

SDS 384 11: Theoretical Statistics

Lecture 14: Uniform Law of Large Numbers- Covering number

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Definitions

- Recall that a metric space (\mathcal{T}, ρ) consists of a nonempty set \mathcal{T} and a mapping $\rho : \mathcal{T} \times \mathcal{T} \rightarrow \mathbb{R}$ that satisfies:
 - Non-negative: $\rho(\theta, \theta') \geq 0$ for all (θ, θ') with equality iff $\theta = \theta'$.
 - Symmetric: $\rho(\theta, \theta') = \rho(\theta', \theta)$ for all pairs (θ', θ) , and
 - Triangle ineq holds: $\rho(\theta, \theta') + \rho(\theta', \theta'') \geq \rho(\theta, \theta'')$
- Examples:
 - $\mathcal{T} = \mathbb{R}^d$, $\rho(\theta, \theta') = \|\theta - \theta'\|_2$
 - $\mathcal{T} = \{0, 1\}^d$ with $\rho(\theta, \theta') = \frac{1}{d} \sum_i 1(\theta_i \neq \theta'_i)$

Covering numbers

Definition

A δ cover of a set \mathcal{T} w.r.t to a metric ρ is a set $\{\theta^1, \dots, \theta^N\}$ such that for every $\theta \in \mathcal{T}$, $\exists i \in [N]$, s.t. $\rho(\theta, \theta^i) \leq \delta$. The δ covering number $N(\delta; \mathcal{T}, \rho)$ is the cardinality of the smallest δ cover.

- We will consider metric spaces which are totally bounded, i.e. $N(\delta; \mathcal{T}, \rho) < \infty$ for all $\delta > 0$.
- The covering number is non-increasing in δ , i.e. $N(\delta) \geq N(\delta')$ for all $\delta < \delta'$
- We are interested in something called Metric entropy, which is the logarithm of the covering number.

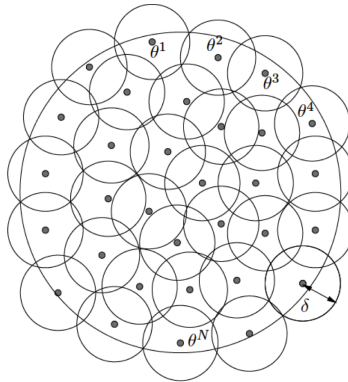


Figure 1: [courtesy: Martin Wainwright's book]

- A δ covering can be thought of as a union of balls with radius δ .

Covering number of a unit cube

Example

Consider the interval $[-1, 1]$ with $\rho(\theta, \theta') = |\theta - \theta'|$. We have
$$N(\delta; [-1, 1], |\cdot|) \leq \frac{1}{\delta} + 1$$

- Divide the interval into L sub-intervals centered at $\theta^i := -1 + (2i - 1)\delta$ for $i \in [L]$ and each of length at most 2δ .
- By construction this is a δ covering.
- So $L \leq 1 + 1/\delta$

Covering the binary hypercube

Example

Consider a d dimensional binary hypercube $\mathcal{T} = \{0, 1\}^d$ with the Hamming metric defined before.

$$\frac{\log N(\delta; \mathcal{T}, \rho)}{\log 2} \leq \lceil d(1 - \delta) \rceil$$

- Let $S = \{1, 2, \dots, \lceil \delta d \rceil\}$
- Consider the set of binary vectors $\mathcal{S}(\delta) := \{\theta \in \mathcal{T} : \theta_j = 0, j \in S(\delta)\}$.
- By construction, for every binary vector $\theta' \in \mathcal{T}$, we can find a vector $\theta \in \mathcal{S}(\delta)$ such that $\rho(\theta, \theta') \leq \delta$
- $N(\delta; \mathcal{T}, \rho) \leq |\mathcal{S}(\delta)| = 2^{\lceil d(1-\delta) \rceil}$

Lower bound on Covering number of the binary hypercube

- Let $\delta \in (0, 1/2)$
- If $\{\theta^1, \dots, \theta^N\}$ is a δ covering, then the (unrescaled) Hamming balls of radius $s = \delta d$ around each θ^ℓ must contain all 2^d vectors.
- Let $s = \lfloor \delta d \rfloor$
- For each θ^i there are exactly $\sum_{j=0}^s \binom{d}{j}$ vectors within δd distance.
- So $N \sum_{j=0}^s \binom{d}{j} \geq 2^d$

