

LINFO 1121 DATA STRUCTURES AND ALGORITHMS



TP3: Arbres de recherche

	Sequential	Binary
Insertion	O(n)	O(n) O(n/2) moyen
Recherche	O(n)	O(log n)

n : nombre d'éléments dans l'array

seq	=

bin =

	Sequential	Binary
Insertion	O(n)	O(n) O(n/2) moyen
Recherche	O(n)	O(log n)

P: Nombre de put

G: Nombre de get

Coûts liés aux puts

$$seq = (\sum_{i=1}^{P} i)$$

$$bin = (\sum_{i=1}^{P} i)$$

	Sequential	Binary
Insertion	O(n)	O(n) O(n/2) moyen
Recherche	O(n)	O(log n)

P: Nombre de put

G: Nombre de get

Coûts liés aux puts

+ coûts liés aux gets

$$seq = (\sum_{i=1}^{P} i) + GP$$

$$bin = (\sum_{i=1}^{P} i) + G \log P$$

	Sequential	Binary
Insertion	O(n)	O(n) O(n/2) moyen
Recherche	O(n)	O(log n)

P: Nombre de put

G: Nombre de get

$$seq = (\sum_{i=1}^{P} i) + GP$$

$$bin = (\sum_{i=1}^{P} i) + G \log P$$

$$GP > G \log P$$

Rappels sur les BSTS

Propriété d'un Binary Search Tree

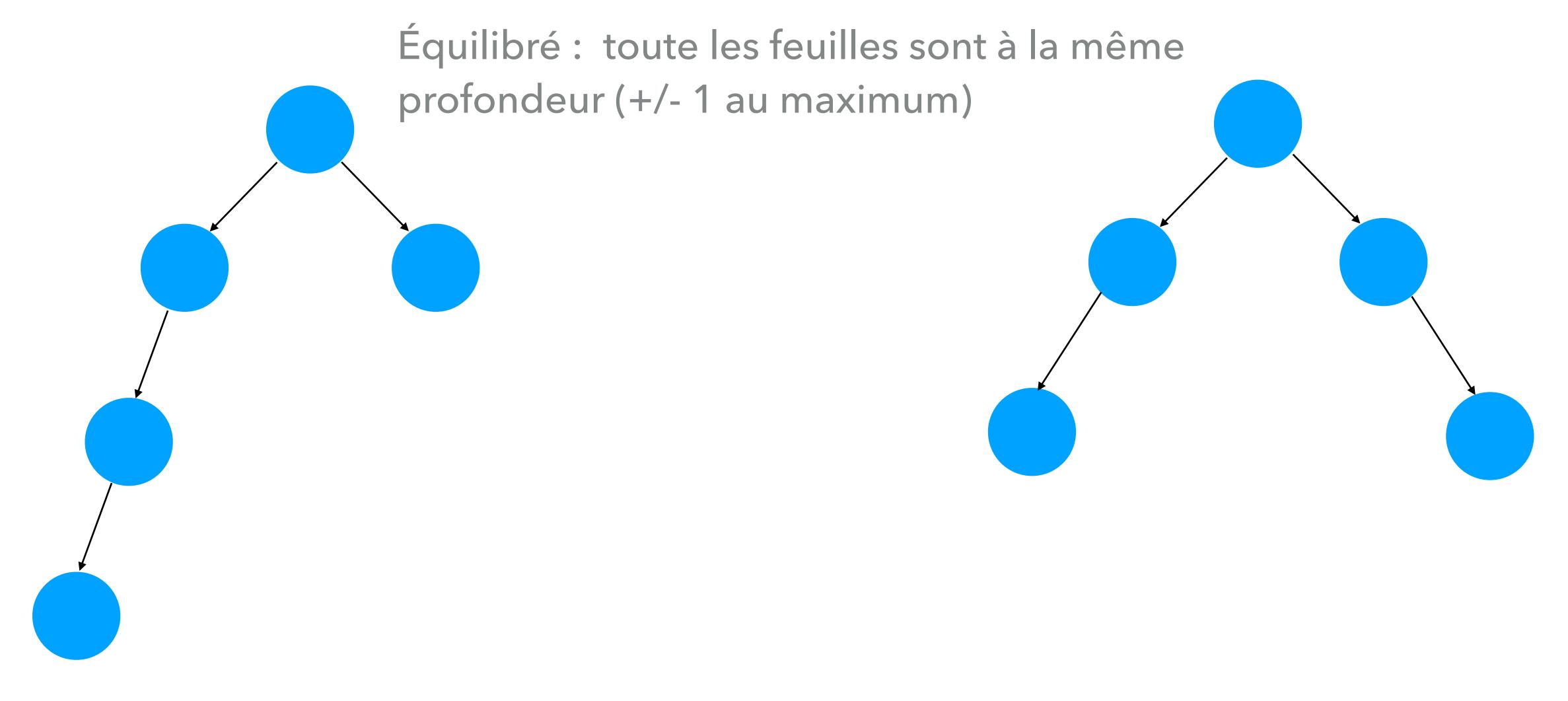
C'est un arbre :

- C'est un graphe
- Connexe
- Possédant n-1 arêtes pour n noeuds
- Pas de cycles
- Au plus un chemin entre toute paire de noeuds

Il est binaire:

- Deux enfants par noeuds maximum
- Les noeuds en dessous, et à gauche doivent avoir des clés plus petites
- Les noeuds en dessous, et à droite doivent avoir des clés plus grandes

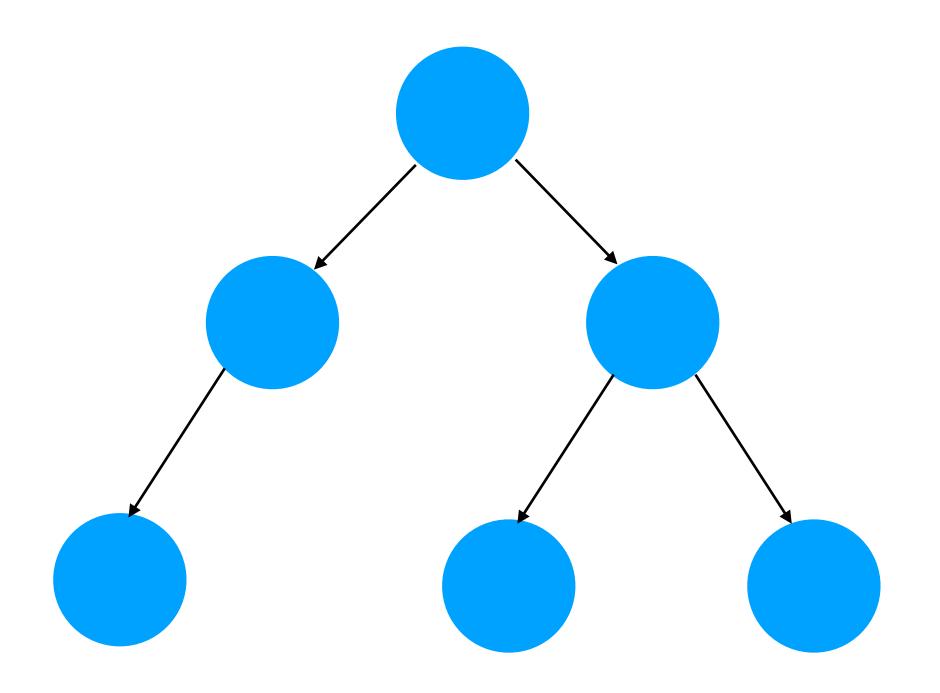
Équilibre d'un arbre



Non équilibré

Équilibré

Équilibre parfait d'un arbre



Parfaitement équilibré si taille des ces chemins est de $log_2(n)(+/-1)$.

3.1.3 Interpolation Search

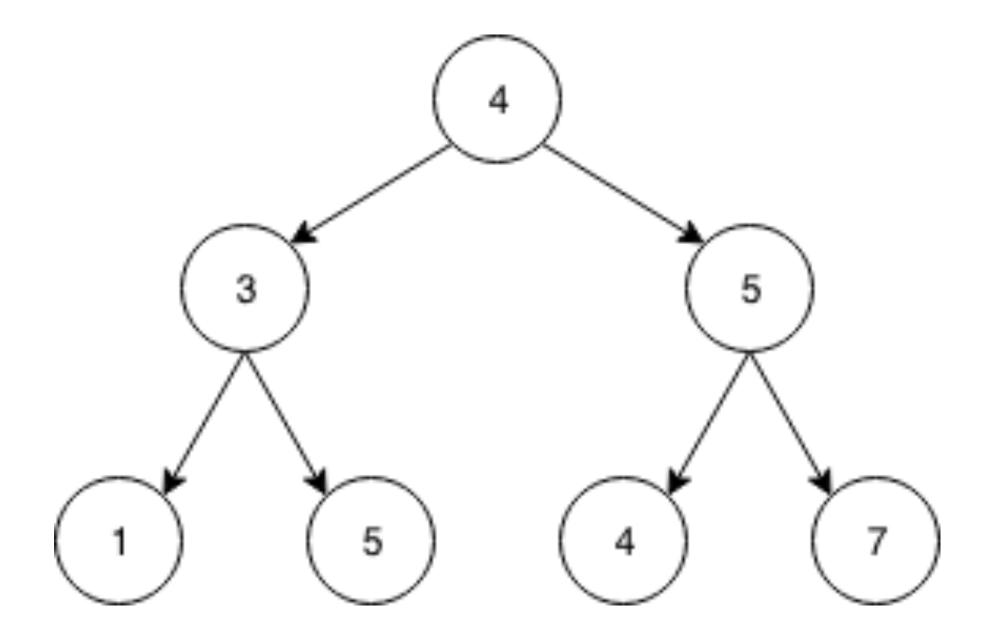
```
public int rank(Key key) {
    Key lo = min(Keys)
    Key hi = max(Keys)
     while (lo <= hi) {
       int index = floor((key - lo)/(hi - lo));
       int cmp = key.compareTo(keys[index]);
             (cmp < 0) hi = index - 1;
       else if (cmp > 0) lo = index + 1;
       else return mid;
     return lo;
```

3.1.4 Caching key

```
public class BinarySearchST <Key extends Comparable<Key>, Value> {
    private Key[] keys;
    private Value[] vals;
    private int N;
    private Key cacheKey;
    private int cacheRank;
```

```
public int rank(Key key) {
  if (key.compareTo(cacheKey) == Q)
     return cacheRank;
  int lo = 0, hi = N-1;
  while (lo \leq hi) {
    int mid = lo + (hi - lo) / 2;
    int cmp = key.compareTo(keys[mid]);
         (cmp < 0) hi = mid - 1;
    else if (cmp > 0) lo = mid + 1;
    else {
       cacheKey = key;
       cacheRank = mid;
       return mid;
  return lo;
```

Question 3.1.5 isBST()



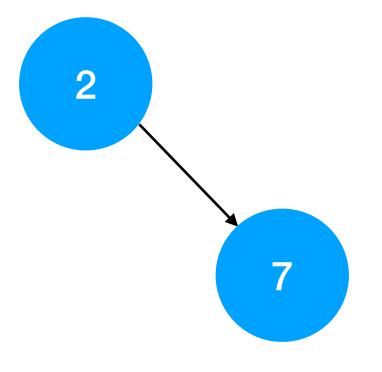
Question 3.1.5 isBST()

```
private boolean isBST() {
    return isBST(root, null, null);
}

private boolean isBST(Node x, Key min, Key max) {
    if (x == null) return true;
    if (min != null && x.key.compareTo(min) <= 0) return false;
    if (max != null && x.key.compareTo(max) >= 0) return false;
    return isBST(x.left, min, x.key) && isBST(x.right, x.key, max);
}
```

- 10,9,8,7,6,5
- •4,10,8,6,5
- 1,10,2,9,3,8,4,7,6,5
- 2,7,3,8,4,5
- 1,2,10,4,8,5



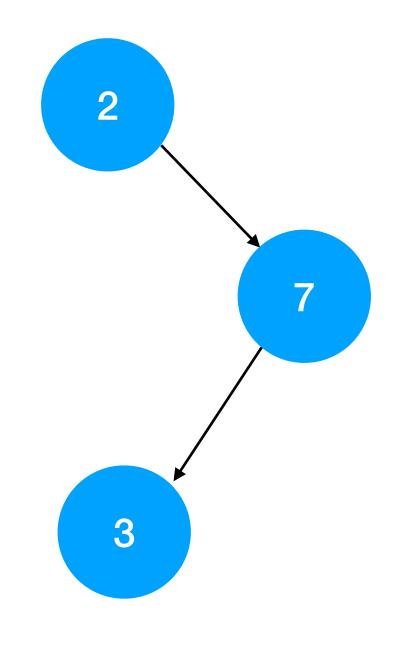


$$LB = -\infty$$

$$UB = +\infty$$

$$7 > 2 \rightarrow LB = 2$$



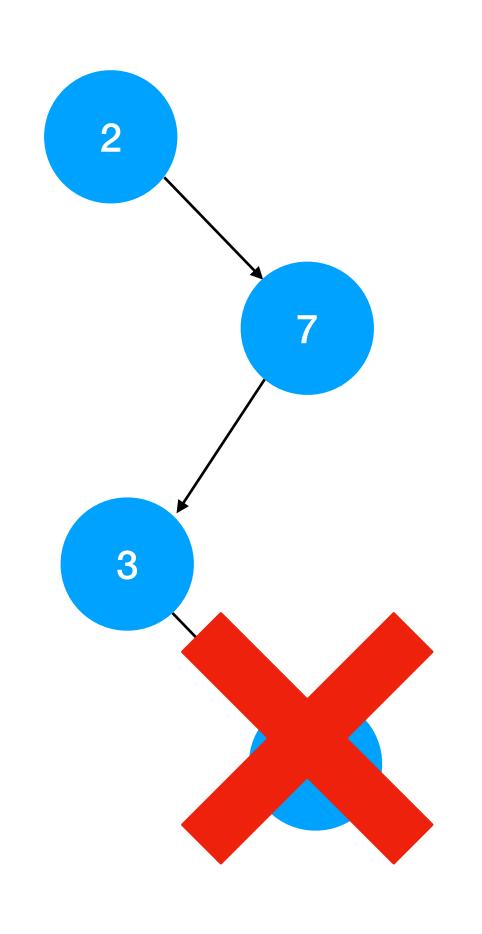


$$LB = 2$$

$$UB = +\infty$$

$$3 < 7 \rightarrow UB = 7$$





$$LB = 2$$

$$UB = 7$$

$$8 > 7 \rightarrow 8 > UB$$

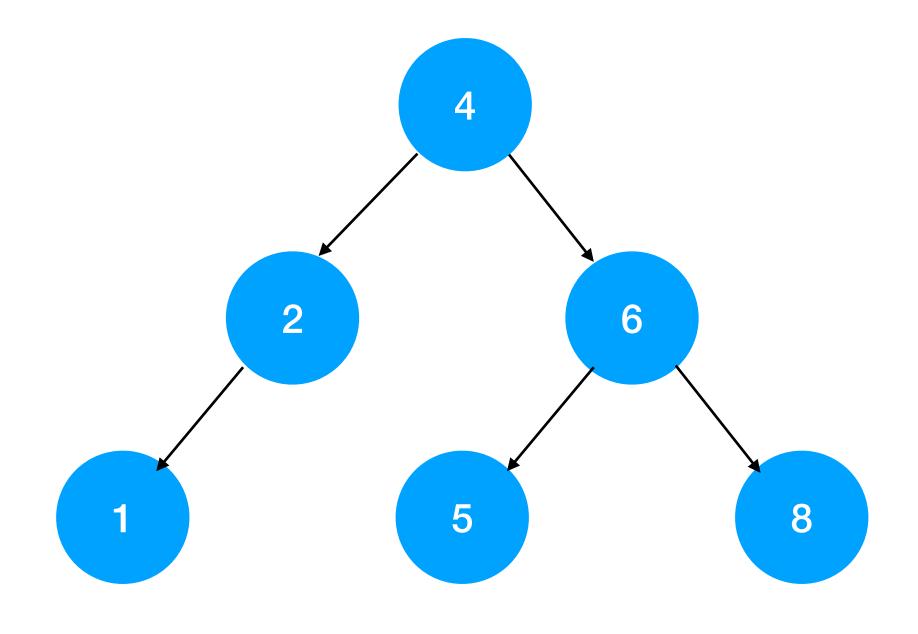
- 10,9,8,7,6,5
- 4,10,8,6,5
- 1,10,2,9,3,8,4,7,6,5
- 2,7,3,8,4,5
- 1,2,10,4,8,5

Question 3.1.8 Énumerer en ordre croissant les clés

Parcours infixe

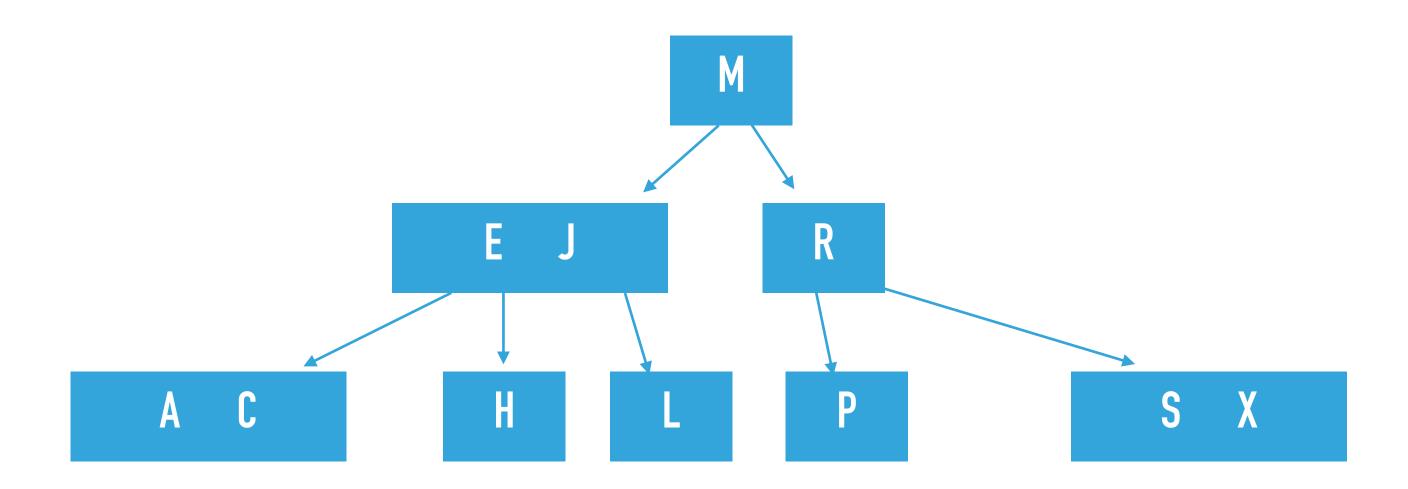
Complexité : O(nombre de noeuds)

```
public static void enumerate(Node parent) {
    enumerate(parent.leftChild)
    System.out.println(parent.key)
    enumerate(parent.rightChild)
}
```



