From graphical clients:

Host: `minerva-login1.npl.co.uk`

Username: `npl\abc123` (your NPL AD account name)

- ▶ From Linux/MacOSX/UNIX (or Cygwin):

```
ssh -Y npl\\abc12@minerva-login1.npl.co.uk
```

Note the double backslash — this is because UNIX command interpreters treat `\` as special.

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Note the double backslash — this is because UNIX command interpreters treat `\` as special.

Connecting: First time login

- ▶ The first connection to a particular hostname produces the following:

```
The authenticity of host 'minerva-login1.npl.co.uk (139.143.201.10)' can't be established.
```

```
ECDSA key fingerprint is SHA256:k/eB+LjcAfQW56XCzK9QptT0wVWF7j3a/CPxPRd7+lE.
```

```
ECDSA key fingerprint is MD5:18:9a:97:e2:87:4c:07:60:cb:43:46:f2:bb:d8:3d:01.
```

```
Are you sure you want to continue connecting (yes/no)? yes
```

```
Warning: Permanently added 'minerva-login1.npl.co.uk (139.143.201.10)' (ECDSA) to the list of known hosts.
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- ▶ One should always check the fingerprint before typing “yes”.
- ▶ Graphical SSH clients *should* ask a similar question.
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Connecting: File Transfer

- ▶ With graphical clients, connect as before and drag and drop.

- ▶ From Linux/MacOSX/UNIX (or Cygwin):

```
rsync -av old_directory/
```

```
npl\\abc12@minerva-login1.npl.co.uk:hpc-work/new_directory
```

copies contents of `old_directory` to `~/hpc-work/new_directory`.

```
rsync -av old_directory
```

```
npl\\abc12@minerva-login1.npl.co.uk:hpc-work/new_directory
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copies `old_directory` (and contents) to

`~/hpc-work/new_directory/old_directory`.

- * Rerun to update or resume after interruption.
- * All transfers are checksummed.
- * For transfers in the opposite direction, place the remote machine as the first argument.

- ▶ Exercise 3 - File transfer.

Connecting: File Transfer

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`rsync -av old_directory/`

`npl\\abc12@minerva-login1.npl.co.uk:hpc-work/new_directory`

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- ▶ Exercise 3 - File transfer.

Connecting: Remote Desktop

- First time starting a remote desktop:

```
[sjr20@login-a-1 ~]$ vncserver
```

You will require a password to access your desktops.

Password:

Verify:

Would you like to enter a view-only password (y/n)? n

New 'login-a-1:99 (sjr20)' desktop is login-a-1:99

Starting applications specified in /home/sjr20/.vnc/xstartup

Log file is /home/sjr20/.vnc/login-a-1:99.log

- NB Choose a **different** password for VNC.
- The VNC password protects your desktop from other users.
- Remember the unique display number (**99** here) of your desktop.

Connecting: Remote Desktop

- ▶ Remote desktop already running:

```
[sjr20@login-a-1 ~]$ vncserver -list
```

TigerVNC server sessions:

X DISPLAY #	PROCESS ID
:99	130655

- ▶ Kill it:

```
[sjr20@login-a-1 ~]$ vncserver -kill :99  
Killing Xvnc process ID 130655
```

- ▶ Typically you only need **one** remote desktop.
- ▶ Keeps running until killed, or the node reboots.

Connecting: Remote Desktop

- ▶ To connect to the desktop from Linux:

```
vncviewer -via npl\\abc12@minerva-login1.npl.ad.local localhost:99
```

- ▶ The display number 99 will be different in general and unique to each desktop.
- ▶ You will be asked firstly for your AD login password, and secondly for your VNC password.
- ▶ Press F8 to bring up the control panel.
- ▶ Exercise 4 - Remote desktop (from Windows)

