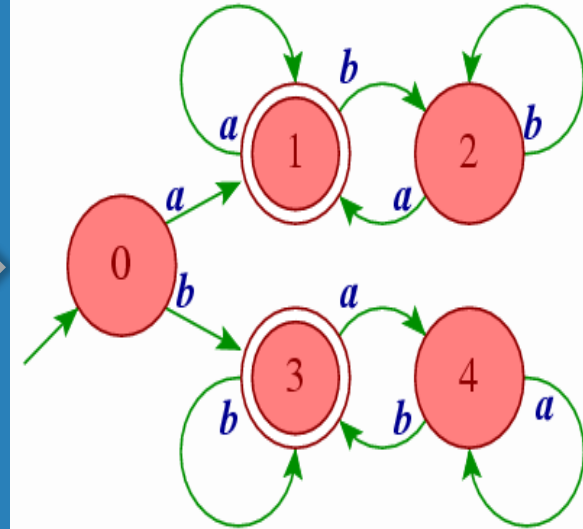


UCS701

Theory of Computation

Chapter 1: Regular Languages & Finite Automata



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Regular Expression

Regular expression

- A regular expression is a sequence of characters that **define a pattern**.
- **Notational shorthand's**
 1. One or more occurrences: $+$
 2. Zero or more occurrences: $*$
 3. Alphabets: Σ

Regular expression

L = Zero or More Occurrences of a = a^*



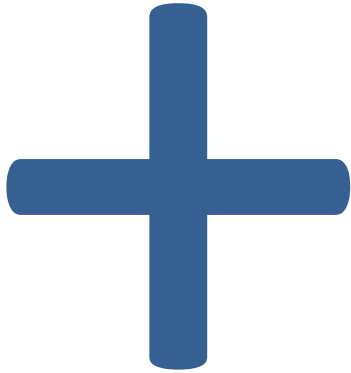
ϵ
a
aa
aaa
aaaa
aaaaa.....

Infinite

A diagram illustrating the strings generated by the regular expression a^* . A vertical list of strings is shown: ϵ , a, aa, aaa, aaaa, and aaaaa...... A blue bracket on the right side of the list groups these strings together, with an arrow pointing from the bracket to the text "Infinite" in red. An arrow also points from the top of the list, specifically to the a^* in the text above, indicating that these strings are the result of applying the Kleene star operation to the character 'a'.

Regular expression

L = One or More Occurrences of a = a^+



a
aa
aaa
aaaa
aaaaa.....

A blue bracket is positioned to the right of the list of 'a' strings, grouping them together. A curved blue arrow originates from the top of the bracket and points upwards towards the a^+ notation in the text above.

Infinite

Precedence and associativity of operators

Operator	Precedence	Associative
Kleene *	1	left
Concatenation	2	left
Union	3	left

Regular expression examples

1. 0 or 1

Strings: 0, 1

R.E. = 0 | 1

2. 0 or 11 or 111

Strings: 0, 11, 111

R.E. = 0 | 11 | 111

3. String having zero or more a .

Strings: ϵ , a , aa , aaa , $aaaa$

*R.E. = a^**

4. String having one or more a .

Strings: a , aa , aaa , $aaaa$

R.E. = a^+

5. Regular expression over $\Sigma = \{a, b, c\}$ that represent all string of length 3.

Strings: abc , bca , bbb , cab , aba

R.E. = $(a|b|c)(a|b|c)(a|b|c)$

6. All binary string.

Strings: 0, 11, 101, 10101, 1111 ...