Arbres 2-3

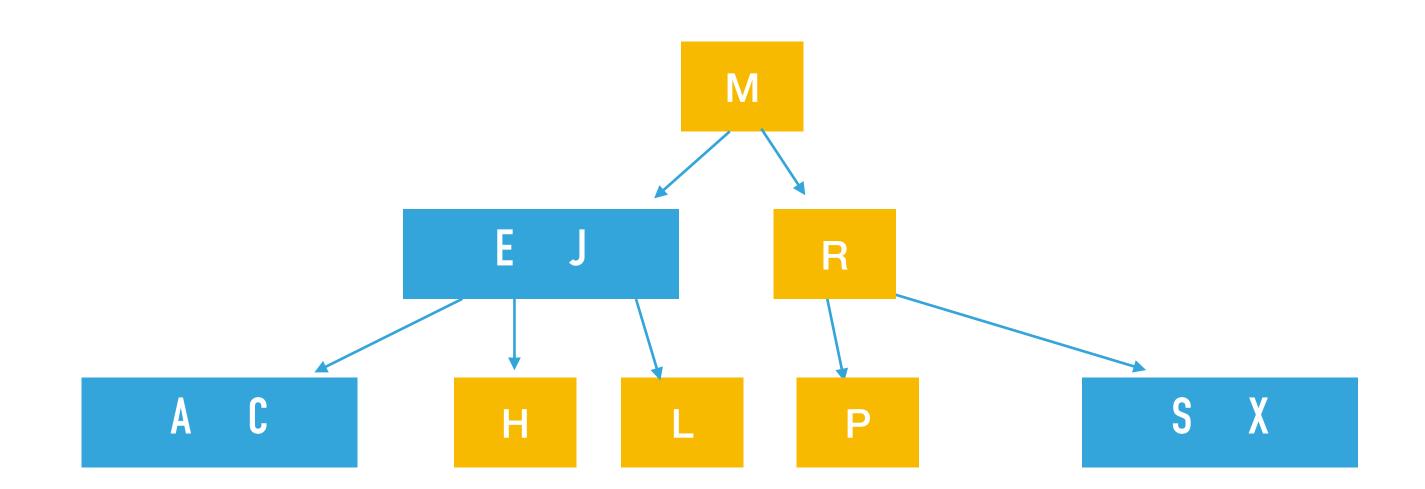
Arbre 2-3 contient soit:



Arbres 2-3

Arbre 2-3 contient soit:

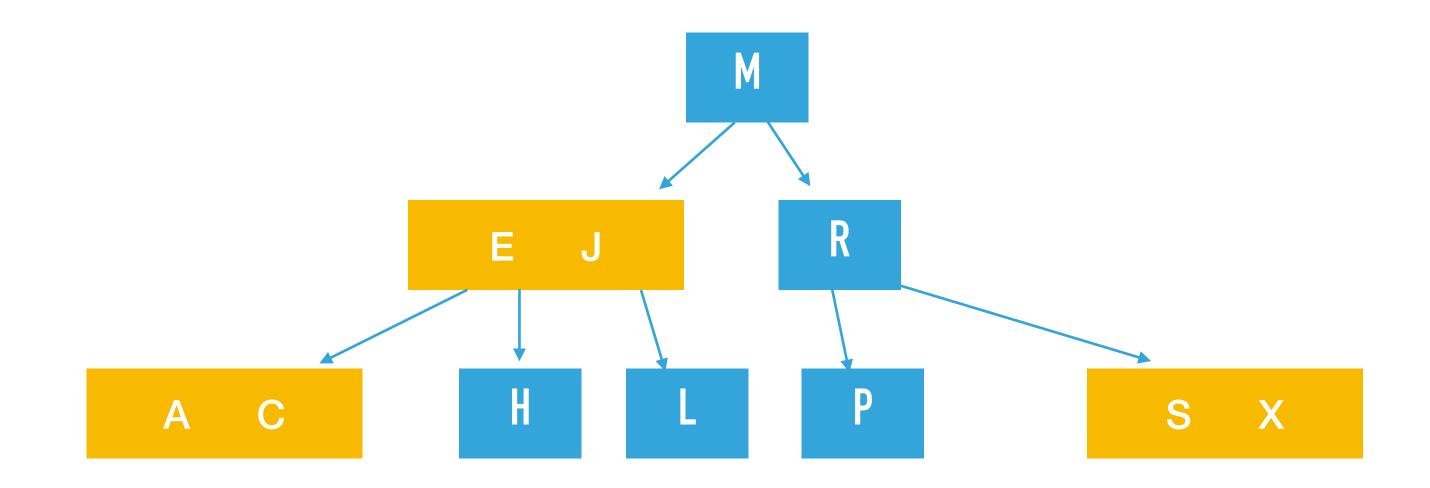
- 2-noeud, avec 1 clé et deux liens
- À gauche : clés plus petites
- À droite : clés plus grandes



Arbres 2-3

Arbre 2-3 contient soit:

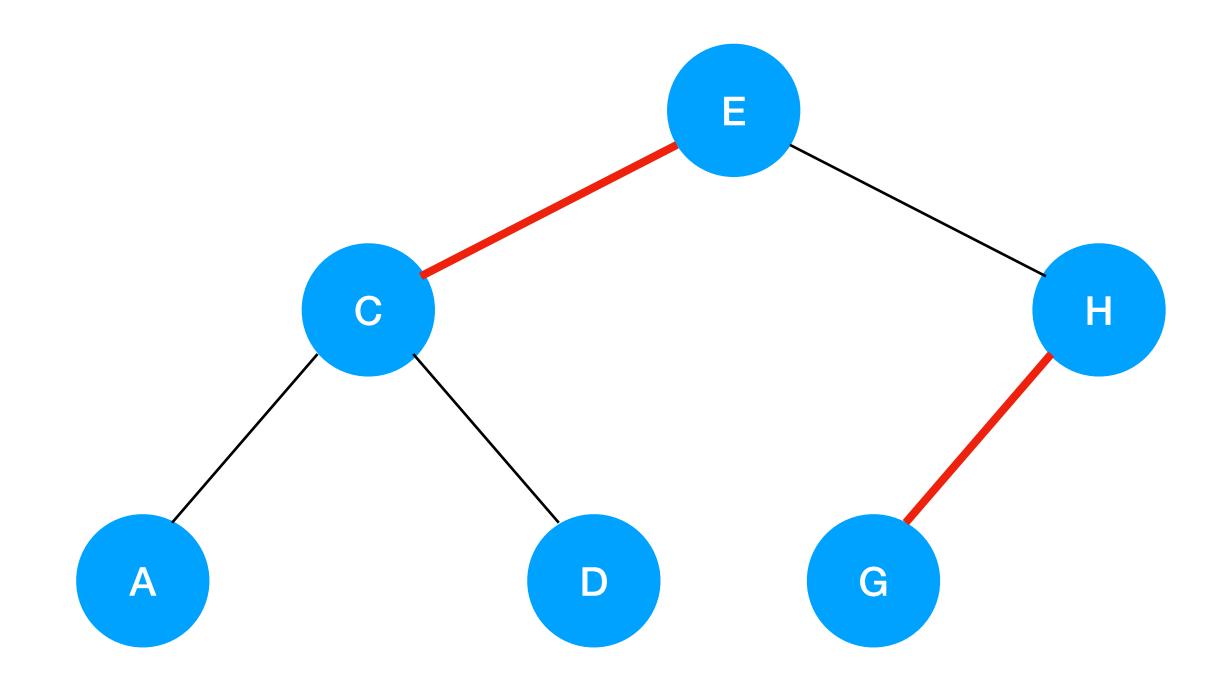
- 2-noeud, avec 1 clé et deux liens
- À gauche : clés plus petites
- À droite : clés plus grandes
- 3-noeud, avec 2 clés et 3 liens
 - À gauche : clés plus petites
 - À droite : clés plus grandes
 - Au milieu : clés entre



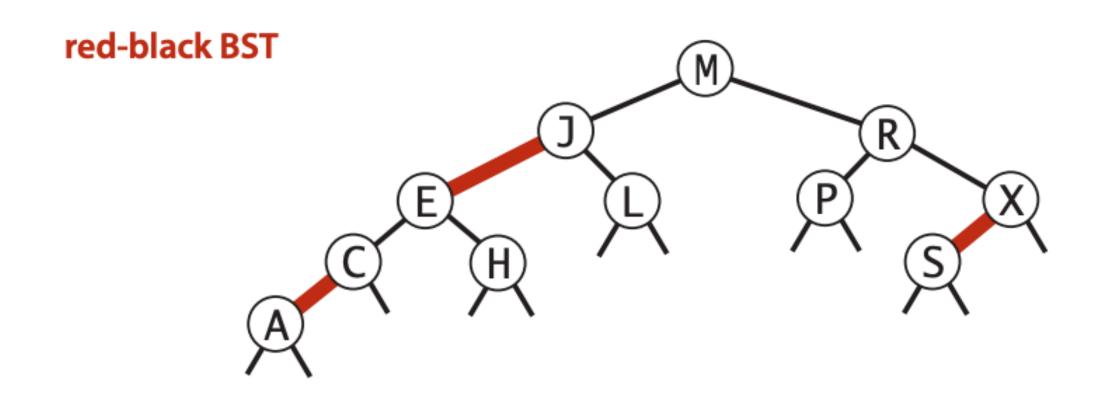
Arbre Red-Black

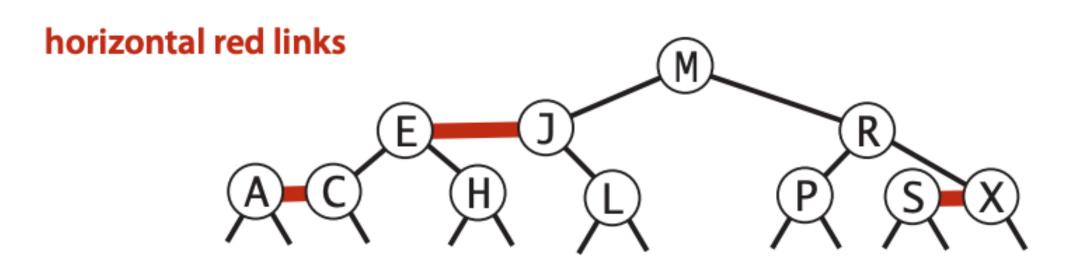
Arbre Red-Black:

- 2 types de liens (rouge ou noir)
- ·Liens rouges tjs à gauche
- ·Pas deux liens rouges d'affilées
- Equilibre parfait sur les liens noirs

