

## FORMULAZIONE COMPLETA

VARIABLES:  $X_{pq}^i \rightarrow$  node  $i$  used for VNF  $i$  on SFC  $p$   
 $B \rightarrow$  bandwidth  $D \rightarrow$  delay  
 $R^w \rightarrow$  node resource  $w$

MINIMIZE:

$$\begin{aligned} \text{cost} = & \sum_i \sum_p \sum_{\substack{q: q \neq \text{usr} \\ q \neq \text{mtr}}} \sum_w X_{pq}^i \cdot R_{pq}^w \cdot C_i^w + \quad \# \text{ nodes} \\ & + \sum_{\substack{i, j \\ i \neq j}} \sum_p \sum_{\substack{q \\ q \neq \text{usr}}} \sum_{\substack{q+1 \\ q \neq \text{mtr}}} X_{pq}^i \cdot X_{pq+1}^j \cdot B_p \cdot C_{ij} \quad \# \text{ links} \\ & + \sum_p \sum_q C_{pq} \quad (?) \quad \# \text{ embedding cost} \end{aligned}$$

## CONSTRAINTS

$$1) \quad \forall i \quad \sum_p \sum_{\substack{q: q \neq \text{usr} \\ q \neq \text{mtr}}} \sum_w X_{pq}^i \cdot R_{pq}^w \leq \max(R_i^w) \quad \# \text{ node res}$$

$$2) \quad \forall i, j \mid i \neq j \quad \sum_p \sum_{\substack{q \\ q \neq \text{usr}}} \sum_{\substack{q+1 \\ q \neq \text{mtr}}} X_{pq}^i \cdot X_{pq+1}^j \cdot B_p \leq \max(B_{ij}) \quad \# \text{ link BW}$$

$$3) \quad \forall p \quad \sum_{\substack{i, j \\ i \neq j}} \sum_{\substack{q \\ q \neq \text{usr}}} \sum_{\substack{q+1 \\ q \neq \text{mtr}}} X_{pq}^i \cdot X_{pq+1}^j \cdot D_{ij} \leq D_p \quad \# \text{ link delay}$$

$$4) \quad \forall p, q \quad \sum_i X_{pq}^i = 1 \quad \# \text{ one node for every function}$$

$$5) \quad \forall p \quad \sum_q \sum_{\substack{i, j \\ i \neq j}} \sum_{\substack{q+1 \\ q \neq \text{mtr}}} X_{pq}^i \cdot X_{pq+1}^j = \text{LEN}(p) - 1 \quad \# \text{ all links of chain are chosen}$$

$$6) \quad \forall_{p,q} \quad \sum_{\substack{i,j \\ i \neq j}} X'_{pq} X'_{pq+1} = 1$$

NOTES

1)  $B_{ij} = 0$  IF  $i, j$  NOT CONNECTED

2)  $D_{ij} \rightarrow \infty$  IF  $i, j$  NOT CONNECTED

QUADRATIC DISCONTINUOUS  
CANNOT BE EXPRESSED  
AS QUBO !!!

# FORMULATIONS INCOMPATIBLE (LINK BASED)

VARIABLES:  $y_{pq}^{ij}$   $\rightarrow$  LINK  $ij$  IS USED FOR GIVING FLOW  
 VALUE  $q$  TO  $q+1$

$B \rightarrow$  BANDWIDTH

$D \rightarrow$  DELAY

$R^w \rightarrow$  NODE RESOURCE  $w$

MINIMIZE:

$\rightarrow$  CONSTANT PER UNIT OF  
 AN EDGE FLOW

$$\begin{aligned} \text{COST} = & \sum_{i,j} \sum_{p,q} \sum_{\substack{q \neq q_{\text{max}} \\ q \neq q_{\text{min}}}} y_{pq}^{ij} \cdot R_{pq}^w \cdot C_{ij}^w + \quad \# \text{ WIRELESS} \\ & + \sum_{i,j} \sum_{p,q} \sum_{\substack{q \neq q_{\text{max}} \\ q \neq q_{\text{min}}}} y_{pq}^{ij} \cdot B_p \cdot C_{ij} \quad \# \text{ LINKS} \\ & + \sum_p \sum_q C_{pq} \quad (?) \quad \# \text{ EMBOSSING COST} \end{aligned}$$

CONSTRAINTS

$$1) \forall_{i,w} \sum_{j \neq i} \sum_{p,q} \sum_{\substack{q \neq q_{\text{max}} \\ q \neq q_{\text{min}}}} y_{pq}^{ij} R_{pq}^w \leq \text{MAX}(R_i^w) \quad \# \text{ NODE RES}$$

$$2) \forall_{i,j} \mid i \neq j \sum_{p,q} \sum_{\substack{q \neq q_{\text{max}} \\ q \neq q_{\text{min}}}} y_{pq}^{ij} B_p \leq \text{MAX}(B_{i,j}) \quad \# \text{ LINK BW}$$

$$3) \forall_p \sum_{i,j} \sum_{\substack{j \neq i \\ q \neq q_{\text{max}} \\ q \neq q_{\text{min}}}} y_{pq}^{ij} D_{ij} \leq D_p \quad \# \text{ LINK DELAY}$$

$$4) \forall_{p,q} \sum_{\substack{q \neq q_{\text{max}} \\ q \neq q_{\text{min}}}} \sum_{i,j \neq i} y_{pq}^{ij} = 1 \quad \# \text{ ONE NODE FOR GIVING FUNCTION}$$

$$5) \forall_p \sum_q \sum_{\substack{q \neq q_{\text{max}} \\ q \neq q_{\text{min}}}} \sum_{i,j \neq i} y_{pq}^{ij} = \text{LEN}(p) - 1 \quad \text{ALL LINKS OF CHAIN ARE PRESENT}$$

$$6) \quad \forall_{p,q} \quad \sum_{i \in S} \sum_{j \in T} y_{pq}^{ij} = 1 \quad \text{ONLY ONE LINK BETWEEN TWO VARS}$$

### NOTES

- EQUAZIONE 5) NON FORZA UNA DUE  
STESSA CATEGORIA AD ESSERE COMPLETI !!

↳ DIFFICILE AGGIUNGERE COME  
CONSTRAINT CHE SIA QUADRATICO