



VIT

Vellore Institute of Technology
(Deemed to be University under section 3 of UGC Act, 1956)
CHENNAI

Continuous Assessment Test II - October 2024

Programme	B.Tech.(CSE)	Semester	Fall 2024-25
Course	Design and Analysis of Algorithms	Code	BCSE 204L
Faculty	Dr.L.Jeganathan, Dr M Janaki Meena, Dr M Raja, Dr R Sivakami, Dr B Indra , Dr G Kavipriya, Dr N.Jeiprathap	Slot/Class No.	A1/CH2024250101354 /CH2024250101360 /CH2024250102306 /CH2024250100952 /CH2024250100957 /CH2024250100961 /CH2024250100543
Time	90 Minutes	Max. Marks	50

Instructions:

- Answer all the FIVE questions.
- If any assumptions are required, assume the same and mention those assumptions in the answer script.
- Use of intelligence is highly appreciated.
- In this open-book exam, you are expected have a maximum-thinking and answer the question with just the necessary required information.
- You are requested to answer the five questions in the first five pages of your answer book, with one question each in a page.
- You are requested to do all the rough works from the page six. Contents from page six onwards, will not be evaluated.

1. You are given two sequences which $X = \langle A_1 A_2, \dots, A_m \rangle$, $Y = \langle B_1 B_2, \dots, B_n \rangle$ where A_i 's and B_j 's are matrices, $0 < i \leq m$, $0 < j \leq n$. The order of the matrix A_i , is represented as a 2-tuple (rA_i, cA_i) , where rA_i is the number of rows in A_i and cA_i represents the number of columns of A_i . Similarly, the order of the matrices B_i , $i = 0$ to n is represented as a 2-tuple (rB_i, cB_i) . The task is to compute the Longest-Common-Subsequence of X and Y , in the sense of the order of the matrices, denoted by $LCS(X, Y)$. For example, if $X = \langle A_1 A_2 A_3 A_4, A_5 \rangle$, with sizes $(5, 10), (7, 3), (3, 12), (12, 5), (5, 7)$ respectively, $Y = \langle B_1 B_2 B_3 B_4 \rangle$ with sizes $(7, 3), (3, 12), (5, 7), (4, 7)$, $LCS(X, Y)$ is $A_2 A_3 A_5$. Here $A_2 \cong B_1$, $A_3 \cong B_2$, $A_5 \cong B_3$. Here $A_i \cong B_j$, for any i, j , means that the matrix A_i and B_j are of the same order. We can also write $LCS(X, Y)$ as $B_1 B_2 B_3$. For the purpose of this problem, $LCS(X, Y)$ should be written in terms of A 's alone. You are required to use the dynamic programming strategy for this problem. The sequence $\langle X, Y \rangle$ is represented as $[m, n]$, where the sequence X has m matrices and Y has n matrices. Length of the $LCS(X, Y)$ is represented as $L[m, n]$,

- (a) Write the recurrence relation used to compute the $L[m, n]$ [2 marks]
- (b) What will be the entry in the $[3, 2]^{th}$ cell of the bottom-up memoization table used for the dynamic programming strategy in the computation of $L[m, n]$, if the order of matrices of X are : $(2, 3), (7, 9), (9, 6), (6, 8), (8, 18)$ and the order of the matrices of Y are: $(9, 6), (6, 8), (7, 8), (8, 18)$. [2 marks]
- (c) Mention the direction (left or top or diagonal) in the $(2, 3)^{th}$ cell for the input given in question 1(b). [2 marks]

Algorithm 1 LCS(X,Y)

(d) $m = X.length$
 $n = Y.length$
Let $b[0,...,m:0,...,n]$ and $c[0,...,m:0,...,n]$ be two new tables, initialised with zero and 'space' respectively
.....
.....
for $i = 0 \rightarrow m$ do
 for $j = 0 \rightarrow n$ do
 if _____ (1) then
 $c[i,j] =$ _____ (2)
 $b[i,j] = \text{'diagonal'}$
 else if _____ (3) then
 $c[i,j] =$ _____ (4)
 $b[i,j] = \text{'top'}$
 else
 $c[i,j] =$ _____ (5)
 $b[i,j] = \text{'left'}$
 end if
 end for
end for

There are four 'fill-in the blanks' in the Algorithm 1, which solves the given problem. Please write the response for 1.d.1 (1 mark), 1.d.2 (1 mark), 1.d.3 (1 mark), 1.d.4 (0.5 mark), 1.d.5 (0.5 mark) respectively, in separate lines. [4 marks]

2. Consider the 'Maximum Event Participation' (MEP problem) described as follows : You are organizing a series of community events : E_1, E_2, \dots, E_n , n is a positive integer. Each event runs for specified duration. Let the d_1, d_2, \dots, d_n be the duration (in minutes) of E_1, E_2, \dots, E_n respectively. A Participant can only attend events in the order they are listed and a participant can attend the events for a maximum duration D minutes. Given, $E_1, E_2, \dots, E_n, d_1, d_2, \dots, d_n$ and D , Task is to determine the maximum number of events that can be attended by a participant without exceeding D . Algorithm 2 is designed to solve the MEP problem.

