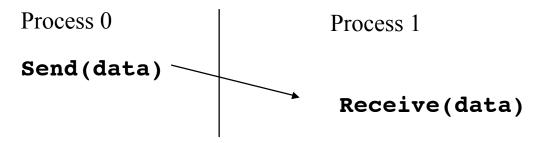
MPI Basic Send/Receive

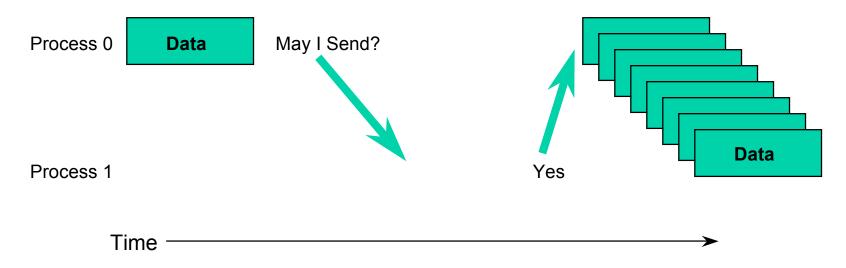
• We need to fill in the details in



- Things that need specifying:
 - How will "data" be described?
 - How will processes be identified?
 - How will the receiver *recognize/screen* messages?
 - What will it mean for these operations to complete?

What is message passing?

Data transfer plus synchronization



- Requires cooperation of sender and receiver
- Cooperation not always apparent in code

MPI Data Types

- The data in a message to sent or received is described by a triple (address, count, datatype), where
- An MPI *datatype* is recursively defined as:
 - predefined, corresponding to a data type from the language (e.g., MPI_INT, MPI_DOUBLE_PRECISION)
 - a contiguous array of MPI datatypes
 - a strided block of datatypes
 - an indexed array of blocks of datatypes
 - an arbitrary structure of datatypes
- There are MPI functions to construct custom datatypes, such an array of (int, float) pairs, or a row of a matrix stored columnwise.

MPI Tags

- Messages are sent with an accompanying userdefined integer *tag*, to *assist* the receiving process in identifying the message.
- Messages can be screened at the receiving end by specifying a specific tag, or not screened by specifying MPI_ANY_TAG as the tag in a receive.
- Some non-MPI message-passing systems have called tags "message types". MPI calls them tags to avoid confusion with datatypes.

MPI Basic (Blocking) Send

MPI_SEND (start, count, datatype, dest, tag, comm)

- The message buffer is described by (start, count, datatype).
- The target process is specified by **dest**, which is the rank of the target process in the communicator specified by **comm**.
- When this function returns, the data has been *delivered* to the system and the buffer can be reused. The message may *not* have been received by the target process.

MPI Basic (Blocking) Receive

MPI_RECV(start, count, datatype, source, tag, comm, status)

- Waits until a matching (on **source** and **tag**) message is received from the system, and the buffer can be used.
- source is rank in communicator specified by comm, or MPI ANY SOURCE.
- status contains further information
- Receiving fewer than **count** occurrences of **datatype** is OK, but receiving more is an error.

Retrieving Further Information

- **Status** is a data structure allocated in the user's program.
- In C:

```
int recvd tag, recvd from, recvd count;
   MPI Status status;
   MPI Recv(..., MPI ANY SOURCE, MPI ANY TAG, ..., &status)
   recvd tag = status.MPI TAG; //actual tag
   recvd from = status.MPI SOURCE; //actual source
   MPI Get count( &status, datatype, &recvd count );
• In Fortran:
   integer recvd tag, recvd from, recvd count
   integer status(MPI STATUS SIZE)
   call MPI RECV(..., MPI ANY SOURCE, MPI ANY TAG, .. status, ierr)
   tag recvd = status(MPI TAG)
   recvd from = status(MPI SOURCE)
   call MPI GET COUNT(status, datatype, recvd count, ierr)
```

Why Data Types?