Lower bound on Covering number of the binary hypercube

- Let $\delta \in (0, 1/2)$
- So $N \sum_{j=0}^s \binom{d}{j} \geq 2^d$
- Now take a Binomial $(d, 1/2)$ random variable X .
- $P(X \leq \delta d) = \sum_{j=0}^s \binom{d}{j} / 2^d$
- So $N \geq \frac{1}{P(X \leq \delta d)}$
- Using the Hoeffding bound gives: $N \geq \exp(\frac{d}{2}(1/2 - \delta)^2)$
- Using the refined version in your homework gives:
 $N \geq \exp(dKL(\delta||1/2))$

Packing numbers

Definition

An δ -packing of \mathcal{T} w.r.t a metric ρ is a set $\{\theta^1, \dots, \theta^M\}$ such that $\rho(\theta^i, \theta^j) > \delta$ for every distinct pair $i, j \in [M]$. The δ packing number $M(\delta; \mathcal{T}, \rho)$ is the cardinality of the largest δ packing.

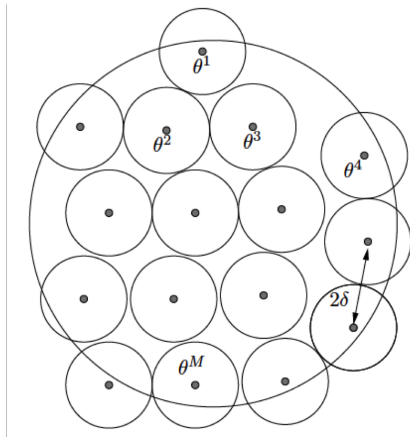


Figure 2: [courtesy: Martin Wainwright's book]

- A 2δ packing can be thought of as a union of balls with radius δ such that no two balls touch.

Relationship between packing and covering numbers

Theorem

For all $\delta > 0$,

$$M(2\delta; \mathcal{T}, \rho) \leq N(\delta; \mathcal{T}, \rho) \leq M(\delta; \mathcal{T}, \rho)$$

- This is saying that packing and covering numbers exhibit the same scaling behavior as $\delta \rightarrow 0$.

- **Upper bound:** Let $V = \{x_1, \dots, x_M\}$ be a δ packing of \mathcal{T} . So for each $y \in \mathcal{T} \setminus V$, $\exists i, \|y - x_i\| \leq \delta$. Otherwise we could have added this point and increased the packing number. So, V is also a δ cover. But since the covering number is the size of the smallest δ covering, the lower bound holds.
- **Lower bound:** Say there is a 2δ packing $\{y_1, \dots, y_M\}$ and a δ covering $\{v_1, \dots, v_n\}$ with $M > n$. Now by pigeonhole, there must be two y_i, y_j who both are in the δ ball around some v_k . But using triangle, we will have $|y_i - y_j| \leq 2\delta$, which is a contradiction. So we must have $m \leq n$.

Covering and Packing numbers-example

Theorem

Let ρ be the Euclidean norm on \mathbb{R}^d . Let $B_1(0)$ be the unit ball centered at the origin (WLOG).

$$\frac{1}{\epsilon^d} \leq N(\epsilon, B_1, \rho) \leq (1 + 2/\epsilon)^d$$

- Consider an ϵ cover $\{\theta^1, \dots, \theta^N\}$. Now,

$$B_1 \subseteq \bigcup_{i=1}^N B_\epsilon(\theta^i)$$

$$\text{vol}(B_1) \leq N \text{vol}(B_\epsilon(\theta^i)) = N \epsilon^d \text{vol}(B_1)$$

$$N \geq 1/\epsilon^d$$

Proof-upper bound

- Consider a ϵ packing $\{\theta^1, \dots, \theta^M\}$
- This is an union of disjoint balls of radius $\epsilon/2$

$$\bigcup_i B_{\epsilon/2}(\theta^i) \subseteq B_{1+\epsilon/2}$$

$$M \text{vol}(B_{\epsilon/2}(\theta^i)) \leq (1 + \epsilon/2) \text{vol}(B_{1+\epsilon/2})$$

$$M(\epsilon/2)^d \text{vol}(B_1) \leq (1 + \epsilon/2)^d \text{vol}(B_1)$$

$$M \leq (1 + 2/\epsilon)^d$$

Suprema over an infinite space

Theorem

Consider a d dimensional vector of independent $\text{subG}(\sigma^2)$ random variables. Let B_d be the unit ball in $\|\cdot\|_2$ norm. Then the following holds:

$$E\left[\sup_{\theta \in B_d} \theta^T X\right] \leq 4\sigma\sqrt{d}$$

Also, for $\delta \in (0, 1)$, with probability $1 - \delta$,

$$\sup_{\theta \in B_d} \theta^T X \leq 4\sigma\sqrt{d} + \sqrt{2\sigma \log(1/\delta)}.$$

