

운영체제 보고서

과제 #07: 8장 연습문제

과목 명: 운영체제

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8.12

$$1KB = 1024 \text{ byte} = 2^{10}$$

단계 1, 주소를 십진수 \rightarrow 이진수

단계 2, page number 부분과 offset 부분으로 나누기

단계 3, offset과 page no 를 이진수 \rightarrow 십진수

(a) 3085

$$1. \begin{array}{r} 2 \overline{) 3085} \\ \underline{1542} \\ 1542 \\ \underline{971} \\ 971 \\ \underline{385} \\ 385 \\ \underline{192} \\ 192 \\ \underline{96} \\ 96 \\ \underline{48} \\ 48 \\ \underline{24} \\ 24 \\ \underline{12} \\ 12 \\ \underline{6} \\ 6 \\ \underline{3} \\ 3 \\ \underline{1} \\ 1 \end{array} \Rightarrow 110000001101$$

$$2. \begin{array}{c} 11 / 0000001101 \\ \hline \text{page no.} \quad \text{offset} \end{array}$$

$$3. 011 = 3, 0000001101 = 13$$

$$\begin{array}{l} \text{page no.} = 3 \\ \text{page offset} = 13 \end{array}$$

(b) 42095

$$1. 42095 \Rightarrow 1010010001101111$$

$$2. \begin{array}{c} 1010010001101111 \\ \hline \text{page no.} \quad \text{offset} \end{array}$$

$$3. 101001 = 41, 0001101111 = 111$$

$$\begin{array}{l} \text{page no.} = 41 \\ \text{page offset} = 111 \end{array}$$

(c) 215201

$$1. 215201 \Rightarrow 110100100010100001$$

$$2. \begin{array}{c} 110100100010100001 \\ \hline \text{page no.} \quad \text{offset} \end{array}$$

$$3. 11010010 = 210, 0010100001 = 161$$

$$\begin{array}{l} \text{page no.} = 210 \\ \text{page offset} = 161 \end{array}$$

(d) 650000

$$1. 650000 \Rightarrow 10011110101100010000$$

$$2. \begin{array}{c} 10011110101100010000 \\ \hline \text{page no.} \quad \text{page offset} \end{array}$$

$$3. 1001111010 = 634, 1100010000 = 784$$

$$\begin{array}{l} \text{page no.} = 634 \\ \text{page offset} = 784 \end{array}$$

(e) 2000001

$$1. 2000001 \Rightarrow 111101000010010000001$$

$$2. \begin{array}{c} 111101000010010000001 \\ \hline \text{page no.} \quad \text{offset} \end{array}$$

$$3. 11110100001 = 1953, 0010000001 = 129$$

$$\begin{array}{l} \text{page no.} = 1953 \\ \text{page offset} = 129 \end{array}$$

8.13

Virtual address : 21-bit

Physical address : 16-bit

Page size : 2-KB = 2^{11}

$$(a) \quad 2^m = \text{no. of pages} \times \text{page size}$$

$$2^{21} = \text{no. of pages} \times 2^{11}$$

$$\therefore \text{no. of pages} = \frac{2^{21}}{2^{11}} = 2^{10}$$

$$\therefore \text{no. of pages} = 2^{10} \text{ pages}$$

$$(b) \quad 2^m = \text{entries} \times \text{page size}$$

$$2^{16} = \text{entries} \times 2^{11}$$

$$\text{entries} = 2^{16} / 2^{11} = 2^5$$

$$\therefore 32 \text{ entries}$$

8.15

(a) logical address space $2^m = \text{no. of pages} \times \text{page size}$
 $= 256 \times 4 \text{ KB}$
 $= 2^8 \times 2^{12} = 2^{20}$

$\therefore m = 20$

(b) number of physical address $\frac{\text{physical address space}}{\text{frame size}}$

physical address space $= 2^x$

$= \text{no. of frames} \times \text{frame size}$

$= 64 \times 4 \text{ KB}$

$= 2^6 \times 2^{12} = 2^{18}$

$\therefore x = 18$

8.17

(a) 2 memory accesses

$\Rightarrow 50 \times 2 = 100$

$\therefore 100 \text{ nanoseconds}$

(b) effective access time $= 75\% \times \text{TLB hit-time} + 25\% \times \text{TLB miss-time}$

$= 0.75 \times (50 \text{ ns} + 2 \text{ ns}) + 0.25 (100 \text{ ns} + 2 \text{ ns})$

$= 64.5 \text{ ns}$

$\therefore 64.5 \text{ nanoseconds}$

8.20

Segment	Base	Length
0	219	600
1	2300	14
2	90	100
3	1327	580
4	1952	96

(a) 0, 430

segment 0은 600의 길이로 430보다 크다.

base 219와 offset 430이므로 $\text{physical address} = 219 + 430 = 649$

(b) 1, 10

segment 1은 14이므로 10보다 작다.

base 2300과 offset 10이므로 $\text{physical address} = 2300 + 10 = 2310$

(c) 2, 500

segment 2의 길이는 100으로 500보다 작다.

참조 불가능

(d) 3, 400

segment 3의 길이는 580으로 400보다 크다.

base 1327과 offset 400이므로 $\text{physical address} = 1327 + 400 = 1727$

(e) 4, 112

segment 4의 길이는 96으로 112보다 작다.

따라서 참조 불가능