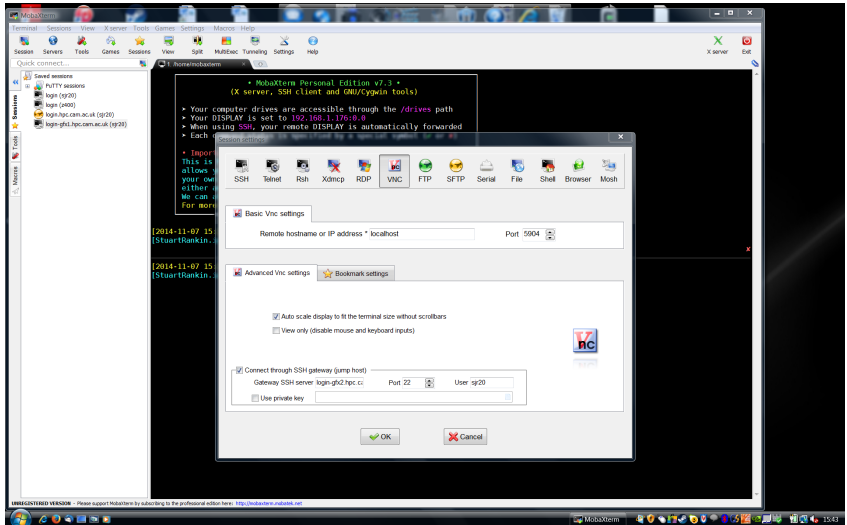
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- ▶ Exercise 4 - Remote desktop (from Windows)

Connecting: Remote Desktop (MobaXterm)



Using HPC: User Environment

- ▶ CentOS Linux 7.4 (Red Hat Enterprise Linux 7.4 rebuild)
 - ▶ bash shell
 - ▶ Gnome or XFCE4 desktop (if you want)
 - ▶ GCC compilers and other development software.
- ▶ But you don't need to know that.

Red Hat Enterprise Linux 7

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Red Hat Enterprise Linux 7

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User Environment: Filesystems

- ▶ `/home/abc12@npl.ad.local`
 - ▶ 50GB quota.
 - ▶ Visible equally from all nodes.
 - ▶ Single storage server.
 - ▶ Regular backups.
 - ▶ Not intended for job outputs or large/many input files.
- ▶ `~/hpc-work`
 - ▶ Visible equally from all nodes.
 - ▶ Larger (1TB initial quota).
 - ▶ Intended for job inputs and outputs.
 - ▶ **Not backed up by default.**

Filesystems: Quotas

► quota

```
[sjr20@login-a-1 ~]$ quota -s
Disk quotas for user sjr20 (uid 1004):
```

Filesystem	space	quota	limit	grace	files	quota	limit	grace
10.44.82.252:/hpc-work	OK	1024G	1126G		1	0	0	
10.44.82.252:/home	13272K	51200M	56320M		345	0	0	

- Aim to stay below the soft limit (*quota*).
- Once over the soft limit, you have 7 days grace to return below.
- When the grace period expires, or you reach the hard limit (*limit*), no more data can be written.
- It is important to rectify an out of quota condition ASAP.

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Filesystems: Permissions

- ▶ Be careful and if unsure, please ask support.
 - ▶ Can lead to accidental destruction of your data or account compromise.
- ▶ Avoid changing the permissions on your home directory.
 - ▶ Files under /home are particularly security sensitive.
 - ▶ Easy to break passwordless communication between nodes.

User Environment: Software

- ▶ Free software accompanying Red Hat Enterprise Linux is (or can be) provided.
- ▶ Other software (free and non-free) is available via modules.
- ▶ Proprietary software currently available includes Matlab and COMSOL.
- ▶ New software may be possible to provide on request.
- ▶ Self-installed software should be properly licensed.
- ▶ *sudo will not work.* (You should be worried if it did.)

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User Environment: Environment Modules

- ▶ Modules load or unload additional software packages.
- ▶ Some are **required** and automatically loaded on login.
- ▶ Others are optional extras, or possible replacements for other modules.
- ▶ **Beware** unloading default modules in `~/.bashrc`.
- ▶ **Beware** overwriting environment variables such as `PATH` and `LD_LIBRARY_PATH` in `~/.bashrc`. If necessary append or prepend.

User Environment: Environment Modules

► Currently loaded:

```
module list
```

```
Currently Loaded Modulefiles:
```

1) dot	3) centos7/global
2) slurm	4) centos7/default-basic

► Available:

```
module av
```

User Environment: Environment Modules

► Whatis:

```
module whatis openmpi-3.0.0-gcc-4.8.5-n2hvjgm  
openmpi-3.0.0-gcc-4.8.5-n2hvjgm: The Open MPI Project is an open source...
```

► Load:

```
module load openmpi-3.0.0-gcc-4.8.5-n2hvjgm
```

► Unload:

```
module unload openmpi-3.0.0-gcc-4.8.5-n2hvjgm
```

- ▶ Matlab

```
module load matlab/r2018a
```

- ▶ Invoking matlab in batch mode:

```
matlab -nodisplay -nojvm -nosplash command
```

where the file `command.m` contains your matlab code.