R.E. = $(0 | 1)^+$

Regular expression examples

7. 0 or more occurrence of either a or b or both

Strings: $\epsilon, a, aa, abab, bab \dots$ *R.E.* = $(a \mid b)^*$

8. 1 or more occurrence of either a or b or both

Strings: $a, aa, abab, bab, bbbaaa \dots$ *R.E.* = $(a \mid b)^+$

9. Binary no. ends with 0

Strings: $0, 10, 100, 1010, 11110 \dots$ *R.E.* = $(0 \mid 1)^* 0$

10. Binary no. ends with 1

Strings: $1, 101, 1001, 10101, \dots$ *R.E.* = $(0 \mid 1)^* 1$

11. Binary no. starts and ends with 1

Strings: $11, 101, 1001, 10101, \dots$ *R.E.* = $1 (0 \mid 1)^* 1$

12. String starts and ends with same character

Strings: $00, 101, aba, baab \dots$ *R.E.* = $1 (0 \mid 1)^* 1$ or $0 (0 \mid 1)^* 0$
 $a (a \mid b)^* a$ or $b (a \mid b)^* b$

Regular expression examples

13. All string of a and b starting with a

Strings: a, ab, aab, abb...

R.E. = $a(a \mid b)^$*

14. String of 0 and 1 ends with 00

Strings: 00, 100, 000, 1000, 1100...

R.E. = $(0 \mid 1)^ 00$*

15. String ends with abb

Strings: abb, babb, ababb...

R.E. = $(a \mid b)^ abb$*

16. String starts with 1 and ends with 0

Strings: 10, 100, 110, 1000, 1100...

R.E. = $1(0 \mid 1)^ 0$*

17. All binary string with at least 3 characters and 3rd character should be zero
Strings: 000, 100, 1100, 1001...

R.E. = $(0 \mid 1)(0 \mid 1)0(0 \mid 1)^$*

18. Language which consist of exactly two b's over the set $\Sigma = \{a, b\}$

Strings: bb, bab, aabb, abba...

R.E. = $a^ b a^* b a^*$*

Regular expression examples

19. The language with $\Sigma = \{a, b\}$ such that 3rd character from right end of the string is always

Strings: aaa, aba, aaba, abb...

R.E. = $(a | b)^ a(a|b)(a|b)$*

20. Any no. of a followed by any no. of b followed by any no. of c

Strings: ϵ , abc, aabbcc, aabc, abb... *R.E. = $a^* b^* c^*$*

21. String should contain at least three 1

Strings: 111, 01101, 0101110...

R.E. = $(0|1)^ 1 (0|1)^* 1 (0|1)^* 1 (0|1)^*$*

22. String should contain exactly two 1

Strings: 11, 0101, 1100, 010010, 100100...

R.E. = $0^ 1 0^* 1 0^*$*

23. Length of string should be at least 1 and at most 3

Strings: 0, 1, 11, 01, 111, 010, 100... *R.E. = $(0|1) | (0|1)(0|1) | (0|1)(0|1)(0|1)$*

24. No. of zero should be multiple of 3

Strings: 000, 010101, 110100, 000000, 100010010... *R.E. = $(1^* 01^* 01^* 01^*)^*$*

Regular expression examples

25. The language with $\Sigma = \{a, b, c\}$ where a should be multiple of 3

Strings: aaa, baaa, bacaba, aaaaaa.. ***R.E. = $((b|c)^* a(b|c)^* a(b|c)^* a(b|c)^*)^*$***

26. Even no. of 0

Strings: 00, 0101, 0000, 100100.... ***R.E. = $(1^* 01^* 01^*)^*$***

27. String should have odd length

Strings: 0, 010, 110, 000, 10010.... ***R.E. = $(0|1) ((0|1)(0|1))^*$***

28. String should have even length

Strings: 00, 0101, 0000, 100100.... ***R.E. = $((0|1)(0|1))^*$***

29. String start with 0 and has odd length

Strings: 0, 010, 010, 000, 00010.... ***R.E. = $(0) ((0|1)(0|1))^*$***

30. String start with 1 and has even length

Strings: 10, 1100, 1000, 100100.... ***R.E. = $1(0|1)((0|1)(0|1))^*$***

Regular expression examples

31. All string begins or ends with 00 or 11

Strings: 00101, 10100, 110, 01011 ... *R.E.* = $(00|11)(0|1)^*|(0|1)^*(00|11)$

32. Language of all string containing both 11 and 00 as substring

Strings: 0011, 1100, 100110, 010011 ...