Proof of first half

- Let $\mathcal{N}_{1/2}$ be a half covering of B_d . So $N(1/2, B_d, \|\cdot\|_2) \leq 5^d$
- So for each $\theta \in B_d$, $\exists z_\theta \in \mathcal{N}_{1/2}$ such that

$$\theta = z_\theta + x, \quad \|x\| \leq 1/2$$

- So,

$$Y := \sup_{\theta \in B_d} \theta^T X \leq \max_{z_\theta \in \mathcal{N}_{1/2}} z_\theta^T X + \underbrace{\sup_{x \in 1/2 B_d} x^T X}_{Y/2}$$

- Thus, we have:

$$EY \leq 2E \left[\max_{z_\theta \in \mathcal{N}_{1/2}} z_\theta^T X \right] \leq 2\sigma \sqrt{2 \log |\mathcal{N}_{1/2}|} \leq \sigma \sqrt{8d \log 5} \leq 4\sigma \sqrt{d}$$

- We used the same result as last time.

Proof of part 2

$$\begin{aligned} P(Y \geq t) &\leq P(\max_{z \in \mathcal{N}_{1/2}} z^T X \geq t/2) \\ &\leq |\mathcal{N}_{1/2}| P(z^T X \geq t/2) \\ &\leq 5^d \exp(-t^2/8\sigma^2 \|z\|^2) \leq 5^d \exp(-t^2/4\sigma^2) \leq \delta \end{aligned}$$

Solving for t gives, $P\left(Y \geq 2\sigma\sqrt{\log 5} + 2\sigma\sqrt{\log(1/\delta)}\right) \leq \delta$

Example-smoothly parametrized problems

- Consider the following function class parametrized by $\theta \in \Theta$.

$$\mathcal{F} := \{f_\theta(\cdot) : \theta \in \Theta\}$$

- Let $\|\cdot\|_\Theta$ be the norm for θ and $\|\cdot\|_{\mathcal{F}}$ be the norm for \mathcal{F} .
- Say $\|f_\theta(\cdot) - f_{\theta'}(\cdot)\|_{\mathcal{F}} \leq L\|\theta - \theta'\|_\Theta$
- Then $N(\epsilon; \mathcal{F}, \|\cdot\|_{\mathcal{F}}) \leq N(\epsilon/L; \Theta, \|\cdot\|_\Theta)$

Example-smoothly parametrized problems

- A Lipschitz parametrization allows us to go from cover of the Θ space to cover of the f_θ space with a loss of L .
- If \mathcal{F} is parametrized by a compact set of d parameters then $N(\epsilon, \mathcal{F}) = O(1/\epsilon^d)$

A parametric class of Lipschitz continuous functions

Example

For any fixed θ , define the real-valued function $f_\theta(x) := \exp(-\theta x)$, and consider the function class

$$\mathcal{F} = \{f_\theta : [0, 1] \rightarrow \mathbb{R} \mid \theta \in [0, 1]\}$$

Using the uniform norm as a metric, i.e.

$\|f - g\|_\infty := \sup_{x \in [0, 1]} |f(x) - g(x)|$. Prove that

$$\left\lfloor \frac{1 - 1/e}{2\delta} \right\rfloor + 1 \leq N(\delta; \mathcal{F}, \|\cdot\|_\infty) \leq \frac{1}{2\delta} + 2.$$

Proof-upper bound

- First note that $\|f_\theta - f_{\theta'}\|_\infty \leq |\theta - \theta'|$
- For any $\delta \in (0, 1)$, let $T = \lfloor \frac{1}{2\delta} \rfloor$
- Consider $S = \{\theta^0, \dots, \theta^{T+1}\}$ where $\theta^i = 2\delta i$ for $i \leq T$ and $\theta^{T+1} = 1$.
- $\{f_{\theta^i} : \theta^i \in S\}$ is a δ cover for \mathcal{F} .
- For any $\theta \in [0, 1]$ we can find $\theta^i \in S$ such that $|\theta^i - \theta| \leq \delta$
- Indeed we have,

$$\begin{aligned}\|f_{\theta^i} - f_\theta\|_\infty &= \sup_{x \in [0, 1]} |\exp(-\theta^i x) - \exp(-\theta x)| \\ &\leq |\theta^i - \theta| \leq \delta\end{aligned}$$

$$\text{So } N(\delta; \mathcal{F}, \|\cdot\|_\infty) \leq 2 + T \leq 2 + \frac{1}{\delta}$$

