

# Task II.1

Before extending real code, we first extend the MiniTriangle type system for the four new constructs we added in CW1.

- **Repeat-until-loops:**

$$\frac{\Gamma \vdash e : \text{Boolean} \quad \Gamma \vdash c}{\Gamma \vdash \text{repeat } c \text{ until } e}$$

- **Conditional expressions:**

$$\frac{\Gamma \vdash e_1 : \text{Boolean} \quad \Gamma \vdash e_2 : T \quad \Gamma \vdash e_3 : T}{\Gamma \vdash (e_1 ? e_2 : e_3) : T}$$

- **Extended syntax for the if-command:**

$$\frac{\Gamma \vdash e_1 : \text{Boolean} \quad \Gamma \vdash \bar{e}_2 : \text{Boolean} \quad \Gamma \vdash c_1 \quad \Gamma \vdash \bar{c}_2 \quad \Gamma \vdash c_3^{\leq 1}}{\Gamma \vdash \text{if } e_1 \text{ then } c_1 \text{ elsif } \bar{e}_2 \text{ then } \bar{c}_2 \text{ else } c_3^{\leq 1}}$$

- **Character literals:**

$$\Gamma \vdash c : \text{Character}$$

# TaskII.2

After extending the type system, we now need to modify real code to enable checker to analyse context grammar after parser's lexical analysis.

I implemented the four constructs one at a time in the following order:

1. **Character Literals:**

- *Type.hs:*

```
71 | Character -- ^ The Character type
```

Adding Character as a new type in Type representation.

```
97 Character == Character = True
```

```
185 showsPrec _ Character = showString "Character"
```

At the same time, extended Eq and Show instance for this type.

- *MTStdEnv.hs*:

```
68      ("Character", Character)]
```

Add type `Character` to `[Types:]` in `MiniTriangle` initial environment.

```
90      ("getchr", Arr [Snk Character] Void,      ESVLbl "getchr"),
91      ("putchr", Arr [Character] Void,          ESVLbl "putchr"),
```

Add `'getchr'` and `'putchr'` to `[Procedures:]` in `MiniTriangle` initial environment.

`'getchr'` reads user input from terminal and has a write-only reference.

In `MTIR.hs` and `PPMTIR.hs`, `Character` has already been implemented. `LibMT.hs` will be modified in TaskII.3 to agree with the code we modified here.

- *TypeChecker.hs*:

```
353 -- T-LITCHR
354 infTpExp env e@(A.ExpLitChr {A.elcVal = c, A.expSrcPos = sp}) = do
355   c' <- toMTChr c sp
356   return (Character,                                     -- env |- n : Character
357          ExpLitChr {elcVal = c', expType = Character, expSrcPos = sp})
```

Check type of the expression.

Return the expected type and expression in Monad Type D.

## 2. Repeat-until-loops:

- *MTIR.hs*:

```
78      -- | Repeat-loop
79      | CmdRepeat {
80          crBody    :: Command,          -- ^ Loop-body
81          crCond    :: Expression,       -- ^ Loop-condition
82          cmdSrcPos :: SrcPos
83      }
```

Add repeat-until-loops to `Command` in `MiniTriangle` internal representation.

In `MTIR`, `Command` has the same structure as in `AST`.

- *PPMTIR.hs*:

```
64 ppCommand n (CmdRepeat {crBody = c, crCond = e, cmdSrcPos = sp}) =
65   indent n . showString "CmdRepeat" . spc . ppSrcPos sp . nl
66   . ppCommand (n+1) c
67   . ppExpression (n+1) e
```

Same to the pretty printer for `MTIR`.

- *TypeChecker.hs*:

```

121 -- T-REPEAT
122 chkCmd env (A.CmdRepeat {A.crBody = c, A.crCond = e, A.cmdSrcPos = sp}) = do
123   c' <- chkCmd env c                                -- env |- c
124   e' <- chkTpExp env e Boolean                       -- env |- e : Boolean
125   return (CmdRepeat {crBody = c', crCond = e', cmdSrcPos = sp})

```

Checks command and type of expression (Loop-condition should be a Boolean).

Transfer from AST Type to Monad Type D and return them for further manipulation.

### 3. Conditional Expressions:

- *MTIR.hs*:

```

168 | ExpCond {
169     ecCond    :: Expression,      -- ^ Condition
170     ecTrue    :: Expression,      -- ^ Value if condition true
171     ecFalse   :: Expression,      -- ^ Value if condition false
172     expType   :: Type,
173     expSrcPos :: SrcPos
174 }

```

Add conditional expression to Expression in MiniTriangle internal representation.

