## The method we used to show $K_{3,3}$ isn't planar generalises:

## Take any cycle C in a graph G

- ▶ If the *G* is planar, *C* will be drawn as a "circle"
- ▶ Any other vertex or edge must lie inside or outside circle
- ► Handle possibilities case by case

#### Stereographic projection:

Don't have to treat inside/outside the circle as separate cases.

But for a complicated graph, could still be a lot of cases...

Best case: graph is Hamiltonian.

Don't have to put vertices inside/outside circle

## Another example: $K_5$ isn't planar

## Use our method with Hamiltonian cycle ABCDEA:

- ▶ WLOG, edge *AC* drawn inside...
- ▶ Then B cut off from DE inside, so BD, BE outside
- ▶ BD cuts off C from E on outside, so CE inside

### Only have AD left, but:

- ► *CE* cuts it off inside
- ▶ *BE* cuts it off outside

Therefore,  $K_5$  isn't planar

## The crossing graph packages the case by case analysis

"Edge  $e_1$  is in, so edge  $e_2$  out, so edge  $e_3$  in, so ..." gets tiresome.

For this slide: G a graph with Hamiltonian cycle C

- ightharpoonup Some pairs of edges in  $G \setminus C$  cross if drawn inside C
- ► Some pairs of edges can be drawn on the same side

### Definition

The crossing graph Cross(G, C) has:

Vertices: the edges in  $G \setminus C$ 

Edges e and f are adjacent if they cross inside C

Theorem (Planarity Algorithm for Hamiltonian graphs)

G is planar if and only if Cross(G, C) is bipartite

## Example of planarity algorithm:

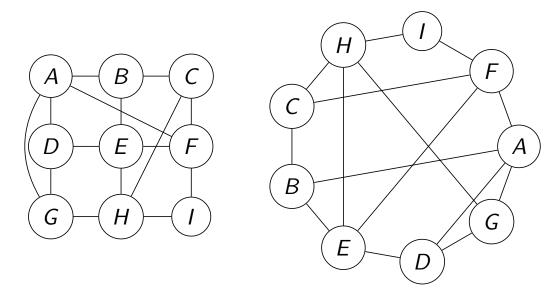


Figure: A graph  $\Gamma$ , then redrawn with Hamiltonian cycle outside

## What is $Cross(\Gamma, AFIHCBEDGA)$ ?

Vertices = edges in middle

Edges = crossings in middle

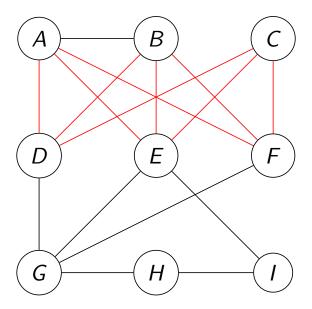
## What if G isn't Hamiltonian? Two lemma suffice.

#### Lemma

If G is a subgraph of H, and G is nonplanar, then H is nonplanar.

### Proof.

To draw H, we're drawing G and then adding some things.  $\square$ 



## Another tool for showing G isn't planar

#### Definition

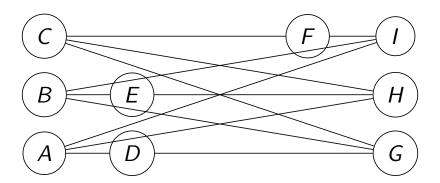
A graph H is a subdivision of G if it can be created from G by adding some vertices of degree two in the middle of edges.

#### Lemma

If H is a subdivision of G and, and G isn't planar, then H isn't planar.

#### Proof.

To draw H, we're drawing G and then adding some dots on the edges.



## Kuratowski's theorem - proves a general G not planar

#### Theorem

A graph G is not planar if and only if it has a subgraph that is a subdivision of  $K_{3,3}$  or  $K_5$ 

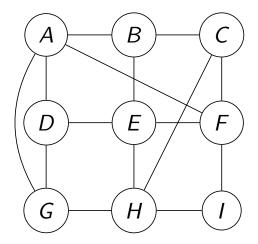
### Proof of the "if" direction:

- $ightharpoonup K_{3,3}$  and  $K_5$  aren't planar
- ▶ So subdivisions of  $K_{3,3}$  or  $K_5$  aren't planar
- ▶ So graphs having subdivisions a  $K_{3,3}$  or  $K_5$

#### Remarks on the "only if" direction

- ► Harder to prove and we won't even sketch
- ▶ Won't *explicitly* use if *G* is planar, prove it by drawing it!
- ▶ Will use implicitly if G isn't planar, we know we can find a  $K_{3,3}$  or  $K_5$  hidden in it

# Example of using Kuratowski's theorem



Give another proof that  $\Gamma$  isn't planar