

PYQs - Part I

Course on Data Structure



CS & IT Engineering

Data Structure
Hashing





Topics

to be covered

1 Hashing-II



Quadratic Probing

24, 2, 13

$$\left. \begin{array}{l} h(24) = 2 \\ h(2) = 2 \\ h(13) = 2 \end{array} \right\} i=1$$

$$\begin{array}{l} H(24,1) = (h(24) + 1^2) \bmod 11 = 3 \\ H(2,1) = (h(2) + 1^2) \bmod 11 = 3 \\ H(13,1) = (h(13) + 1^2) \bmod 11 = 3 \end{array}$$

$$\begin{array}{l} i=2 \\ H(24,2) = (h(24) + 2^2) \bmod 11 = 6 \\ H(2,2) = (h(2) + 2^2) \bmod 11 = 6 \\ H(13,2) = (h(13) + 2^2) \bmod 11 = 6 \end{array}$$

$$\begin{array}{l} i=3 \\ H(24,3) = (h(24) + 3^2) \bmod 11 = 0 \\ H(2,3) = 0 \end{array}$$

$$\begin{array}{l} i=4 \\ H(24,4) = (h(24) + 4^2) \bmod 11 = 7 \\ H(2,4) = 7 \\ H(13,4) = 7 \end{array}$$

$$\begin{array}{l} i=5 \\ H(24,5) = (h(24) + 5^2) \bmod 11 = 5 \\ H(2,5) = 5 \\ H(13,5) = 5 \end{array}$$

$$\begin{array}{l} i=6 \\ H(24,6) = 5 \\ H(2,6) = 5 \\ H(13,6) = 5 \end{array}$$

$$\begin{array}{l} i=7 \\ H(24,7) = 7 \\ H(2,7) = 7 \\ H(13,7) = 7 \end{array}$$

$$H(13,2) = 0$$

0	X X X X X X X X X X
1	
2	X X X X X X X X X X
3	X X X X X X X X X X
4	
5	X X X X X X X X X X
6	X X X X X X X X X X
7	X X X X X X X X X X
8	
9	
10	

2, 3, 6, 0, 7, 5, 5, 7, 0, 6, 3, 2, ...

24, 2, 13
↓ ↓ ↓
2 2 2

Keys that are hashed to same locations follow the same resolution path bcz of which we are not able to utilize the table size efficiently.

Inspite of almost 50% available space, we are not able to insert a new element.

Double Hashing

Let $h(k)$ is the function

$$h(k) = k \bmod m \Rightarrow \text{Collision.}$$

$$H(k, i) = (h(k) + i) \bmod m \quad \text{L.P}$$

$$= (h(k) + i^2) \bmod m \quad \text{Q.P}$$

$$H(k, i) = (h(k) + i \cdot h'(k)) \bmod m$$

Primary hash

Secondary hash func.

Q - What if the value generated by

$$h'(k) = 0 \quad ?$$

$$H(k, i) = h(k) \bmod m$$

$h'(k)$ never generate 0

Double Hashing

keys: 13, 17, 21, 2, 57, 28, 30, 27

$$h(x) = x \bmod 11$$

$$h'(x) = 7 - (x \bmod 7)$$

$$h(13) = 2$$

$$h(17) = 6$$

$$h(21) = 10$$

$$h(2) = 2^x$$

$$H(2, 1) = (h(2) + 1 \cdot h'(2)) \bmod 11$$

$$h'(2) = 7 - (2 \bmod 7) = 5$$

$$H(2, 1) = (2 + 1 \cdot 5) \bmod 11 = 7$$

$$h(57) = 2^x$$

$$H(57, 1) = (h(57) + 1 \cdot h'(57)) \bmod 11$$

$$h'(57) = 7 - 57 \bmod 7 = 7 - 1 = 6$$

$$H(57, 1) = (2 + 6) \bmod 11 = 8$$

$$h(28) = 6^x$$

$$H(28, 1) = (h(28) + h'(28)) \bmod 11$$

$$= (6 + 7) \bmod 11$$

$$= 2^x$$

$$H(28, 2) = (h(28) + 2 \cdot h'(28)) \bmod 11$$

$$H(28, 2) = (6 + 2 \cdot 7) \bmod 11 = 9$$

$$h(30) = 8^x$$

$$H(30, 1) = (8 + 1 \cdot h'(30)) \bmod 11 = (8 + 5) \bmod 11 = 2$$

$$H(30, 2)$$

$$= (8 + 2 \cdot 5) \bmod 11$$

$$= 7^x$$

$$H(30, 3) = (8 + 3 \cdot 5) \bmod 11 = 1$$

$$h(27) = 5$$

$$\left. \begin{array}{l} h(2) = 2 \\ h(57) = 2 \end{array} \right\} \begin{array}{l} H(2,1) = 7 \\ H(57,1) = 8 \end{array} \quad \begin{array}{l} \text{Problem} \\ \text{*overhead} \end{array}$$

\Rightarrow 2 Hash func.