Algorithm 2 MEP(E,d,D)

a) 1: $E = [E_1, E_2, \dots, E_n]$
2: $d = [d_1, d_2, \dots, d_n]$
3: $SE = \{\}$ // {Set of Selected events}
4: $CT = 0$ // {Total duration of the selected events}
5: for _____ (1) do
6: _____ (2)
7: _____ (3)
8: _____ (4)
9: end for
10: return SE

There are four 'fill-in the blanks' in the Algorithm 2. Please write the response for 2.a.1, 2.a.2, 2.a.3, 2.a.4 respectively, in separate lines.

- (b) What will be returned by the MEP pseudocode if $d = [5, 7, 4, 3, 6, 1]$ and $D = 17$. [4 marks]
(c) Given $D = 15$, compute n and $d = (d_1, d_2, \dots, d_n)$ such that the following conditions are satisfied. [2 Marks]

- n is maximum (in the sense that d cannot be of size greater than n that satisfies the other required conditions)
- d is in decreasing order of d_i 's and all d_i 's are distinct integers.

(d) Describe the time-complexity of the above MEP pseudocode.

[2 marks]

3. Given a text T and a pattern P , Compute the length of the Longest Common Subsequence of T and P such that $LCS(T, P)$ occurs in T from the leftmost character to the rightmost and the $LCS(T, P)$ Occurs in P from the rightmost to the leftmost. You are provided with two algorithms : Algorithm 3, Algorithm 4 to compute $LCS(T, P)$.

Algorithm 3 $LCS(T, P)$

```

1:  $m = \text{length}(T)$ 
2:  $n = \text{length}(P)$ 
3:  $PR = \underline{\hspace{2cm}}$  (1)
4: Let  $L[0, \dots, m; 0, \dots, n]$  initialized to 0
5: for  $i = 1$  to  $m$  do
6:   for  $j = 1$  to  $n$  do
7:     if  $\underline{\hspace{2cm}}$  (2) then
8:        $\underline{\hspace{2cm}}$  (3)
9:     else
10:       $\underline{\hspace{2cm}}$  (4)
11:    end if
12:  end for
13: end for
14: return  $L[m][n]$ 
```

There are four 'fill-in the blanks' in the Algorithm 3. Please write the response for 3.a.1, 3.a.2, 3.a.3, 3.a.4 respectively, in separate lines.

[4 marks]

Algorithm 4 BruteForce-LCS

```

1: Input: Text  $T$ , Pattern  $P$ 
2: Output: Length of Longest Common Subsequence
3:  $m \leftarrow \text{length}(T)$ 
4:  $n \leftarrow \text{length}(P)$ 
5:  $RP \leftarrow \underline{\hspace{2cm}}$  (1)
6:  $\text{max\_LCS\_length} \leftarrow 0$ 
7:  $AST \leftarrow \text{GenerateSubsequences}(T)$ 
8:  $ASP \leftarrow \underline{\hspace{2cm}}$  (2)
9: for  $\underline{\hspace{2cm}}$  (3) do
10:  for  $\underline{\hspace{2cm}}$  (4) do
11:    if  $A = B$  then
12:       $\underline{\hspace{2cm}}$  (5)
13:    end if
14:  end for
15: end for
16: return  $\text{max\_LCS\_length}$ 
```

Assume that the Algorithm 4, is supported by a function $\text{GenerateSubsequences}()$. If you want any other functions, you can assume the same. There are six 'fill-in the blanks' in the Algorithm 4. Please write the response for 3.b.1(1 mark), 3.b.2(1 mark) 3.b.3 (1 mark), 3.b.4 (1 mark), 3.b.5(2 marks) respectively in separate lines.