- Since all data is labeled by type, an MPI implementation can support communication between processes on machines with very different memory representations and lengths of elementary datatypes (heterogeneous communication from PowerPC processor to Intel's).
- Specifying application-oriented layout of data in memory
 - reduces memory-to-memory copies in the implementation
 - allows the use of special hardware (scatter/gather) when available

Tags and Contexts

- Separation of messages used to be accomplished by use of tags, but
 - this requires libraries to be aware of tags used by other libraries.
 - this can be defeated by use of "wild card" tags.
- Contexts are different from tags
 - no wild cards allowed
 - allocated dynamically by the system when a library sets up a communicator for its own use.
- User-defined tags still provided in MPI for user convenience in organizing application
- Use MPI Comm split to create new communicators

Sources of Deadlocks

- Send a large message from process 0 to process 1
 - If there is insufficient storage at the destination, the send must wait for the user to provide the memory space (through a receive)
- What happens with

| Process 0 | Process 1 |
|-------------|-------------|
| Send(dst=1) | Send(dst=0) |
| Recv(src=1) | Recv(src=0) |

 This is called "unsafe" because it depends on the availability of system buffers

Some Solutions to the "unsafe" Problem

• Order the operations more carefully:

| Process 0 | Process 1 |
|-----------|-----------|
| Send(1) → | Recv(0) |
| Recv(1) ← | Send(0) |

• Use non-blocking operations:

| Process 0 | Process 1 | |
|-----------|-----------|--|
| Isend(1) | Isend(0) | |
| Irecv(1) | Irecv(0) | |
| Waitall | Waitall | |

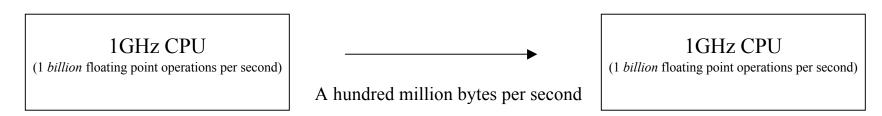
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Nonblocking Send/Receive

- How to send a large message without knowing the limitation of a computer system?
 - Use nonblocking send, MPI_Isend, and follow later with MPI_Wait
 - This can send an arbitrarily large message
 - Improve performance by not making an internal copy

Nonblocking Send/Receive (continue)

- MPI_SEND/MPI_RECV are *blocking* operations
 - A"Send" operation do not complete until matching "Receive" is issued on the destination process
- Moving and manipulating data within a single process is much faster than moving data from one process to another
 - MPI solution: nonblocking send and receive
 - Permit a program to continue to execute or compute instead of waiting for communications to complete



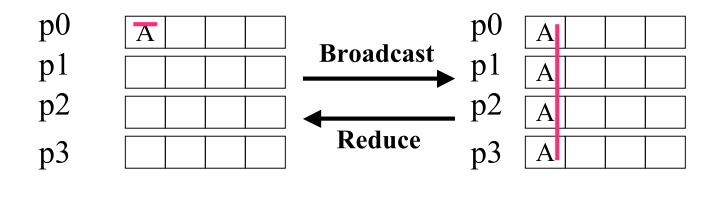
Nonblocking Send/Receive (continue)

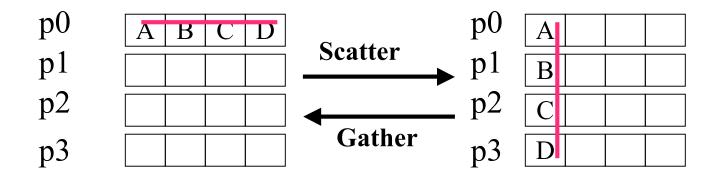
- Library for nonblocking send/receive
 - MPI_ISEND(buffer, count, datatype, dest, tag, comm, request, ierr)
 - MPI_IRECV(buffer, count, datatype, source, tag, comm, request, ierr)
 - MPI_WAIT(request, status, ierr)
 - MPI WAITALL(count, array of requests, array of status, ierr)
- Example:

```
call MPI_IRECV(..., requests(1), ierr)
call MPI_IRECV(..., requests(2), ierr)
...
call MPI_WAITALL(2, requests, status, ierr)
where "status" is declared as "integer status (MPI_STATUS_SIZE, 2)"
```

