

# CSCI-4116

## Assignment 2

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January 24, 2025

### Question 1

#### Part a)

We know that a residue class  $a + m\mathbb{Z}$  is invertible iff  $\gcd(a, m) = 1$ . Thus only the classes represented by  $\{1, 3, 5, 7, 9, 11, 13, 15\}$  are invertible.

Given a residue class  $a + m\mathbb{Z}$  the inverse is a residue class  $a^{-1} + m\mathbb{Z}$  such that  $a \times a^{-1} \equiv 1 \pmod{16}$ .<sup>1</sup> So we have it that:

- $1 \times 1 \equiv 1 \pmod{16}$
- $3 \times 11 \equiv 1 \pmod{16}$
- $5 \times 13 \equiv 1 \pmod{16}$
- $7 \times 7 \equiv 1 \pmod{16}$
- $9 \times 9 \equiv 1 \pmod{16}$
- $15 \times 15 \equiv 1 \pmod{16}$

**aside:** The format  $x \times y \equiv 1 \pmod{m}$  just means that  $y + m\mathbb{Z}$  is the inverse of  $x + m\mathbb{Z}$  and vice versa. Since  $(\mathbb{Z}/m\mathbb{Z}, \cdot)$  is a commutative monoid we know that  $a \cdot a^{-1} = a^{-1} \cdot a$ .

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<sup>1</sup>I'm using the euclidean algorithm to find  $a^{-1}$  but immiting the steps

**Part b)****Group of Units**

The group of units is the group of all invertible residue classes, and thus, is the group of all residue classes  $a + 15\mathbb{Z}$  where  $\gcd(a, 15) = 1$ . Then given the inverses(similar to partA):

- $1 \times 1 \equiv 1(\text{mod } 15)$
- $2 \times 8 \equiv 1(\text{mod } 15)$
- $4 \times 4 \equiv 1(\text{mod } 15)$
- $7 \times 13 \equiv 1(\text{mod } 15)$
- $11 \times 11 \equiv 1(\text{mod } 15)$
- $14 \times 14 \equiv 1(\text{mod } 15)$

We have it that  $(\mathbb{Z}/15\mathbb{Z})^* = \{1, 2, 4, 7, 8, 11, 13, 14\}$

**Zero Divisors:**

We know that the zero divisors of  $\mathbb{Z}/m\mathbb{Z}$  are the residue classes where  $1 < \gcd(a, m) < m$ . This means that the zero divisors is the set of all residue classes that are not in  $(\mathbb{Z}/m\mathbb{Z})^*$  and not the class  $0 + m\mathbb{Z}$ . For  $m = 15$  the zero divisors are:

- $3 + 15\mathbb{Z}$
- $5 + 15\mathbb{Z}$
- $6 + 15\mathbb{Z}$
- $9 + 15\mathbb{Z}$
- $10 + 15\mathbb{Z}$
- $12 + 15\mathbb{Z}$

## Question 2

b) is cryptosystem while a) is not.

### Reason:

Recall that one property of a cryptosystem is that for all encryption keys in the keyspace there must exist a decryption key such that decrypting the encryption returns the original plain text.

In the context of the provided scheme, the mapping from  $x(\text{mod } m) \mapsto kx(\text{mod } m)$  can be reversed by multiplying  $k$  with its inverse in modulo  $m$ .

Since scheme a) puts no restrictions on  $k$  this means that all keys in the keyspace must be invertible in  $m$ , which is only true when  $m$  is a prime number, and 26 is clearly not. Thus there exists a key where its inverse,  $k^{-1}$ , is not uniquely identifiable. Therefore there is no guarantee that  $\mathcal{D}_{k^{-1}}(\mathcal{E}_k(p)) = p$ , thus the property is not satisfied. (The restriction in scheme b guarantees that all keys are invertible).

