HOMEWORK 4

APOSTOLOPOULOU IOANNA
TOLOUDIS PANAGIOTIS



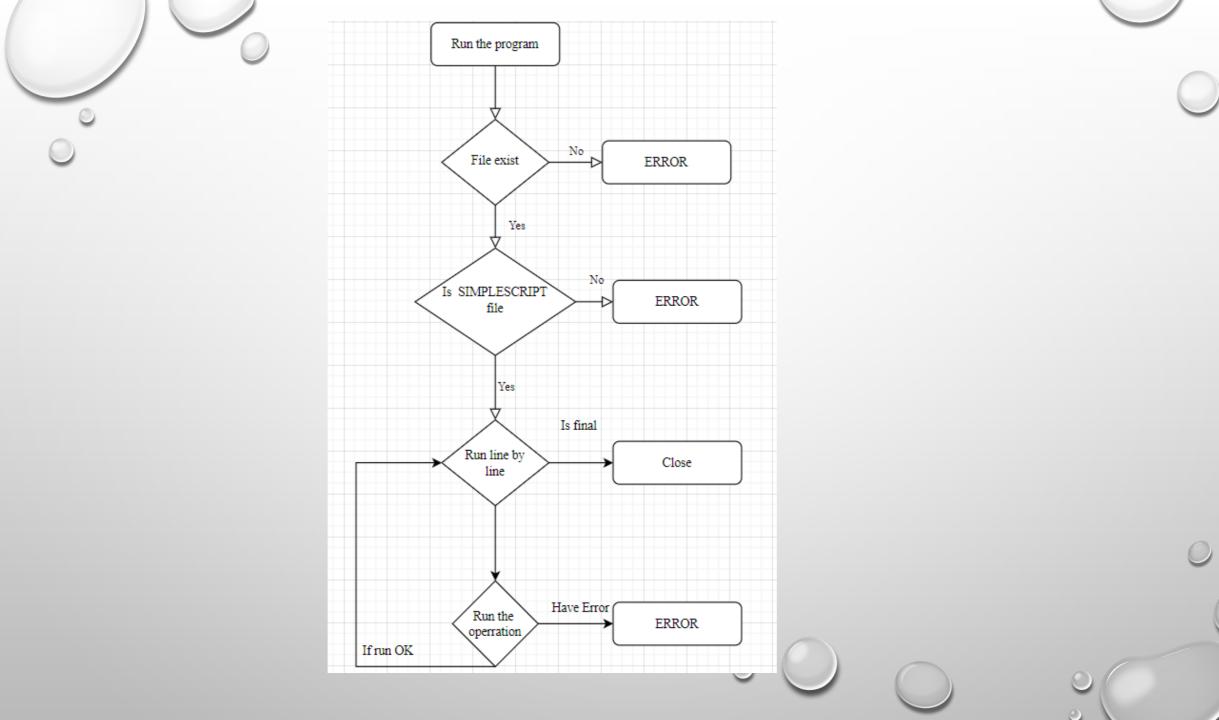
ENVIRONMENT

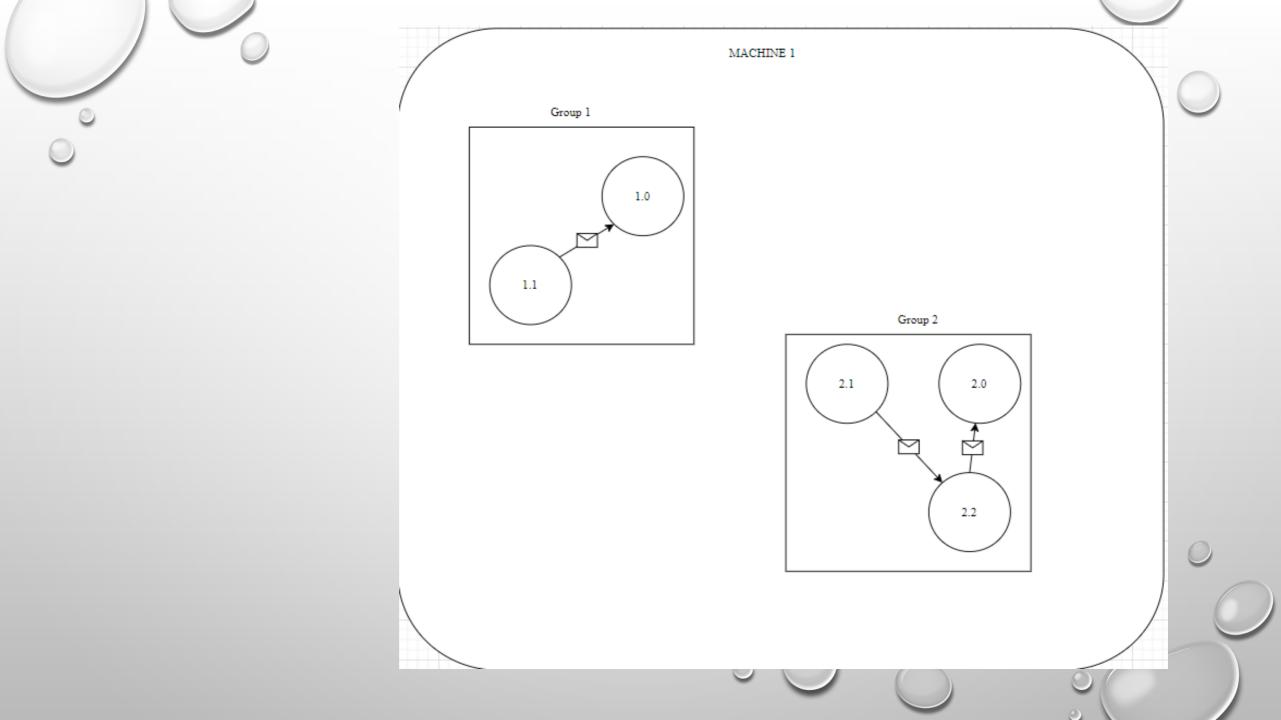
- RUN
 - RUNS PROGRAMS REQUESTED AND PLACES THEM IN GROUPS
- LIST
 - PRINTS ALL RUNNING PROCESSES AND THREADS AT THE MOMET IN THE WORKSPACE
- KILL
 - KILLS ALL PROCESS IN GROUP REQUESTED
- MIGRATE
 - TAKES A PROCESS FORM 1 MACHINE TO ANOTHER BY SENDING ITS DATA OVER NETWORK
- HELP
 - PRINTS INSTRUCTIONS ON HOW TO USE OUR WORKSPACE



PARSER

- OPENS FILES REQUESTED TO RUN
- CHECKS FROM LINE 0 TO END OF FILE FOR SYNTAX ERRORS
 - IF SYNTAX ERRORS ARE FOUND TERMINATES THE PROCEDURE.
 - IF LOOP IN LABELS IS DETECTED, EXITS THE RUNING PROGRAM WITH ERROR CODE.