[6 marks]

4. Given $S = \{a_1, a_2, \dots, a_n\}$ where a_i 's are activities with the respective start-time and finish-time as (s_i, f_i) , and a pattern $P = \{(s_\alpha, f_\alpha)\}$, Task is to compute the maximum subset S' of S which has mutually compatible activities such that S' has at the least one activity which has the time interval described in $P = (s_\alpha, f_\alpha)$ within itself. For example, the activity $(2, 7)$ has the time interval $(3, 5)$, within itself. If $S = \{(2, 7), (3, 5), (6, 8), (8, 12), (7, 17)\}$, $P = \{(3, 5)\}$, then $S' = \{(3, 5), (6, 8), (8, 12)\}$. One

approach to solve this problem is as follows: From S and P , we first compute a set F which will have the activities that has the time interval of P , with in itself. If F is non-empty, that means that S has activities which have the time interval of P with in itself. In the above example, $F = \{(3, 5), (2, 7)\}$. In that case, we can find the maximum compatible subset of S , which will ensure that S' has atleast one activity which will have the time interval of P with in itself. Based on this method, Algorithm 5 is designed to compute S' .

Algorithm 5 Maximum Compatible Subset with Overlapping Activity

a) 1: **Input:** Set of activities $S = \{(s_i, f_i)\}$, Pattern $P = (s_\alpha, f_\alpha)$
 2: **Output:** Maximum subset S' of mutually compatible activities that overlap with P
 3: $F \leftarrow \{\}$
 4: $counter \leftarrow 0$
 5: **for** each activity (s_i, f_i) in S **do**
 6: **if** _____ (1) **and** _____ **then**
 7: $F \leftarrow$ _____ (2)
 8: $counter \leftarrow counter + 1$
 9: **end if**
 10: **end for**
 11: **if** $counter > 0$ **then**
 12: F _____ (3)
 13: **else**
 14: **return** F
 15: **end if**
 16: **if** F is not empty **then**
 17: Sort the activities in F by finish time
 18: $S' \leftarrow \{\}$
 19: $A \leftarrow 0$
 20: **for** _____ (4) **do**
 21: **if** $s_i \geq A$ **then**
 22: $S' \leftarrow$ _____ (5)
 23: _____ (6)
 24: **end if**
 25: **end for**
 26: **end if**
 27: **return** S'

There are six 'fill-in the blanks' in the Algorithm 5 . Please write the response for 4.a.1, 4.a.2, 4.a.3, 4.a.4, 4.a.5, 4.5.6 respectively in separate lines. [6 marks]

(b) Compute the time-complexity of Algorithm 5. [2 Marks]

(c) For the input $S = \{(2, 7), (3, 5), (8, 12), (1, 4), (7, 17)\}$ and the pattern $P = \{(3, 5)\}$, compute the output returned by the Algorithm 5. [2 marks]

5. Understand the Algorithm 6 for the N-queen's problem and answer the following.

(a) In the Algorithm 6, first queen is placed in which row? [2 marks]

(b) If the board is of size $n \times m$, $n > m$, how many queens can be placed in the board satisfying all the constraints. [2 marks]

(c) How many diagonals are checked before placing a queen? [2 marks]

(d) If $n=4$, how many recursion calls are made? [2 marks]

(e) Compute the time-complexity of the recursive function alone in Algorithm 6. [2 marks]

Algorithm 6 N-Queens Problem: Finding All Solutions

```
1: Input: Integer  $n$  (size of the board, number of queens)
2: Output: All possible solutions where queens are non-attacking
3:  $board \leftarrow$  empty board of size  $n \times n$ 
4:  $solutions \leftarrow [ ]$ 
5: PlaceQueen( $board, n, n - 1$ )
6: return  $solutions$ 
7: Procedure PlaceQueen( $board, n, row$ )
8: if  $row < 0$  then
9:   Add  $board$  to  $solutions$ 
10:  return
11: end if
12: for  $col = 0$  to  $n - 1$  do
13:   if IsSafe( $board, row, col$ ) then
14:     Place queen at  $(row, col)$ 
15:     PlaceQueen( $board, n, row - 1$ )
16:     Remove queen from  $(row, col)$ 
17:   end if
18: end for
19: return
20: Function IsSafe( $board, row, col$ )
21: for  $i = row + 1$  to  $n$  do
22:    $j = col - 1$ 
23:   if  $j < 0$  then
24:     Continue
25:   end if
26:   if  $board[i][j] = \text{true}$  then
27:     return false
28:   end if
29:    $j = j - 1$ 
30: end for
31: for  $i = row - 1$  to  $1$  step  $-1$  do
32:    $j = col - 1$ 
33:   if  $j < 0$  then
34:     Continue
35:   end if
36:   if  $board[i][j] = \text{true}$  then
37:     return false
38:   end if
39:    $j = j - 1$ 
40: end for
41: return true
```