- ▶ The current site license contains the Parallel Computing Toolbox.

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- ▶ Purge:

```
module purge
```

- ▶ Defaults:

```
module show centos7/default-basic  
module load centos7/default-basic
```

- ▶ Run time environment must match compile time environment.

► GCC

```
gcc -O3 -mtune=native code.c -o prog  
gfortran -O3 -mtune=native code.f90 -o prog
```

```
module load openmpi-3.0.0-gcc-4.8.5-n2hvjgm  
mpicc -O3 -mtune=native mpi_code.c -o mpi_prog  
mpif90 -O3 -mtune=native mpi_code.f90 -o mpi_prog
```

► Exercise 5: Modules and Compilers

- ▶ GCC

```
gcc -O3 -mtune=native code.c -o prog  
gfortran -O3 -mtune=native code.f90 -o prog
```

```
module load openmpi-3.0.0-gcc-4.8.5-n2hvjgm  
mpicc -O3 -mtune=native mpi_code.c -o mpi_prog  
mpif90 -O3 -mtune=native mpi_code.f90 -o mpi_prog
```

- ▶ Exercise 5: Modules and Compilers

Using HPC: Job Submission



Using HPC: Job Submission

- ▶ Compute resources are managed by a scheduler:
SLURM/PBS/SGE/LSF/...
- ▶ Jobs are submitted to the scheduler
 - analogous to submitting jobs to a print queue
 - a file (*submission script*) is copied and queued for processing.

Using HPC: Job Submission

- ▶ Jobs are submitted from the **login node**
 - not itself managed by the scheduler.
- ▶ Jobs may be either non-interactive (**batch**) or **interactive**.
- ▶ Batch jobs run a shell script on the first of a list of allocated nodes.
- ▶ Interactive jobs provide a command line on the first of a list of allocated nodes.

Using HPC: Job Submission

- ▶ Jobs are submitted from the **login node**
 - not itself managed by the scheduler.
- ▶ Jobs may be either non-interactive (**batch**) or **interactive**.
- ▶ **Batch** jobs run a shell script on the first of a list of allocated nodes.
- ▶ **Interactive** jobs provide a command line on the first of a list of allocated nodes.

Using HPC: Job Submission

- ▶ Jobs may use **part** or **all** of one or more nodes
 - the owner can specify `--exclusive` to force exclusive node access.
- ▶ Template submission scripts are available under `~/job_templates`.

Job Submission: Using SLURM

- Prepare a shell script and submit it to SLURM:

```
[abc123@login-a-1]$ sbatch slurm_submission_script  
Submitted batch job 790299
```

Job Submission: Show Queue

- Submitted job scripts are copied and stored in a queue:

```
[abc123@login-a-1]$ squeue -u abc123
```

JOBID	PARTITION	NAME	USER	ST	TIME	NODES	NODELIST(REASON)
790299	skylake	Test3	abc123	PD	0:00	2	(Priority)
790290	skylake	Test2	abc123	R	27:56:10	2	cpu-a-[1,10]

Job Submission: Show Queue

- Submitted job scripts are copied and stored in a queue:

```
[abc123@login-a-1]$ squeue -u abc123
```

JOBID	PARTITION	NAME	USER	ST	TIME	NODES	ODELIST(REASON)
790299	skylake	Test3	abc123	PD	0:00	2	(Resources)
790290	skylake	Test2	abc123	R	27:56:10	2	cpu-a-[1,10]

Job Submission: Show Queue

- Submitted job scripts are copied and stored in a queue:

```
[abc123@login-a-1]$ squeue -u abc123
```

JOBID	PARTITION	NAME	USER	ST	TIME	NODES	NODELIST(REASON)
790299	skylake	Test3	abc123	PD	0:00	2	(AssocGrpCPUMinsLimit)
790290	skylake	Test2	abc123	R	27:56:10	2	cpu-a-[1,10]

Job Submission: Monitor Job

- Examine a particular job:

```
[abc123@login-a-1]$ scontrol show job=790290
```

Job Submission: Cancel Job

- ▶ Cancel a particular job:

```
[abc123@login-a-1]$ scancel 790290
```

Job Submission: Scripts

► SLURM

In `~/job_templates`, see examples: `slurm_submit.skylake.generic`,
`slurm_submit.skylake.matlab`.

```
#!/bin/bash
#! Name of the job:
#SBATCH -J myjob
#! Which project should be charged:
#SBATCH -A NPL-GENERAL-CPU
#! How many whole nodes should be allocated?
#SBATCH --nodes=1
#! How many tasks will there be in total? (<= nodes*32)
#SBATCH --ntasks=1
#! How much wallclock time will be required?
#SBATCH --time=02:00:00
#! Select partition:
#SBATCH -p skylake
...
```

- `#SBATCH` lines are *structured comments*
 - correspond to sbatch command line options.
- The above job will be given 1 cpu on 1 node for 2 hours (by default there is 1 task per node, and 1 cpu per task).