R.E. = $((0|1)^*00(0|1)^*11(0|1)^*)|((0|1)^*11(0|1)^*00(0|1)^*)$

33. String ending with 1 and not contain 00

Strings: 011, 1101, 1011 ...

R.E. = $(1|01)^+$

34. Language of C identifier

Strings: area, i, redious, grade1 ... *R.E.* = $(_ + L)(_ + L + D)^*$

where L is Letter & D is digit

Finite Automata

Finite Automata

- A **finite automaton**, or **finite state machine** is a 5-tuple $(Q, \Sigma, q_0, A, \delta)$ where
 - Q is finite set of states;
 - Σ is finite alphabet of *input symbols*;
 - $q_0 \in Q$ (*initial state*);
 - $A \subseteq Q$ (the set of *accepting* states);
 - δ is a function from $Q \times \Sigma$ to Q (the *transition* function).
- For any element q of Q and any symbol $a \in \Sigma$, we interpret $\delta(q, a)$ as the state to which the FA moves, if it is in state q and receives the input a .

Example: Finite Automata

- $M = (Q, \Sigma, q_0, A, \delta)$

- $Q = \{q_0, q_1, q_2\}$

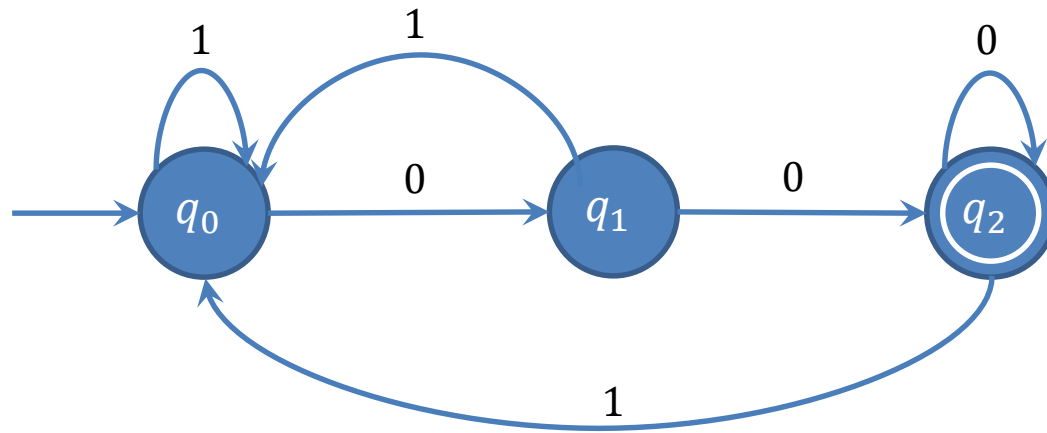
- $\Sigma = \{0,1\}$

- $q_0 = q_0$

- $A = \{q_2\}$

- δ is defined as

δ	Input	
State	0	1
q_0	q_1	q_0
q_1	q_2	q_0
q_2	q_2	q_0

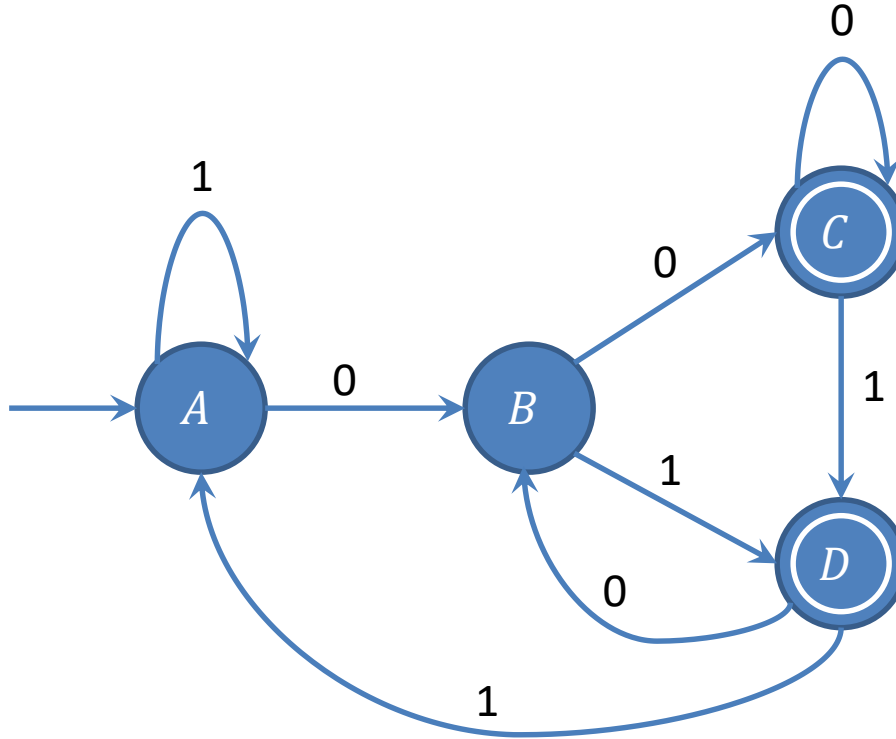


Applications of FA

- Lexical analysis phase of a compiler.
- Design of digital circuit.
- String matching.
- Communication Protocol for information exchange.

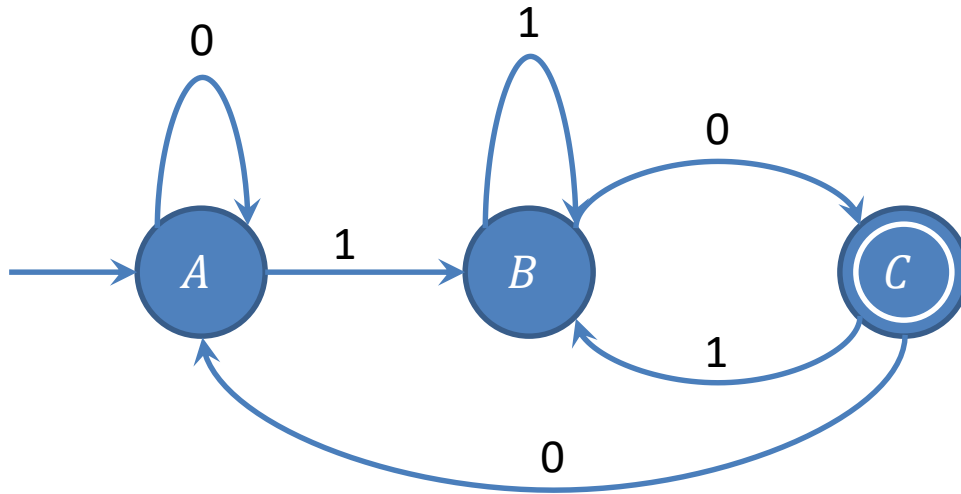
FA Examples

- The string with next to last symbol as 0.



