

Trees : Level 3

Lesson 1

[minimum-element-in-bst](#)

[search-a-node-in-bst](#)

[array-to-bst](#)

[Binary Tree to BST](#)

[Find Common Nodes in two BSTs](#)

[print-bst-elements-in-given-range](#)

Binary Search Tree

Binary Search Tree is a binary tree data structure which has the following properties:

- The left subtree of a node contains only nodes with keys lesser than the node's key.
- The right subtree of a node contains only nodes with keys greater than the node's key.
- The left and right subtree each must also be a binary search tree.

Minimum element in BST

Given a **Binary Search Tree**. The task is to find the minimum element in this given BST.

Expected Time Complexity: O(Height of the BST)

Expected Auxiliary Space: O(Height of the BST).

Constraints:

$1 \leq N \leq 10^4$

```
1. int minValue(Node* root)
2. {
3.     if(root->left == NULL)
4.         return root->data;
5.     minValue(root->left);
6.
7.     // Code here
8. }
```

Search a node in BST

Given a **Binary Search Tree** and a node value X, find if node with value X is present in the BST or not.

Expected Time Complexity: $O(\text{Height of the BST})$

Expected Auxiliary Space: $O(1)$.

Constraints:

$1 \leq \text{Number of nodes} \leq 105$

```
1. bool search(Node* root, int x)
2. {
3.     if (root == NULL )
4.         return false;
5.     if( root->data == x)
6.         return true;
7.
8.     // Key is greater than root's key
9.     if (root->data < x)
10.        return search(root->right, x);
11.
12.    // Key is smaller than root's key
13.    return search(root->left, x);
14.}
```

Print BST elements in given range

Given a Binary Search Tree and a range. Find all the numbers in the BST that lie in the given range.

Note: Elements greater than or equal to root go to the right side.

Example 1:

Input:

```
    17
   /  \
  4    18
 /  \
2    9
```

l = 4, h = 24

Output: 4 9 17 18

Example 2:

Input:

```
    16
   /  \
  7    20
 /  \
1    10
```

l = 13, h = 23

Output: 16 20

Expected Time Complexity: $O(N)$.

Expected Auxiliary Space: $O(\text{Height of the BST})$.

Constraints:

$1 \leq \text{Number of nodes} \leq 104$

$1 \leq l < h < 105$

```
1. void nodes(Node*root,int low,int high,vector<int>&v)
2. {
3.     if(root==NULL)
4.         Return;
5.     //inorder traversal
6.     nodes(root->left,low,high,v);
7.     if(root->data>=low&&root->data<=high)
8.         v.push_back(root->data);
9.     nodes(root->right,low,high,v);
10.}
11.vector<int> printNearNodes(Node *root, int low, int high)
12.{
13.    vector<int>v;
14.    nodes(root,low,high,v);
15.    //sort(v.begin(),v.end());
16.    return v;
17. // your code goes here
18.}
```

Array to BST

Given a sorted array. Convert it into a Height balanced Binary Search Tree (BST).

Height balanced BST means a binary tree in which the depth of the two subtrees of every node never differ by more than 1.

Example 1:

Input: nums = {1, 2, 3, 4}

Output: {3, 2, 1, 4}

Explanation:

The preorder traversal of the following BST formed is {3, 2, 1, 4}:



Example 2:

Input: nums = {1,2,3,4,5,6,7}

Output: {4,2,1,3,6,5,7}

Explanation:

The preorder traversal of the following BST formed is {4,2,1,3,6,5,7} :

4

/\

2 6

/\ /\

1 3 5 7

Expected Time Complexity: $O(n)$

Expected Space Complexity: $O(n)$

Constraints:

$1 \leq |\text{nums}| \leq 10^4$

$-10^4 \leq \text{nums}[i] \leq 10^4$

```
1. void array(int left , int right ,vector<int>&nums,vector<int>&v)
2. {
3.     if(left<=right)
4.     {
5.         int mid=(left+right)/2;
6.         v.push_back(nums[mid]);
7.         array(left,mid-1,nums,v);
8.         array(mid+1,right,nums,v);
9.     }
10. }
11. vector<int> sortedArrayToBST(vector<int>& nums)
12. {
13.     int n=nums.size();
14.     vector<int>v;
15.     array(0,n-1,nums,v);
16.     return v;
17. }
```


Binary Tree to BST

Given a Binary Tree, convert it to Binary Search Tree in such a way that keeps the original structure of Binary Tree intact.

Example 1:

Input:

1

/ \

2 3

Output: 1 2 3

Example 2:

Input:

1

/ \

2 3

/

4

Output: 1 2 3 4

Explanation:

The converted BST will be

3

/ \

2 4

/

1

Expected Time Complexity: $O(N \log N)$.

Expected Auxiliary Space: $O(N)$.

Constraints:

1 <= Number of nodes <= 1000

```
1. void inorder(Node*root,vector<int>&v)
2. {
3.     if(root==NULL)
4.         return;
5.     inorder(root->left,v);
6.     v.push_back(root->data);
7.     inorder(root->right,v);
8. }
9. void push_el(Node*root, vector<int>&v,int &i)
10. {
11.     if(root==NULL)
12.         return;
13.     push_el(root->left,v,i);
14.     root->data=v[i];
15.     i++;
16.     push_el(root->right,v,i);
17. }
18. Node *binaryTreeToBST (Node *root)
19. { int i=0;
20.     vector<int>v;
21.     inorder(root,v);
22.     sort(v.begin(),v.end());
23.     push_el(root,v,i);
24.     return root; }
```

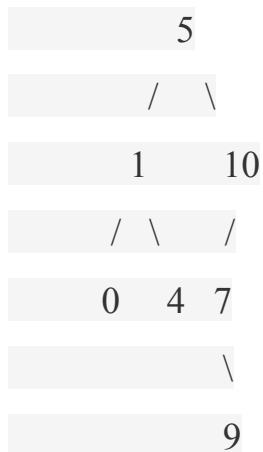
Find Common Nodes in two BSTs

Given two Binary Search Trees. Find the nodes that are common in both of them,
ie- find the intersection of the two BSTs.

Example 1:

Input:

BST1:



BST2:



Output: 4 7 9 10

Example 2:

Input:

BST1:

```
10
 / \
2   11
 / \
1   3
```

BST2:

```
2
 / \
1   3
```

Output: 1 2 3

Expected Time Complexity: $O(N1 + N2)$ where $N1$ and $N2$ are the sizes of the 2 BSTs.

Expected Auxiliary Space: $O(H1 + H2)$ where $H1$ and $H2$ are the heights of the 2 BSTs.

Constraints:

$1 \leq \text{Number of Nodes} \leq 10^5$