Job Submission: Scripts

► SLURM

In `~/job_templates`, see examples: `slurm_submit.skylake.generic`, `slurm_submit.skylake.matlab`.

```
#!/bin/bash
#! Name of the job:
#SBATCH -J myjob
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#SBATCH -A NPL-GENERAL-CPU
#! How many whole nodes should be allocated?
#SBATCH --nodes=1
#! How many tasks will there be in total? (<= nodes*32)
#SBATCH --ntasks=1
#! How much wallclock time will be required?
#SBATCH --time=02:00:00
#! Select partition:
#SBATCH -p skylake
...
```

- **#SBATCH** lines are *structured comments*
 - correspond to sbatch command line options.
- The above job will be given 1 cpu on 1 node for 2 hours (by default there is 1 task per node, and 1 cpu per task).

Job Submission: Scripts

► SLURM

In `~/job_templates`, see examples: `slurm_submit.skylake.generic`, `slurm_submit.skylake.matlab`.

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#!/bin/bash
#! Name of the job:
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#SBATCH -A NPL-GENERAL-CPU
#! How many whole nodes should be allocated?
#SBATCH --nodes=1
#! How many tasks will there be in total? (<= nodes*32)
#SBATCH --ntasks=1
#! How much wallclock time will be required?
#SBATCH --time=02:00:00
#! Select partition:
#SBATCH -p skylake
...
```

- **#SBATCH** lines are *structured comments*
 - correspond to sbatch command line options.
- The above job will be given **1 cpu** on 1 node for 2 hours (by default there is 1 task per node, and 1 cpu per task).

Job Submission: Scripts

► SLURM

In `~/job_templates`, see examples: `slurm_submit.skylake.generic`,
`slurm_submit.skylake.matlab`.

```
#!/bin/bash
#! Name of the job:
#SBATCH -J myjob
#! Which project should be charged:
#SBATCH -A NPL-GENERAL-CPU
#! How many whole nodes should be allocated?
#SBATCH --nodes=1
#! How many tasks will there be in total? (<= nodes*32)
#SBATCH --ntasks=16
#! How much wallclock time will be required?
#SBATCH --time=02:00:00
#! Select partition:
#SBATCH -p skylake
...
```

- **#SBATCH** lines are *structured comments*
 - correspond to sbatch command line options.
- The above job will be given **16 cpus** on 1 node for 2 hours (by default there is 1 task per node, and 1 cpu per task).

Job Submission: Accounting Commands

- How many core hours available do I have?

```
mybalance
```

User	Usage	Account	Usage	Account Limit	Available (hours)
-----	-----	-----	-----	-----	-----
	+		+		
sjr20	3	SUPPORT-CPU	2,929	22,425,600	22,422,671
sjr20	0	SUPPORT-GPU	0	87,600	87,600

- How many core hours does some other project or user have?

```
gbalance -p SUPPORT-CPU
```

User	Usage	Account	Usage	Account Limit	Available (hours)
-----	-----	-----	-----	-----	-----
	+		+		
pfb29	2,925	SUPPORT-CPU	2,929	22,425,600	22,422,671
sjr20 *	3	SUPPORT-CPU	2,929	22,425,600	22,422,671
...					

(Use -u for user.)

- List all jobs charged to a project/user between certain times:

```
gstatement -p NPL-GENERAL-CPU -u xyz10 -s "2018-04-01-00:00:00" -e "2018-04-30-00:00:00"
```

JobID	User	Account	JobName	Partition	End	ExitCode	State	CompHrs
-----	-----	-----	-----	-----	-----	-----	-----	-----
263	xyz10	support-c+	_interact+	skylake	2018-04-18T19:44:40	0:0	TIMEOUT	1.0
264	xyz10	support-c+	_interact+	skylake	2018-04-18T19:48:07	0:0	CANCELLED+	0.1
275	xyz10	support-c+	_interact+	skylake	Unknown	0:0	RUNNING	0.3
...								

Job Submission: Single Node Jobs

- Serial jobs requiring large memory, or OpenMP codes.

