

# Parser - Documentation:

Language used: Python

=== Grammar ===

non\_terminals: List<String> => Contains the set of nonterminal symbols  
terminals: List<String> => Contains the set of terminal symbols  
S: List<Pair<Pair<String, String>, String>> => Contains the set of states  
productions: List<Production> => Contains the set of production rules

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get\_grammar\_from\_file(String): Void => Reads the data from the input file and stores it

get\_productions\_for\_non\_terminal(String): returns the set of productions for a give nonterminal string

CFG\_check(): Boolean => Returns True if the grammar is Context-free, False otherwise

expand(): =>

advance(): =>

momentary\_insuccess(): =>

back(): =>

another\_try(): =>

success(): =>

=== Production ===

start: String => the "head" of the rule, contains the string of symbols with at least one nonterminal

rules: List<String> => the "body" of the rule, contains the strings of symbols that can result from the start

=== ParserConfiging ===

grammar: Grammar => Contains the grammar used for parsing

s: String => Contains the current state of the parser

i: Integer => Contains the current position of the symbol in input sequence

alpha: List => Represents the working stack ( contains the way the parse is built)

beta: List => Represents the input stack, part of the tree to be built

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expand(): => For a nonterminal head of input stack it expands it according to the rule

advance(): => For a terminal head of input stack that equals the currenty symbol from input, i is incremented and alpha stack advances with the top of beta stack

momentary\_insuccess(): => For a terminal head of input stack that is diferent from the current symbol from input, parser state is set to back state ('b')

back(): => For a terminal head of working stack, i is decremented and alpha stack advances

another\_try(): => For a nonterminal head of working stack, either the program goes into error state ('e') or normal state ('q')

success(): => Sets the final state ('f') corresponding to success  $w \in L(G)$

gl.txt:

S A B

a b epsilon

S  
 S → epsilon | a B | b A  
 A → a  
 A → a A | b B  
 B → b | b B  
 A B → a

g2.txt:

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program cmpdstmt decllist declaration type type1 arraydecl stmt stmtlist simplstmt
assignstmt expression iostmt structstmt ifstmt forstmt condition RELATION
IDENTIFIER "Boolean" "Integer" "String" "List" "<" type1 ">" "=" "+" "-" "*" "/" "("
")" "[" "]" "," "read" "print" CONSTANT : { } "if" "else" "for" < <= == != >= >
program
program -> decllist cmpdstmt
decllist -> declaration | declaration decllist
declaration -> type IDENTIFIER
type -> type1 | arraydecl
type1 -> "Boolean" | "Integer" | "String"
arraydecl -> "List" "<" type1 ">"
cmpdstmt -> "BEGIN" stmtlist "END"
stmtlist -> stmt | stmt ";" stmtlist
stmt -> simplstmt | structstmt
simplstmt -> assignstmt | iostmt
assignstmt -> IDENTIFIER "=" expression
expression -> expression "+" | "-" term | term
iostmt -> "read" "(" IDENTIFIER ")" | "write" "(" IDENTIFIER | CONSTANT ")"
structstmt -> cmpdstmt | ifstmt | forstmt
ifstmt -> "if" condition ":" stmt ["else" ":" stmt]
forstmt -> "for" assignstmt condition assignstmt ":" stmt
condition -> expression RELATION expression
RELATION -> < | <= | == | != | >= | >

```