

# HASHING FUNCTIONS

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# Hashing Function

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- Hashing function is a function which is applied on a key by which it produces an integer, which can be used as an address in hash table.
- A simple hashing function:  $h(k) = k \bmod m$

# Properties of Hashing Functions

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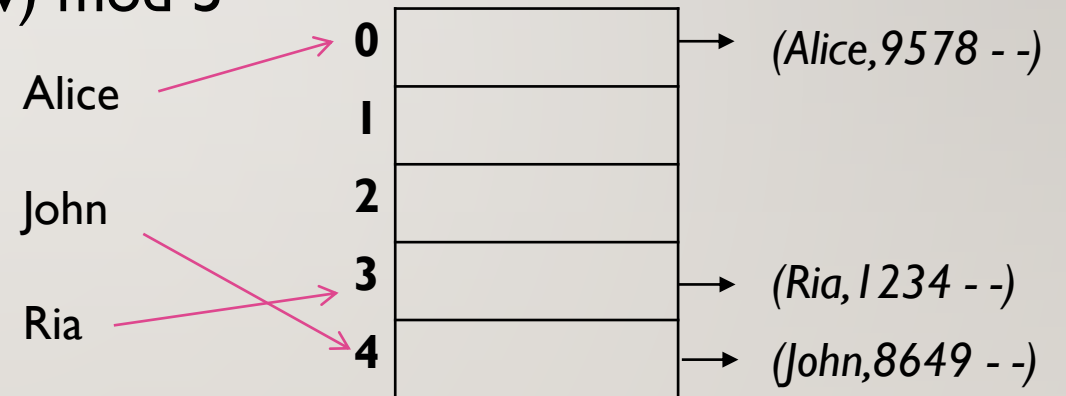
- Easy to compute
- Uniform distribution
- Less collisions

# Hash Table: Example

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- **Example:** phone book with table size  $N = 5$
- **hash function**  $h(w) = (\text{length of the word } w) \bmod 5$

- **Problem:** collisions
- Where to store Joe (collides with Ria)



# Collisions

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- Collisions occur when different elements are mapped to the same cell.
- Keys  $k_1, k_2$  with  $h(k_1) = h(k_2)$  are said to collide

What should we do now?

- Find a better hashing algorithm
- Use a bigger table
- Need a system to deal with collisions



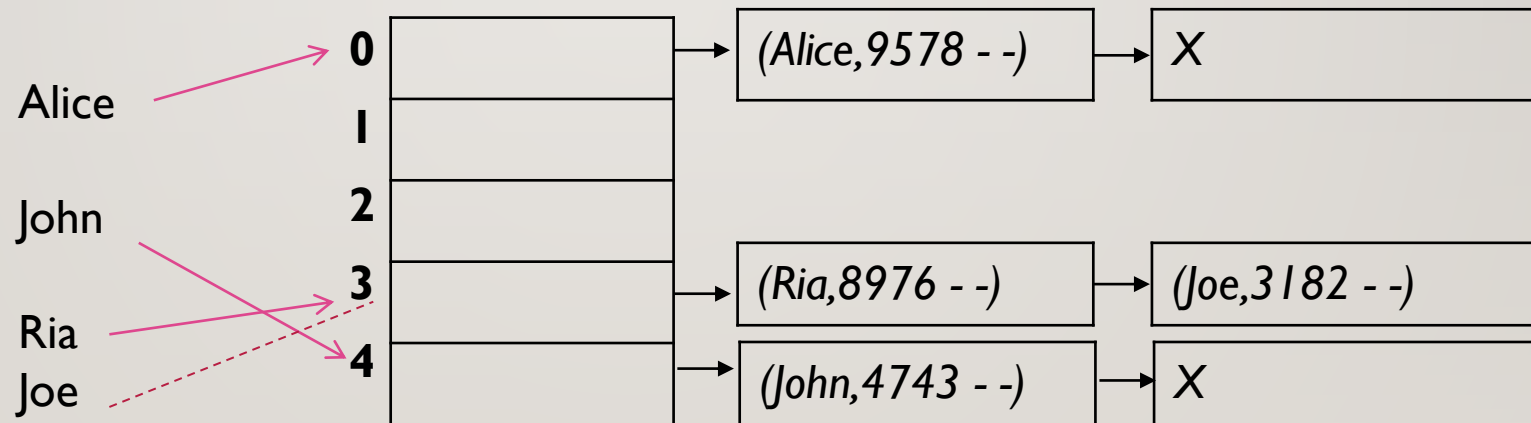
# Resolving Collisions

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- Two different methods for collision resolution:
  - **Separate Chaining:** Use a dictionary data structure (such as a linked list) to store multiple items that hash to the same slot.
  - **Closed Hashing (or *Open Addressing*):** search for empty slots using a second function and store item in first empty slot that is found.