```
#!/bin/bash
...
#SBATCH --nodes=1
#SBATCH --ntasks=1
# Default is 1 task per node
#SBATCH --cpus-per-task=

#SBATCH --mem=5990
# Memory per node in MB - default is pro rata by cpu number
# Increasing --mem or --cpus-per-task implicitly increases the other
...
export OMP_NUM_THREADS= # For OpenMP across cores
$application $options
...
```

Job Submission: Single Node Jobs

- Serial jobs requiring large memory, or OpenMP codes.

```
#!/bin/bash
...
#SBATCH --nodes=1
#SBATCH --ntasks=1
# Default is 1 task per node
#SBATCH --cpus-per-task=

#SBATCH --mem=5990
# Memory per node in MB - default is pro rata by cpu number
# Increasing --mem or --cpus-per-task implicitly increases the other
...
export OMP_NUM_THREADS= # For OpenMP across cores
$application $options
...
```

Job Submission: Single Node Jobs

- Serial jobs requiring large memory, or OpenMP codes.

```
#!/bin/bash
...
#SBATCH --nodes=1
#SBATCH --ntasks=1
# Default is 1 task per node
#SBATCH --cpus-per-task=1
# Default is 1 cpu (core) per task
#SBATCH --mem=5990
# Memory per node in MB - default is pro rata by cpu number
# Increasing --mem or --cpus-per-task implicitly increases the other
...
export OMP_NUM_THREADS= # For OpenMP across cores
$application $options
...
```

Job Submission: Single Node Jobs

- Serial jobs requiring large memory, or OpenMP codes.

```
#!/bin/bash
...
#SBATCH --nodes=1
#SBATCH --ntasks=1
# Default is 1 task per node
#SBATCH --cpus-per-task=1
# Default is 1 cpu (core) per task
#SBATCH --mem=5990
# Memory per node in MB - default is pro rata by cpu number
# Increasing --mem or --cpus-per-task implicitly increases the other
...
export OMP_NUM_THREADS= # For OpenMP across cores
$application $options
...
```

Job Submission: Single Node Jobs

- Serial jobs requiring large memory, or OpenMP codes.

```
#!/bin/bash
...
#SBATCH --nodes=1
#SBATCH --ntasks=1
# Default is 1 task per node
#SBATCH --cpus-per-task=1
# Default is 1 cpu (core) per task
#SBATCH --mem=5990
# Memory per node in MB - default is pro rata by cpu number
# Increasing --mem or --cpus-per-task implicitly increases the other
...
export OMP_NUM_THREADS= # For OpenMP across cores
$application $options
...
```


Job Submission: Single Node Jobs

- Serial jobs requiring large memory, or OpenMP codes.

```
#!/bin/bash
...
#SBATCH --nodes=1
#SBATCH --ntasks=1
# Default is 1 task per node
#SBATCH --cpus-per-task=32 # Whole node

#SBATCH --mem=5990
# Memory per node in MB - default is pro rata by cpu number
# Increasing --mem or --cpus-per-task implicitly increases the other
...
export OMP_NUM_THREADS=32 # For OpenMP across 32 cores
$application $options
...
```

Job Submission: Single Node Jobs

- Serial jobs requiring large memory, or OpenMP codes.

```
#!/bin/bash
...
#SBATCH --nodes=1
#SBATCH --ntasks=1
# Default is 1 task per node
#SBATCH --cpus-per-task=16 # Half node

#SBATCH --mem=5990
# Memory per node in MB - default is pro rata by cpu number
# Increasing --mem or --cpus-per-task implicitly increases the other
...
export OMP_NUM_THREADS=16 # For OpenMP across 16 cores
$application $options
...
```

Job Submission: Single Node Jobs

- Serial jobs requiring large memory, or OpenMP codes.

```
#!/bin/bash
...
#SBATCH --nodes=1
#SBATCH --ntasks=1
# Default is 1 task per node
#SBATCH --cpus-per-task=32 # Whole node

#SBATCH --mem=5990
# Memory per node in MB - default is pro rata by cpu number
# Increasing --mem or --cpus-per-task implicitly increases the other
...
export OMP_NUM_THREADS=16 # For OpenMP across 16 cores (using all memory)
$application $options
...
```

Job Submission: MPI Jobs

- ▶ Parallel job across multiple nodes.

