

Subdividing Alpha Complex

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Basic Definitions

Let B be a set of balls and D_B be its Delaunay complex.

- A Delaunay simplex δ is an alpha simplex if its corresponding balls intersect with the dual of δ .

An Example of alpha complex

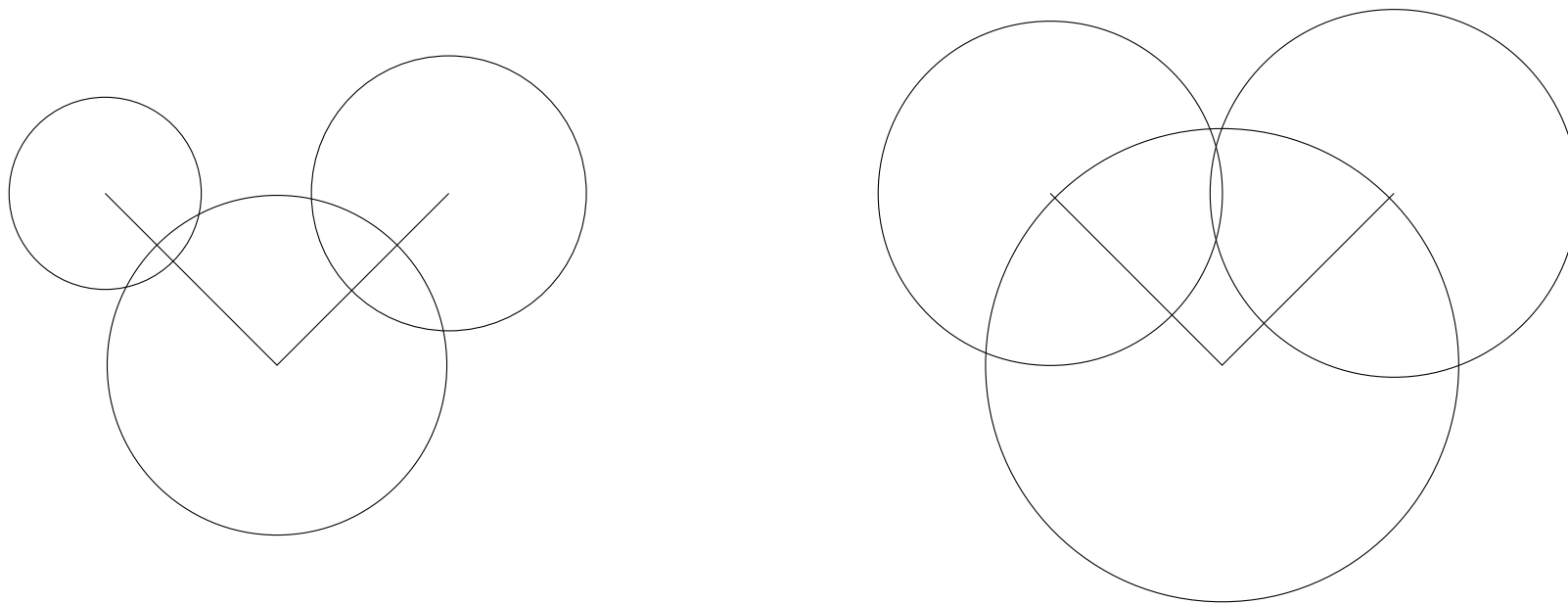


Figure 1:

Subdividing Alpha Complex

An alpha complex \mathcal{K}_B *subdivides* a simplicial complex \mathcal{C} if:

- $|\mathcal{K}_B| = |\mathcal{C}|$, and
- every simplex in \mathcal{K}_B is contained in a simplex in \mathcal{C}

Problem: Given a simplicial complex \mathcal{C} , construct a set of balls B such that \mathcal{K}_B subdivides $\text{set}\mathcal{C}$

Assumption: We are given the constraint triangulation of \mathcal{C} , i.e. the triangulation of the convex hull of \mathcal{C} which contains \mathcal{C} .

Main Theorem

Theorem. *Let B be a set of balls and \mathcal{C} be a simplicial complex. If B satisfies the following Conditions C1 and C2:*

C1. *for a subset $X \subseteq B$, if $\bigcap X \neq \emptyset$ then $z(X) \subseteq \sigma$ for some $\sigma \in \mathcal{C}$, and,*

C2. *for each $\sigma \in \mathcal{C}$, define $B(\sigma) = \{b \in B \mid b \cap \sigma \neq \emptyset\}$.*

Then we have: $z(B(\sigma)) \subseteq \sigma \subseteq \bigcup B(\sigma)$,

then \mathcal{K}_B subdivides \mathcal{C} .