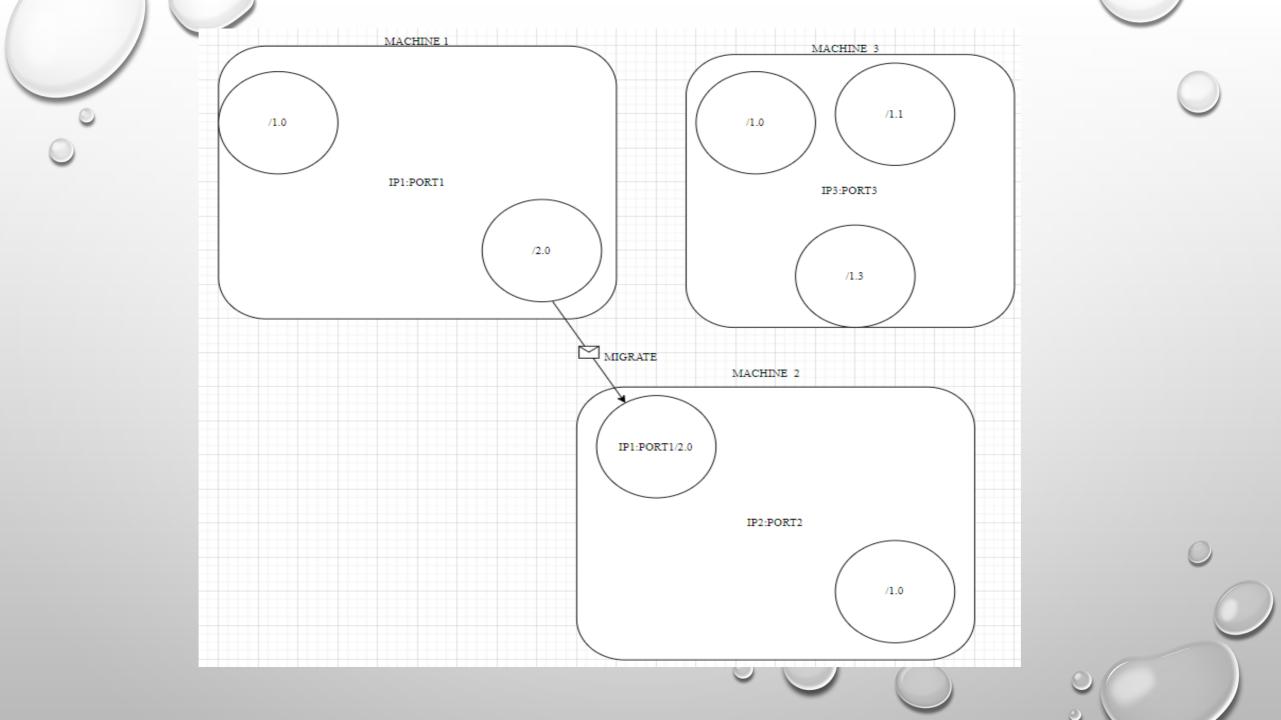






MIGRATION

- THE MIGRATION FROM ONE MACHINE TO OTHER IS DONE BY :
 - TCP/IP PROTOCOL
 - USER DEFINES THE IP AND PORT FOR THE NEW HOST MACHINE
- MESSAGES SENT BETWEEN MENBERS OF THE SAME GROUP IN WHICH SOME PROCESSES MIGRATED INTO OTHER MACHINE:
 - FIND THE MACHINE THAT HOSTS THE PROCESS BY USING MULTICAST
 - SENDS MESSAGE OR DATA OVER TCP/IP PROTOCOL



LOAD BALANCE (THEORETICALLY)

- 1 MACHINE IS DEFINED AS THE CO-ORDINATOR
 - SENDS A MULTICAST REQUEST TO EVERY ACTIVE MACHINE RUNING THE WORKSPACE ASKING FOR THEIR LOAD
 - WHEN DATA ARE COLLECTED PLACES THEM INTO A GLOBAL DISTRIBUTED BUFFER THAT EVERY MACHINE HAS ACCESS TO
 - BUFFER IS CONSISTED OF THE NAME_OF_MACHINE (IP.PORT) AND THE COUNTER OF LOAD
 - SUM OF LOADS IS CALCULATED
 - AND THEN THE AVERAGE NUMBER OF PROCESSES PER MACHINE IS CALCULATED
- THEN THE CO-ORDINATOR INFORMS EACH MACHINE OF THE MIGRATIONS NEEDED TO BE DONE SO IT IS BALANCED

PROBLEMS DETECTED IN DISTRIBUTED LOAD BALANCE

- WHO CAN BE THE CO-ORIDINATOR?
- HOW CAN WE ELECT CO-ORDINATOR BETWEEN MACHINES THAT CAN BE ALL AROUND THE GLOBAL NETWORK AND NOT IN A SPECIFIC GROUP?
- IF A RANDOM MACHINE DECIDES TO SEND MULTICAST REQUEST, BECAUSE THE SYSTEM IS NOT SYNCHRONOUS 2 OR MORE MULTICASTS CAN BE HAPPENING SIMULTANEOUSLY AND DICTATED MIGRATIONS MIGHT VARY
- WHAT IS THE T(MAX) WE SHOULD WAIT FOR REPLIES OF MULTICAST?
- HUGE OVERHEAD DUE TO CONTINIOUS UPDATES OF LOAD_BUFFER

PROBLEMS DETECTED IN DISTRIBUTED LOAD BALANCE

- WHAT IF WE DO NOT UPDATE THE LOAD_BUFFER AND INFO DO NOT CORRESPOND AT THE SPECIFIC TIME?
- ASYNCHRONOUS CLOCKS == WRONG LOAD ?



References:

https://www.researchgate.net/publication/342709018_LOad_balancing_research_paper

https://www.sciencedirect.com/science/article/pii/S131915782100046X

https://arxiv.org/abs/2007.07515