```
#!/bin/bash
...
#SBATCH --nodes=4
#SBATCH --ntasks=128      # i.e. 32x4 MPI tasks in total.
#SBATCH --cpus-per-task=2
...
mpirun -np 128 $application $options
...
```

- ▶ SLURM-aware MPI launches remote tasks via SLURM.
- ▶ The template script uses `$SLURM_TASKS_PER_NODE` to set PPN.

Job Submission: MPI Jobs

- ▶ Parallel job across multiple nodes.

```
#!/bin/bash
...
#SBATCH --nodes=4
#SBATCH --ntasks=64      # i.e. 16x4 MPI tasks in total.
#SBATCH --cpus-per-task=2
...
mpirun -ppn 16 -np 64 $application $options
...
```

- ▶ SLURM-aware MPI launches remote tasks via SLURM.
- ▶ The template script uses `$SLURM_TASKS_PER_NODE` to set PPN.

Job Submission: MPI Jobs

- ▶ Parallel job across multiple nodes.

```
#!/bin/bash
...
#SBATCH --nodes=4
#SBATCH --ntasks=64      # i.e. 16x4 MPI tasks in total.
#SBATCH --cpus-per-task=2
...
mpirun -ppn 16 -np 64 $application $options
...
```

- ▶ SLURM-aware MPI launches remote tasks via SLURM.
- ▶ The template script uses `$SLURM_TASKS_PER_NODE` to set PPN.

Job Submission: Hybrid Jobs

- ▶ Parallel jobs using both MPI and OpenMP.

```
#!/bin/bash
...
#SBATCH --nodes=4
#SBATCH --ntasks=64      # i.e. 16x4 MPI tasks in total.
#SBATCH --cpus-per-task=2
...
export OMP_NUM_THREADS=2  # i.e. 2 threads per MPI task.
mpirun -ppn 16 -np 64 $application $options
...
```

- ▶ This job uses 128 CPUs (each MPI task splits into 2 OpenMP threads).

Job Submission: Hybrid Jobs

- ▶ Parallel jobs using both MPI and OpenMP.

```
#!/bin/bash
...
#SBATCH --nodes=4
#SBATCH --ntasks=64      # i.e. 16x4 MPI tasks in total.
#SBATCH --cpus-per-task=2
...
export OMP_NUM_THREADS=2  # i.e. 2 threads per MPI task.
mpirun -ppn 16 -np 64 $application $options
...
```

- ▶ This job uses 128 CPUs (each MPI task splits into 2 OpenMP threads).

Job Submission: High Throughput Jobs

- ▶ Multiple serial jobs across multiple nodes.
- ▶ Use `srun` to launch tasks (job steps) within a job.

```
#!/bin/bash
...
#SBATCH --nodes=2
...
cd directory_for_job1
srun --exclusive -N 1 -n 1 $application $options_for_job1 > output 2> err &
cd directory_for_job2
srun --exclusive -N 1 -n 1 $application $options_for_job2 > output 2> err &
...
cd directory_for_job64
srun --exclusive -N 1 -n 1 $application $options_for_job64 > output 2> err &
wait
```

- ▶ Exercise 6 - Submitting Jobs.

Job Submission: High Throughput Jobs

- ▶ Multiple serial jobs across multiple nodes.
- ▶ Use `srun` to launch tasks (job steps) within a job.

```
#!/bin/bash
...
#SBATCH --nodes=2
...
cd directory_for_job1
srun --exclusive -N 1 -n 1 $application $options_for_job1 > output 2> err &
cd directory_for_job2
srun --exclusive -N 1 -n 1 $application $options_for_job2 > output 2> err &
...
cd directory_for_job64
srun --exclusive -N 1 -n 1 $application $options_for_job64 > output 2> err &
wait
```

- ▶ Exercise 6 - Submitting Jobs.

Job Submission: High Throughput Jobs

- ▶ Multiple serial jobs across multiple nodes.
- ▶ Use `srunch` to launch tasks (job steps) within a job.

```
#!/bin/bash
...
#SBATCH --nodes=2
...
cd directory_for_job1
srunch --exclusive -N 1 -n 1 $application $options_for_job1 > output 2> err &
cd directory_for_job2
srunch --exclusive -N 1 -n 1 $application $options_for_job2 > output 2> err &
...
cd directory_for_job64
srunch --exclusive -N 1 -n 1 $application $options_for_job64 > output 2> err &
wait
```

- ▶ Exercise 6 - Submitting Jobs.

Job Submission: High Throughput Jobs

- ▶ Multiple serial jobs across multiple nodes.
- ▶ Use `srun` to launch tasks (job steps) within a job.

