## Colouring Graphs



#### Started with 4 colour theorem:

Any map can be coloured with four colours so that adjacent countries have different colours.

- ▶ Posed by Guthrie 1852
- Proved by Appel and Haken 1976

#### **Definition**

The chromatic number  $\chi(G)$  of a graph G is the least number of colours needed to colour the vertices so that adjacent vertices have different colours.

### Theorem (The four colour theorem)

A planar graph G has  $\chi(G) \leq 4$ .

## First examples of chromatic number

### The empty graph $E_n$

The only graphs with  $\chi(G) = 1$  are the empty graphs  $E_n$ .

## The complete graph $K_n$

The complete graph has  $\chi(K_n) = n$ 

### Bipartite graphs

A graph G has  $\chi(G) = 2$  if and only if G is bipartite.

The wheel graph  $W_n$ 

$$\chi(W_n) = \begin{cases} 3 & n \text{ even} \\ 4 & n \text{ odd} \end{cases}$$

Why?

# How to find $\chi(G)$ : sandwich it!

Start by finding rough upper and lower bounds.

#### Upper bound: colour it

If you can colour the vertices of G with k colours so that adjacent vertices don't share a vertex, then  $\chi(G) \leq k$ .

#### Lower bound: often case by case

A few trivial tricks:

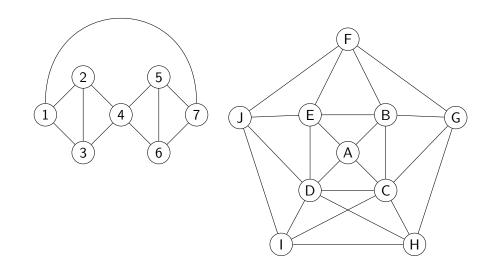
- ▶ If a vertex is adjacent to everything, as in  $W_n$
- ▶ If H is a subgraph of G, then  $\chi(G) \ge \chi(H)$ .

If lower bound isn't equal to upper bound, refine.

### Finding the $\chi(G)$ is NP-hard

So no beautiful answer. But for small graphs not that bad.

# Example: find $\chi(G)$ for the graphs shown below



## General upper bounds

#### **Definition**

 $\Delta(G)$  is the maximum degree of any vertex of G.

#### Lemma

$$\chi(G) \leq \Delta(G) + 1$$

#### Proof.

Colour the vertices one by one in any order.

#### The bound is tight, but for very few graphs:

- $\chi(K_n) = n = \Delta(K_n) + 1$
- For n odd,  $\chi(C_n) = 3 = \Delta(C_n) + 1$

### Theorem (Brooks)

If G isn't a complete graph or an odd cycle, then  $\chi(G) \leq \Delta(G)$ .

# A story problem for $\chi(G)$

Suppose you have some things you want to separate into groups, but certain things can't be in the same group. How many groups do you need?

- ▶ Group vacation, several cottages, some people don't get on
- Storing chemicals, some react dangerously with each other

## Make a graph G with:

- Vertices are the things
- Edges mean the vertices can't be in the same group

The groups are the colours.

 $\chi(G)$  is the lowest feasible number of groups