There is an expected return type in conditional expression (different from AST).

- *PPMTIR.hs*:

```

120 ppExpression n (ExpCond {ecCond = c, ecTrue = t, ecFalse = f, expType = et, expSrcPos = sp}) =
121   indent n . showString "ExpCond" . spc . ppSrcPos sp . nl
122   . ppExpression (n+1) c
123   . ppExpression (n+1) t
124   . ppExpression (n+1) f
125   . indent n . showString ": " . shows et . nl

```

Same to the pretty printer for MTIR.

- *TypeChecker.hs*:

```

427 -- T-COND
428 infTpExp env (A.ExpCond {A.ecCond = e1, A.ecTrue = e2, A.ecFalse = e3, A.expSrcPos = sp}) = do
429   e1' <- chkTpExp env e1 Boolean
430   (t, e2') <- infNonRefTpExp env e2
431   (t, e3') <- infNonRefTpExp env e3
432   return (t, ExpCond {ecCond = e1', ecTrue = e2', ecFalse = e3', expType = t, expSrcPos = sp})

```

Since conditional expression is an expression, it should have a return type.

However, e2 and e3 allow to have variables in them, which are reference types.

Therefore, if we need checker to do contextual analysis, we first need to use `infNonRefTpExp` function to resolve the type of e2 and e3 before checker's contextual analysis.

When we use `infNonRefTpExp` function, we assume that e2 and e3 have the same type. Thus, the whole conditional expression will only have one kind of return type.

### 4. Extension for if-then-else-commands:

- *MTIR.hs*:

```

65     -- | Conditional command
66     | CmdIf {
67         ciCondThens    :: [(Expression,
68                             Command)],           -- ^ Conditional branches
69         ciMbElse       :: Maybe Command,         -- ^ Optional else-branch
70         cmdSrcPos      :: SrcPos
71     }

```

Extend if-then-else commands to Command in MiniTriangle internal representation (same as AST).

- *PPMTIR.hs*:

```

56 ppCommand n (CmdIf {ciCondThens = ecs, ciMbElse = mc, cmdSrcPos = sp}) =
57     indent n . showString "CmdIf" . spc . ppSrcPos sp . nl
58     . ppSeq (n+1) (\n (e,c) -> ppExpression n e . ppCommand n c) ecs
59     . ppOpt (n+1) ppCommand mc

```

Same to the pretty printer for MTIR (same as PPAST).

- *TypeChecker.hs*:

```

96 -- T-IF
97 chkCmd env (A.CmdIf {A.ciCondThens = ecs, A.ciMbElse = mc2,
98                     A.cmdSrcPos=sp}) = do
99     ecs' <- mapM (chkOptIf env) ecs           -- env |- ecs
100     mc2' <- case mc2 of                       -- env |-
101         Just c2 -> do
102             c2' <- chkCmd env c2
103             return (Just c2')
104         Nothing -> do
105             return Nothing
106     return (CmdIf {ciCondThens = ecs', ciMbElse = mc2', cmdSrcPos = sp})

133 chkOptIf :: Env -> (A.Expression, A.Command) -> D (Expression, Command)
134 chkOptIf env (e, c) = do
135     e' <- chkTpExp env e Boolean             -- env |- e : Boolean
136     c' <- chkCmd env c                       -- env |- c
137     return (e', c')

```

ecs represents a list of tuples, consist of an expression and a command (ecs : [(e, c)]).

mc2 represents a command wrapped with a Maybe type. Maybe type allows the else-branch can be missing.

chkOptIf is a helper function which checks a single tuple consisted of an expression (should be a Boolean) and a command and transfer them into Monad Type D and return the new tuple.

MapM function enables chkOptIf to operate on every element of ecs list.