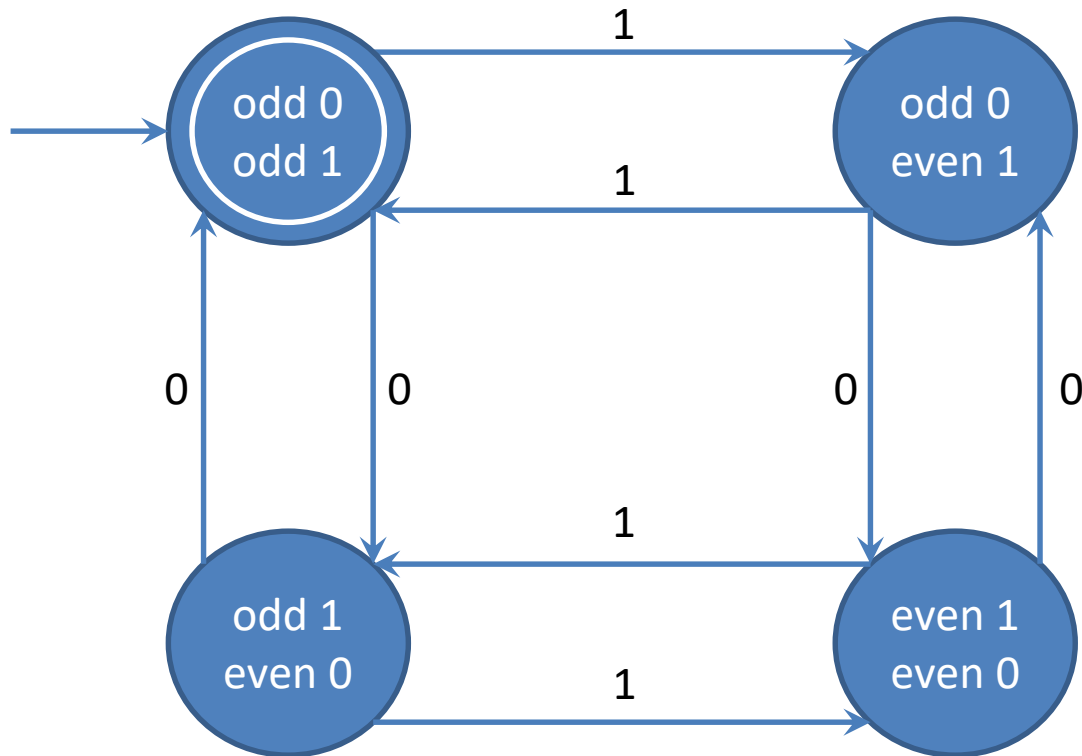
FA Examples

- The strings ending with 10.