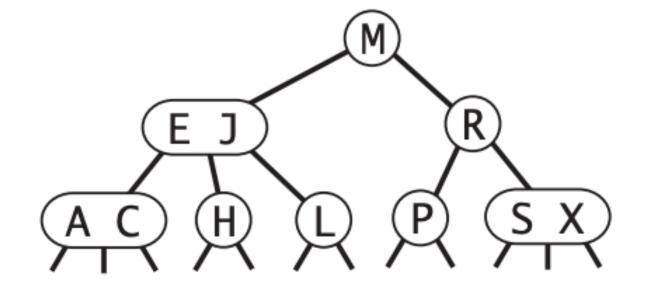


Correspondance 2-3 et Red-Black

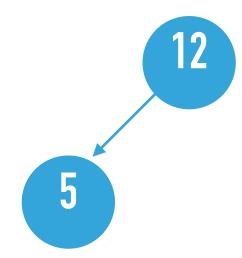


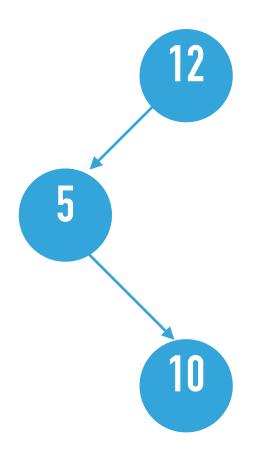


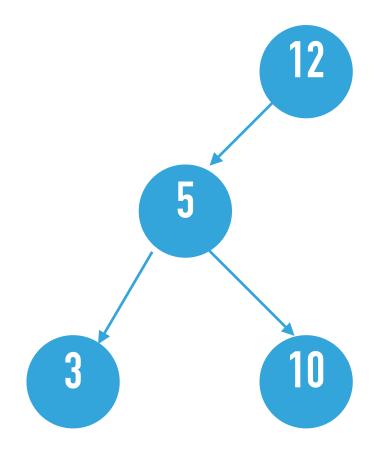
2-3 tree

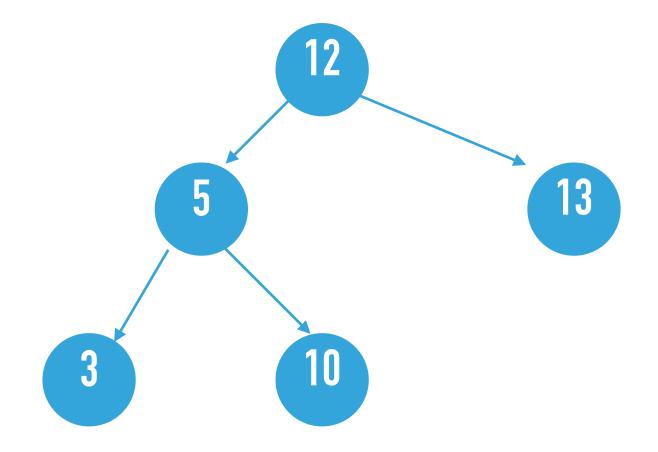


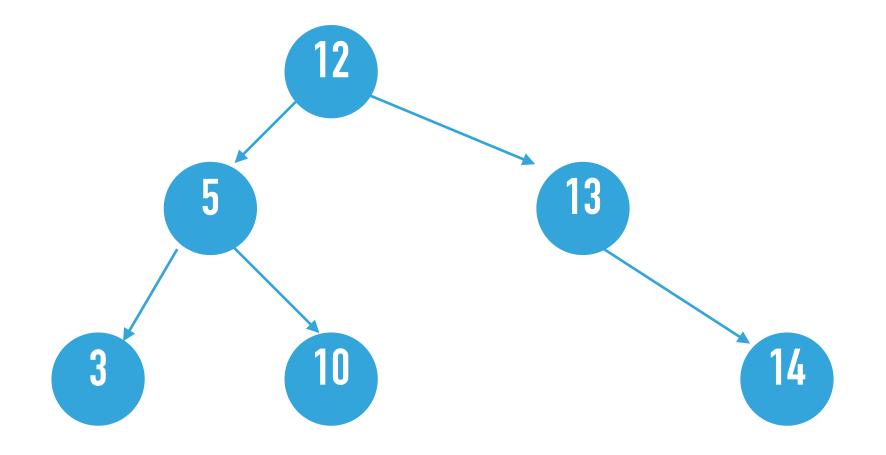
12

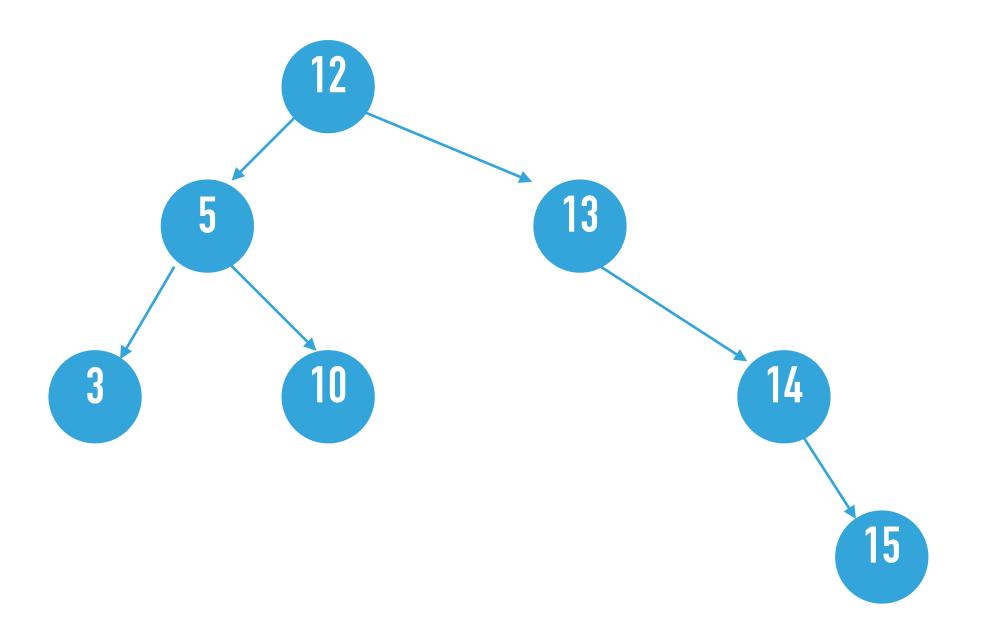


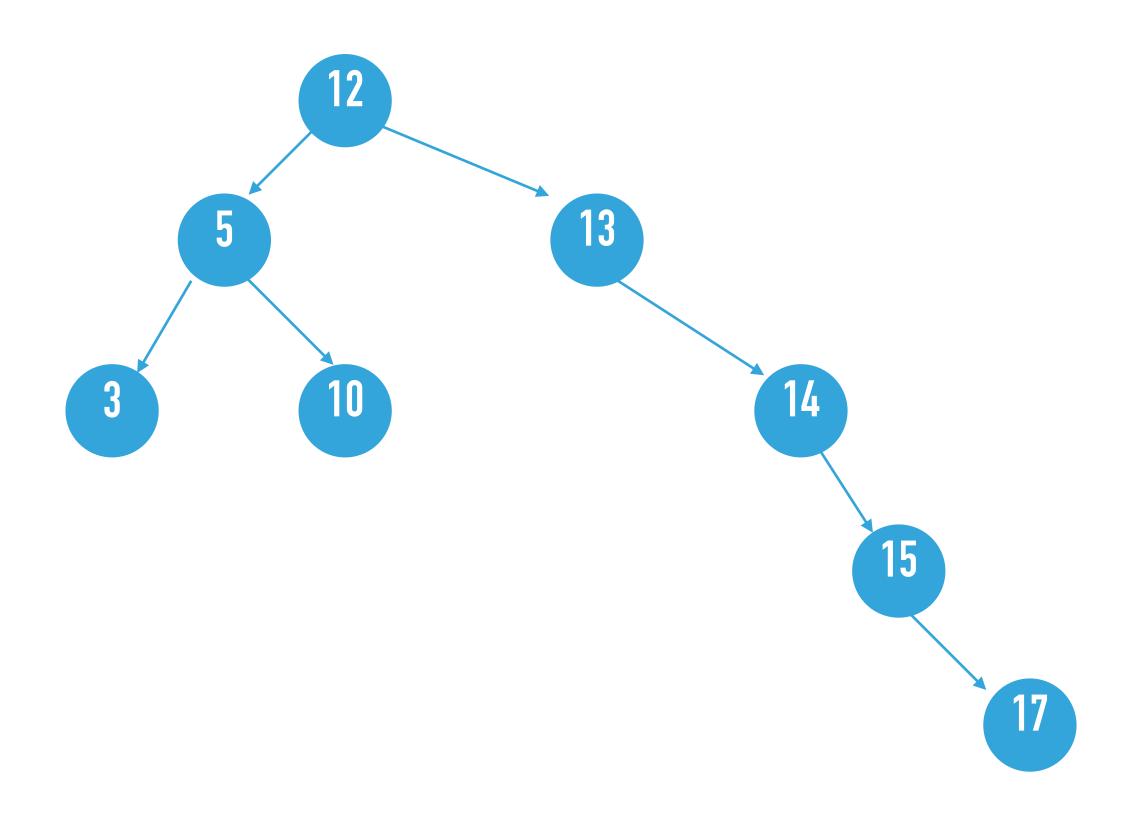


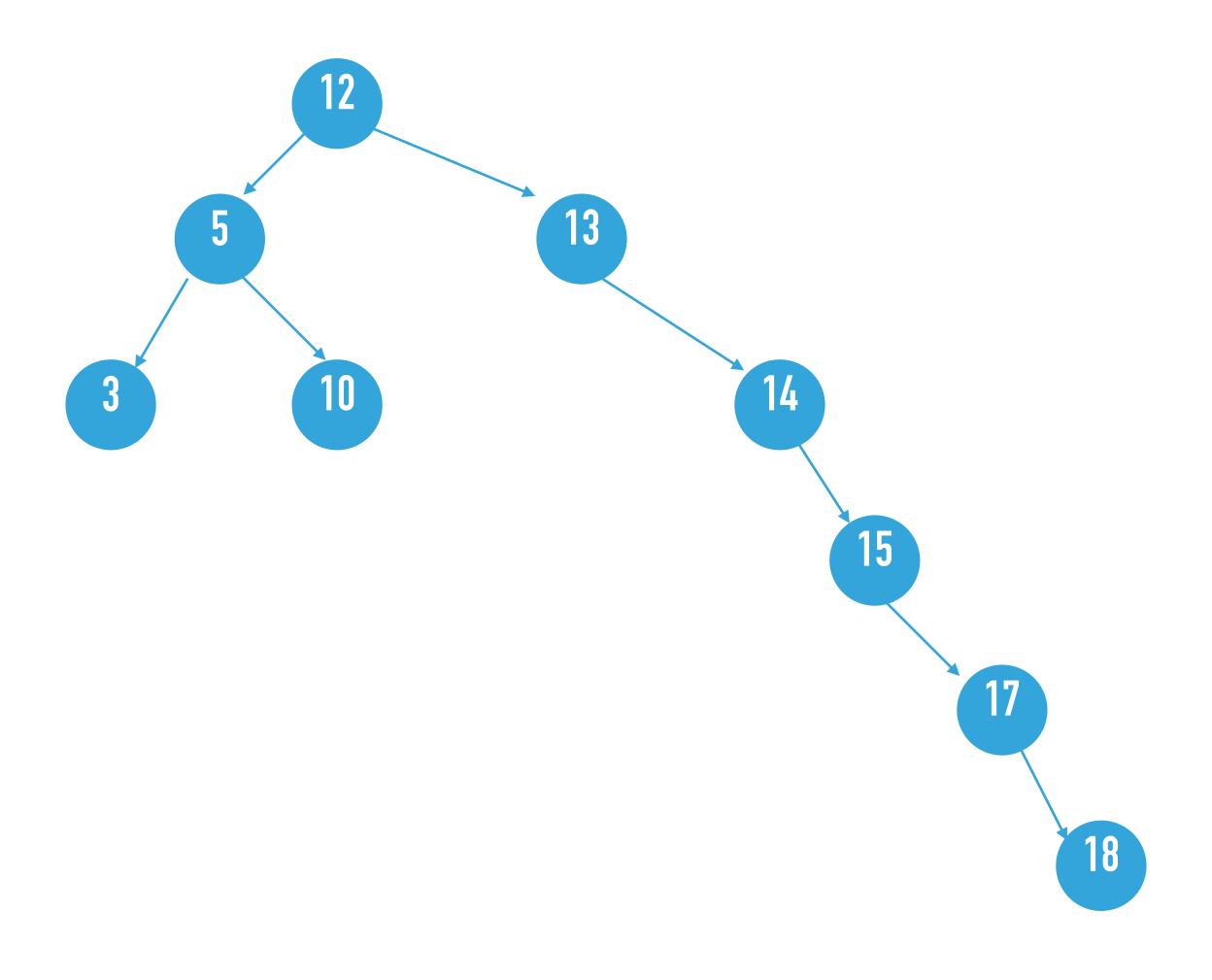




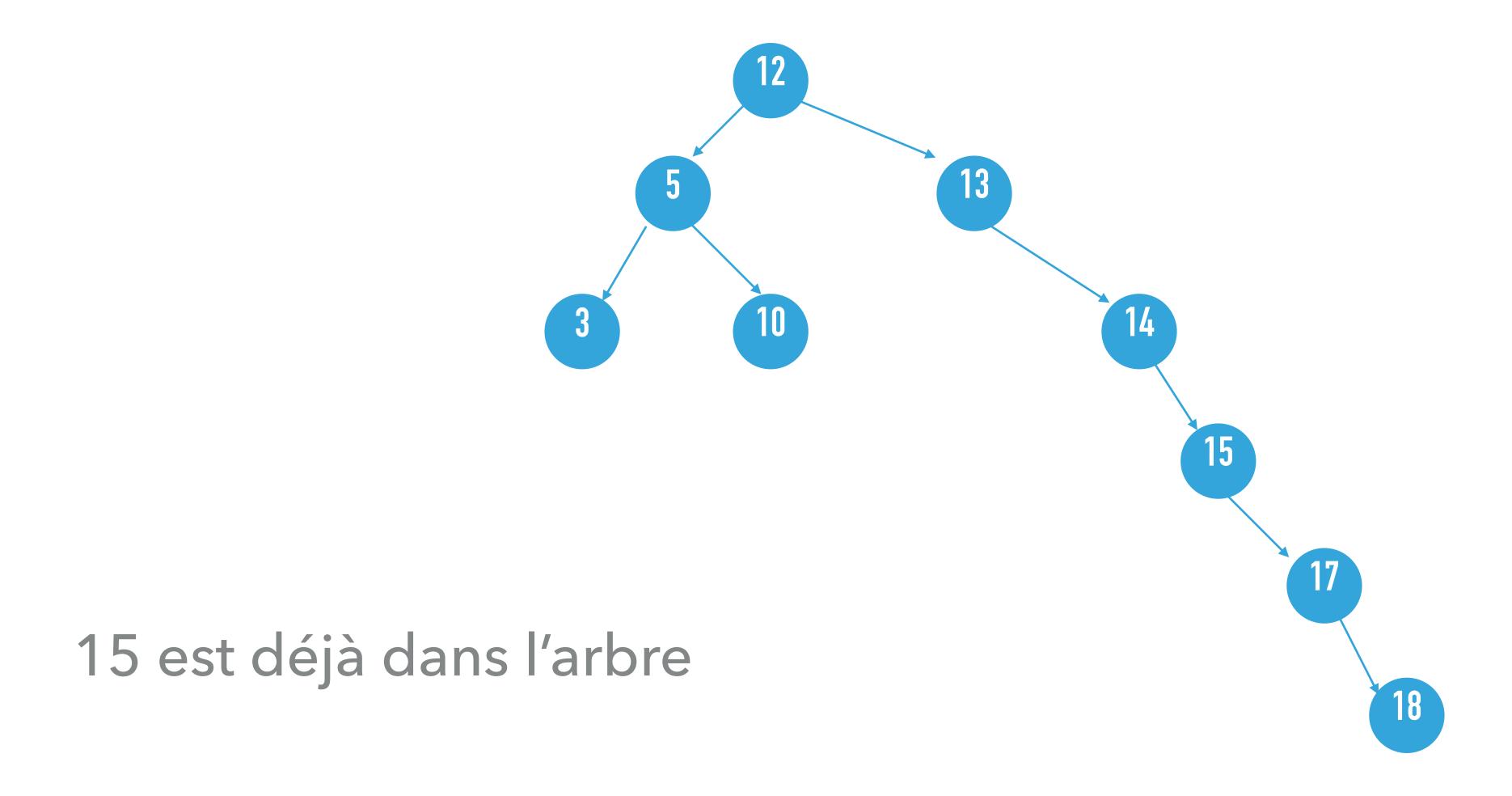








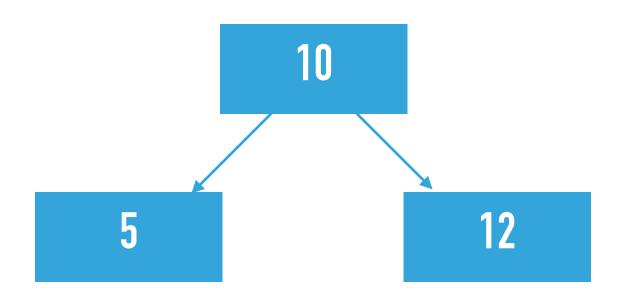
12, 5, 10, 3, 13, 14, 15, 17, 18, 15

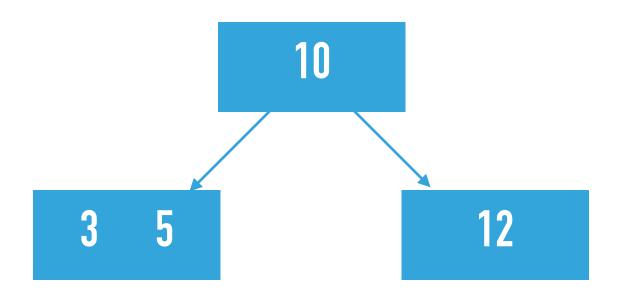


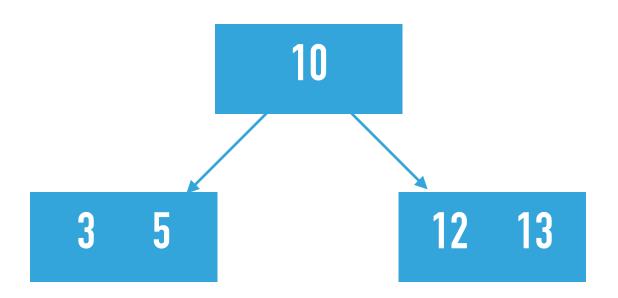
12, 5, 10, 3, 13, 14, 15, 17, 18, 15

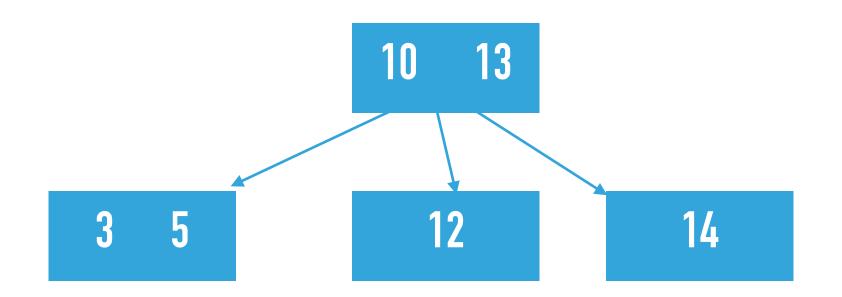
2

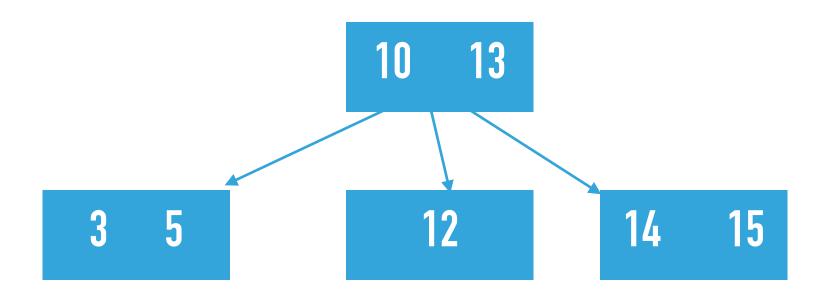
5 12

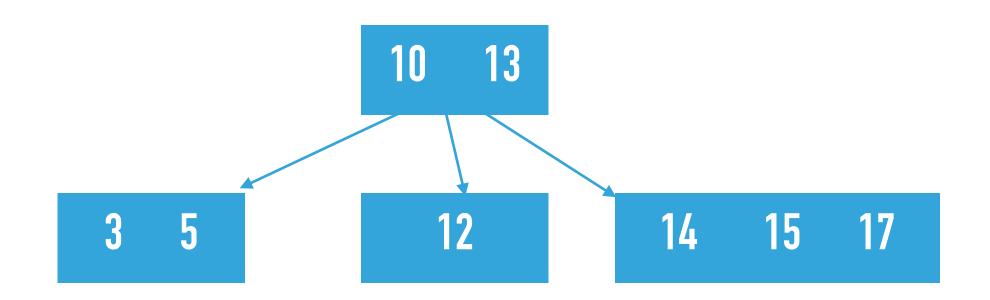


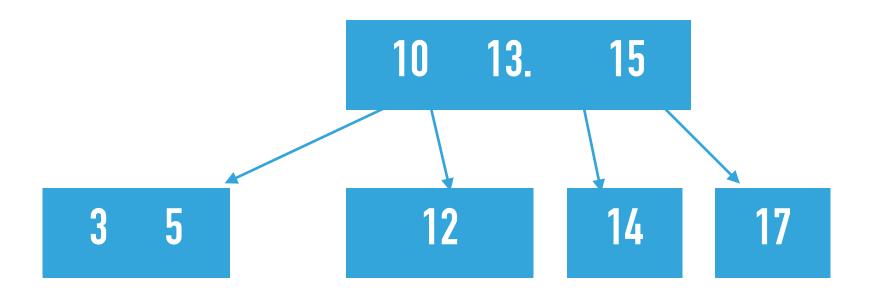


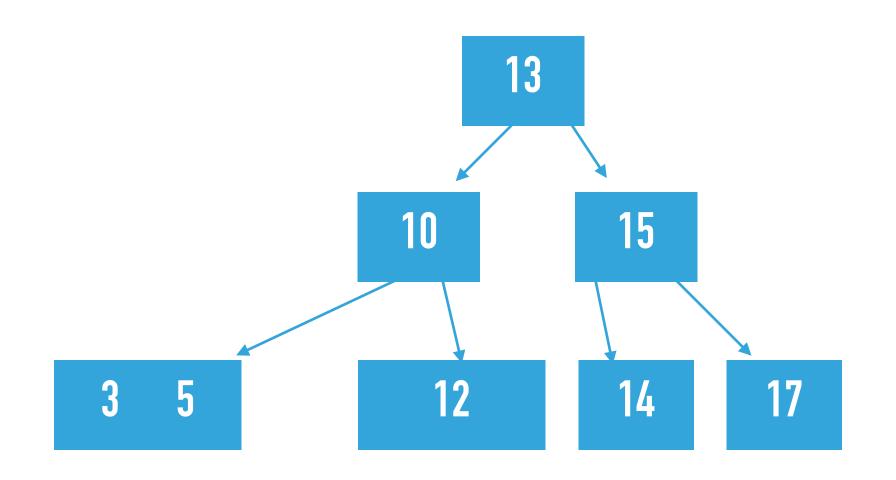


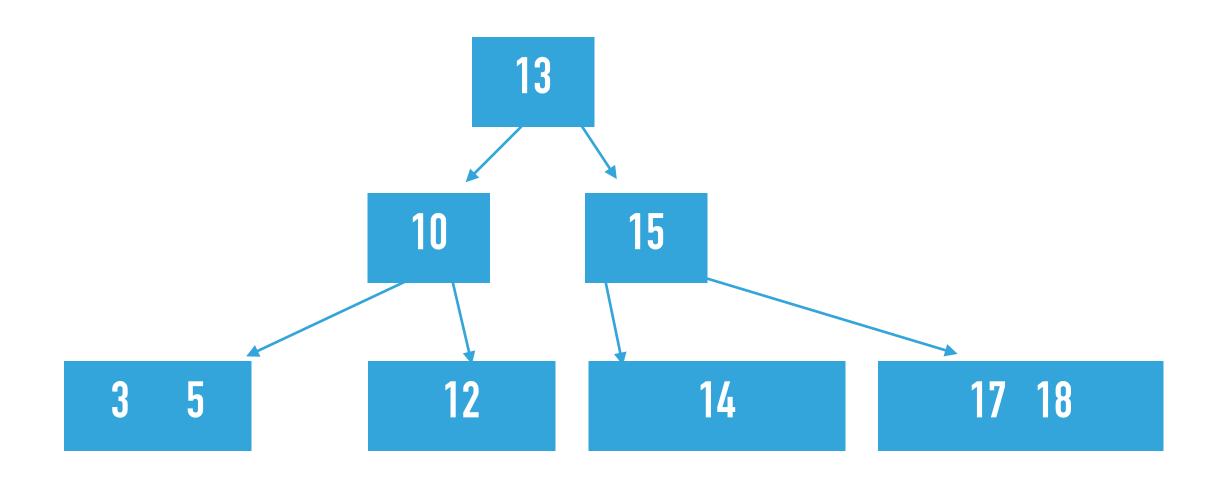


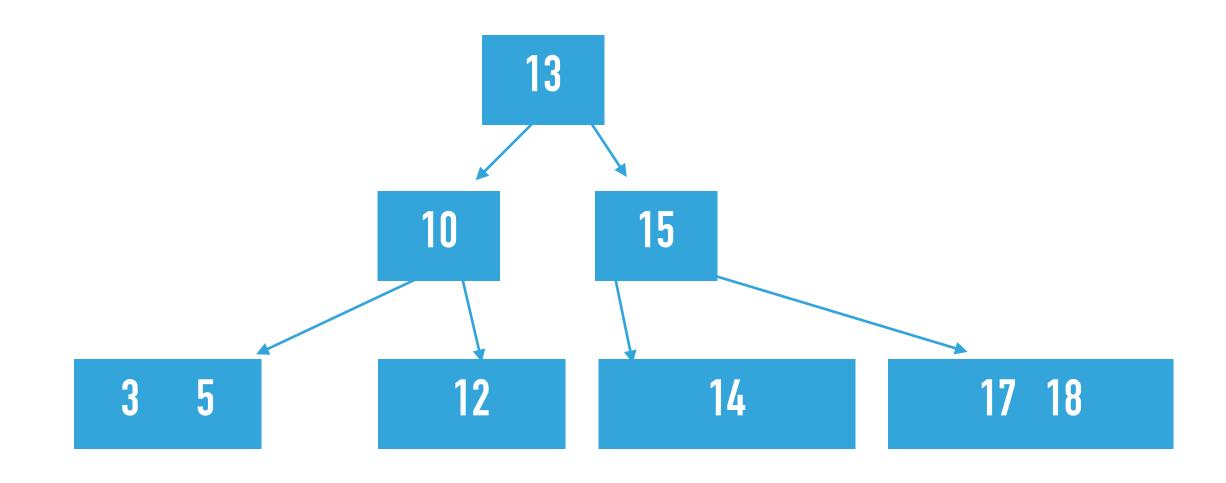






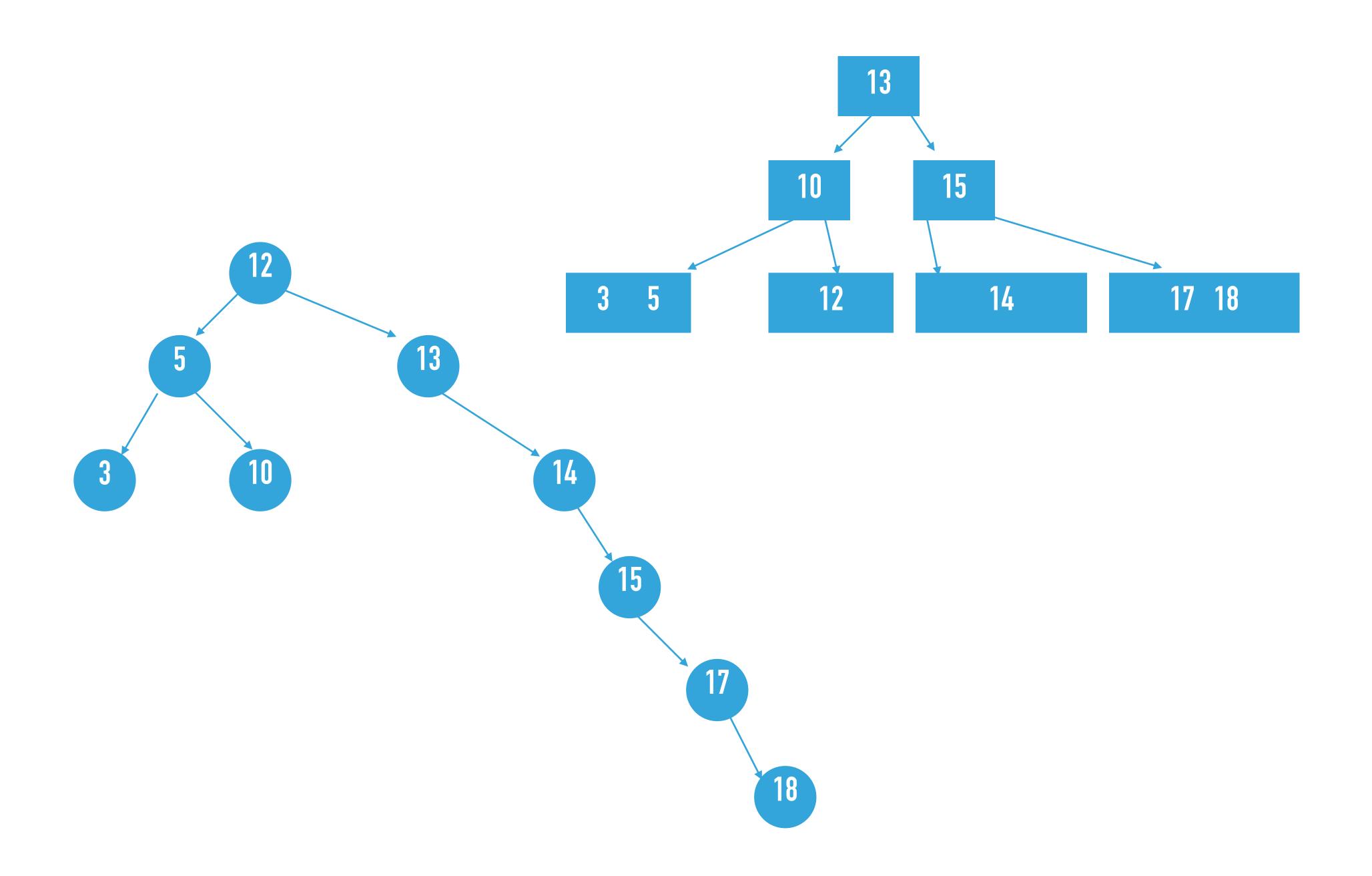




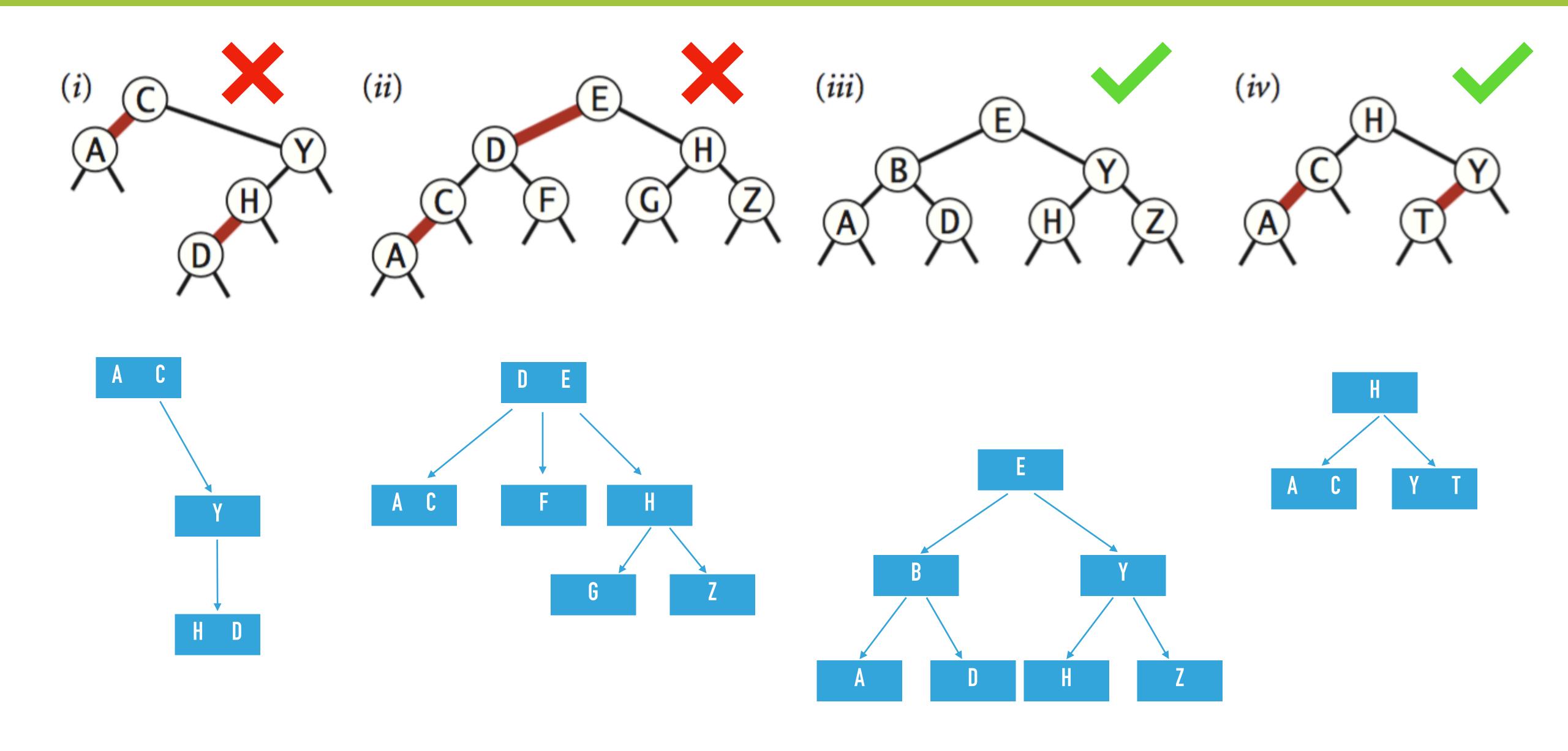