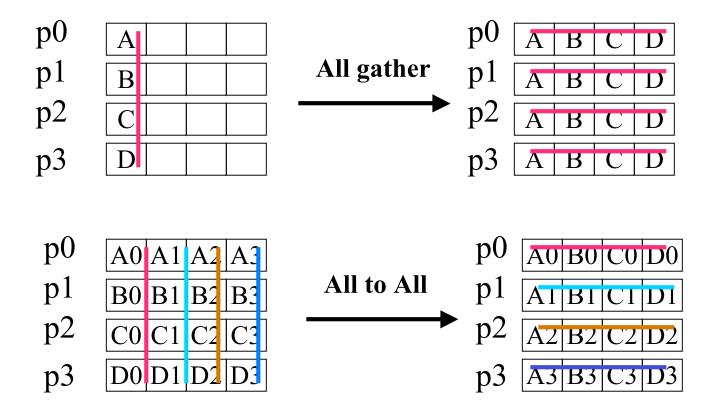
• More in the example of solving Poisson problems

Collective Operations





Advanced Collective Operations (continue)



Distributed Computing

- MPI is not suitable for distributed computing
 - The intent of MPI is to support tightly-coupled parallel computing,
 where all processes are working on the same or a similar problem
 - MPI-1 contains no support for the client-server model available in CORBA and DCOM
 - MPI-2 provide limited support for client-server model
 - OpenMPI implementation appears to support this feature
- Grid-enabled MPI is intended to perform message-passing operations via Grid
 - Appear to be inactive now

Monte Carlos Method

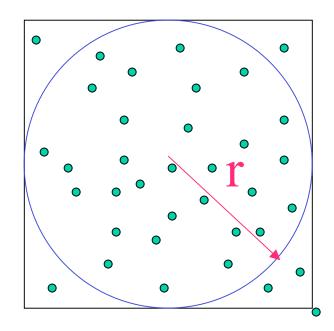
- A class of computational algorithms that rely on repeated random sampling to compute their results
- The term was coined in the 1940s by physicists working on nuclear weapon projects in the Los Alamos National Laboratory
- Tend to be used when it is infeasible or impossible to compute an exact result with a deterministic algorithm such as molecular dynamics.
- Especially useful in studying systems with a large number of coupled degrees of freedom, such as fluids, disordered materials, strongly coupled solids, and cellular structures.
- More broadly, useful for modeling phenomena with significant uncertainty in inputs, such as the calculation of risk in business.
- In finance and mathematical finance, used to value and analyze (complex) instruments, portfolios and investments by simulating the various sources of uncertainty affecting their value, and then determining their average value over the range of resultant outcomes.
 - Its advantage increases as the dimensions (sources of uncertainty) of the problem increase.

Monte-Carlo Method (Algorithm)

- 1) Specify an initial points x0 in phase space.
- 2) Generate a new state x'.
- 3) Compute the transition probability W(x,x')
- 4) Generate a uniform random number R between [0,1]
- 5) If w<R, then return to step 2
- 6) Otherwise, accept the new state and return to step 2

Monte Carlo computation of pi

- Square area = 2r * 2r
- Circle area = pi r*r
- Pi = 4*circle area/square area
- Selection criterion: A particle inside the circle or not
- The area value is represented with number of dots
- Radius = 1 for convenience
- In computation, "dimensionless variables are typically used



Monte Carlo Computation of pi (1)