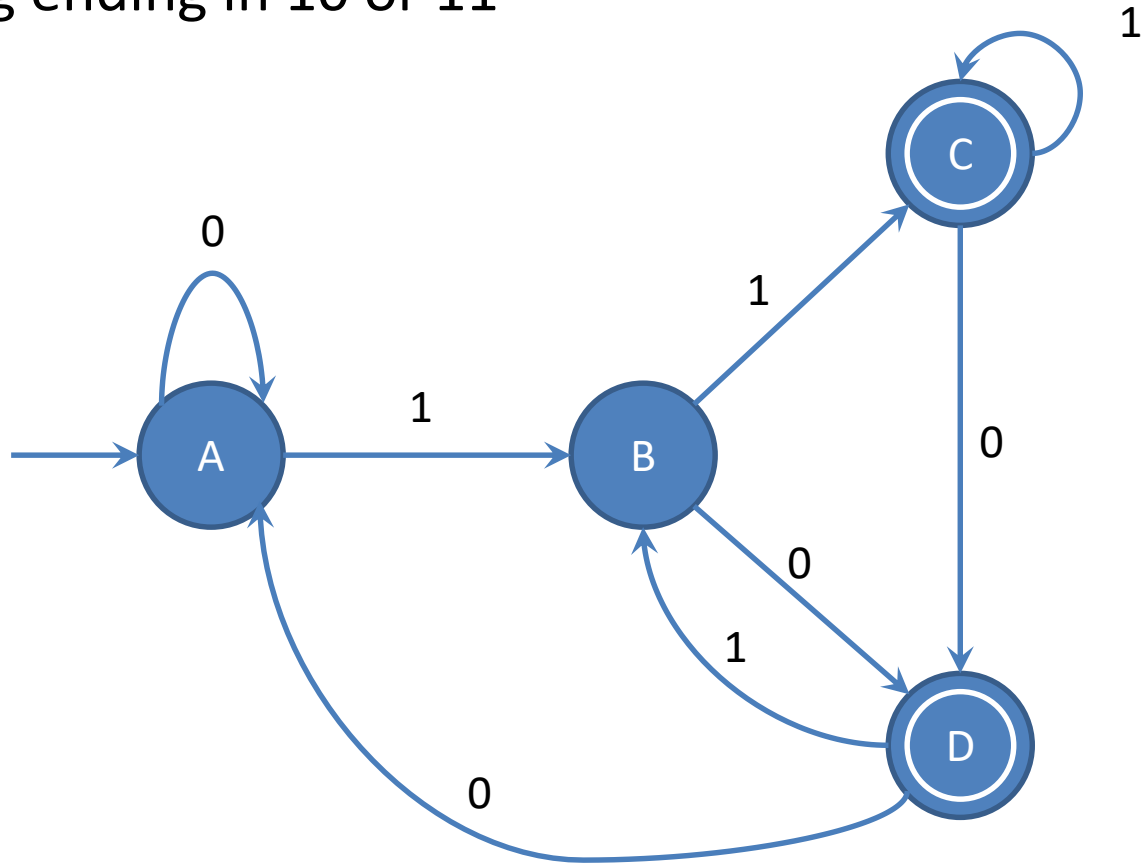
FA Examples

- The string with number of 0's and number of 1's are odd



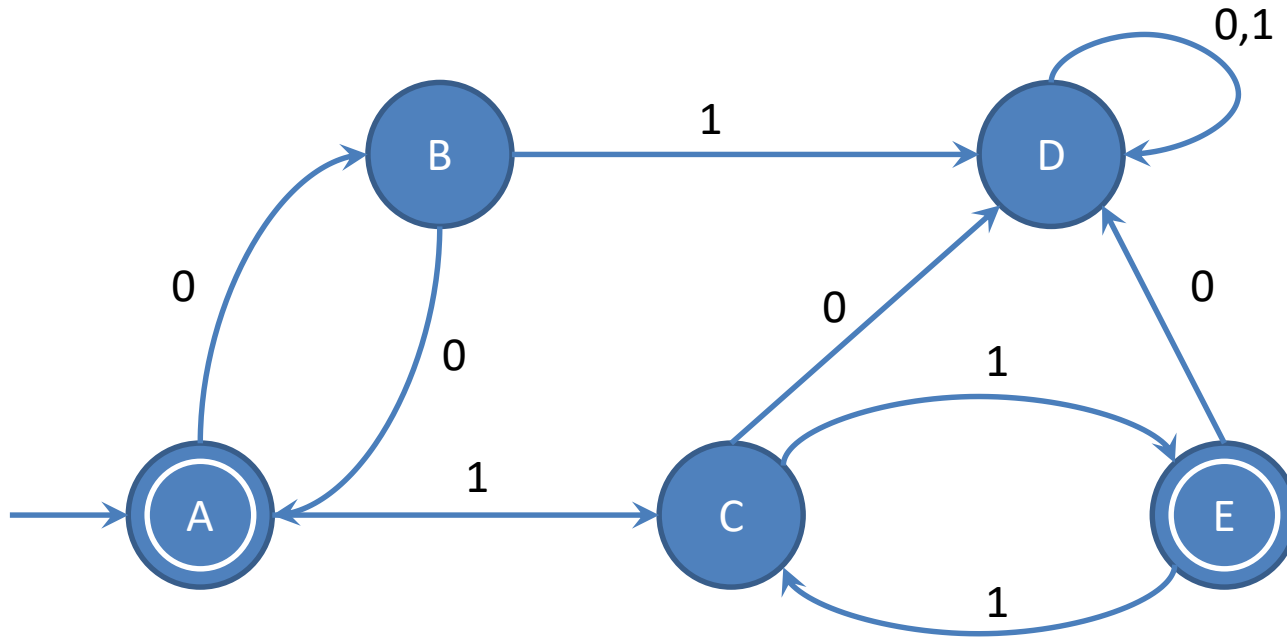
FA Examples

- The string ending in 10 or 11



FA Examples

- The string corresponding to Regular expression $\{00\}^*\{11\}^*$



FA Examples

- $(a+b)^*baaaa$

