

# Colouring Graphs



## Started with 4 colour theorem:

Any map can be coloured with four colours so that adjacent countries have different colours.

- ▶ Posed by Guthrie 1852
- ▶ Proved by Appel and Haken 1976

## Definition

The *chromatic number*  $\chi(G)$  of a graph  $G$  is the least number of colours needed to colour the vertices so that adjacent vertices have different colours.

## Theorem (The four colour theorem)

A planar graph  $G$  has  $\chi(G) \leq 4$ .

# First examples of chromatic number

The empty graph  $E_n$

The only graphs with  $\chi(G) = 1$  are the empty graphs  $E_n$ .

The complete graph  $K_n$

The complete graph has  $\chi(K_n) = n$

Bipartite graphs

A graph  $G$  has  $\chi(G) = 2$  if and only if  $G$  is bipartite.

The wheel graph  $W_n$

$$\chi(W_n) = \begin{cases} 3 & n \text{ even} \\ 4 & n \text{ odd} \end{cases}$$

Why?

# How to find $\chi(G)$ : sandwich it!

Start by finding rough upper and lower bounds.

## Upper bound: colour it

If you can colour the vertices of  $G$  with  $k$  colours so that adjacent vertices don't share a vertex, then  $\chi(G) \leq k$ .

## Lower bound: often case by case

A few trivial tricks:

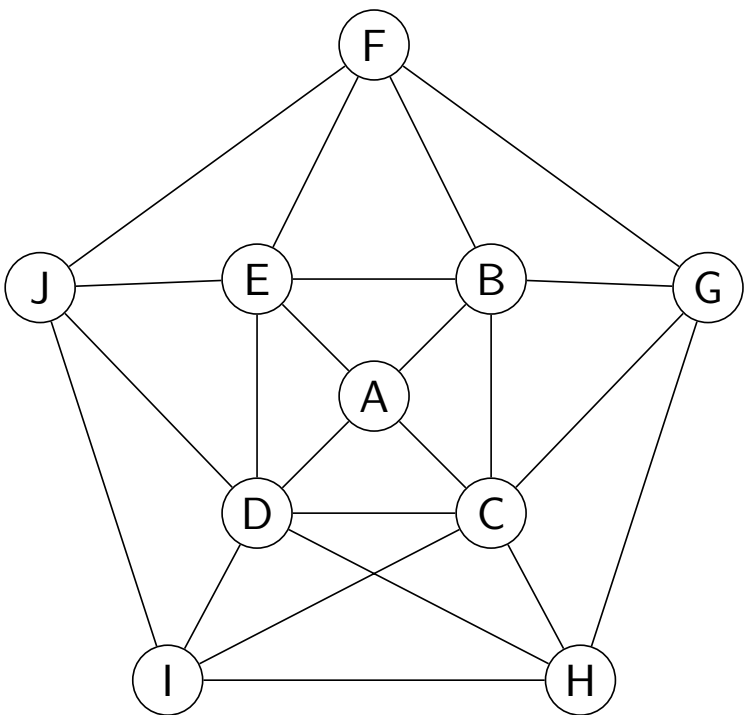
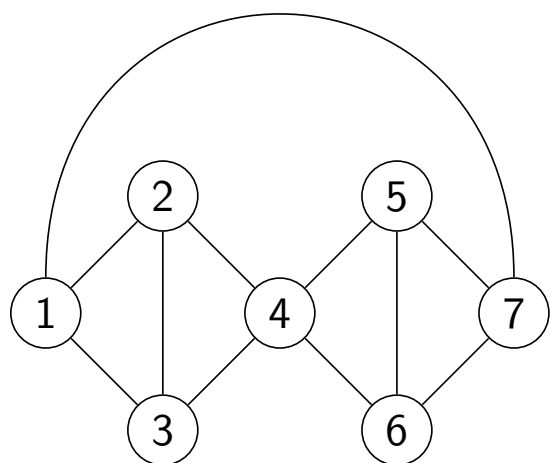
- ▶ If a vertex is adjacent to everything, as in  $W_n$
- ▶ If  $H$  is a subgraph of  $G$ , then  $\chi(G) \geq \chi(H)$ .

If lower bound isn't equal to upper bound, refine.

## Finding the $\chi(G)$ is NP-hard

So no beautiful answer. But for small graphs not that bad.

Example: find  $\chi(G)$  for the graphs shown below



# General upper bounds

## Definition

$\Delta(G)$  is the maximum degree of any vertex of  $G$ .

## Lemma

$$\chi(G) \leq \Delta(G) + 1$$

## Proof.

Colour the vertices one by one in any order. □

The bound is tight, but for very few graphs:

- ▶  $\chi(K_n) = n = \Delta(K_n) + 1$
- ▶ For  $n$  odd,  $\chi(C_n) = 3 = \Delta(C_n) + 1$

## Theorem (Brooks)

*If  $G$  isn't a complete graph or an odd cycle, then  $\chi(G) \leq \Delta(G)$ .*

## A story problem for $\chi(G)$

Suppose you have some things you want to separate into groups, but certain things can't be in the same group. How many groups do you need?

- ▶ Group vacation, several cottages, some people don't get on
- ▶ Storing chemicals, some react dangerously with each other

Make a graph  $G$  with:

- ▶ Vertices are the things
- ▶ Edges mean the vertices can't be in the same group

The groups are the colours.

$\chi(G)$  is the lowest feasible number of groups