### Cryptosystem's details: (scheme b)

- Plaintext Space:  $\mathbb{Z}_{26}$
- Ciphertext Space:  $\mathbb{Z}_{26}$
- Key Space:  $\{1, 3, 5, 7, 9, 11, 13, 15, 17, 19, 21, 23, 25\}$ .<sup>2</sup>

## Question 3

1.  $\varphi(2024) = 880$
2.  $\varphi(2025) = 1080$
3.  $\varphi(8958) = 2984$

3

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<sup>2</sup>any value in  $(\mathbb{Z}/26\mathbb{Z})^*$  is a valid key for the mapping.

<sup>3</sup>I wrote a C++ program to calculate Question 3 for me

## Question 4

### Plain and Cipher Spaces:

The plain and cipher text can be any string over  $\Sigma$ . Thus the Plain text and Cipher text space is the set of all possible strings over the alphabet, denoted as  $\Sigma^*$

aside:

The alphabet  $\Sigma = \{A..Z\}$  can be mapped to integers  $\mathbb{Z}_{26} = \{0..25\}$  in lexicographical order.

### Encryption and Decryption Functions

The described encryption and decryption procedures can each be represented as the composition of 2 functions:

- Encryption: function  $\mathcal{E}'$  that shifts symbols and a function  $R$  that reverses a sequence
- Decryption: function  $R$  that reverses a sequence and a function  $\mathcal{D}'$  that shifts symbols

First, let a string/sequence  $X = \{x_i\}_{i=1}^n$

We now define the functions  $\mathcal{E}'$ ,  $\mathcal{D}'$  and  $R$  as follows :

$$\mathcal{E}'_{k_1, k_2}(X): \quad \begin{array}{ll} x_i + k_2 \pmod{26}, & i|2 \\ x_i + k_1 \pmod{26}, & \text{otherwise} \end{array}$$

$$\mathcal{D}'_{k_1, k_2}(X): \quad \begin{array}{ll} x_i - k_2 \pmod{26}, & i|2 \\ x_i - k_1 \pmod{26}, & \text{otherwise} \end{array}$$

$$R(X): \quad x_i = x_{n-i+1}$$

We can now write the Encryption and Decryption functions as:

$$\mathcal{E} = \mathcal{E}'_{k_1, k_2} \circ R$$

$$\mathcal{D} = R \circ \mathcal{D}'_{k_1, k_2}$$

### Key Space:

Observe that the symbols and the operations  $\mathcal{E}'$  and  $\mathcal{D}'$  that we perform on them is just the group  $(\mathbb{Z}/m\mathbb{Z}, +)$  with  $m = 26$ . So the inverse of any  $k \pmod{26}$  is  $-k \pmod{26}$ . This means that any integer value of  $k$  is a valid key. Thus the keyspace is  $\mathbb{Z}_{26}$ .

### Conclusion:

Since the procedure has:

1. a Plain text space
2. a Cipher text space
3. a Key space

4. an Encryption Function
5. a Decryption Function
6. Satisfies that  $\mathcal{D}(\mathcal{E}(p)) = p$  (since all keys are invertible)

This means that it is a valid cryptosystem as it satisfies all the properties of one.

## Question 5

Given the encryption function:

$$\mathcal{E}(x) = ax + b \pmod{m}$$

And the restriction that  $a$  needs to have a multiplicative inverse in  $m$ , we have it that for:

- **m=30:**

$b \in \mathbb{Z}_{30}$  so  $b$  can have 30 different values

$a \in (\mathbb{Z}/30\mathbb{Z})^*$  so  $a$  can have  $\varphi(30) = 8$

Thus there are  $30 * 8 = 240$  possible keys

- **m=29:**

$b \in \mathbb{Z}_{29}$  so  $b$  can have 29 different values

$a \in (\mathbb{Z}/29\mathbb{Z})^*$  so  $a$  can have  $\varphi(29) = 28$

Thus there are  $29 * 28 = 812$  possible keys