Satisfying Condition 1

Subdivision

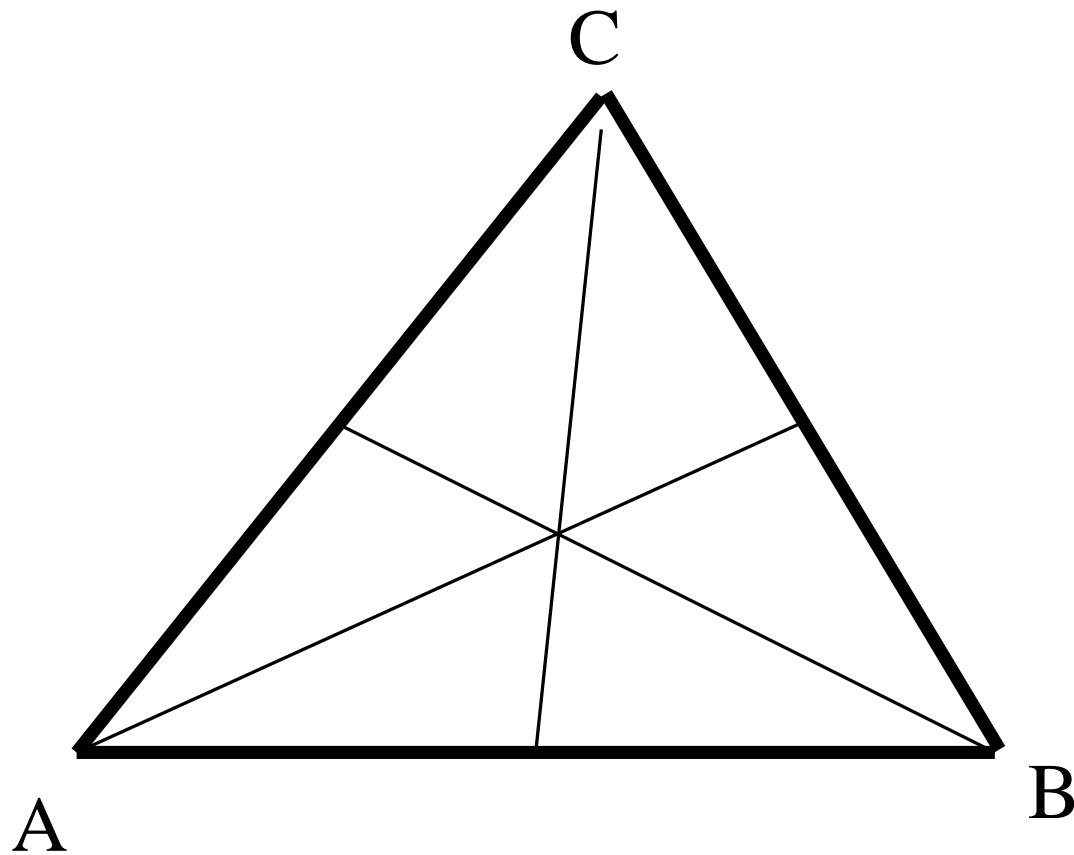


Figure 2:

Satisfying Condition 1

Protecting cells

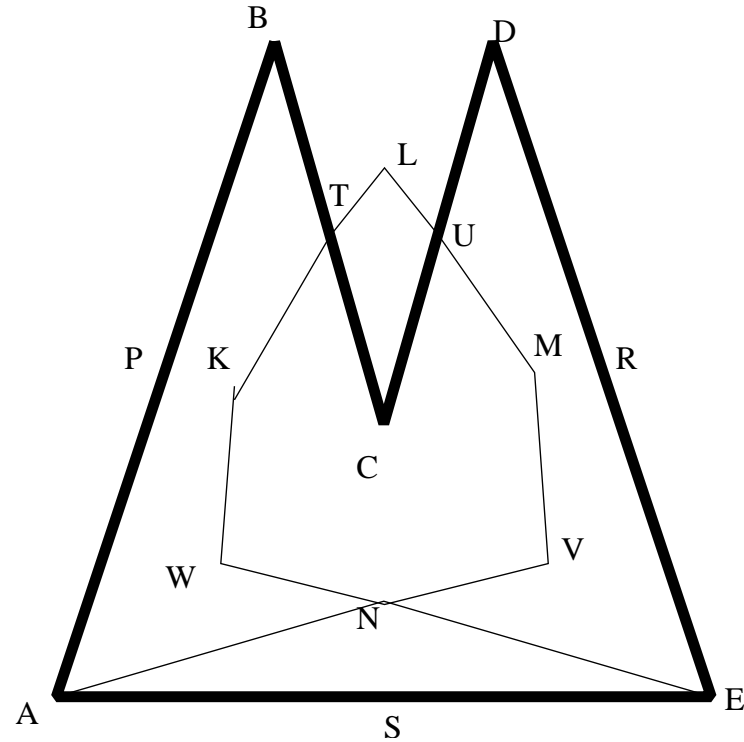
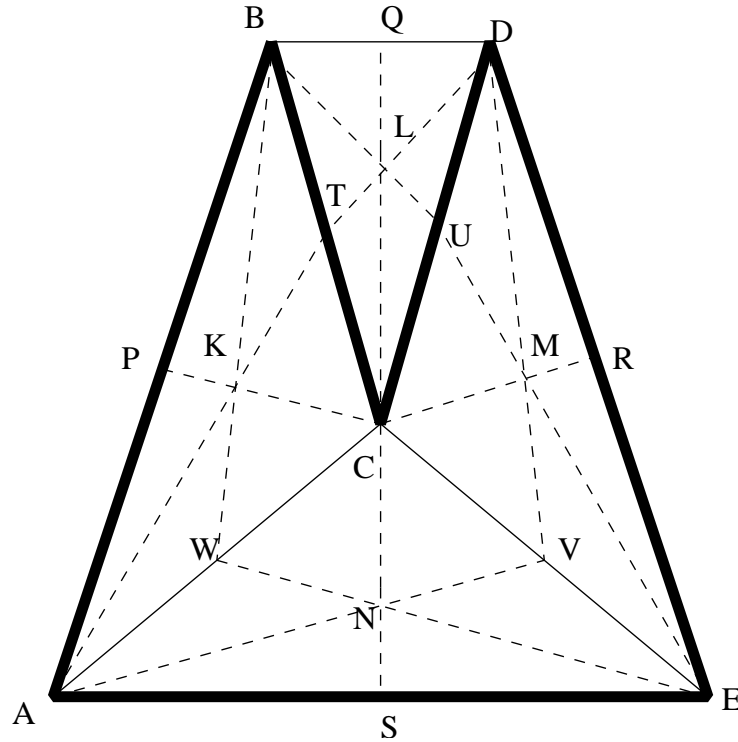


Figure 3:

Satisfying Condition 1

For a point $p \in \rho$, $\text{MaxWeight}(p)$ is the weight of the smallest ball centered on p that intersects the protecting cell of ρ .

The Algorithm

Restricted Voronoi complex:

Let $\sigma \in \mathcal{C}$ and X a set of balls.

- Restricted Voronoi complex $V_X(\sigma)$ is the restriction of V_X on ρ .
- A Voronoi vertex v is called a *positive* vertex, if v is outside every ball in X .

The Main Algorithm:

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1:  for  $i = 0, \dots, d$ 
2:    for all  $\rho \in \mathcal{C}$  of dimension  $i$ 
3:      ConstructBalls( $\rho$ )
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The Procedure ConstructBalls(σ):

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1:  if  $\dim \sigma = 0$ 
2:     $B(\sigma) := (\sigma, \gamma \cdot \text{MaxWeight}(\sigma))$ 
3:  else
4:    Let  $l := \dim \sigma$ 
5:    Let  $\tau_1, \dots, \tau_{l+1}$  be the  $(l-1)$ -faces of  $\sigma$ .
6:     $X := B(\tau_1) \cup \dots \cup B(\tau_{l+1})$ 
7:    while  $\exists$  a positive vertex  $u$  in  $V_X(\sigma)$ 
8:       $w := \gamma \cdot \text{MaxWeight}(u)$ 
9:       $X := X \cup \{(u, w)\}$ 
10:   endwhile
11:    $B(\sigma) := X$ 
12:endif
```

Proof of Correctness:

Condition C2 follows from the following proposition:

Proposition. *Let X be a set of balls. Suppose $z(X) \subseteq \sigma$. Then $\sigma \subseteq \bigcup X$ if and only if there is no positive vertex in $V_x(\sigma)$.*

Proof of Termination:

Fact. *Let Λ be a subset of σ whose boundary lies entirely in the interior of σ . Then there exists a constant $c > 0$ such that for all $p \in \Lambda$, $\text{MaxWeight}(p) > c$.*

Termination follows from the fact and the compactness of ρ .

Object Approximation