# TaskII.3

- *myTamCode3a:*

|    |        |         |          |
|----|--------|---------|----------|
| 1  |        | GETINT  |          |
| 2  |        | LOAD    | [SB + 0] |
| 3  |        | LOADL   | 0        |
| 4  |        | GTR     |          |
| 5  |        | JUMPIFZ | #end     |
| 6  |        | LOADL   | 0        |
| 7  | #loop: |         |          |
| 8  |        | LOADL   | 1        |
| 9  |        | ADD     |          |
| 10 |        | LOAD    | [SB + 1] |
| 11 |        | PUTINT  |          |
| 12 |        | LOAD    | [SB + 0] |
| 13 |        | LOAD    | [SB + 1] |
| 14 |        | EQL     |          |
| 15 |        | JUMPIFZ | #loop    |
| 16 | #end:  |         |          |
| 17 |        | HALT    |          |

- *myTamCode3b:*

|    |             |         |            |
|----|-------------|---------|------------|
| 1  |             | GETINT  |            |
| 2  |             | LOADL   | 0          |
| 3  |             | LOAD    | [SB + 0]   |
| 4  |             | LOADA   | [SB + 1]   |
| 5  |             | CALL    | #0_fac     |
| 6  |             | PUTINT  |            |
| 7  |             | HALT    |            |
| 8  | #0_fac:     |         |            |
| 9  |             | LOAD    | [LB - 2]   |
| 10 |             | LOADL   | 2          |
| 11 |             | LSS     |            |
| 12 |             | JUMPIFZ | #recursive |
| 13 |             | LOADL   | 1          |
| 14 |             | LOAD    | [LB - 1]   |
| 15 |             | STOREI  | 0          |
| 16 |             | RETURN  | 0 2        |
| 17 | #recursive: |         |            |
| 18 |             | LOAD    | [LB - 2]   |
| 19 |             | LOADL   | 1          |
| 20 |             | SUB     |            |
| 21 |             | LOAD    | [LB - 1]   |
| 22 |             | CALL    | #0_fac     |
| 23 |             | LOAD    | [LB - 2]   |
| 24 |             | LOAD    | [LB - 1]   |
| 25 |             | LOADI   | 0          |
| 26 |             | MUL     |            |
| 27 |             | LOAD    | [LB - 1]   |
| 28 |             | STOREI  | 0          |
| 29 |             | RETURN  | 0 2        |

- *LibMT.hs:*

```

165 -- getchr
166   Label "getchr",
167   GETCHR,
168   LOAD (LB (-1)),
169   STOREI 0,
170   RETURN 0 1,
171
172 -- putchr
173   Label "putchr",
174   LOAD (LB (-1)),
175   PUTCHR,
176   RETURN 0 1,

```

Add 'getchr' and 'putchr' to agree with codes we modified in TaskII.2.

## TaskII.4

- *Repeat-until-loop:*

```

180 execute maj1 env n (CmdRepeat {crBody = c, crCond = e}) = do
181   lblLoop <- newName
182   emit (Label lblLoop)
183   execute maj1 env n c
184   evaluate maj1 env e
185   emit (JUMPIFZ lblLoop)

```

Repeat-until-loop is similar with While-loop.

The only difference is the Repeat-until-loop will always run for at least one time. Therefore, there is no need to check loop-condition at first.

- *Conditional Expressions:*

```

408 evaluate maj1 env (ExpCond {ecCond = e1, ecTrue = e2, ecFalse = e3, expType = t}) = do
409   lblFalse <- newName
410   lblFinish <- newName
411   evaluate maj1 env e1
412   emit (JUMPIFZ lblFalse)
413   evaluate maj1 env e2
414   emit (JUMP lblFinish)
415   emit (Label lblFalse)
416   evaluate maj1 env e3
417   emit (Label lblFinish)

```

Conditional Expression is similar with simple If-Command. Since in the generator, there is no worry about types. Just add expected return type and appropriate function ('evaluate' not 'execute') to it.

- *If-then-else-command:*

```

162 execute maj1 env n (CmdIf {ciCondThens = ecs, ciMbElse = mc2}) = do
163     lblFinish <- newName
164     mapM (exeOptIf maj1 env n lblFinish) ecs
165     case mc2 of
166         Just c2 -> do
167             execute maj1 env n c2
168         Nothing -> do
169             return ()
170     emit (Label lblFinish)

191 exeOptIf :: MSL -> CGEnv -> MTInt -> Name -> (Expression, Command) -> TAMCG ()
192 exeOptIf maj1 env n lblFinish (e, c) = do
193     lblOptIf <- newName
194     evaluate maj1 env e
195     emit (JUMPIFZ lblOptIf)
196     execute maj1 env n c
197     emit (JUMP lblFinish)
198     emit (Label lblOptIf)

```

Similar with what I have done for the Typechecker.hs.

exeOptIf helper function can generate code for every if-then branch and elsif-then branch.

Wrapped with Maybe type allows else-branch to be optional.