15 est déjà dans l'arbre

12, 5, 10, 3, 13, 14, 15, 17, 18, 15

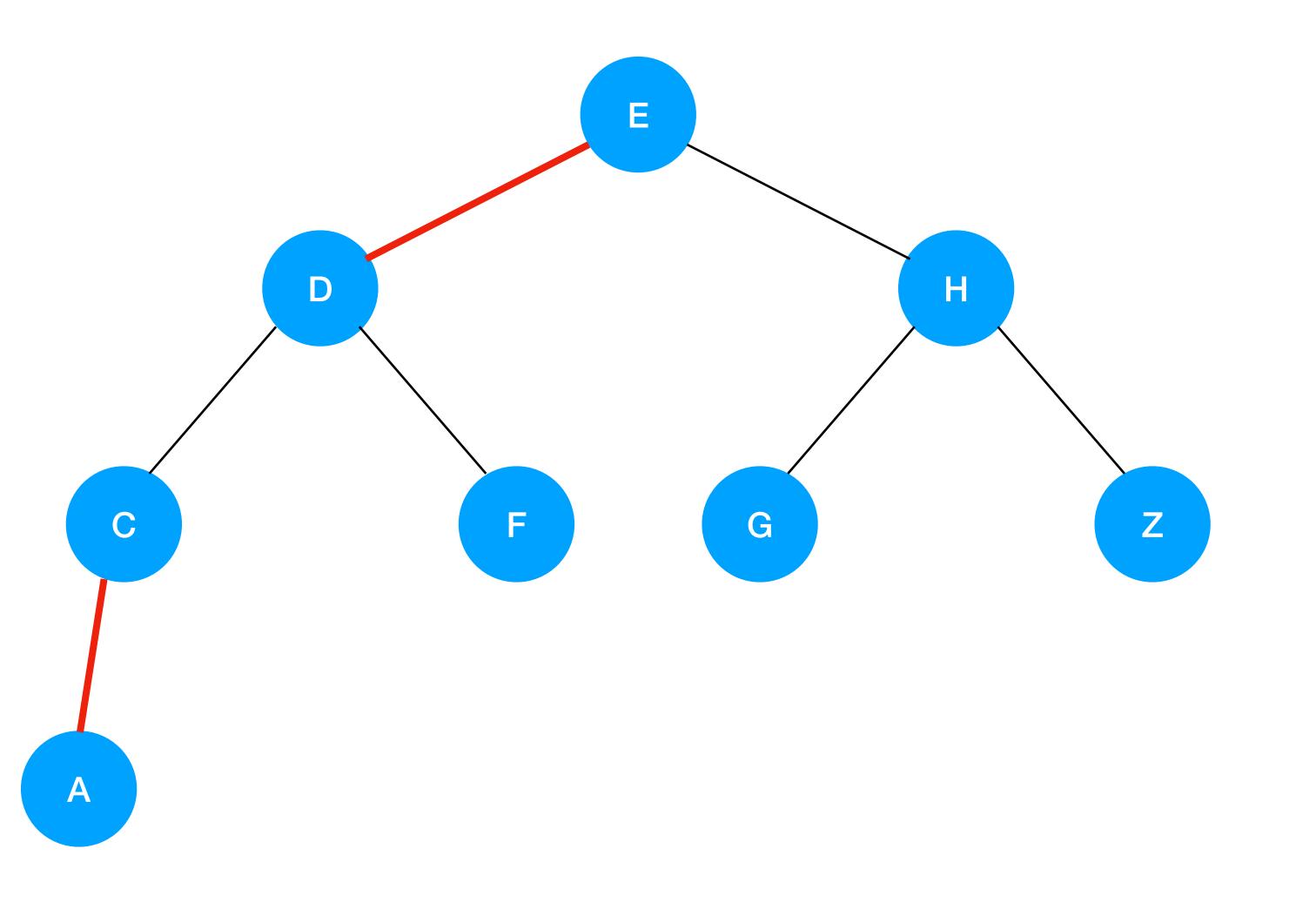


Question 3.1.10 RB?



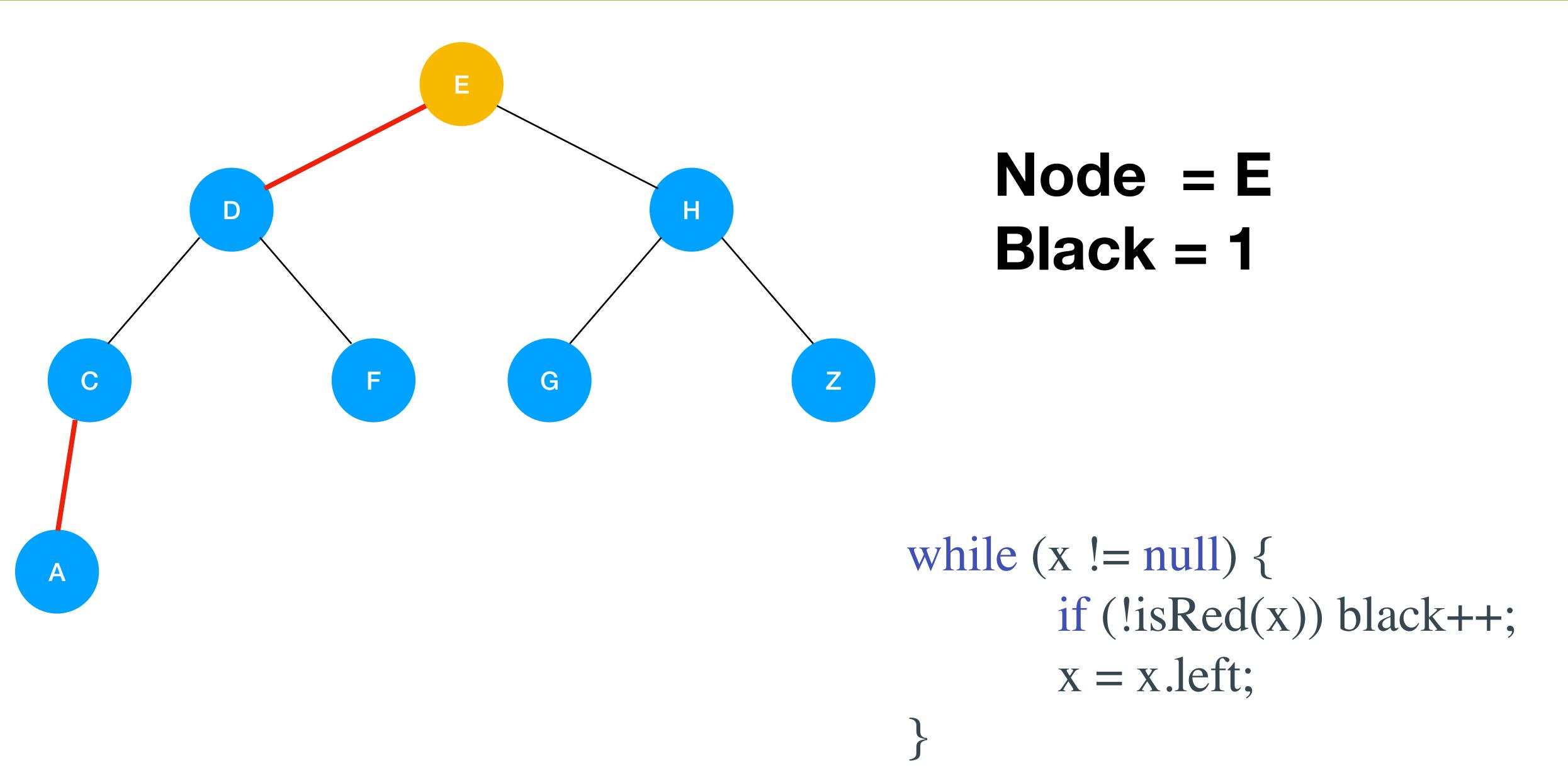
```
private boolean is23() { return is23(root); }
private boolean is23(Node x) {
  if (x == null)
     return true;
  //Est-ce que le lien à droite est rouge ?
  if (isRed(x.right))
     return false;
  //Est-ce que le noeud est connecté à deux liens rouges ?
  if (x != root && isRed(x) && isRed(x.left))
     return false;
  return is23(x.left) && is23(x.right);
```

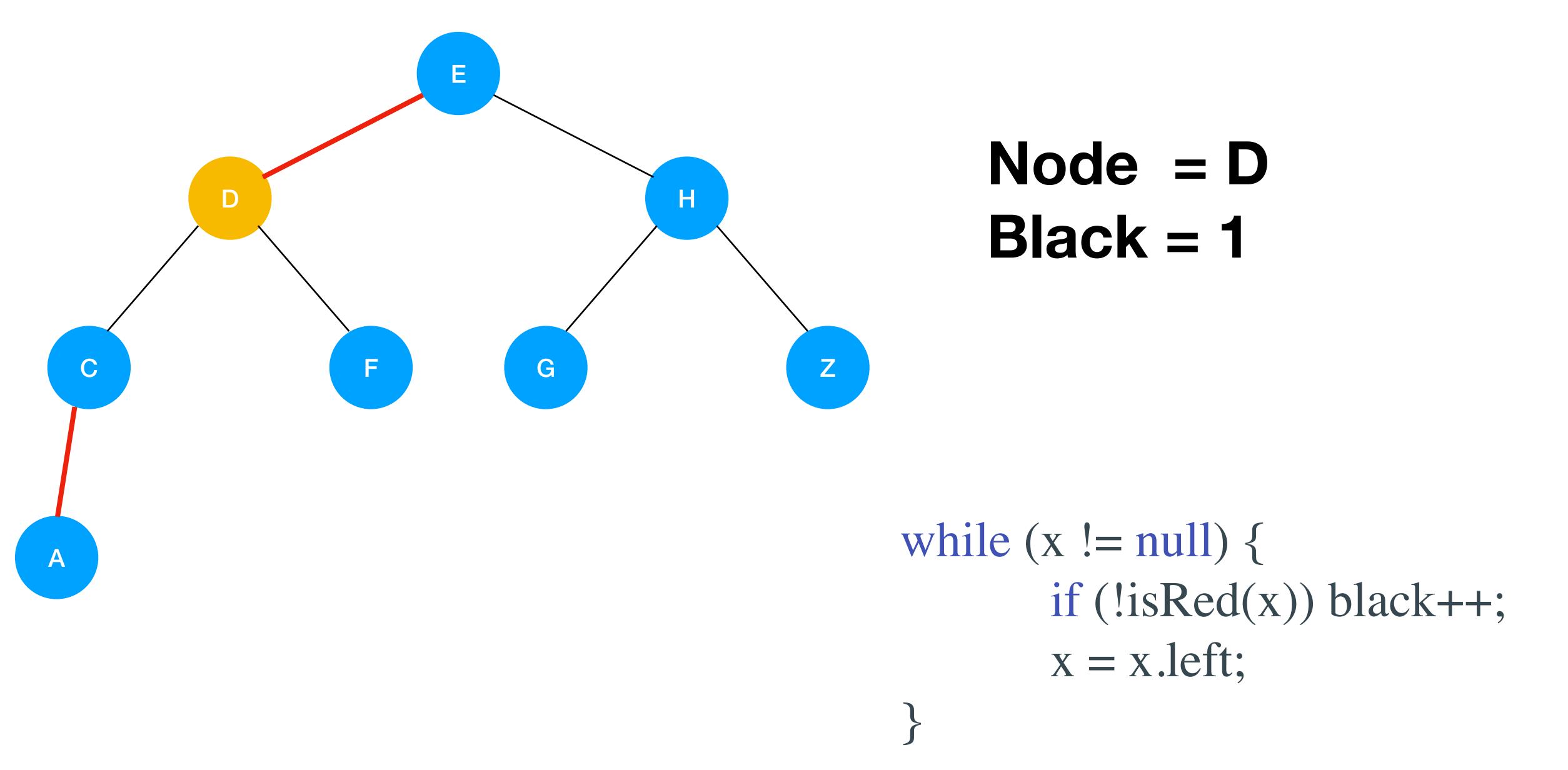
```
private boolean isBalanced() {
     // nombre de liens noirs de la racine à la clé min
     int black = 0;
     Node x = root;
     while (x != null) \{
       if (!isRed(x)) black++;
       x = x.left;
     // Est-ce qu'on retrouve ce nombre de liens pour chaque feuille ?
     return isBalanced(root, black);
private boolean isBalanced(Node x, int black) {
  if (x == null) return black == 0;
  if (!isRed(x)) black--;
  return isBalanced(x.left, black) && isBalanced(x.right, black);
```