Proof-lower bound

- We will do a δ packing.
- Let $\theta^i = -\log(1 - i\delta)$ for $i = 0, \dots, T$
- $-\log(1 - T\delta) = 1$, and so the largest integral value is $T = \lfloor \frac{1 - 1/e}{\delta} \rfloor$
- So $M(\delta; \mathcal{F}, \|\cdot\|_\infty) \geq 1 + \lfloor \frac{1 - 1/e}{\delta} \rfloor$
- $N(\delta; \mathcal{F}, \|\cdot\|_\infty) \geq M(2\delta; \mathcal{F}, \|\cdot\|_\infty) \geq 1 + \lfloor \frac{1 - 1/e}{2\delta} \rfloor$

Example-Lipschitz functions on the unit interval

Example

$$\mathcal{F}_L = \{g : [0, 1] \rightarrow \mathbb{R} \mid g(0) = 0, |g(x) - g(y)| \leq L|x - y|, \forall x, y \in [0, 1]\}$$

Metric entropy scales as $\log N(\delta; \mathcal{F}_L, \|\cdot\|_\infty) \asymp L/\delta$ for small enough $\delta > 0$.

- Its sufficient to consider a sufficiently large packing of \mathcal{F}_L
- For a given ϵ define $M = \lfloor \frac{1}{\epsilon} \rfloor$
- Let $x_i = (i - 1)\epsilon$ for $i = 1, \dots, M + 1$
-

$$\phi(x) := \begin{cases} 0 & x < 0 \\ x & x \in [0, 1] \\ 1 & x > 1 \end{cases} \quad (1)$$

- Define $f_\beta(y) = \sum_{i=1} \beta_i L \epsilon \phi\left(\frac{y - x_i}{\epsilon}\right)$ for $\beta \in \{-1, 1\}^M$

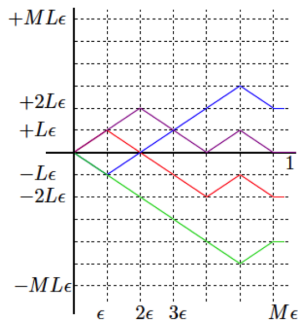


Figure 5-2. The function class $\{f_\beta, \beta \in \{-1, +1\}^M\}$ used to construct a packing of the Lipschitz class \mathcal{F}_L . Each function is piecewise linear over the intervals $[0, \epsilon]$, $[\epsilon, 2\epsilon]$, \dots , $[(M-1)\epsilon, M\epsilon]$ with slope either $+L$ or $-L$. There are 2^M functions in total, where $M = \lceil 1/\epsilon \rceil$.

- For any pair $\beta \neq \beta' \in \{-1, 1\}^M$ there is at least one interval where they have the same starting point.
- So $\|f_\beta(y) - f_{\beta'}(y)\|_\infty \geq 2L\epsilon$
- $f_\beta \in \mathcal{F}_L$ for all $\beta \in \{-1, 1\}^M$
- So f_β forms a $2L\epsilon$ packing.
- Making $\epsilon L = \delta$ we see

$$N(\delta; \mathcal{F}_L, \|\cdot\|_\infty) \geq M(2L\epsilon; \mathcal{F}_L, \|\cdot\|_\infty) = 2^{\lfloor \frac{1}{\epsilon} \rfloor} = 2^{\lfloor \frac{L}{\delta} \rfloor}$$

- Also the set f_β also form a suitable covering of the original functions, and this gives the upper bound.

- The last example can be extended to Lipschitz functions on the Unit cube in higher dimensions, i.e.

$$|f(x) - f(y)| \leq \|x - y\|_{\infty} \quad \text{for all } x, y \in [0, 1]^d$$

- The same method can be used to show that the metric entropy for this class is the same order as $(L/\delta)^d$

Make a comparison

- Recall that for a L Lipschitz continuous functions supported on $[0, 1]$ with $f(0) = 0$, the metric entropy was L/δ
- Also recall that for a L Lipschitz continuous functions supported on $[0, 1]^d$ with $f(0) = 0$, the metric entropy was $(L/\delta)^d$
- However for a given function class like the last one the metric entropy is $\log(1/\delta)$
- Recall that for Unit hypercubes in d dimensions the metric entropy is $d \log(1 + 1/\delta)$
- Note that for Lipschitz continuous functions the dependence on d is exponential. This is a much richer class of functions, so the size is considerably larger and scales poorly with d .

Acknowledgment

This lecture was very much based on Martin Wainwright's book.

