

Homework Assignment 5

Due in class, Monday May 1st

SDS 384-11 Theoretical Statistics

1. (VC dimension) Compute the VC dimension of the following function classes
 - (a) Circles in \mathcal{R}^2
 - (b) Axis aligned rectangles in \mathcal{R}^2
 - (c) Axis aligned squares in \mathcal{R}^2
2. In class, you upper bounded the Rademacher complexity of a function class. Now you will derive a lower bound.
 - (a) For function classes \mathcal{F} with function values in $[0, 1]$, prove that $E\|\hat{P}_n - P\|_{\mathcal{F}} \geq \frac{\mathcal{R}_{\mathcal{F}}}{2} - \sqrt{\frac{\log 2}{2n}}$. *Hint: may be it is easier to start from $\mathcal{R}_{\mathcal{F}}$ and show that $\mathcal{R}_{\mathcal{F}} \leq 2E\|\hat{P}_n - P\|_{\mathcal{F}} + \sqrt{\frac{2\log 2}{n}}$. In order to do this, you would need to add and subtract $E[f(X)]$ and then use triangle inequality.*
 - (b) Now prove that $\|P - \hat{P}_n\|_{\mathcal{F}} \geq E\|P - \hat{P}_n\|_{\mathcal{F}} - \epsilon$ with probability at least $1 - \exp(-cn\epsilon^2)$ for some constant c .
 - (c) Recall the class of all subsets with finite size in $[0, 1]$? Prove that then Rademacher complexity of this class is at least $1/2$. What does this imply?
3. In this exercise, we explore the connection between VC dimension and metric entropy. Given a set class \mathcal{S} with finite VC dimension ν , we show that the function class $\mathcal{F}_{\mathcal{S}} := 1_S, S \in \mathcal{S}$ of indicator functions has metric entropy at most

$$N(\delta; \mathcal{F}_{\mathcal{S}}, L^1(P)) \leq \left(\frac{K \log(3e/\delta)}{\delta} \right)^{\nu} \quad \text{For a constant } K \quad (1)$$

Let $\{1_{S_1}, \dots, 1_{S_N}\}$ be a maximal delta packing in the $L^1(P)$ norm, so that:

$$\|1_{S_i} - 1_{S_j}\|_1 = E[|1_{S_i}(X) - 1_{S_j}(X)|] > \delta \quad \text{for all } i \neq j$$

This is an upper bound on the δ covering number.

- (a) Suppose that we generate n samples $X_i, i = 1, \dots, n$ drawn i.i.d. from P . Show that the probability that every set S_i picks out a different subset of $\{X_1, \dots, X_n\}$ is at least $1 - \binom{N}{2}(1 - \delta)^n$.
- (b) Using part (a), show that for $N \geq 2$ and $n = \lceil 2 \log N / \delta \rceil$, there exists a set of n points from which \mathcal{S} picks out at least N subsets, and conclude that $N \leq \left(\frac{3e \log N}{\nu \delta} \right)^{\nu}$.
- (c) Use part (b) to show that Eq (1) holds with $K := 3e^2/(e - 1)$. *Hint: Note that you have $\frac{N^{1/\nu}}{\log N} \leq \frac{3e}{\nu \delta}$. Let $g(x) = x / \log x$. We are solving for $g(m^{1/\nu}) \leq 3e/\delta$. Prove that $g(x) \leq y$ implies $x \leq \frac{e}{e-1} y \log y$.*