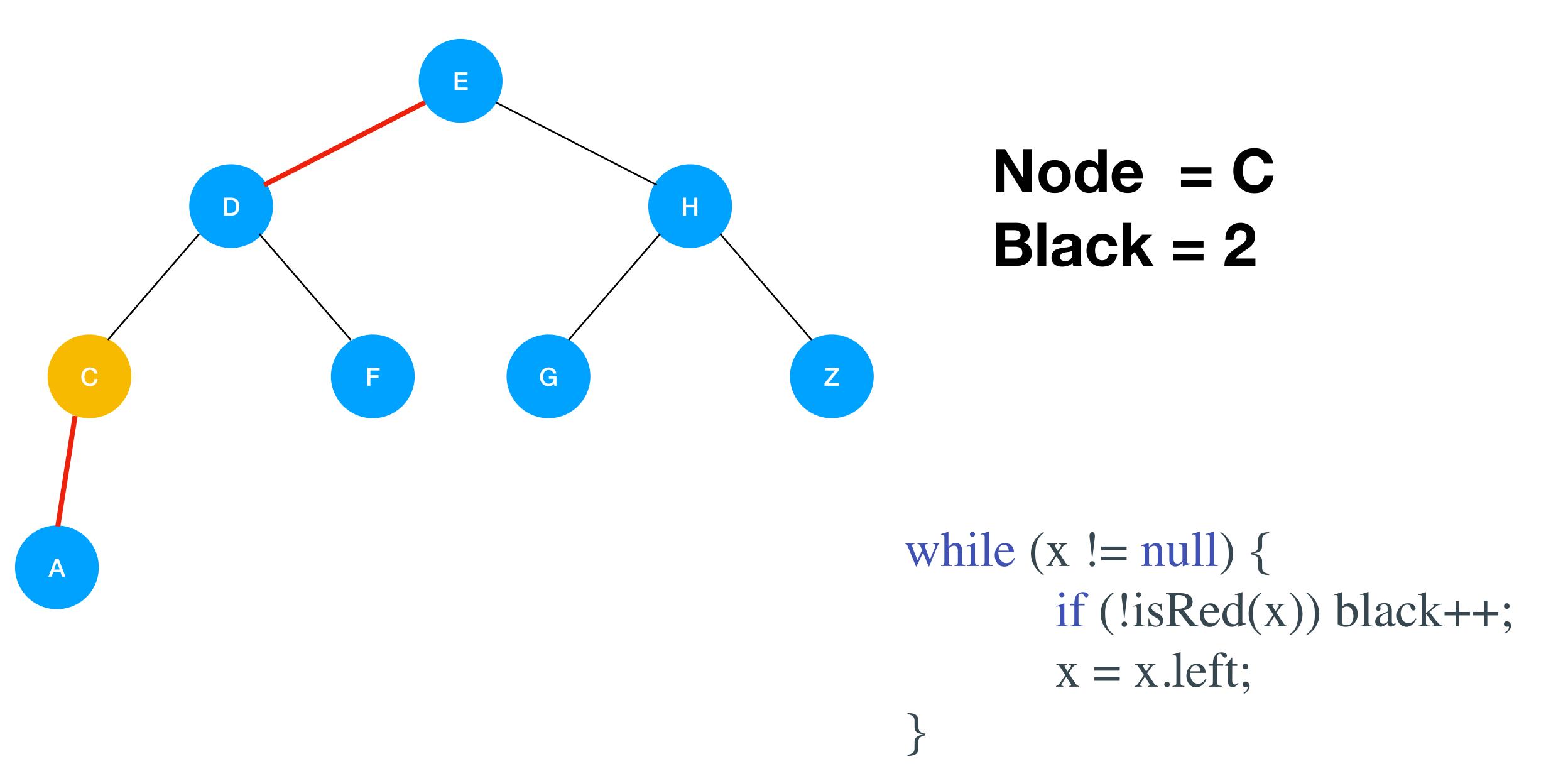


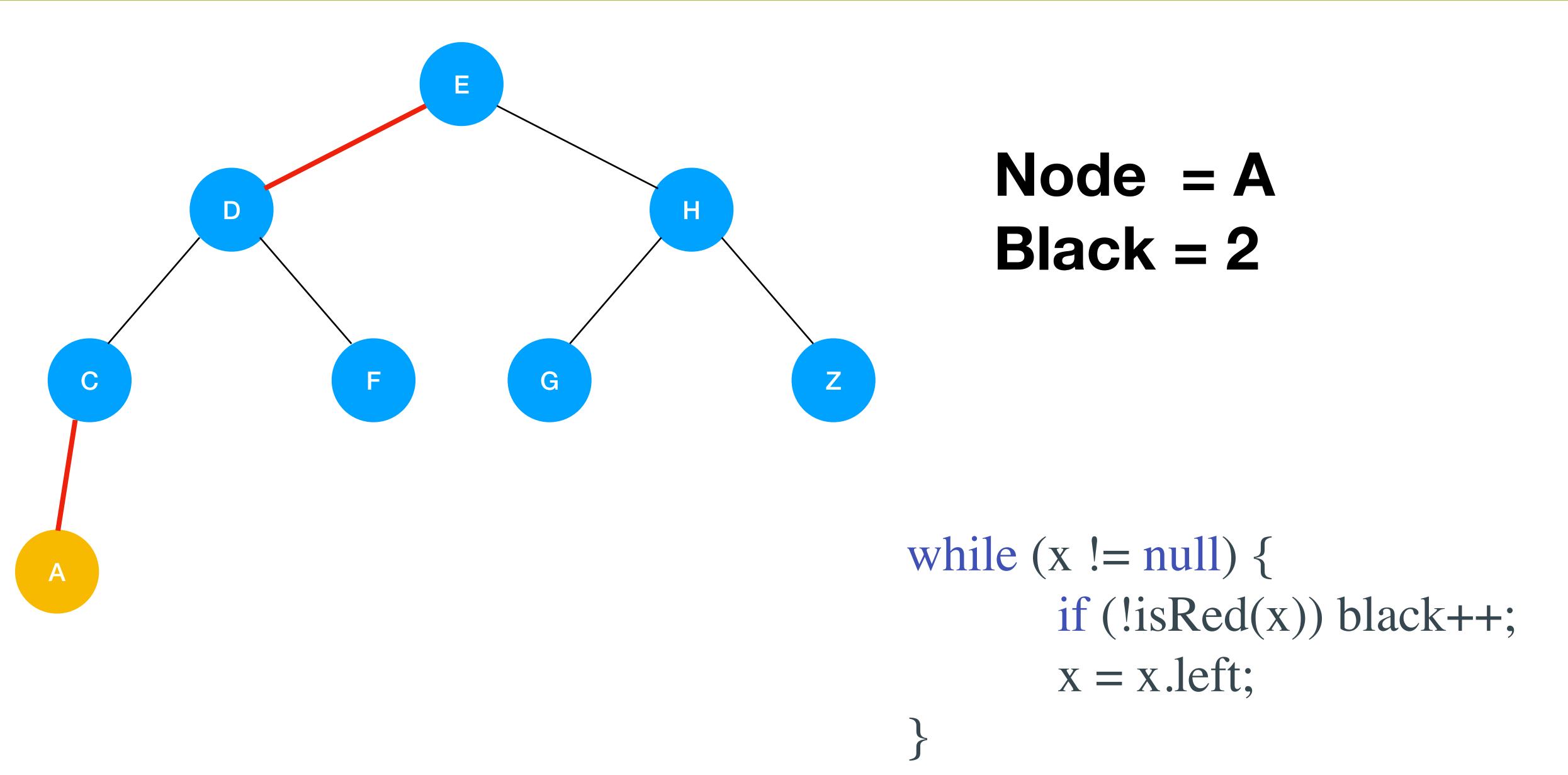
Node =
Black = 0

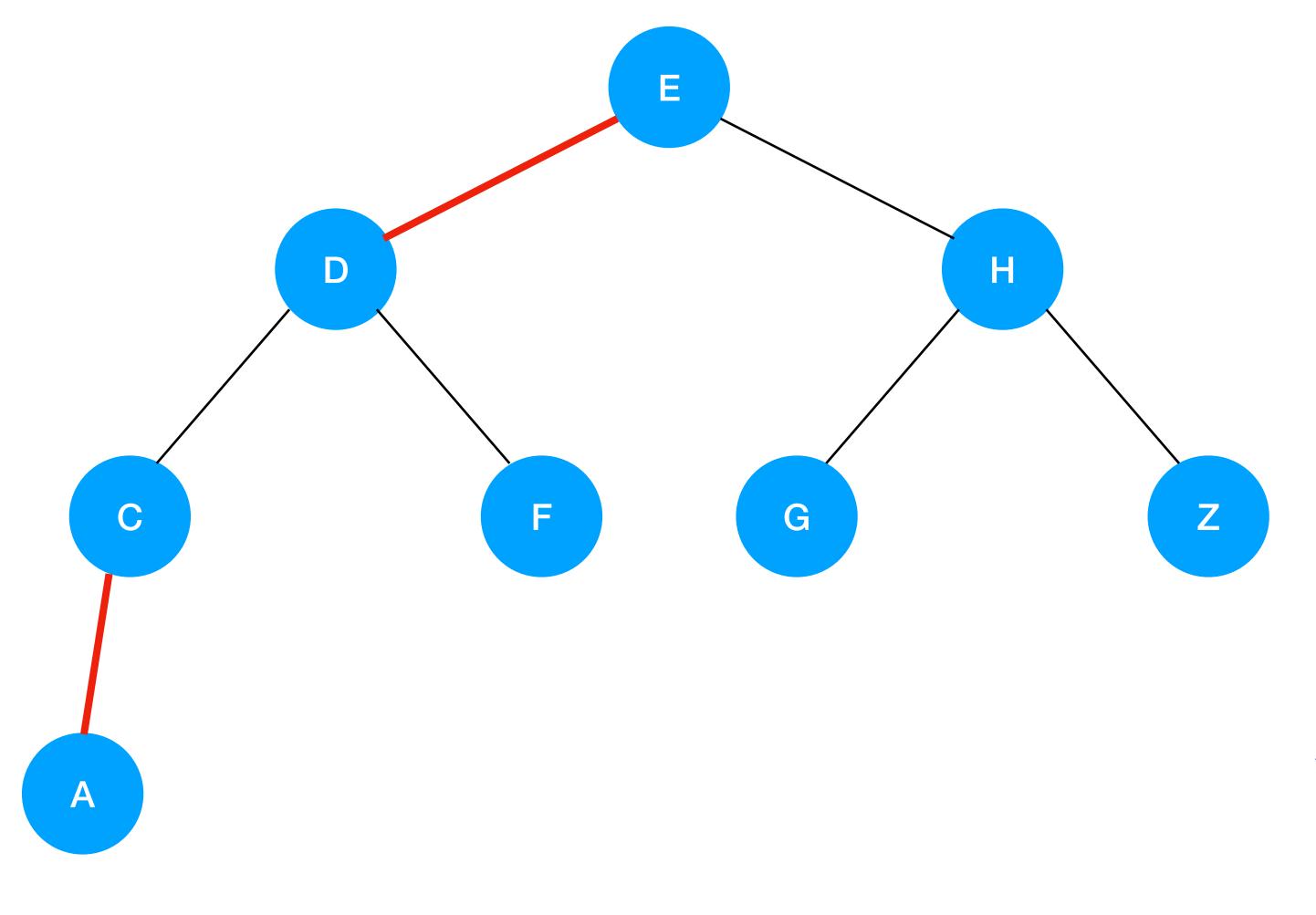
int black = 0;
Node x = root;





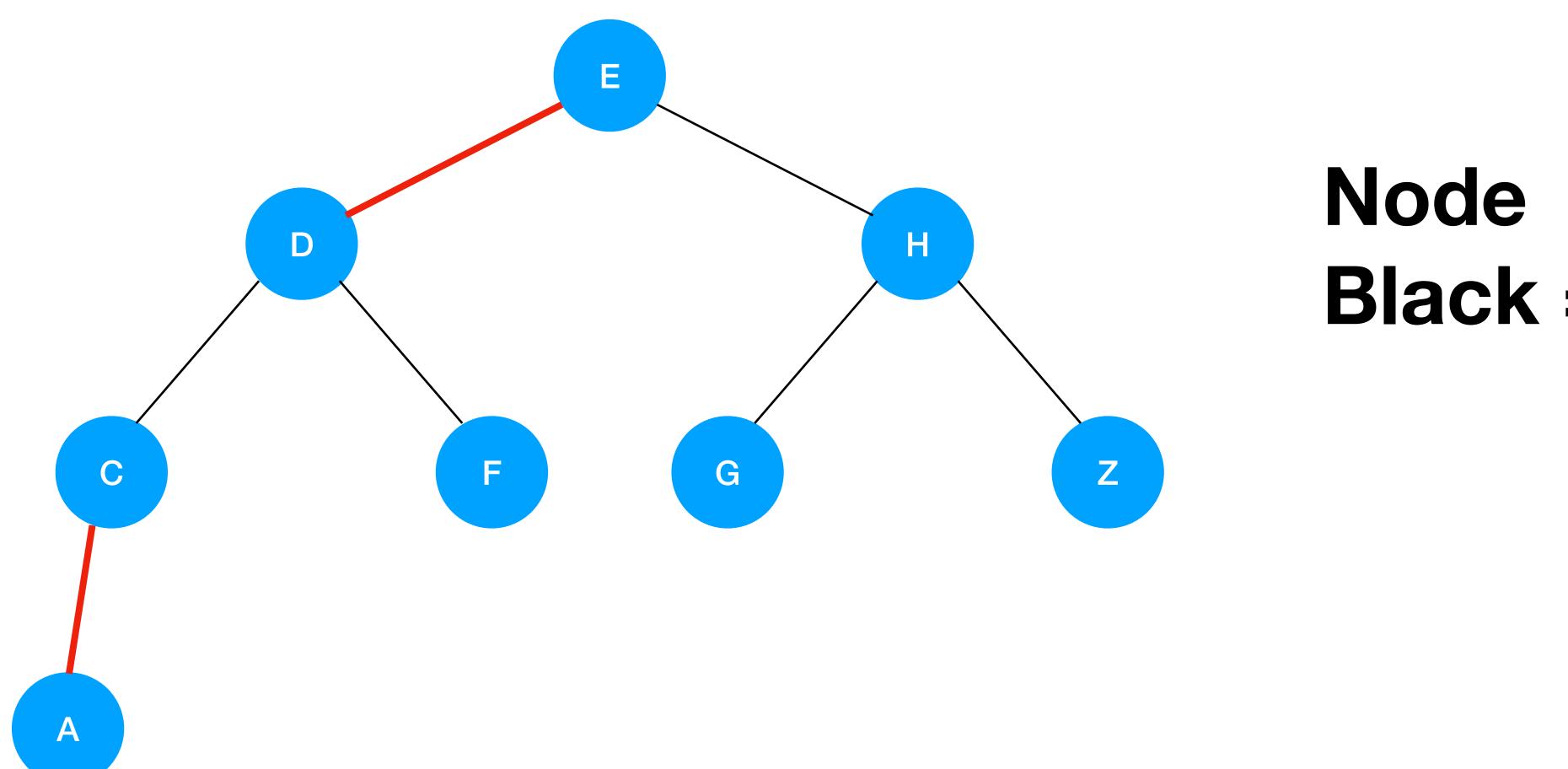






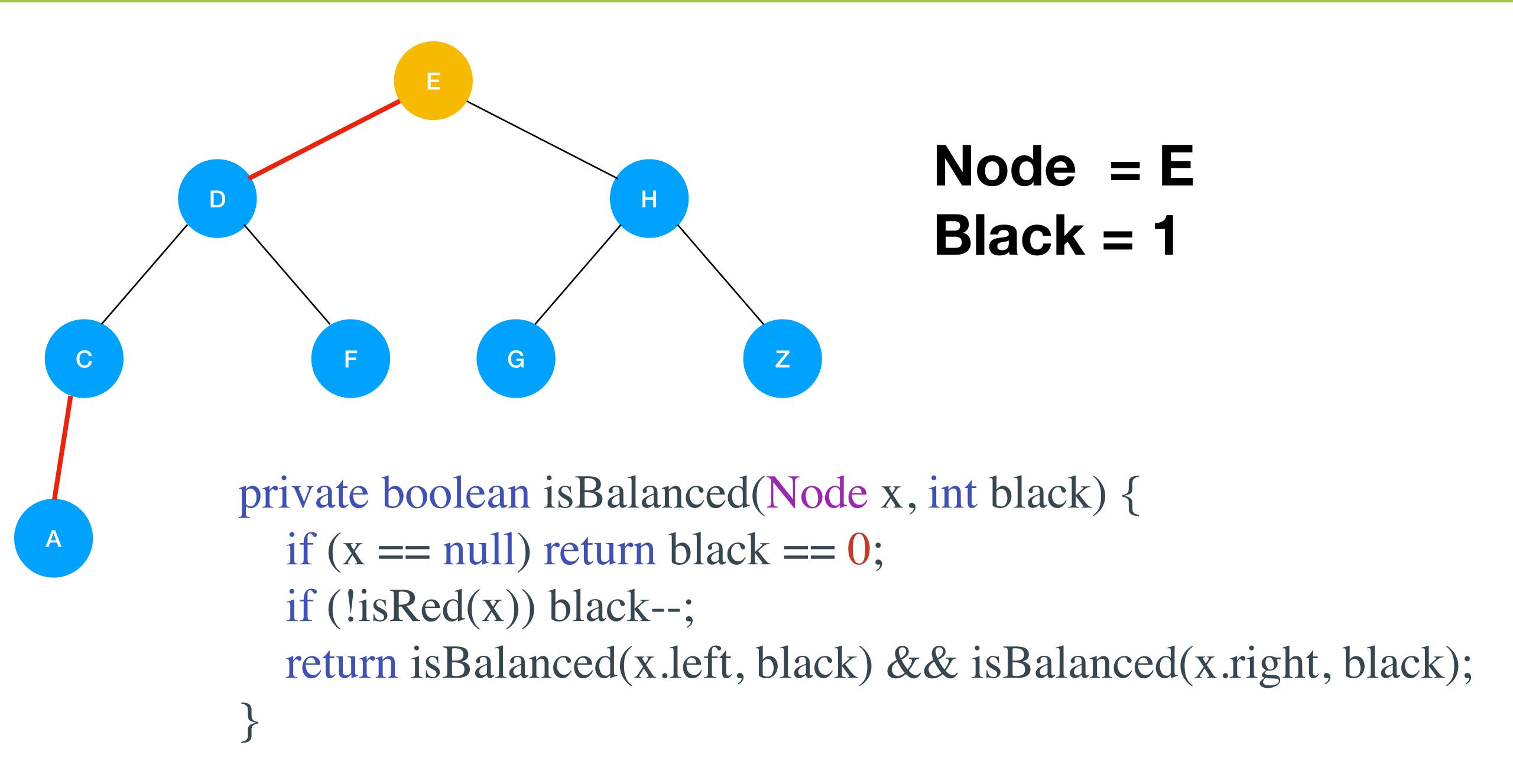
Node = null
Black = 2

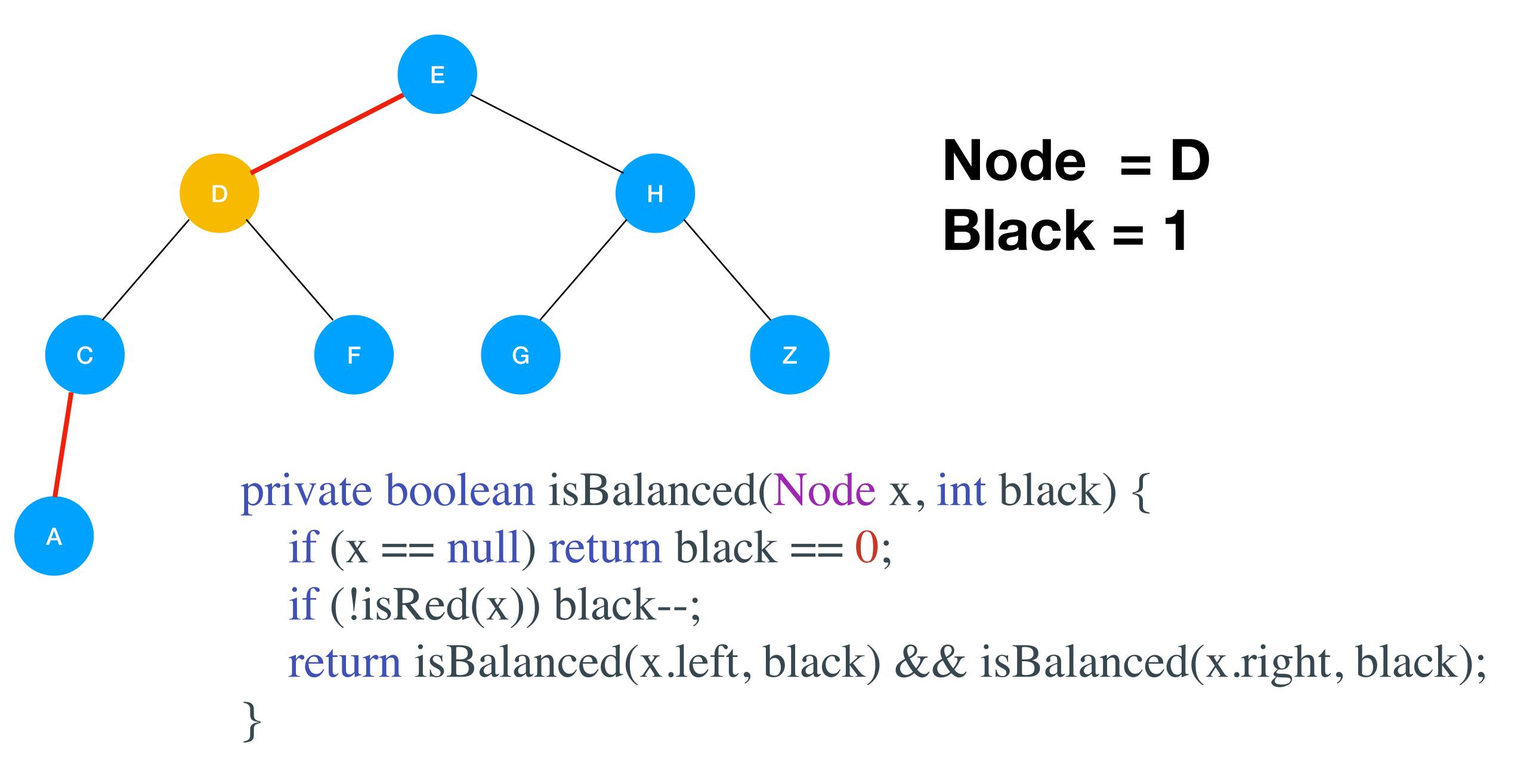
```
while (x != null) {
    if (!isRed(x)) black++;
        x = x.left;
}
```