```
#!/bin/bash
...
#SBATCH --nodes=2
...
cd directory_for_job1
srun --exclusive -N 1 -n 1 $application $options_for_job1 > output 2> err &
cd directory_for_job2
srun --exclusive -N 1 -n 1 $application $options_for_job2 > output 2> err &
...
cd directory_for_job64
srun --exclusive -N 1 -n 1 $application $options_for_job64 > output 2> err &
wait
```

- ▶ Exercise 6 - Submitting Jobs.

Job Submission: High Throughput Jobs

- ▶ Multiple serial jobs across multiple nodes.
- ▶ Use `srunch` to launch tasks (job steps) within a job.

```
#!/bin/bash
...
#SBATCH --nodes=2
...
cd directory_for_job1
srunch --exclusive -N 1 -n 1 $application $options_for_job1 > output 2> err &
cd directory_for_job2
srunch --exclusive -N 1 -n 1 $application $options_for_job2 > output 2> err &
...
cd directory_for_job64
srunch --exclusive -N 1 -n 1 $application $options_for_job64 > output 2> err &
wait
```

- ▶ Exercise 6 - Submitting Jobs.

Job Submission: High Throughput Jobs

- ▶ Multiple serial jobs across multiple nodes.
- ▶ Use `srun` to launch tasks (job steps) within a job.

```
#!/bin/bash
...
#SBATCH --nodes=2
...
cd directory_for_job1
srun --exclusive -N 1 -n 1 $application $options_for_job1 > output 2> err &
cd directory_for_job2
srun --exclusive -N 1 -n 1 $application $options_for_job2 > output 2> err &
...
cd directory_for_job64
srun --exclusive -N 1 -n 1 $application $options_for_job64 > output 2> err &
wait
```

- ▶ Exercise 6 - Submitting Jobs.

Job Submission: Interactive

- ▶ Compute nodes are accessible via SSH while you have a job running on them.

- ▶ Alternatively, submit an interactive job:

```
sintr -A NPL-GENERAL-CPU -N1 -n8 -t 2:0:0
```

- ▶ Within the window (screen session):

- * Launches a shell on the first node (when the job starts).
- * Graphical applications should display correctly.
- * Create new shells with `ctrl-a c`, navigate with `ctrl-a n` and `ctrl-a p`.
- * `ssh` or `srunk` can be used to start processes on any nodes in the job.
- * SLURM-aware MPI will do this automatically.

Job Submission: Interactive

- ▶ Compute nodes are accessible via SSH while you have a job running on them.

- ▶ Alternatively, submit an interactive job:

```
sintr -A NPL-GENERAL-CPU -N1 -n8 -t 2:0:0
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- * Create new shells with `ctrl-a c`, navigate with `ctrl-a n` and `ctrl-a p`.
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Job Submission: Interactive

- ▶ Compute nodes are accessible via SSH while you have a job running on them.

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```

- ▶ Within the window (screen session):

- * Launches a shell on the first node (when the job starts).
- * Graphical applications should display correctly.
- * Create new shells with `ctrl-a c`, navigate with `ctrl-a n` and `ctrl-a p`.
- * `ssh` or `srun` can be used to start processes on any nodes in the job.
- * SLURM-aware MPI will do this automatically.

Job Submission: Array Jobs

- ▶ [http : //slurm.schedmd.com/job_array.html](http://slurm.schedmd.com/job_array.html)
- ▶ Used for submitting and managing large sets of similar jobs.
- ▶ Each job in the array has the same **initial** options.
- ▶ SLURM

```
[abc123@login-a-1]$ sbatch --array=1-7 -A NPL-GENERAL-CPU submit_script
Submitted batch job 791609
```

```
[abc123@login-a-1]$ squeue -u abc123
```

JOBID	PARTITION	NAME	USER	ST	TIME	NODES	NODELIST(REASON)
791609_1	skylake	hpl	abc123	R	0:06	1	cpu-a-6
791609_3	skylake	hpl	abc123	R	0:06	1	cpu-a-16
791609_5	skylake	hpl	abc123	R	0:06	1	cpu-a-7
791609_7	skylake	hpl	abc123	R	0:06	1	cpu-a-7

791609_1, 791609_3, 791609_5, 791609_7

i.e. $\{\text{SLURM_ARRAY_JOB_ID}\}_{-\{\text{SLURM_ARRAY_TASK_ID}\}}$

SLURM_ARRAY_JOB_ID = SLURM_JOBID for the first element.

