

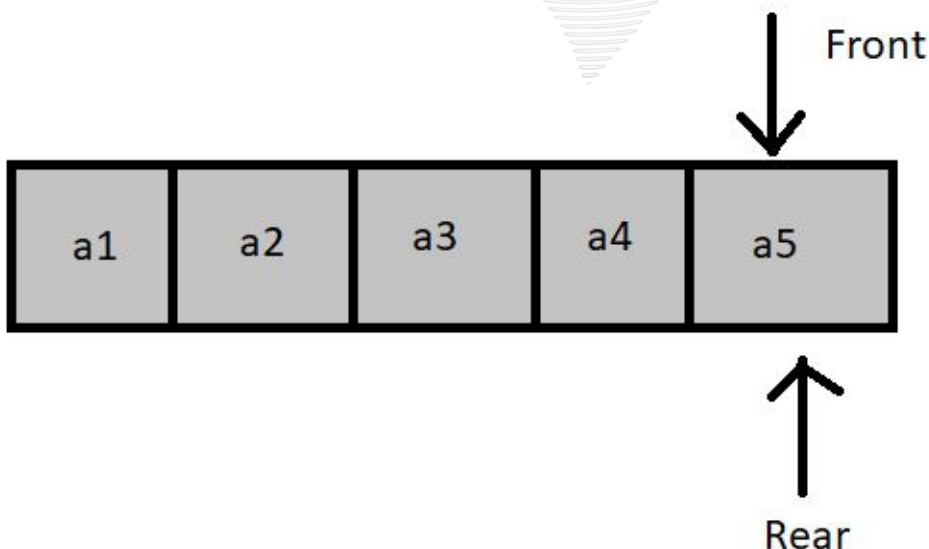
Disadvantage of Implementing Queue using Array

1)Waste of Space:

While deleting elements, we are moving the front pointer to the next index of the array. For example after deleting the first element of array i.e index 0, our front pointer will move to the next index i.e index 1. Now we will never go back to that index of array since there is no element present there and hence that index is useless for us. But still, since that index is part of our array it will occupy some space in our memory which we are never going to use. Similar thing will happen if we remove one more element and hence more space will get unused and hence more space will get wasted. Hence there is a waste of space while implementing a queue using arrays.

2)A condition will arise when a queue will be Empty as well as full

Now let us see this condition with the help of a diagram.



Now in the above case , the queue is empty as well as full and this is very ambiguous . The rear pointer is pointing to the last index of the array and hence our queue is full . On the other hand , as all the elements are deleted from the queue and hence now the front is pointing to the same index as that of the rear . Since , $\text{rear} == \text{front}$, the queue is empty as well . But this is impossible . A queue can't be full and empty at the same time . The problem is arising because we are only using each location only once . In ideal cases , as soon as an element is removed from the queue and that index becomes empty , we should use it again to store any new incoming element . But we are not doing any such thing here , we are just using each location once and hence this problem arises .

Solution of these Disadvantages

CIRCULAR QUEUE

Now ,the biggest problem we are facing while using an array for implementing a queue is that we are not reusing the location from where the element is deleted . As soon as the element is deleted from a particular index , we are abandoning that index and not using that index anymore . But in ideal cases , we should reuse the index to store any new incoming element. This gives rise to the concept of circular queue. In a circular queue we cover up the limitation of a linear queue of not using the same location twice by just adding one single condition . As soon as the queue becomes full (i.e rear pointer value becomes equal to $\text{size}-1$) , instead of terminating the program and saying the queue is full , we redirect our rear pointer to the beginning of the array (i.e index 0) by using modulus operator($\%$) . Now , if the element has been deleted from that index , we can insert our new element there by using the rear pointer as we normally do , hence using the same location twice , we cover up the limitation of the linear queue . Hence here we are avoiding wastage of space in the ambiguous case by just adding one condition.