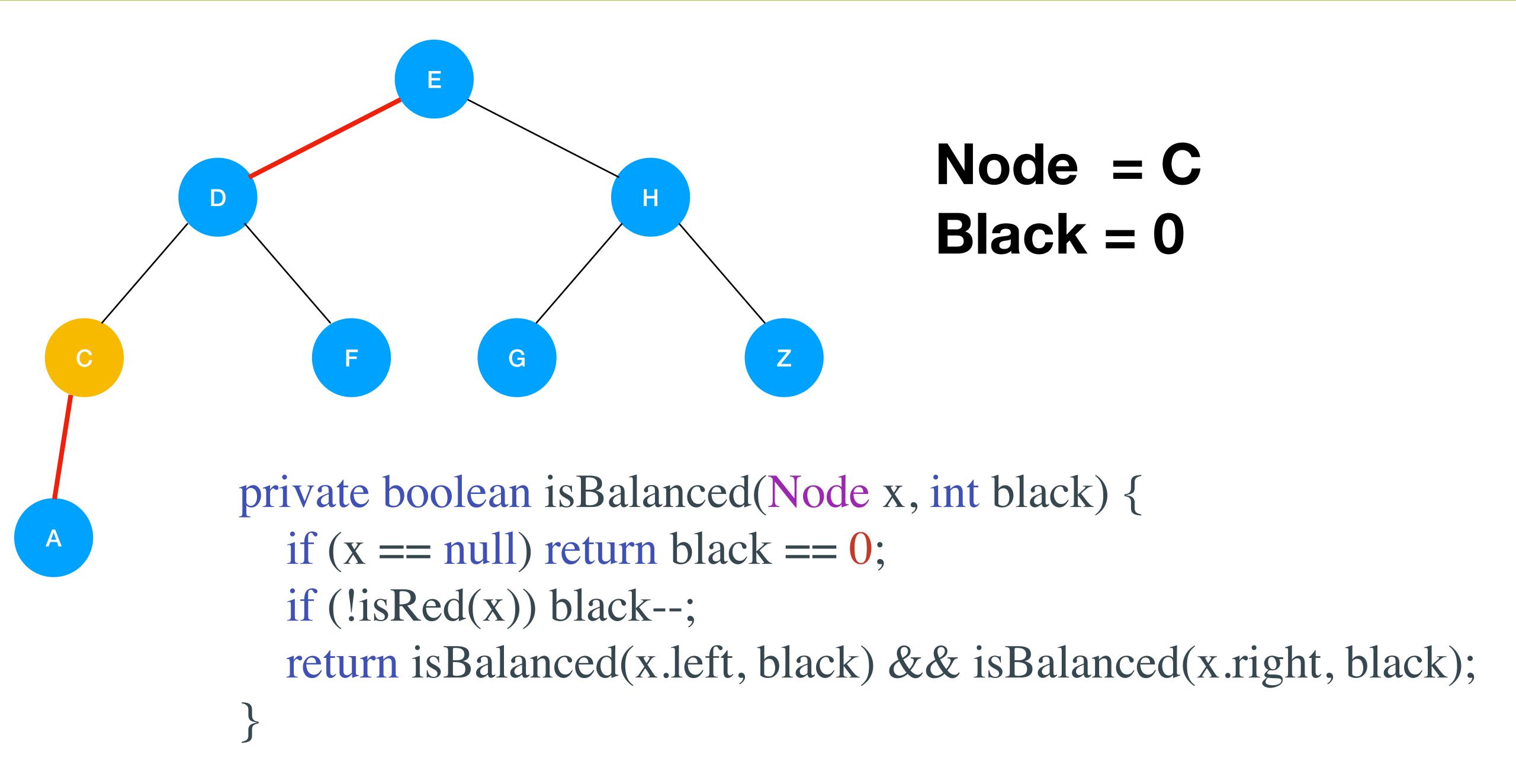


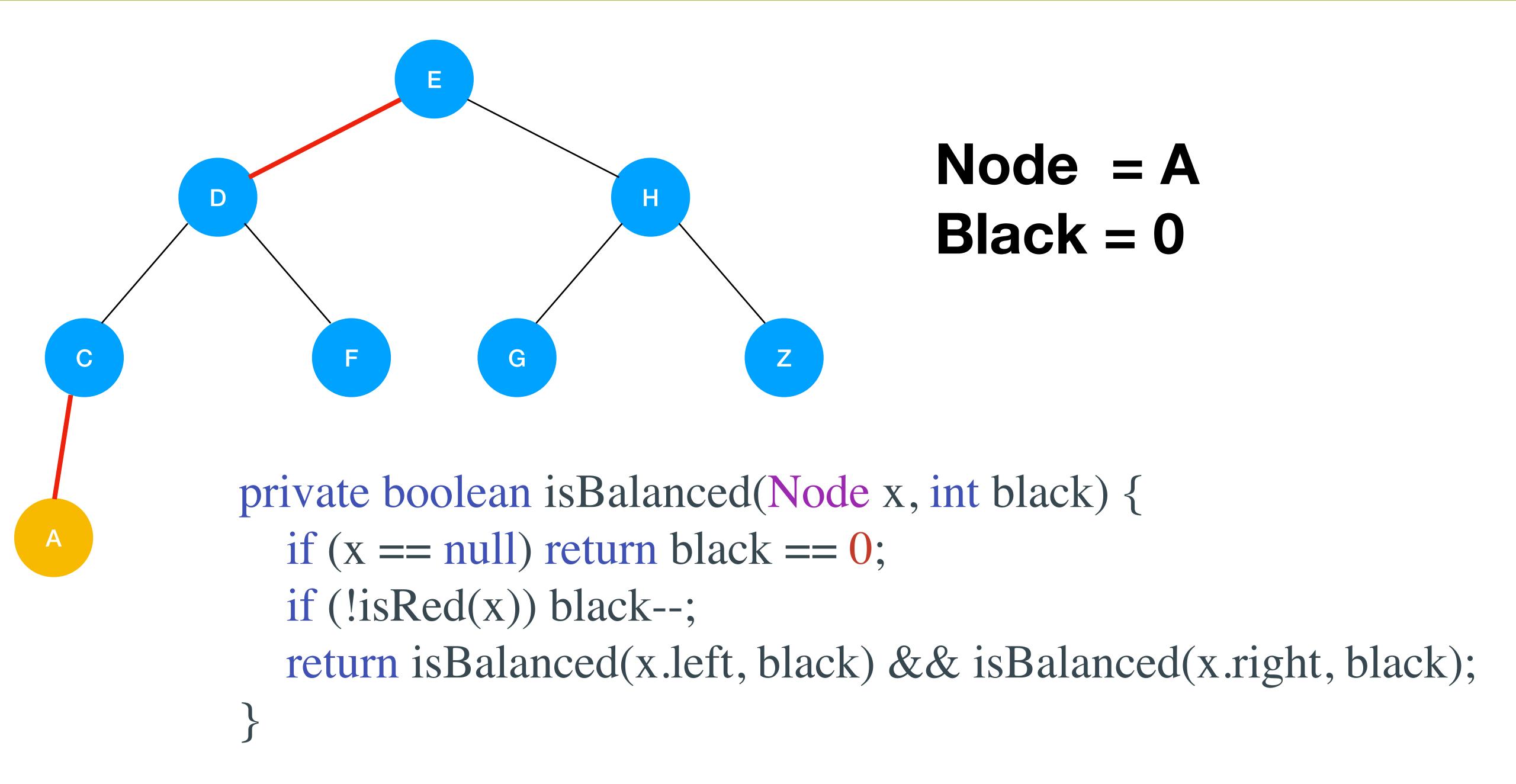
Node = Black = 2

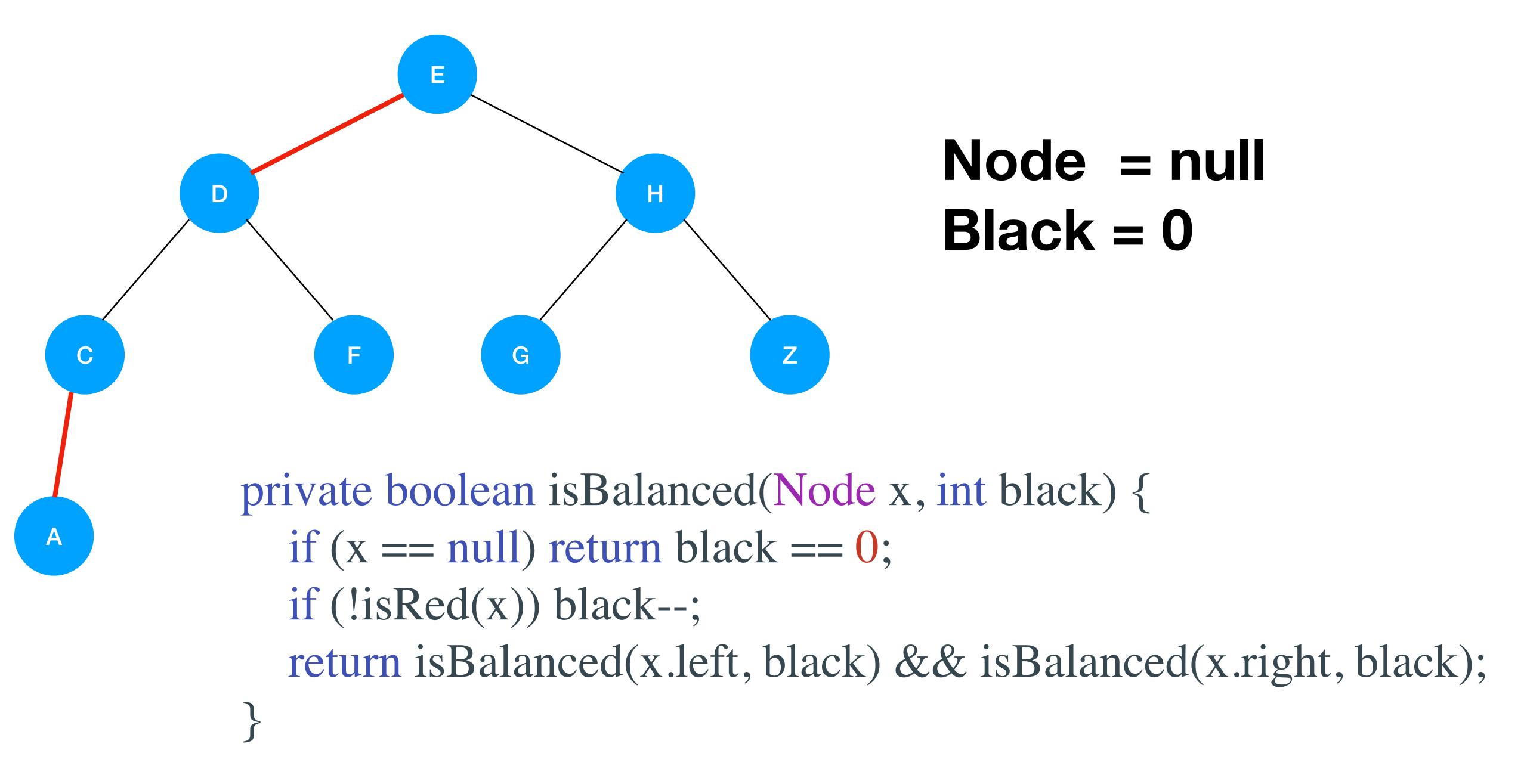
return isBalanced(root, black);