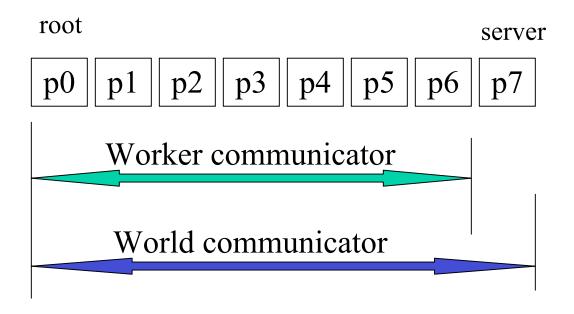
```
/* compute pi using Monte Carlo method */
#include <math.h>
#include "mpi.h"
#include "mpe.h"
#define CHUNKSIZE
                        1000
/* We'd like a value that gives the maximum value returned by the function
 random, but no such value is *portable*. RAND MAX is available on many
 systems but is not always the correct value for random (it isn't for
  Solaris). The value ((unsigned(1)<<31)-1) is common but not guaranteed */
#define INT MAX 1000000000
/* message tags */
#define REQUEST 1
#define REPLY 2
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```

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Monte Carlo Computation of pi (2)

```
int main( int argc, char *argv[])
  int iter;
  int in, out, i, iters, max, ix, iy, ranks[1], done, temp;
  double x, y, Pi, error, epsilon;
  int numprocs, myid, server, totalin, totalout, workerid;
  int rands[CHUNKSIZE], request;
  MPI Comm world, workers;
  MPI_Group world_group, worker_group;
  MPI Status status;
  MPI Init(&argc,&argv);
  world = MPI COMM WORLD;
  MPI Comm size(world, &numprocs);
  MPI Comm rank(world, &myid);
  server = numprocs-1; /* last proc is server */
```

Monte Carlo Computation of pi (0)



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Monte Carlo Computation of pi (3)

Monte Carlo Computation of pi (4)

```
if (myid == server) { /* I am the server process for generating random number*/
         do {
            MPI_Recv(&request, 1, MPI_INT, MPI_ANY_SOURCE, REQUEST,
                      world, &status);
                                                     /* request is valid */
           if (request) {
                   for (i = 0; i < CHUNKSIZE;)
                        rands[i] = random();
                        if (rands[i] <= INT_MAX) i++;</pre>
                    MPI Send(rands, CHUNKSIZE, MPI INT,
             status.MPI_SOURCE, REPLY, world);
         while(request>0);
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```

Monte Carlo Computation of pi (5)

```
/* I am one of the worker processes */
else {
   request = 1;
         done = in = out = 0;
         max = INT MAX; /* max int, for normalization */
   MPI Send( &request, 1, MPI INT, server, REQUEST, world );
   MPI Comm rank( workers, &workerid ); /* get a worker's rank */
         iter = 0:
         while (!done) {
            iter++;
            request = 1;
            MPI Recv( rands, CHUNKSIZE, MPI INT, server, REPLY,
                       world, &status);
           //process the random numbers of CHUNKSIZE
           for (i=0; i<CHUNKSIZE-1; ) {
              x = (((double) rands[i++])/max) * 2 - 1;
              y = (((double) rands[i++])/max) * 2 - 1;
                    if (x*x + y*y < 1.0)
                      in++:
                    else
                      out++;
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```

Monte Carlo Computation of pi (6)

```
MPI Allreduce(&in, &totalin, 1, MPI INT, MPI SUM,
                      workers);
MPI Allreduce(&out, &totalout, 1, MPI INT, MPI SUM,
                      workers);
Pi = (4.0 * totalin)/(totalin + totalout);
error = fabs( Pi-3.141592653589793238462643);
done = (error < epsilon || (totalin+totalout) > 1000000);
request = (done) ? 0 : 1; /*?? */
//continue the operations based on the value of request
if (myid == 0) {
         printf( "\rpi = \%23.20f", Pi );
         MPI Send(&request, 1, MPI INT, server, REQUEST,
                     world);
}
else {
         if (request)
            MPI Send(&request, 1, MPI INT, server, REQUEST,
                        world);
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```

Monte Carlo Computation of pi (7)