Computation time

Time Complexity.

Load factor(λ)

$$\lambda = \frac{3}{3}$$

no. of keys

Table size

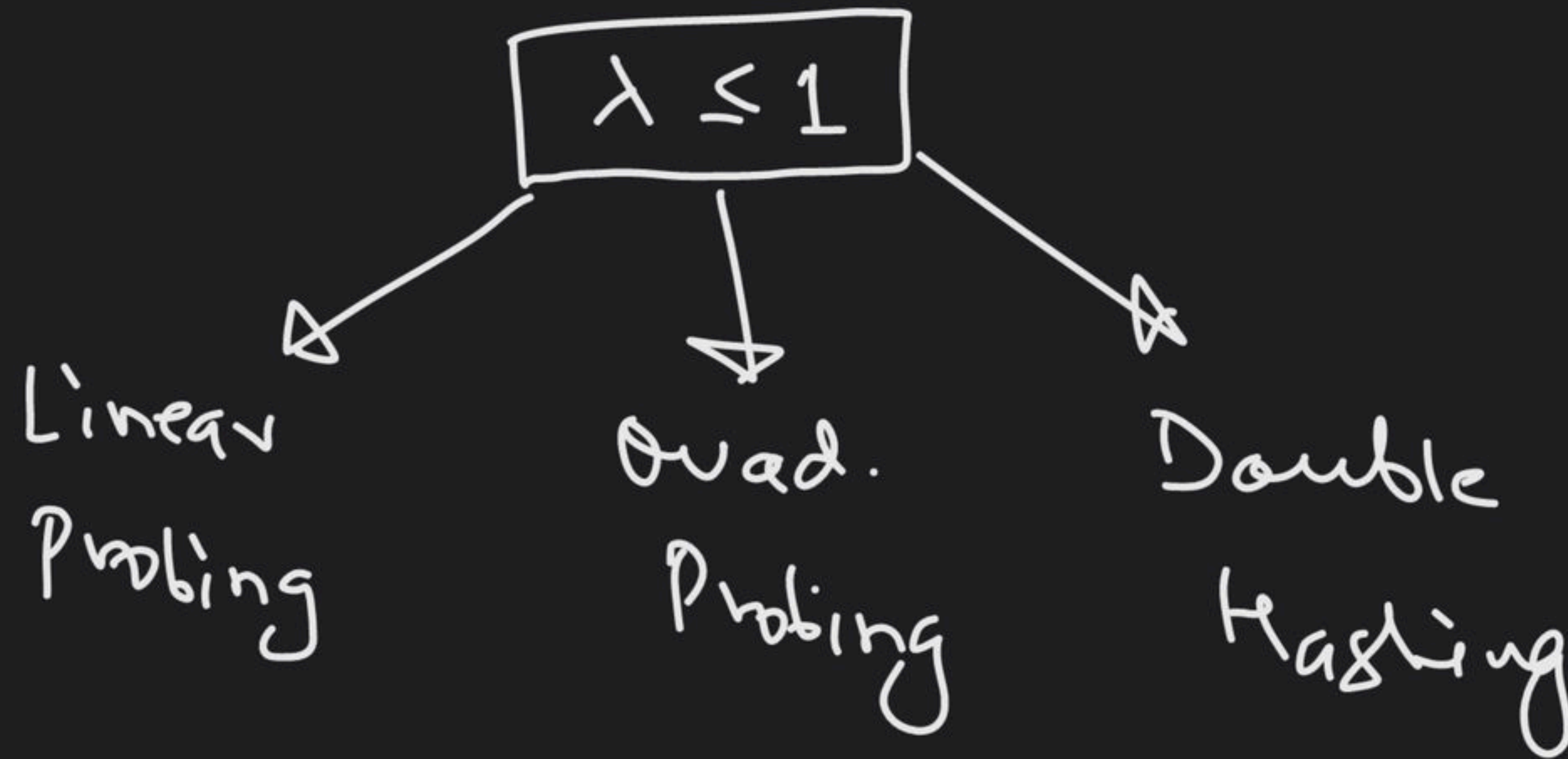
$$n = 20$$

$$m = 40$$

$$\lambda = \frac{20}{40} = \frac{1}{2}$$

$$n = 40$$

$$m = 30$$



Separate chaining $\Rightarrow \lambda > 1$