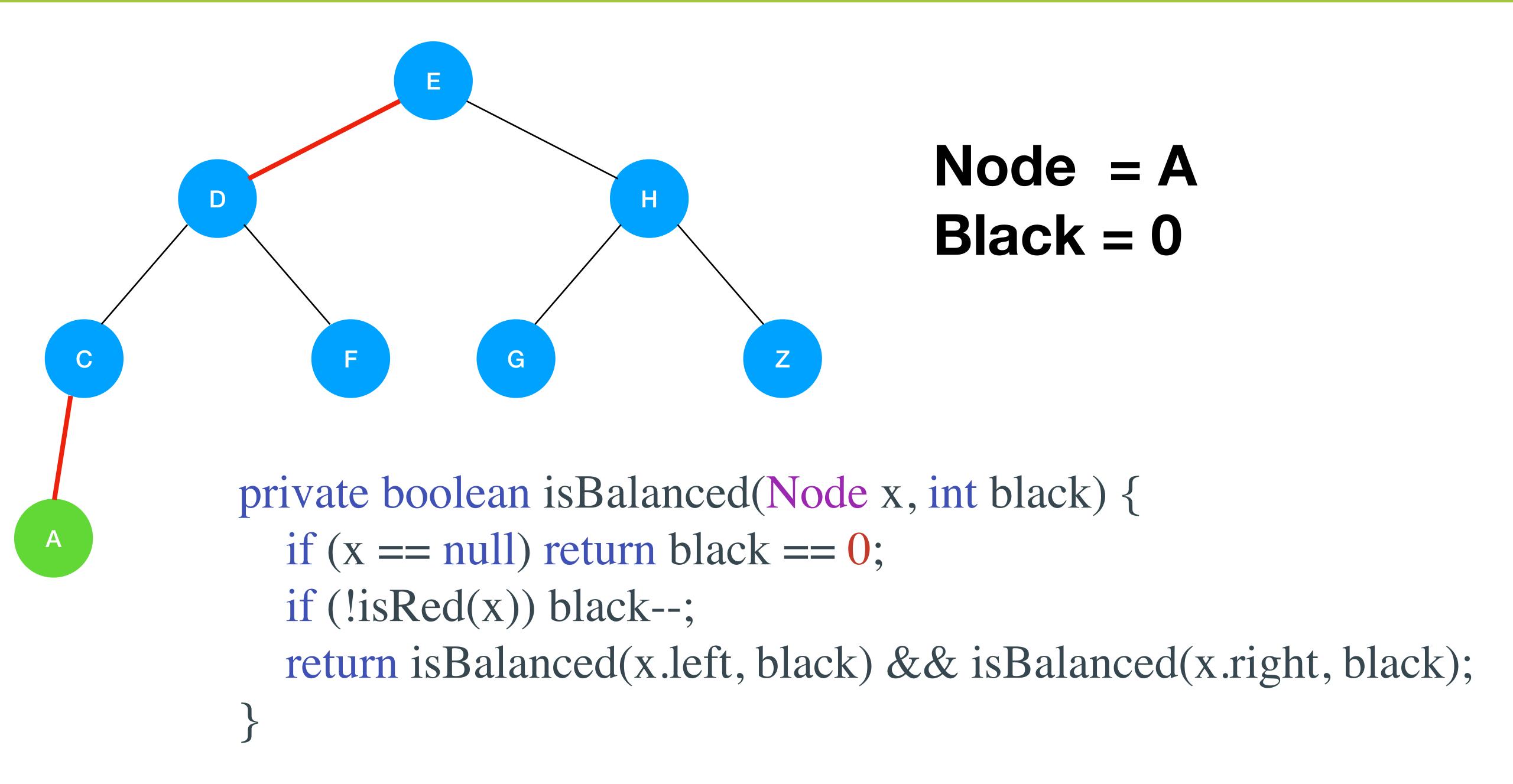


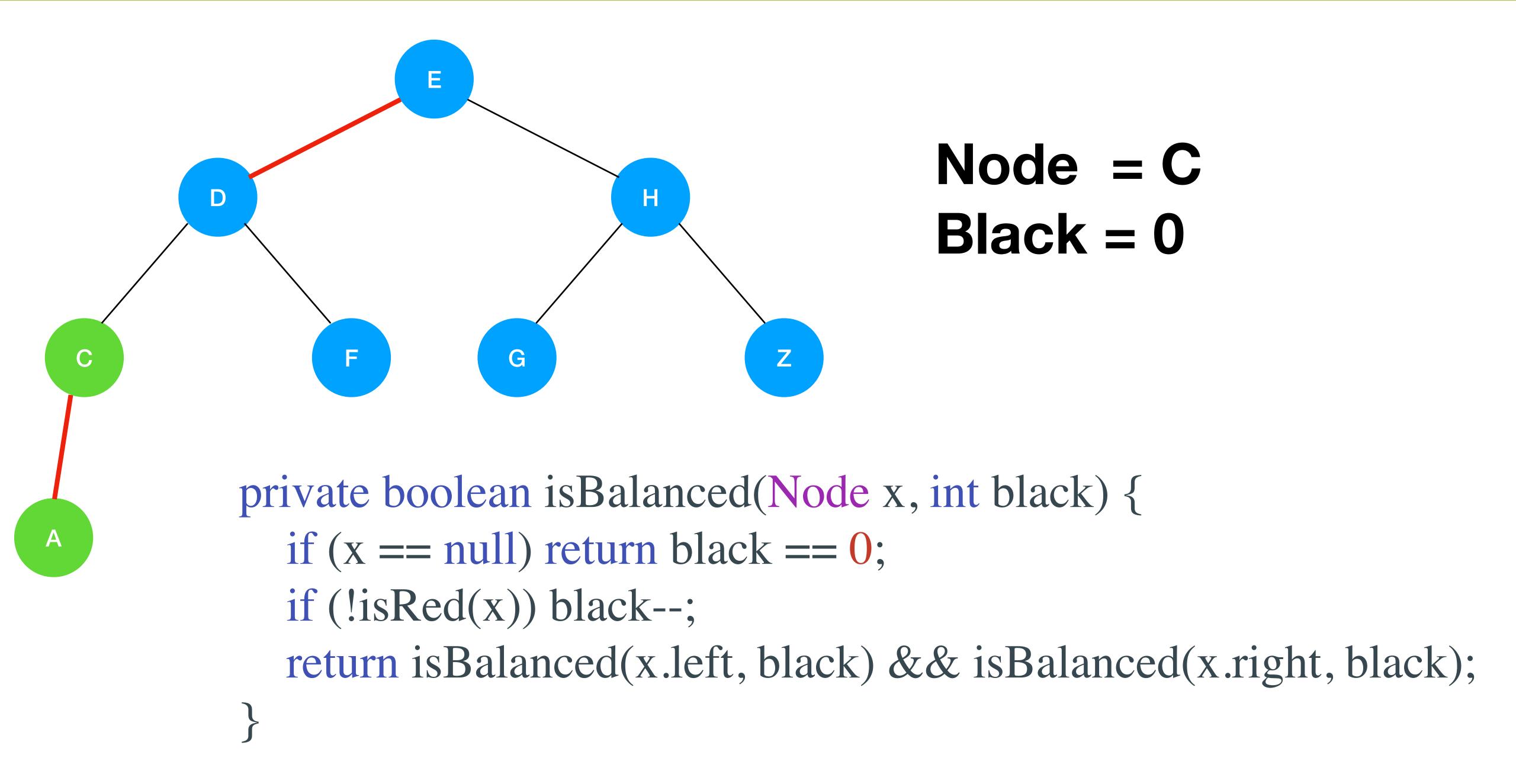


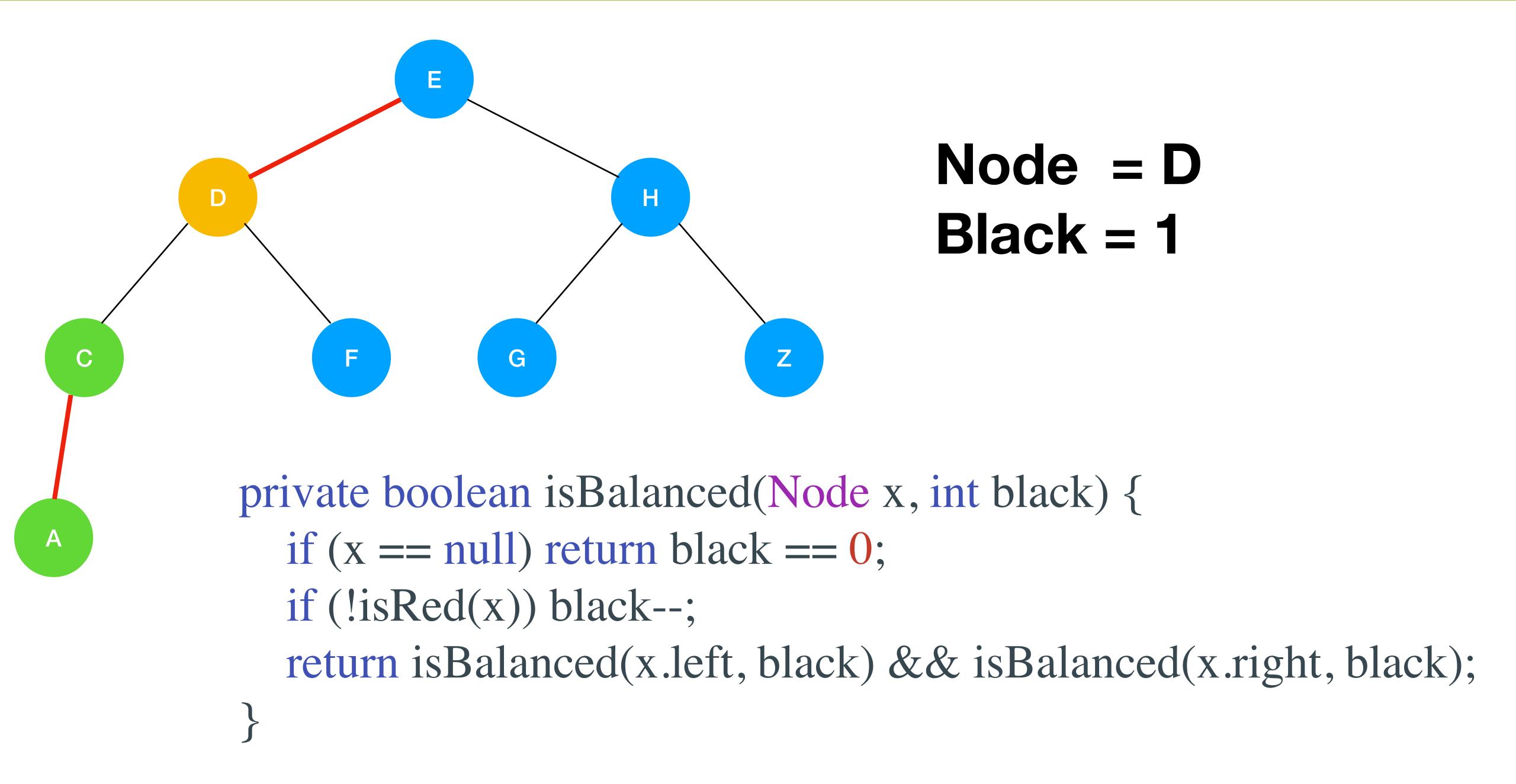


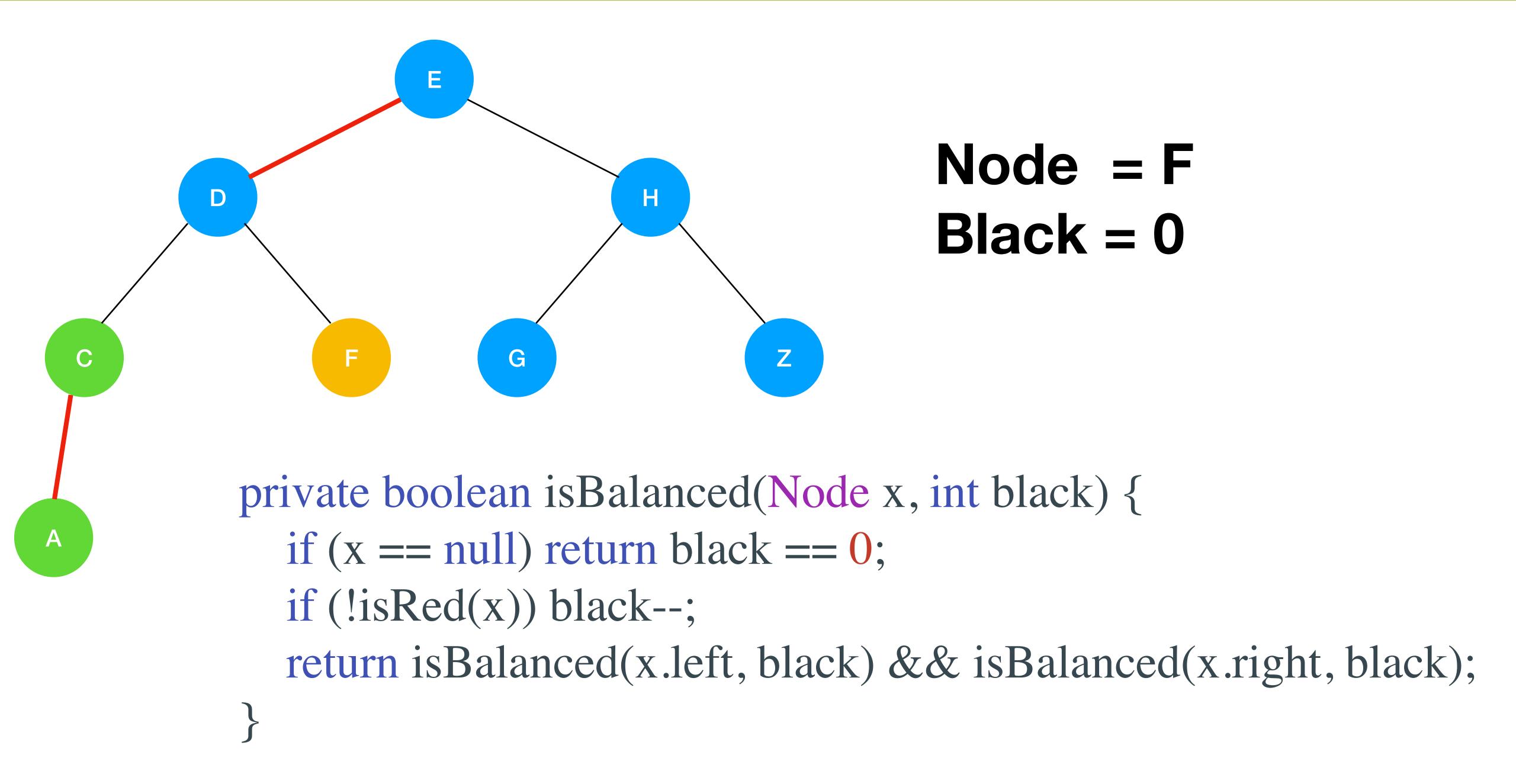


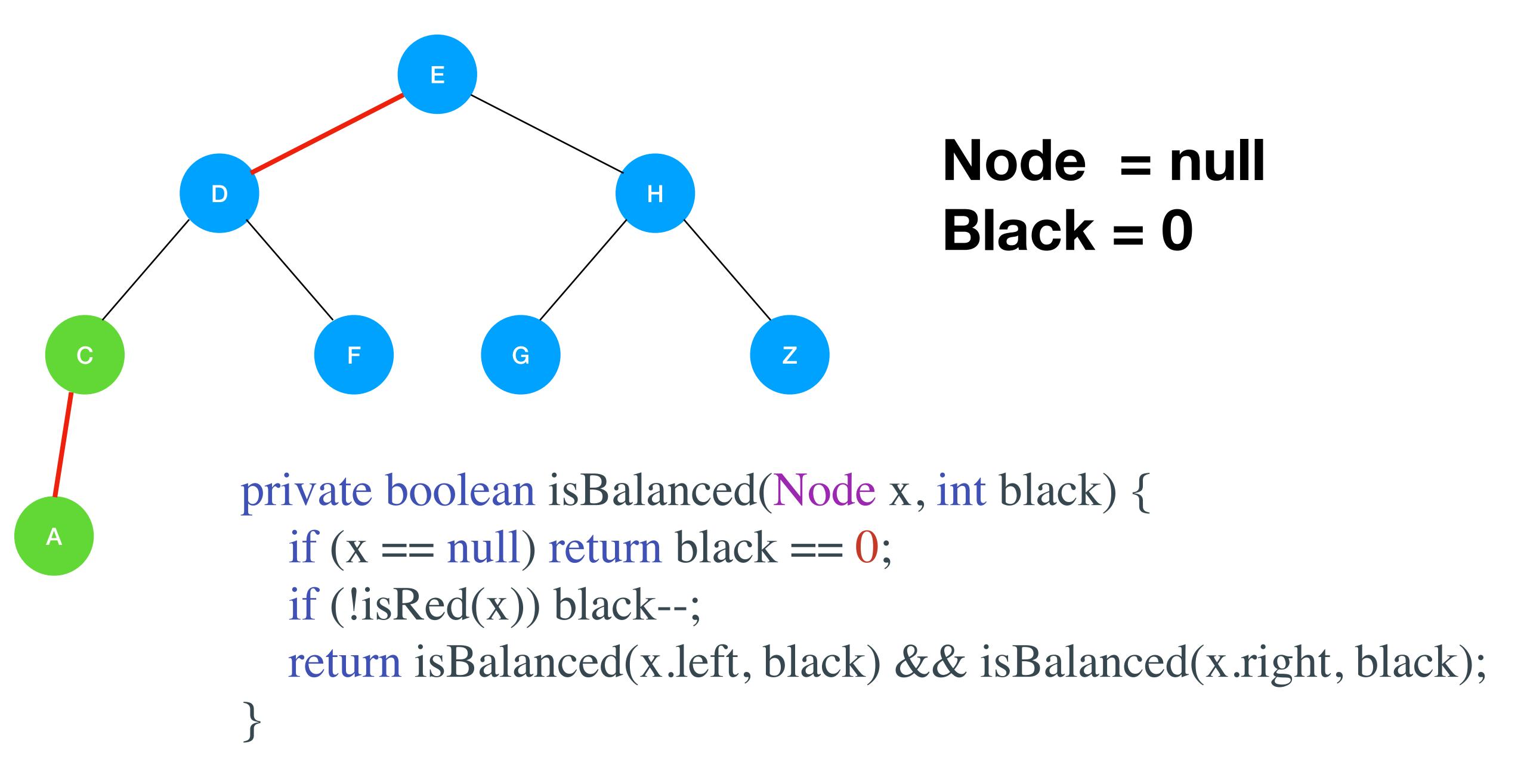


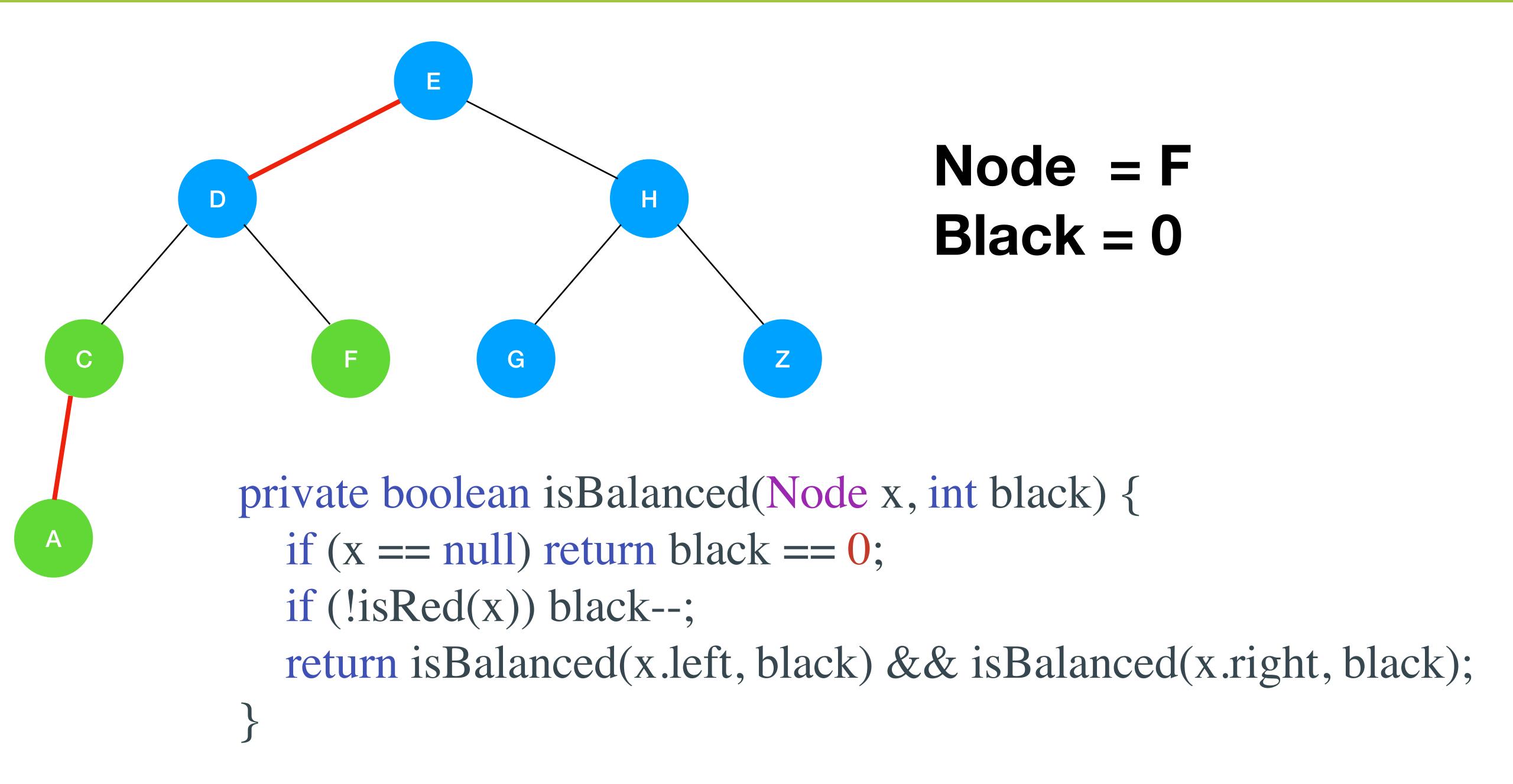


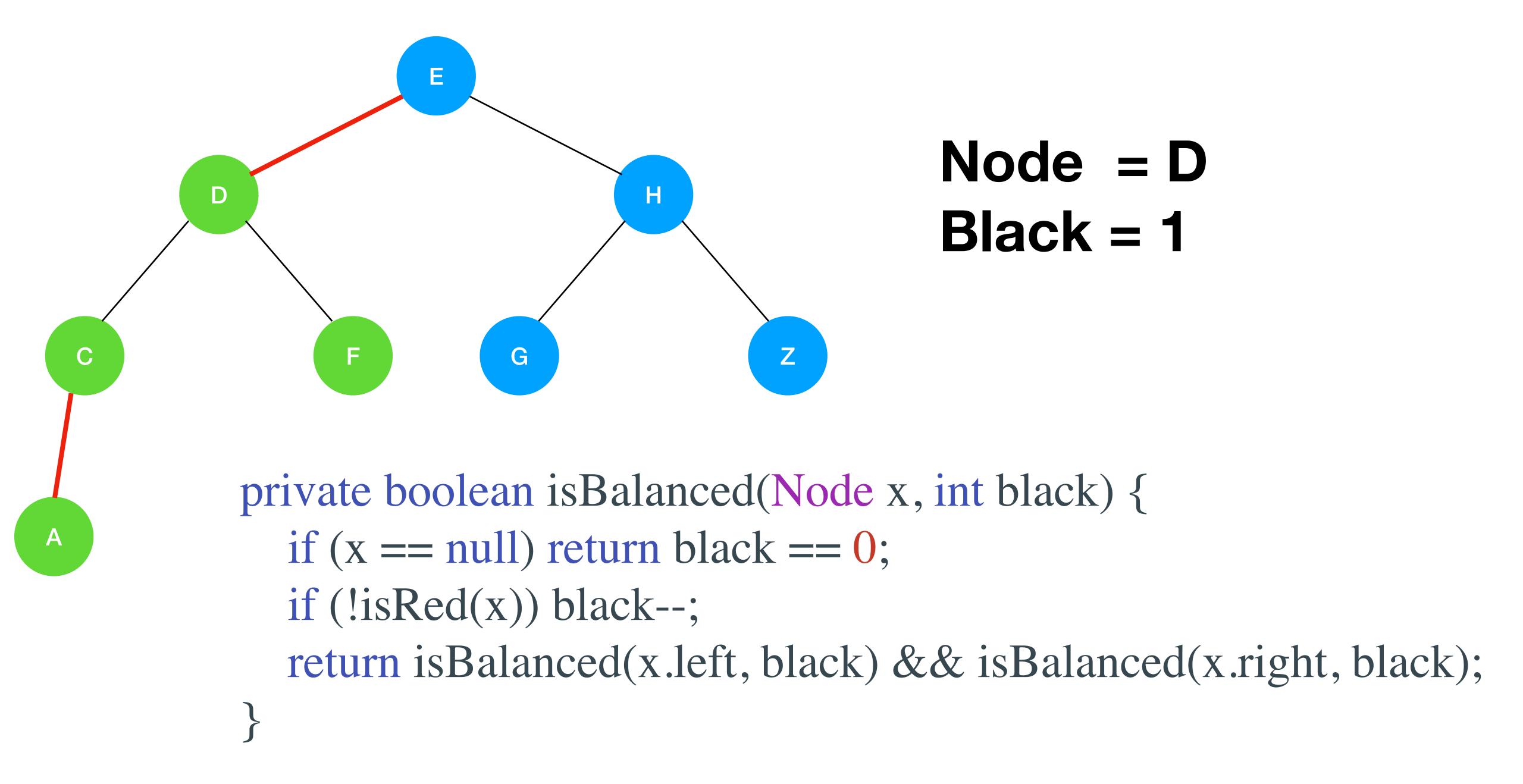


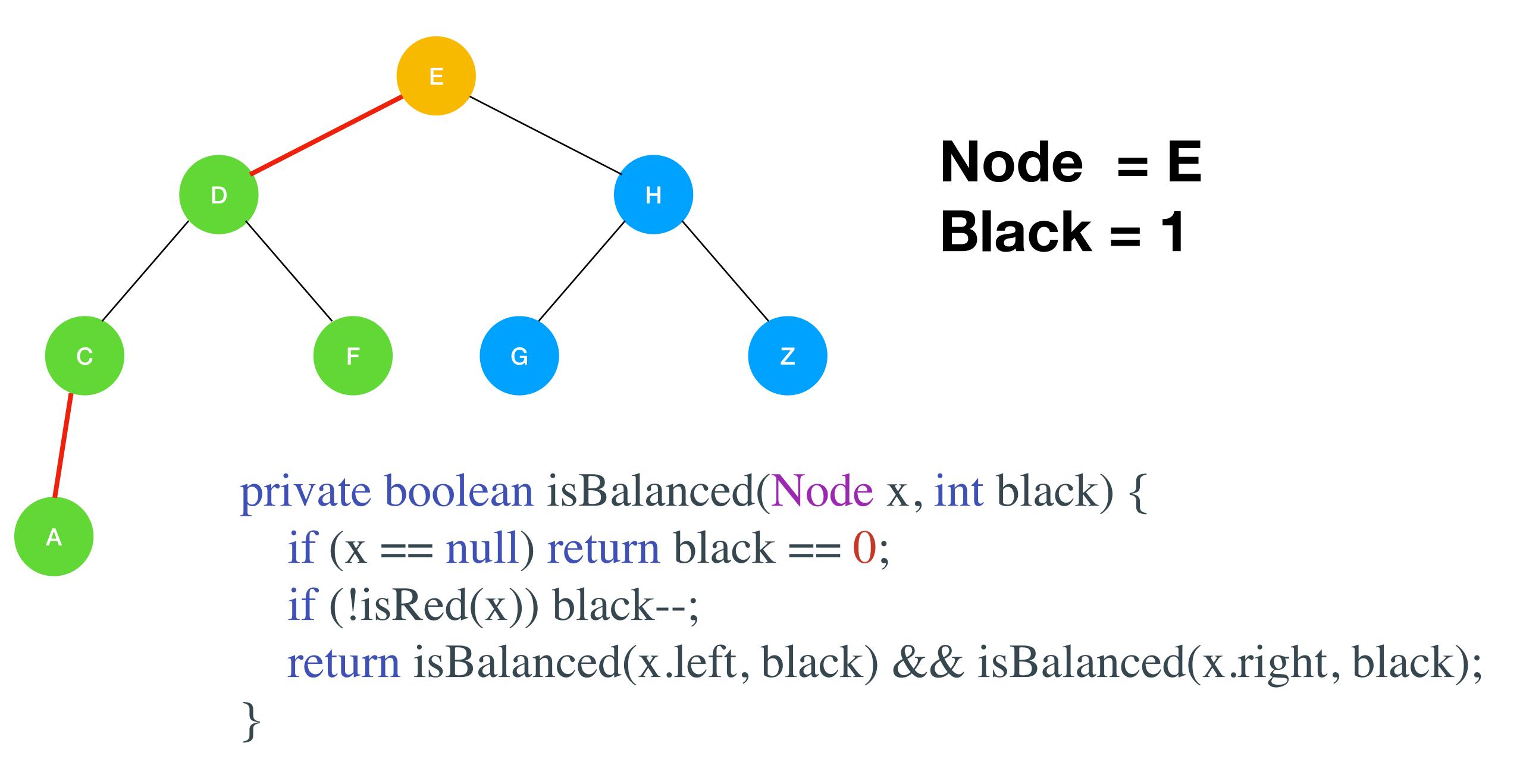


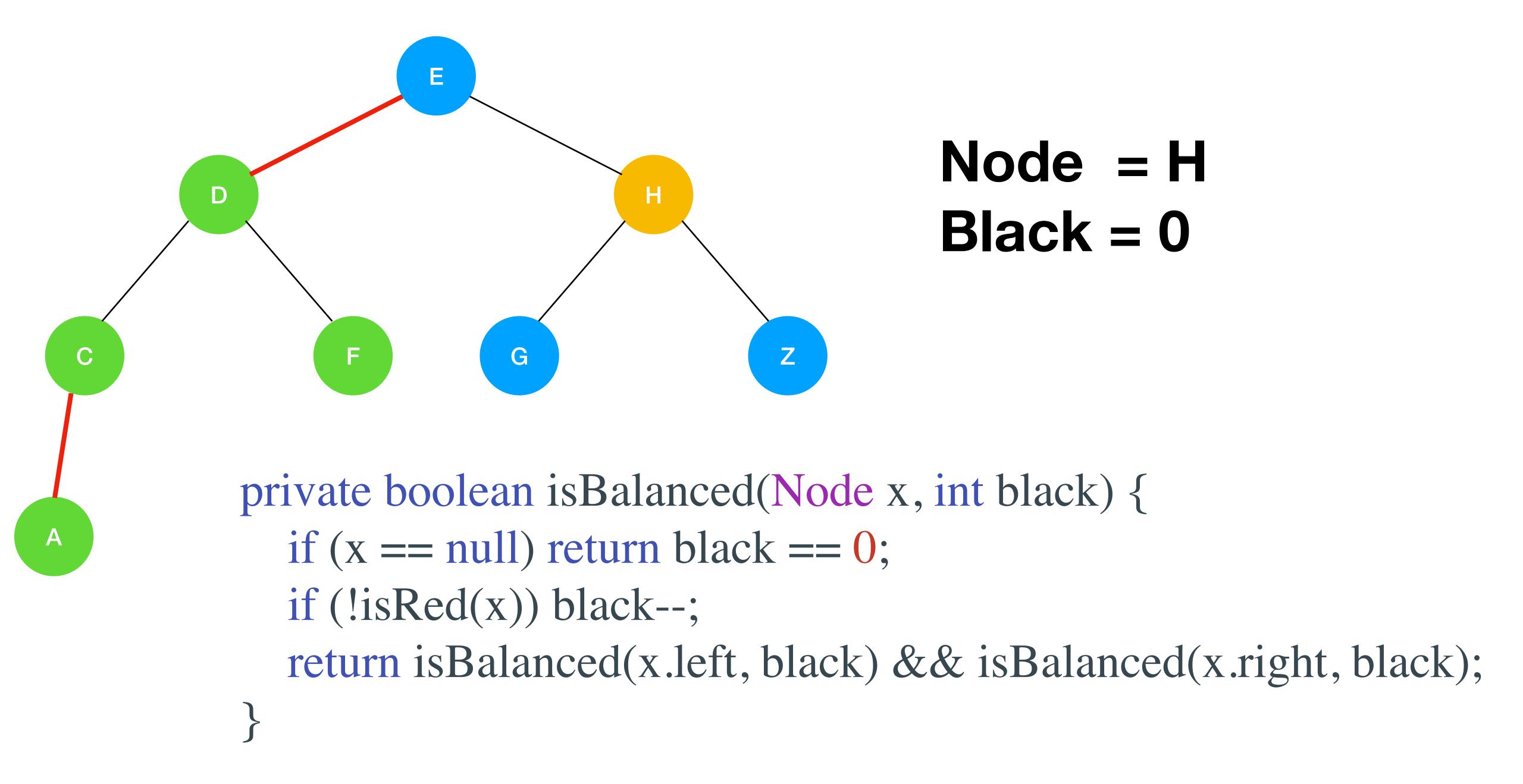


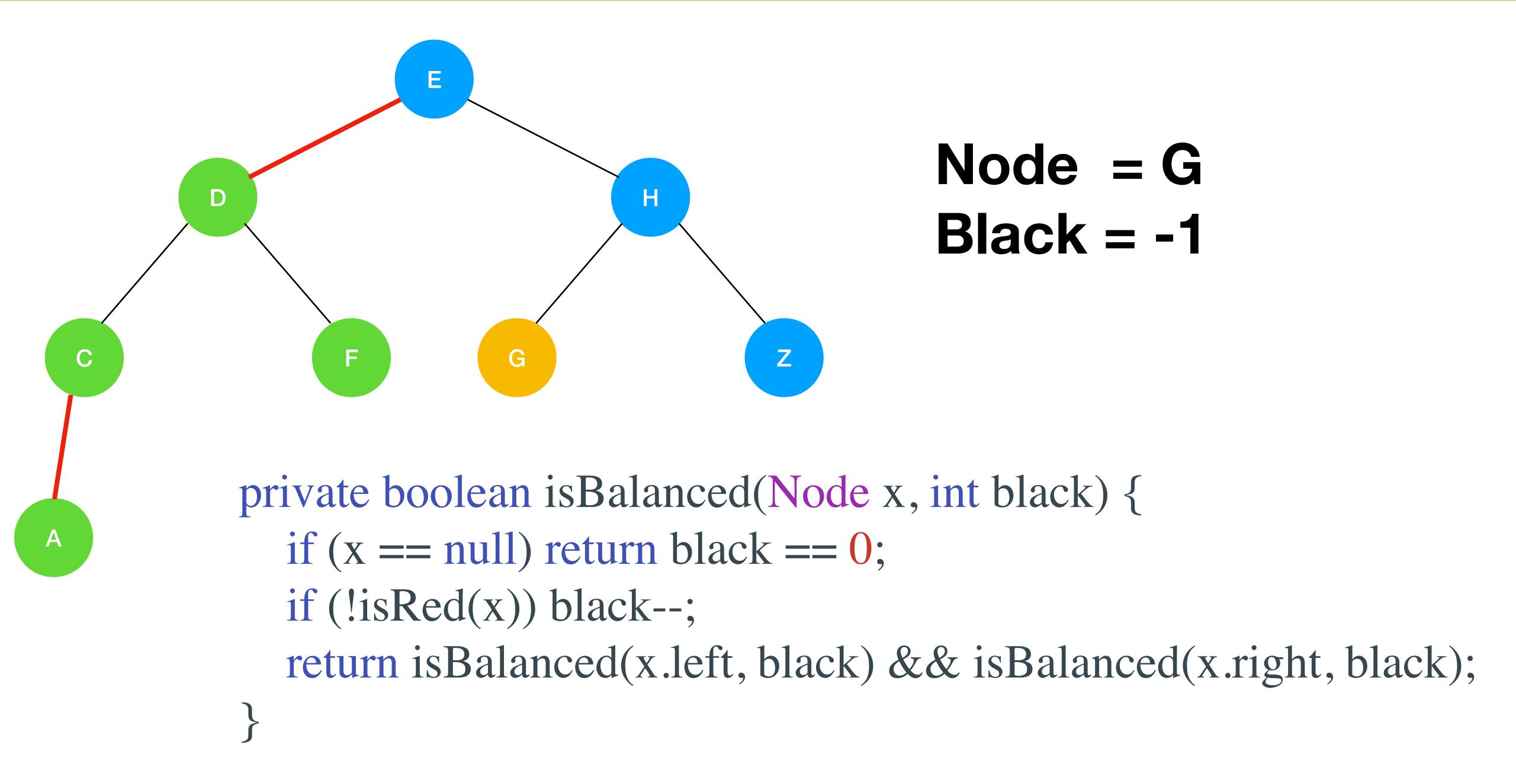


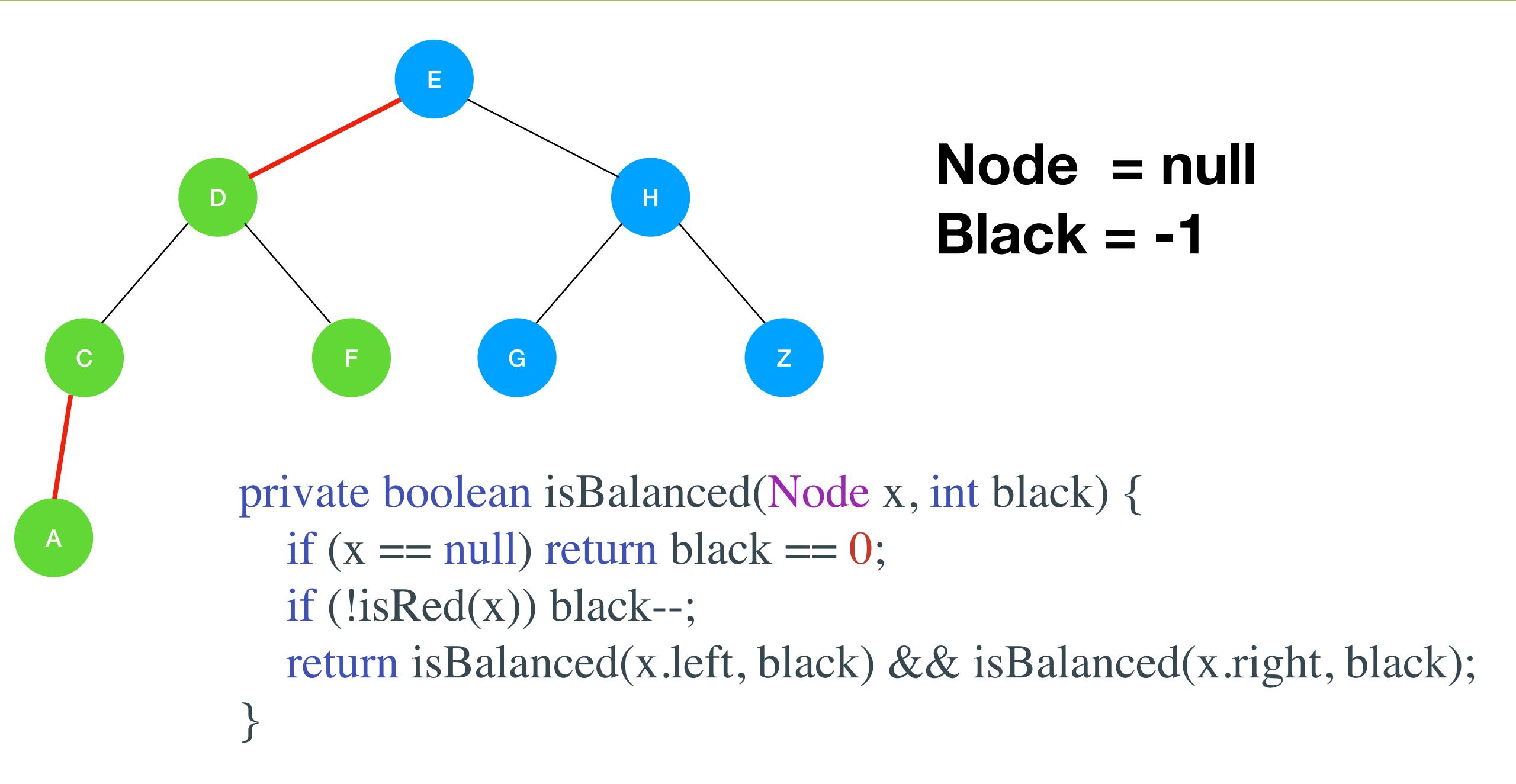


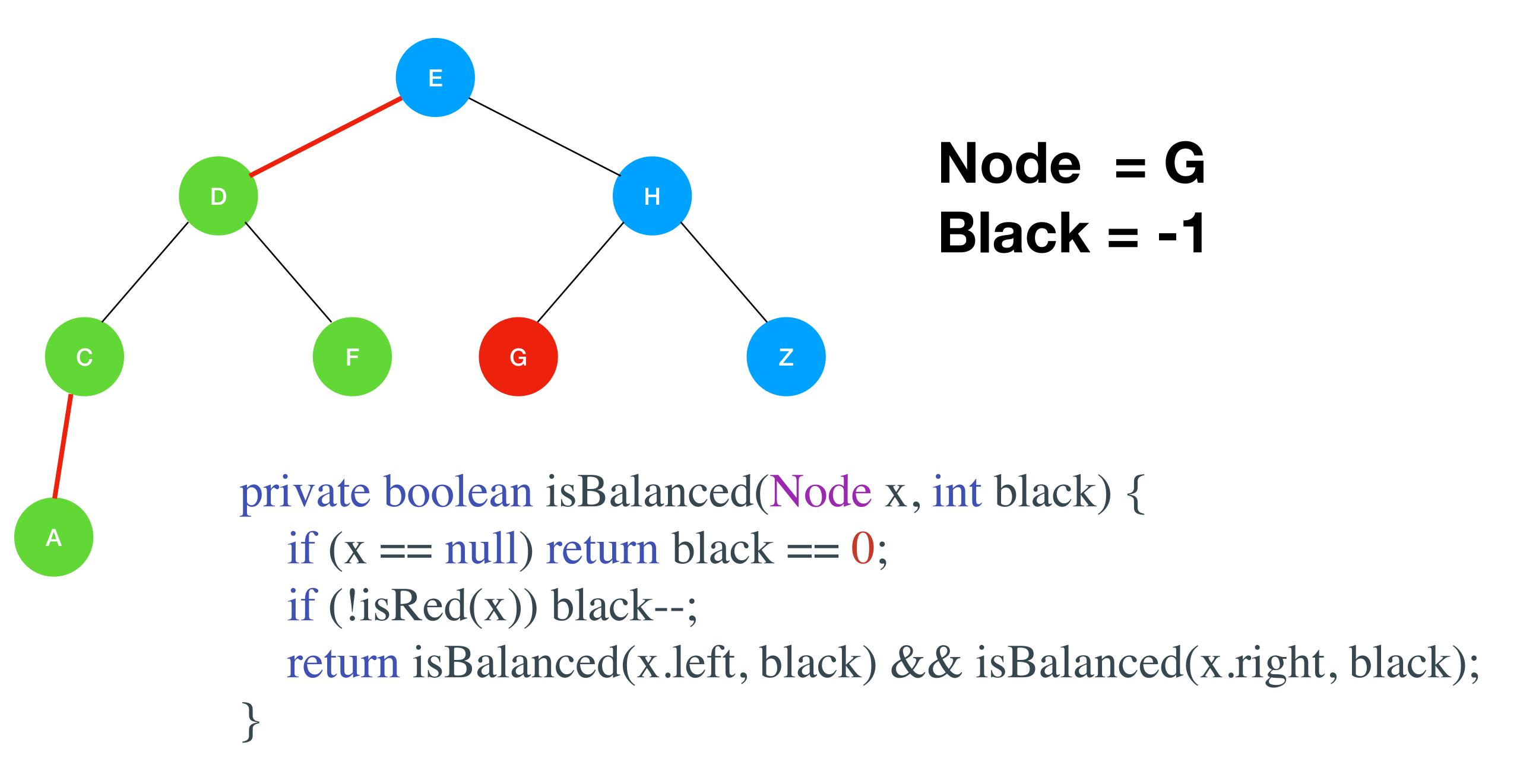












algorithm (data structure)	worst-case cost (after N inserts)		average-case cost (after N random inserts)		efficiently support ordered
	search	insert	search hit	insert	operations?
sequential search (unordered linked list)	N	N	N/2	N	no
binary search (ordered array)	$\lg N$	N	$\lg N$	N/2	yes
binary tree search (BST)	N	N	$1.39 \lg N$	$1.39 \lg N$	yes
2-3 tree search (red-black BST)	$2 \lg N$	$2 \lg N$	$1.00 \lg N$	$1.00 \lg N$	yes

Cost summary for symbol-table implementations (updated)