# Separate Chaining

- Each cell of the hash table points to a linked list of elements that are mapped to this cell.
- Simple, but requires additional memory outside of the table



# Closed Hashing or Open Addressing

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- Open addressing does not introduce a new structure.
- If a collision occurs then we look for availability in the next spot generated by an algorithm.
- There are many implementations of open addressing, using different strategies for where to probe next:
  1. Linear Probing
  2. Quadratic Probing
  3. Double Hashing



## Contd..

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- Given an item  $X$ , try cells  $h_0(X), h_1(X), h_2(X), \dots, h_i(X)$ 
  - $h_i(X) = (\text{Hash}(X) + F(i)) \bmod \text{TableSize}$
  - $F(0) = 0$
- $F$  is the *collision resolution* function. Some possibilities:
  - **Linear**:  $F(i) = i$
  - **Quadratic**:  $F(i) = i^2$
  - **Double Hashing**:  $F(i) = i * \text{Hash}_2(X)$

# Linear Probing Example

insert(14)    insert(8)    insert(21)    insert(2)  
 $14\%7 = 0$      $8\%7 = 1$      $21\%7 = 0$      $2\%7 = 2$

---

0	14
1	
2	
3	
4	
5	
6	

0	14
1	8
2	
3	
4	
5	
6	

0	14
1	8
2	21
3	
4	
5	
6	

0	14
1	8
2	21
3	2
4	
5	
6	

# Quadratic Probing Example

insert(**14**)    insert(**8**)    insert(**21**)    insert(**2**)  
 $14\%7 = 0$      $8\%7 = 1$      $21\%7 = 0$      $2\%7 = 2$

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0	14
1	
2	
3	
4	
5	
6	

0	14
1	8
2	
3	
4	
5	
6	

0	14
1	8
2	
3	
4	21
5	
6	

0	14
1	8
2	2
3	
4	21
5	
6	

# Double Hashing

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- Double hashing can be done using :  
 **$(hash1(key) + i * hash2(key)) \% TABLE\_SIZE$**

- First hash function is typically

$$hash1(key) = key \% TABLE\_SIZE$$

- A popular second hash function is :

**$$hash2(key) = PRIME - (key \% PRIME)$$**

where PRIME is a prime smaller than the TABLE\_SIZE.





# Double Hashing Example

insert(19)  
 $19 \% 13 = 6$

0	
1	
2	
3	
4	
5	
6	19
7	
8	
9	
10	
11	
12	

insert(27)  
 $27 \% 13 = 1$

0	
1	27
2	
3	
4	
5	
6	19
7	
8	
9	
10	
11	
12	

insert(36)  
 $36 \% 13 = 10$

0	
1	27
2	
3	
4	
5	
6	19
7	
8	
9	
10	36
11	
12	

insert(10)  
 $10 \% 13 = 10$

0	
1	27
2	
3	
4	
5	10
6	19
7	
8	
9	
10	36
11	
12	

Collision 2

Let  $\text{Hash2}(\text{key}) = 7 - (\text{key} \% 7)$

$\text{Hash1}(10) = 10 \% 13 = 10$  (Collision 1)

$\text{Hash2}(10) = 7 - (10 \% 7) = 4$

$(\text{Hash1}(10) + 1 * \text{Hash2}(10)) \% 13 = 1$  (Collision 2)

$(\text{Hash1}(10) + 2 * \text{Hash2}(10)) \% 13 = 5$

Collision 1

# PRACTICE QUESTIONS ON GROWTH OF FUNCTIONS

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Q1. Give a big-O notation to estimate the sum of the first  $n$  positive integers.

Q2. Give a big-O estimate for the factorial function.

Q3. Give a big-O estimate for the following function:

$$f(n) = 3n \log(n!) + (n^2 + 3) \log n$$



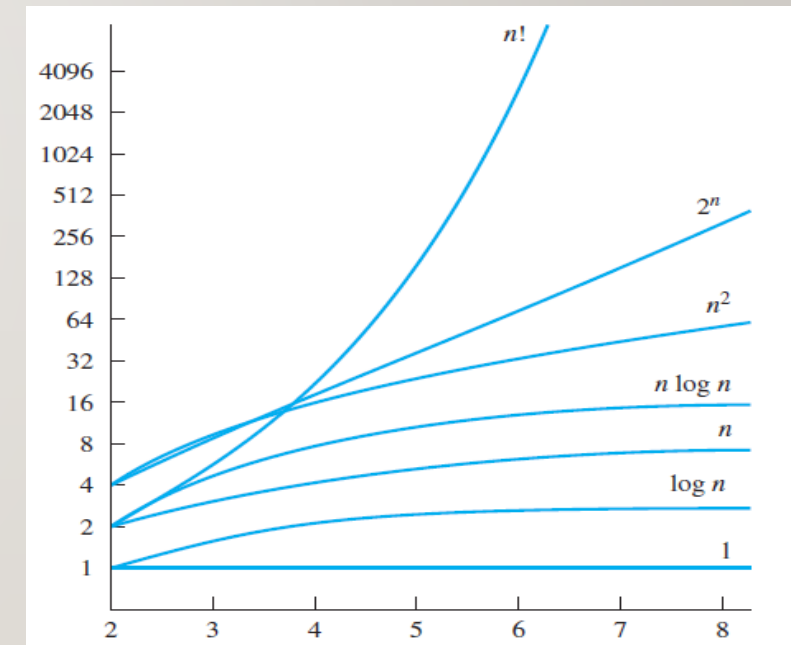
Q4. Give a big-O estimate for the following function:

$$f(x) = (x + 1) \log (x^2 + 1) + 3x^2$$

As mentioned before, *big-O notation is used to estimate the number of operations needed to solve a problem using a specified procedure or algorithm. The functions used in these estimates often include the following:*

*1,  $\log n$ ,  $n$ ,  $n \log n$ ,  $n^2$ ,  $2^n$ ,  $n!$*

Using calculus it can be shown that each function in the list is smaller than the succeeding function, in the sense that the ratio of a function and the succeeding function tends to zero as  $n$  grows without bound. *Figure displays the graphs of these functions, using a scale for the values of the functions that doubles for each successive marking on the graph. That is, the vertical scale in this graph is logarithmic.*



**A Display of the Growth of Functions Commonly Used in Big-O Estimates.**

# Thank You

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