Collision resolve



Keys: 400, 500, 635, 425, 36, 86, 126, 16

$$m = 10$$

$$h(400) = 400 \bmod 10 \\ = 0$$

$$h(500) = 500 \\ \bmod 10 \\ = 0$$

0	2018	2018 400 NULL
1		
2		
3		
4		
5		
6		
7		
8		
9		

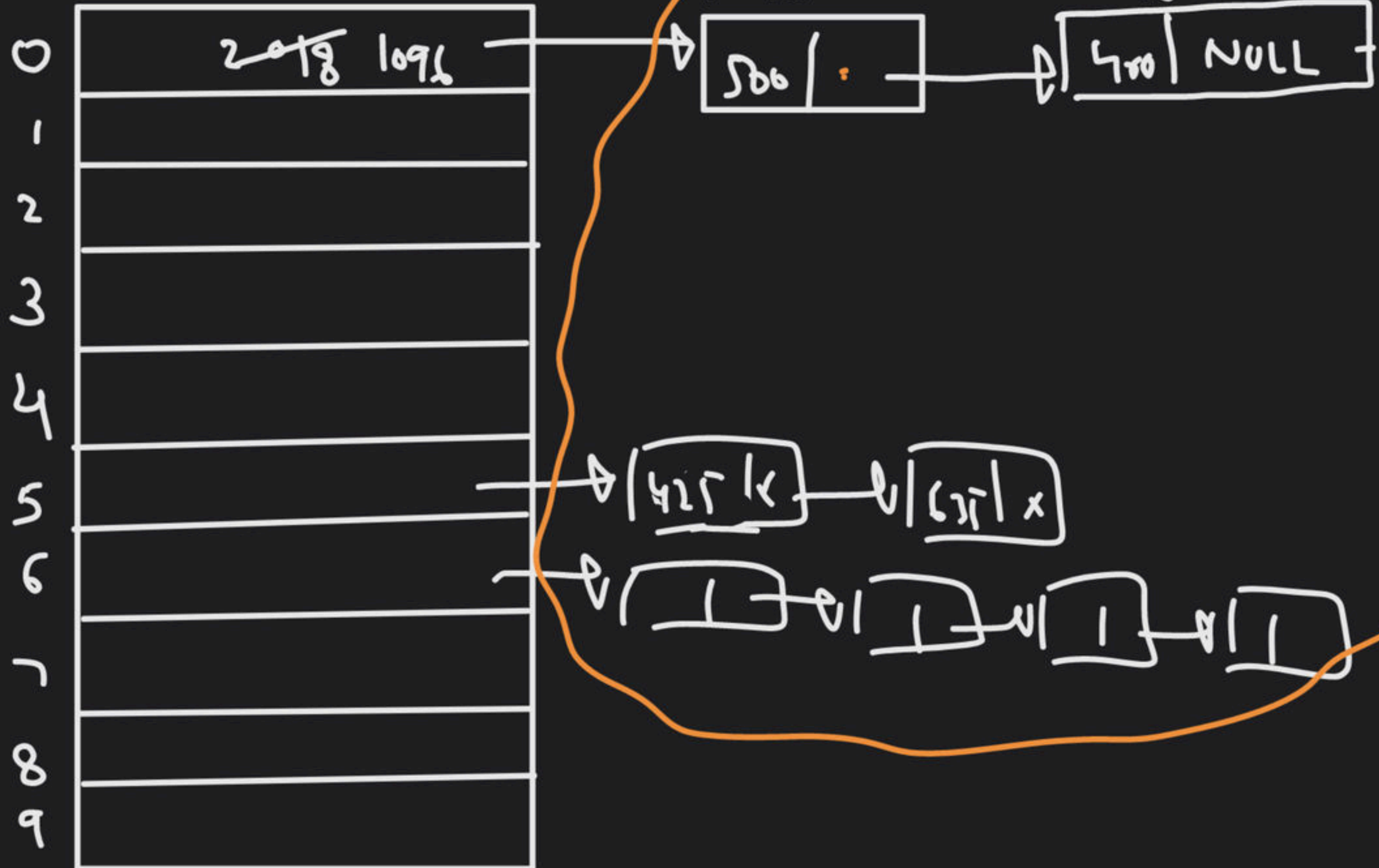
Keys: 400, 500, 635, 425, 36, 86, 126, 16