Job Submission: Array Jobs

- ▶ [http : //slurm.schedmd.com/job_array.html](http://slurm.schedmd.com/job_array.html)
- ▶ Used for submitting and managing large sets of similar jobs.
- ▶ Each job in the array has the same **initial** options.
- ▶ SLURM

```
[abc123@login-a-1]$ sbatch --array=1-7:2 -A NPL-GENERAL-CPU submit_script
Submitted batch job 791609
```

```
[abc123@login-a-1]$ squeue -u abc123
```

JOBID	PARTITION	NAME	USER	ST	TIME	NODES	NODELIST(REASON)
791609_1	skylake	hpl	abc123	R	0:06	1	cpu-a-6
791609_3	skylake	hpl	abc123	R	0:06	1	cpu-a-16
791609_5	skylake	hpl	abc123	R	0:06	1	cpu-a-7
791609_7	skylake	hpl	abc123	R	0:06	1	cpu-a-7

791609_1, 791609_3, 791609_5, 791609_7

i.e. $\{\text{SLURM_ARRAY_JOB_ID}\}_{-\{\text{SLURM_ARRAY_TASK_ID}\}}$

SLURM_ARRAY_JOB_ID = SLURM_JOBID for the first element.

Job Submission: Array Jobs

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- ▶ Used for submitting and managing large sets of similar jobs.
- ▶ Each job in the array has the same **initial** options.
- ▶ SLURM

```
[abc123@login-a-1]$ sbatch --array=1,3,5,7 -A NPL-GENERAL-CPU submit_script
Submitted batch job 791609
```

```
[abc123@login-a-1]$ squeue -u abc123
```

JOBID	PARTITION	NAME	USER	ST	TIME	NODES	NODELIST(REASON)
791609_1	skylake	hpl	abc123	R	0:06	1	cpu-a-6
791609_3	skylake	hpl	abc123	R	0:06	1	cpu-a-16
791609_5	skylake	hpl	abc123	R	0:06	1	cpu-a-7
791609_7	skylake	hpl	abc123	R	0:06	1	cpu-a-7

791609_1, 791609_3, 791609_5, 791609_7

i.e. $\{\text{SLURM_ARRAY_JOB_ID}\}_{\text{SLURM_ARRAY_TASK_ID}}$

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```
[abc123@login-a-1]$ sbatch --array=1,3,5,7 -A NPL-GENERAL-CPU submit_script
Submitted batch job 791609
```

```
[abc123@login-a-1]$ squeue -u abc123
```

JOBID	PARTITION	NAME	USER	ST	TIME	NODES	NODELIST(REASON)
791609_1	skylake	hpl	abc123	R	0:06	1	cpu-a-6
791609_3	skylake	hpl	abc123	R	0:06	1	cpu-a-16
791609_5	skylake	hpl	abc123	R	0:06	1	cpu-a-7
791609_7	skylake	hpl	abc123	R	0:06	1	cpu-a-7

791609_1, 791609_3, 791609_5, 791609_7

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Job Submission: Array Jobs

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- ▶ Used for submitting and managing large sets of similar jobs.
- ▶ Each job in the array has the same **initial** options.
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JOBID	PARTITION	NAME	USER	ST	TIME	NODES	NODELIST(REASON)
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791609_3	skylake	hpl	abc123	R	0:06	1	cpu-a-16
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 - ▶ QoS — payer or non-payer?
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Don't cancel jobs that seem to wait too long.
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Payers with little recent usage receive boost (not implemented yet).
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 - ▶ Promote lower priority jobs into gaps left by higher priority jobs.
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Job Submission: Scheduling Top Dos & Don'ts

▶ Do ...

- ▶ Give reasonably accurate wall times (allows [backfilling](#)).
- ▶ Check your balance occasionally ([mybalance](#)).
- ▶ Test on a small scale first.
- ▶ Implement [checkpointing](#) if possible (reduces resource wastage).

▶ Don't ...

- ▶ Request more than you need
 - you will wait longer and use more credits.
- ▶ Cancel jobs unnecessarily
 - priority increases over time.