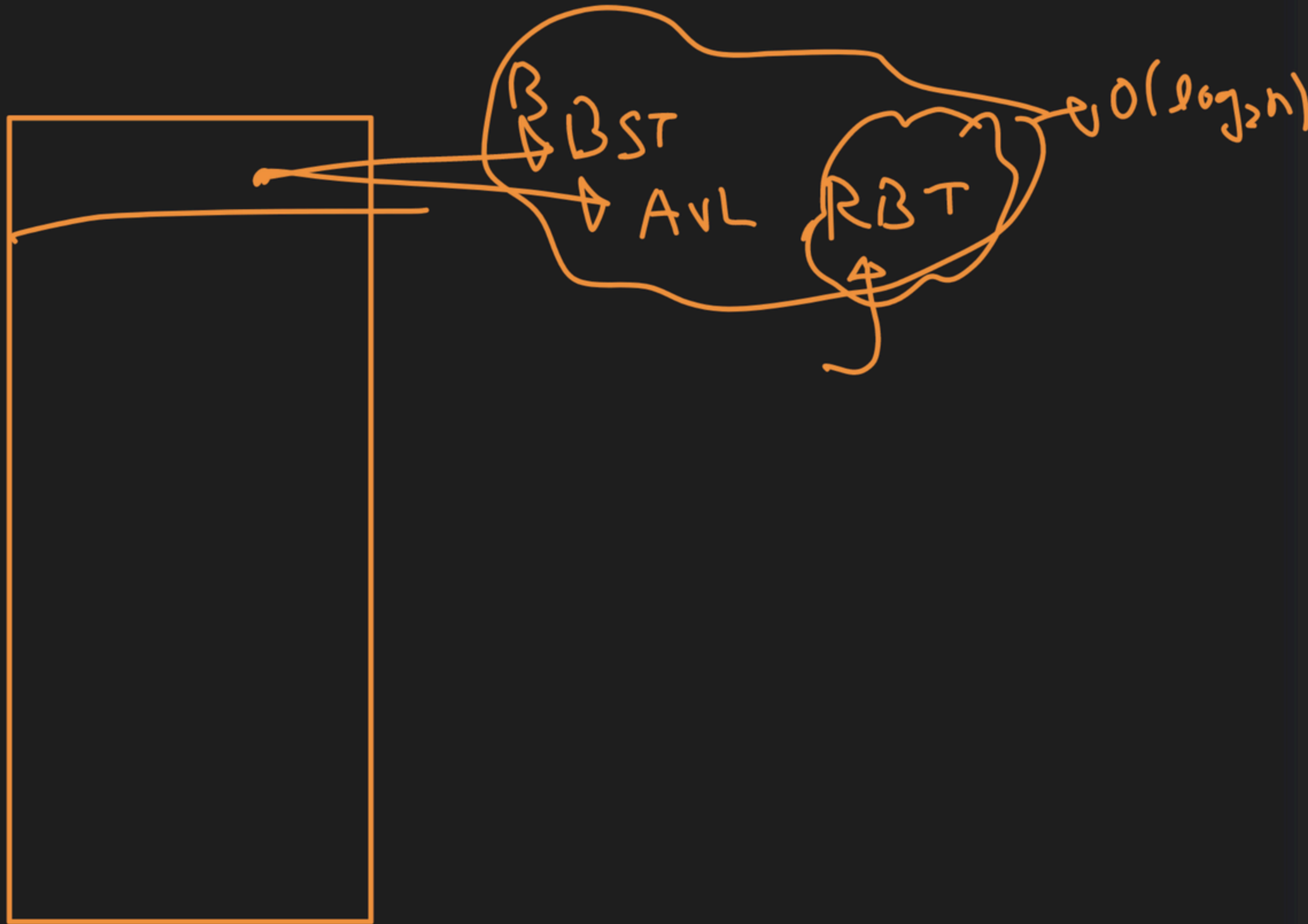
$m = 10$

$$h(400) = 400 \bmod 10 \\ = 0$$

$$h(500) = 500 \\ \bmod 10 \\ = 0$$

$(h+m)$
↓





(L.P)

keys: 31, 20, 43, 27, 34, 12, 46, 14,

58

m=12

delete 20

Search 14

$14 \bmod 12$

$= 2$



after deletion

⇒ Re-hash all remaining keys

0	12
1	58
2	20
3	27
4	14
5	
6	
7	31
8	43
9	
10	34
11	46

Delete



THANK YOU!

Here's to a cracking journey ahead!